

NANYANG TECHNOLOGICAL UNIVERSITY

MIDTERM I (CA1)

MH1812 – Discrete Mathematics

February 2020

TIME ALLOWED: 50 minutes

Name:

Matric. no.:

Tutor group:

INSTRUCTIONS TO CANDIDATES

1. **DO NOT TURN OVER PAPER UNTIL INSTRUCTED.**
2. This midterm paper contains **THREE (3)** questions.
3. Answer **ALL** questions. The marks for each question are indicated at the beginning of each question.
4. Candidates can write anywhere on this midterm paper.
5. This **IS NOT** an **OPEN BOOK** exam.
6. Candidates should clearly explain their reasoning when answering each question.

QUESTION 1.**(30 marks)**

- (a) Which integer $a \in \{0, 1, \dots, 14\}$ is congruent to $2020 + 1010 + 550 + 225$ modulo 15? (10 marks)
- (b) Write down each integer $a \in \{0, 1, 2\}$ for which there exists an integer n such that $a \equiv n^2 + n - 1 \pmod{3}$. (10 marks)
- (c) Let $S = \{\text{integers congruent to 1 modulo 5}\}$ and Δ be multiplication. Is S closed under Δ ? Justify your answer.

Solution:

- (a) We have

$$\begin{aligned}
 2020 + 1010 + 550 + 225 &= 3030 + (2 \cdot 225 + 100) + 225 \\
 &= 2 \cdot 15 \cdot 101 + 3 \cdot 225 + 100 \\
 &= 2 \cdot 15 \cdot 101 + 15 \cdot 45 + 90 + 10 \\
 &= 15 \cdot (2 \cdot 101 + 45 + 6) + 10
 \end{aligned}$$

Hence $2020 + 1010 + 550 + 225 \equiv 10 \pmod{15}$.

- (b) Modulo 3, we have 3 possibilities for n .

- For $n \equiv 0 \pmod{3}$ we have $n^2 + n - 1 \equiv 2 \pmod{3}$.
- For $n \equiv 1 \pmod{3}$ we have $n^2 + n - 1 \equiv 1 \pmod{3}$.
- For $n \equiv 2 \pmod{3}$ we have $n^2 + n - 1 \equiv 4 + 2 - 1 \equiv 2 \pmod{3}$.

So $a = 1$ or $a = 2$.

- (c) Here S is closed under Δ . Indeed, for generic elements $x \in S$ and $y \in S$, we can write $x = 5p + 1$ and $y = 5q + 1$ for some integers p and q . Then

$$x \cdot y = (5p + 1)(5q + 1) = 25pq + 5p + 5q + 1 = 5(5pq + p + q) + 1,$$

which is congruent to 1 modulo 5.

QUESTION 2.**(40 marks)**

(a) Prove or disprove the following logical equivalences:

(i) (10 marks)

$$p \wedge (T \rightarrow p) \equiv p$$

(ii) (10 marks)

$$(p \wedge q \wedge r) \rightarrow (p \vee s) \equiv (p \rightarrow s) \vee (q \rightarrow s) \vee (r \rightarrow s)$$

(b) Decide whether or not the following argument is valid (20 marks):

$$\begin{aligned} &\neg q \vee p; \\ &\neg q \rightarrow F; \\ &p \rightarrow (\neg r \rightarrow s); \\ &q \rightarrow \neg r \\ &\therefore s \end{aligned}$$

Briefly justify your answers.

Solution:

p	$T \rightarrow p$	$p \wedge (T \rightarrow p)$
T	T	T
F	F	F

This proves the logical equivalence.

(ii) For $p = T$, $q = T$, $r = T$, $s = F$ the LHS is true and the RHS is false. This disproves the logical equivalence.

(b) The argument is valid.

(1) $\neg q \vee p$

(2) $\neg q \rightarrow F$

(3) $p \rightarrow (\neg r \rightarrow s)$

(4) $q \rightarrow \neg r$

(5) $\therefore q$

from (2)

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|---------------------------------------|------------------|
| (6) $\therefore p$ | from (5) and (1) |
| (7) $\therefore \neg r \rightarrow s$ | from (6) and (3) |
| (8) $\therefore \neg r$ | from (5) and (4) |
| (9) $\therefore s$ | from (8) and (7) |

Alternatively, one can show that the argument is valid using a truth table.

QUESTION 3.**(30 marks)**

- (a) Let X and Y be domains, and let $P(x)$ and $Q(y)$ be predicates. Which of the following statements is the *negation* of the statement

$\forall x \in X, \exists y \in Y, P(x) \vee \neg Q(y)$? (10 marks)

- (i) $\forall y \in Y, \exists x \in X, \neg P(x) \wedge Q(y)$;
- (ii) $\exists x \in X, \forall y \in Y, P(x) \rightarrow \neg Q(y)$;
- (iii) $\exists y \in Y, \forall x \in X, \neg P(x) \rightarrow \neg Q(y)$;
- (iv) $\exists x \in X, \forall y \in Y, \neg P(x) \wedge Q(y)$;
- (v) none of the above.

Consider the domains $A = \{3, 4\}$ and $B = \{0, 3, 6\}$ and the predicate $P(x, y) = "x^2 - y \geq 9"$.

Determine the truth value of the following statements:

- (i) $\forall x \in A, \exists y \in B, P(x, y)$. (10 marks);
- (ii) $\exists x \in A, \forall y \in B, P(x, y)$. (10 marks).

Briefly justify your answers.

Solution:

- (a) We can write

$$\forall x \in X, \exists y \in Y, P(x) \vee \neg Q(y) \equiv \exists \forall x \in X, R(x),$$

where $R(x)$ is the predicate $R(y) = \exists y \in Y, P(x) \vee \neg Q(y)$. The negation of " $\forall x \in X, R(x)$ " is " $\exists x \in X, \neg R(x)$ ". Next we see that the negation of $R(x)$ is just $\forall y \in Y, \neg(P(x) \vee \neg Q(y))$. Then

$$\begin{aligned} \neg(P(x) \vee \neg Q(y)) &\equiv \neg P(x) \wedge \neg \neg Q(y) && \text{(De Morgan's law)} \\ &\equiv \neg P(x) \wedge Q(y) && \text{(double negation)} \end{aligned}$$

Hence the answer is (iv).

- (b) (i) True. For $x = 3$ take $y = 0$. For $x = 4$ take $y = 0$.
- (ii) True. For $x = 4$ the predicate is true for each $y \in B$.