Seat No. _____

King Mongkut's University of Technology Thonburi

Midterm Exam of Second Semester, Academic Year 2017

Key 23/3/11

CPE 224 Computer Architectures

Computer Engineering Department, 2nd Yr.
Section: ABCD

Tuesday 6th March 2018

13.00-16.00

Instructions

1. This examination contains 6 problems, 14 pages (including this cover page).
2. The answers must be written in this exam paper. Please read the instructions carefully.
3. A calculator and a paper dictionary are allowed.
4. A single A4-sized handwritten note may be taken into the examination room. The note has to be handed in with the exam.

Students will be punished if they violate any examination rules. The highest punishment is dismissal.

This examination is designed by
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Tel. 081-922-8466

Name: Student ID: Section:

Instruction: This exam has 6 questions for a total of 100 points. Write your NAME, ID, Section on EVERY page of the exam, else your question might not be graded. This is a closed-book exam. However, you are allowed to bring with you one A4 sheet of paper of notes. Hand in your note with the exam before leaving the room. Calculator is allowed.

1. (20 points) -- Basic C & MIPS Instructions

For questions 1.1 and 1.2, assume that the variables f, g, h, i and j are assigned to registers $s0, $s1, $s2, $s3 and $s4 respectively, and that the base address of the arrays A and B are in registers $s6 and $s7, respectively.

1.1 (5 pts) For a C statement B[j] = f + A[g+i] ,what is a corresponding MIPS assembly code?

1.2 (5 pts) Write a C code version of the following MIPS assembly code

```
addi $t0, $s6, 4
add  $t1, $s6, $0
sw   $t1, 0($t0)
lw   $t0, 0($t0)
add  $s0, $t1, $t0
```

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For question 1.3 and 1.4, the table below shows 32-bit values of an array stored in memory.

Address	Data
24	2
28	4
32	3
36	6
40	1

1.3 (5 pts) For the memory locations in the table above, write C code to sort the data from lowest to highest, placing the lowest value in the smallest memory location shown in the table. Assume that the data shown represents the C variable called Array, which is an array of type int (32-bit), and that the first number in the array shown is the first element in the array.

1.4 (5 pts) Translate your C code above to the MIPS assembly code. Assume the base address of Array is stored in register $s6.

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2. (25 points) -- Conditionals and Procedures in MIPS Assembly

Given the following recursive C code segment to calculate 4^n

```
int tp(int n) {
    if (n==0) return 1;
    else return 4*tp(n-1);
}
```

2.1 (10 pts) Translate the function tp into MIPS assembly language. Do not use pseudo-instructions and do not forget to comment your code.

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2.2 (15 pts) Assume that the MIPS code starts at location 80000 in a one-word (4-byte) wide instruction memory. Convert your MIPS assembly code in question 2.1 to the MIPS machine language code (with all fields in decimal format). Enter your code in the table below.

Instruction	Mem address	op	rs	rt	rd	shamt	funct
		op	rs	rt	immediate		
		op	address				
	80000						
	80004						
	80008						
	80012						

Name: Student ID: Section:

3. (15 points) -- Computer Arithmetic

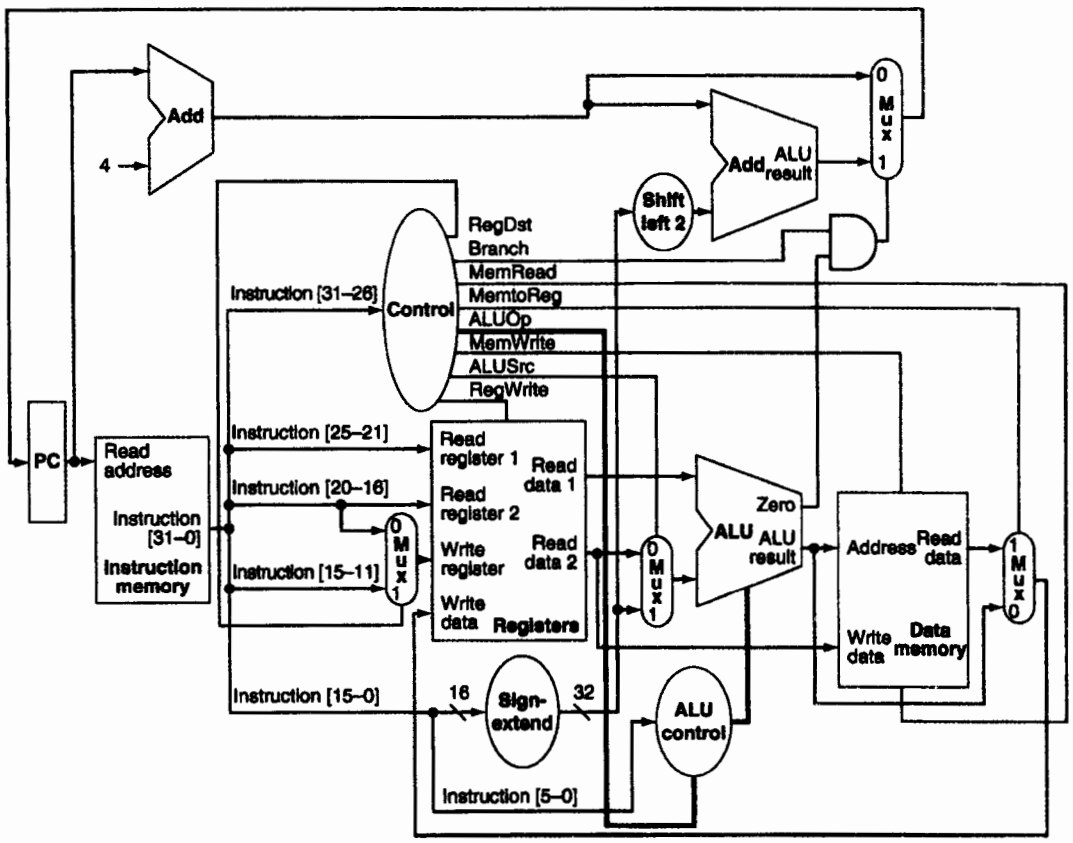
3.1 (5 pts) Assuming single precision IEEE 754 format, what decimal number is represented by this word: 1 01111110 11000000000000000000000

3.2 (5 pts) Show the 32-bit IEEE 754 binary representation of the decimal number -32.50 in single precision.

3.3 (5 pts) IEEE 754-2008 contains a half precision that is only 16 bits wide. The leftmost bit is still the sign bit, the exponent is 5 bits wide and has a bias of 15, and the mantissa is 10 bits long. A hidden 1 is assumed. Write down the bit pattern to represent -1.875×10^{-1} assuming a version of this format, which uses an excess-16 format to store the exponent.

Name: Student ID: Section:

4. (15 points – Processor datapath)
Consider the basic MIPS processor shown in the diagram below.



Assume that the instruction set in this MIPS processor only consist of the following instructions:

- Memory reference: LW and SW
- R-Type: ADD, SUB, AND, OR, SLT
- Control transfer: BEQ

Assume that the delays in the datapath when performing instruction above are as follows:

- Internal Memory and Register Files Access (READ/Write): 50ps
- Register Access (READ/WRITE): 1ps
- ALU and adders: 100ps
- Control logic, Sign Extend, Shift: 20ps
- Logic Gates and Multiplexors: 1ps

Calculate the clock cycle time (clock period) needed for each instructions. Then, calculate the minimal clock cycle time for this MIPS processor. Assume that the processor is a single cycle design (fetch, decode and execute each instruction in one clock cycle).

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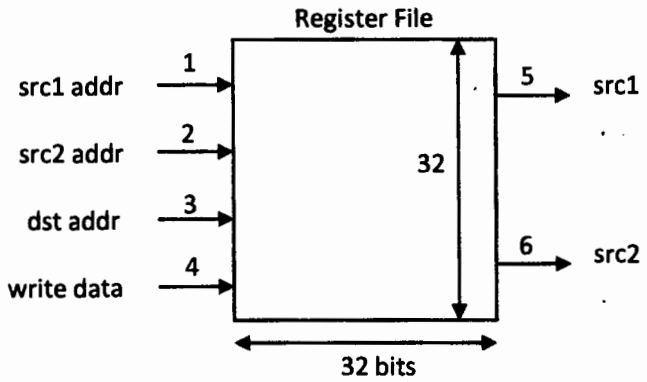
Name: Student ID: Section:

5. (15 points) Short answer (only use the space provided in this exam sheet)

5.1 In MIPS, the execution of instructions can be divided into three steps: Fetch, Decode, Execute. Briefly explain each step.

5.2 MIPS (32-bit architecture) register file contains 32 registers. It has 3 address lines and three data lines (labeled as 1-6 in the figure below). Please specify the width (number of bits) of each line.

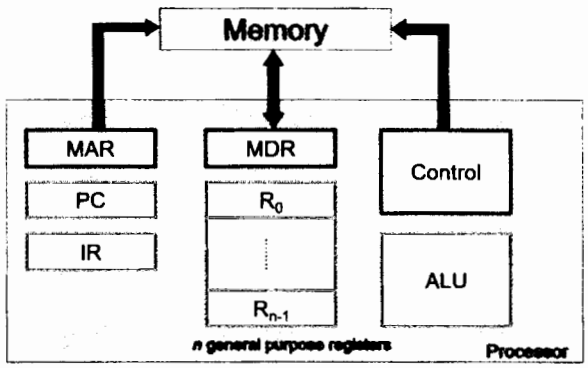


ANS

Line #	Width (Number of bits)
1	
2	
3	
4	
5	
6	

Name: Student ID: Section:

5.3 From the figure below, explain how MDR and MAR are used for READ and WRITE transactions.



5.4 How does the technique of pipelining increase performance? Explain the increased instruction throughput, compared with a multicycle non-pipelined processor. Does pipelining reduce the execution time for individual instructions? Why?

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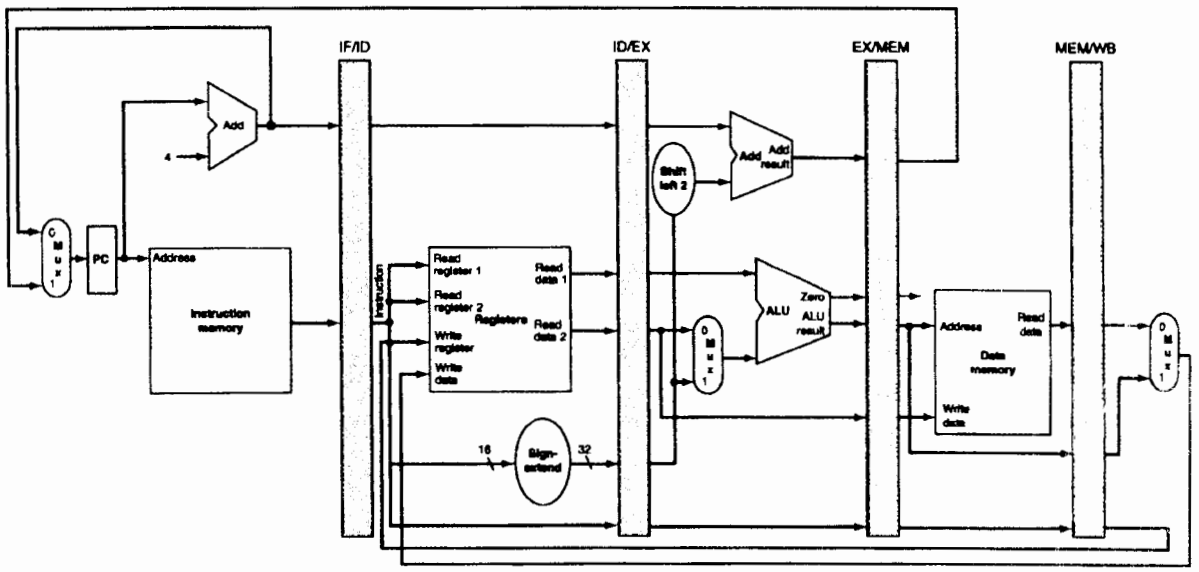
6. (10 pts) Consider the diagram below



6.1 Assume that you have the following sequence of pipelined instructions:

lw $6, 0($7)

add $8, $9, $10

sub $11, $6, $8

Where will the data operands that are processed during the **EX** stage of the subtract (sub) instruction come from?

Name: Student ID: Section:

6.2 Show how the instructions in the sequence given below will proceed through the pipeline:

BEQ $1, $2, X

LW $10, 0($11)

SUB $14, $10, $10

X: add $4, $1, $2

LW $1, 0($4)

SUB $1, $1, $1

ADD $1, $1, $1

Note 1: We will predict that the BEQ instruction is not taken.
Note 2: When the BEQ instruction is executed, the value in $1 is equal to the value in $2
Note 3: Put pipeline stage symbol in the table below (F=Fetch, D=Decode, E=Execute, M=Memory, W=Write)

ANS

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
beq $1, $2, X															
lw $10, 0($11)															
sub $14, $10, $10															
X: add $4, $1, $2															
lw $1, 0($4)															
sub $1, $1, $1															
add $1, $1, $1															

MIPS Reference Data Card ("Green Card") 1. Pull along perforation to separate card 2. Fold bottom side (columns 3 and 4) togcther

MIPS Reference Data

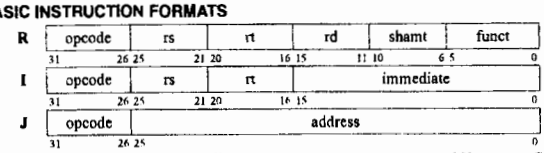
①

CORE INSTRUCTION SET

NAME, MNEMONIC		FOR-MAT	OPERATION (in Verilog)		OPCODE / FUNCT (Hex)
Add	add	R	R[rd] = R[rs] + R[rt]	(1)	0 / 20_{hex}
Add Immediate	addi	I	R[rt] = R[rs] + SignExtImm	(1,2)	8_{hex}
Add Imm. Unsigned	addiu	I	R[rt] = R[rs] + SignExtImm	(2)	9_{hex}
Add Unsigned	addu	R	R[rd] = R[rs] + R[rt]		0 / 21_{hex}
And	and	R	R[rd] = R[rs] & R[rt]		0 / 24_{hex}
And Immediate	andi	I	R[rt] = R[rs] & ZeroExtImm	(3)	c_{hex}
Branch On Equal	beq	I	if(R[rs]==R[rt]) PC=PC+4+BranchAddr	(4)	4_{hex}
Branch On Not Equal	bne	I	if(R[rs]!=R[rt]) PC=PC+4+BranchAddr	(4)	5_{hex}
Jump	j	J	PC=JumpAddr	(5)	2_{hex}
Jump And Link	jal	J	R[31]=PC+8;PC=JumpAddr	(5)	3_{hex}
Jump Register	jr	R	PC=R[rs]		0 / 08_{hex}
Load Byte Unsigned	lbu	I	R[rt]={24'b0,M[R[rs]+SignExtImm](7:0)}	(2)	24_{hex}
Load Halfword Unsigned	lhu	I	R[rt]={16'b0,M[R[rs]+SignExtImm](15:0)}	(2)	25_{hex}
Load Linked	ll	I	R[rt] = M[R[rs]+SignExtImm]	(2,7)	30_{hex}
Load Upper Imm.	lui	I	R[rt] = {imm, 16'b0}		f_{hex}
Load Word	lw	I	R[rt] = M[R[rs]+SignExtImm]	(2)	23_{hex}
Nor	nor	R	R[rd] = ~ (R[rs] | R[rt])		0 / 27_{hex}
Or	or	R	R[rd] = R[rs] | R[rt]		0 / 25_{hex}
Or Immediate	ori	I	R[rt] = R[rs] | ZeroExtImm	(3)	d_{hex}
Set Less Than	slt	R	R[rd] = (R[rs] < R[rt]) ? 1 : 0		0 / $2a_{hex}$
Set Less Than Imm.	slti	I	R[rt] = (R[rs] < SignExtImm)? 1 : 0	(2)	a_{hex}
Set Less Than Imm. Unsigned	sltiu	I	R[rt] = (R[rs] < SignExtImm) ? 1 : 0	(2,6)	b_{hex}
Set Less Than Unsig.	sltu	R	R[rd] = (R[rs] < R[rt]) ? 1 : 0	(6)	0 / $2b_{hex}$
Shift Left Logical	sll	R	R[rd] = R[rt] << shamt		0 / 00_{hex}
Shift Right Logical	srl	R	R[rd] = R[rt] >>> shamt		0 / 02_{hex}
Store Byte	sb	I	M[R[rs]+SignExtImm](7:0) = R[rt](7:0)	(2)	28_{hex}
Store Conditional	sc	I	M[R[rs]+SignExtImm] = R[rt]; R[rt] = (*atomic*) ? 1 : 0	(2,7)	38_{hex}
Store Halfword	sh	I	M[R[rs]+SignExtImm](15:0) = R[rt](15:0)	(2)	29_{hex}
Store Word	sw	I	M[R[rs]+SignExtImm] = R[rt]	(2)	$2b_{hex}$
Subtract	sub	R	R[rd] = R[rs] - R[rt]	(1)	0 / 22_{hex}
Subtract Unsigned	subu	R	R[rd] = R[rs] - R[rt]		0 / 23_{hex}

(1) May cause overflow exception
(2) SignExtImm = { 16{immediate[15]}, immediate }
(3) ZeroExtImm = { 16{1b'0}, immediate }
(4) BranchAddr = { 14{immediate[15]}, immediate, 2'b0 }
(5) JumpAddr = { PC+4[31:28], address, 2'b0 }
(6) Operands considered unsigned numbers (vs. 2's comp.)
(7) Atomic test&set pair; R[rt] = 1 if pair atomic, 0 if not atomic

BASIC INSTRUCTION FORMATS

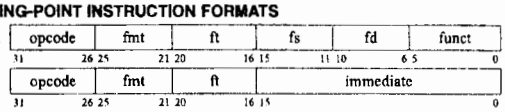


ARITHMETIC CORE INSTRUCTION SET

② OPCODE / FMT /FT / FUNC T (Hex)

NAME, MNEMONIC		FOR-MAT	OPERATION		
Branch On FP True	bc1t	FI	if(FPcond)PC=PC+4+BranchAddr	(4)	11/8/1/--
Branch On FP False	bc1f	FI	if(!FPcond)PC=PC+4+BranchAddr	(4)	11/8/0/--
Divide	div	R	Lo=R[rs]/R[rt]; Hi=R[rs]%R[rt]		0/--/--/1a
Divide Unsigned	divu	R	Lo=R[rs]/R[rt]; Hi=R[rs]%R[rt]	(6)	0/--/--/1b
FP Add Single	add.s	FR	F[fd]= F[fs] + F[ft]		11/10/--/0
FP Add Double	add.d	FR	{F[fd],F[fd+1]} = {F[fs],F[fs+1]} + {F[ft],F[ft+1]}		11/11/--/0
FP Compare Single	c.x.s*	FR	FPcond = (F[fs] *op* F[ft]) ? 1 : 0		11/10/--/*y*
FP Compare Double	c.x.d*	FR	FPcond = ({F[fs],F[fs+1]} *op* {F[ft],F[ft+1]}) ? 1 : 0		11/11/--/*y*
* (*x* is eq, lt, or le) (*op* is ==, <, or <=) (*y* is 32, 3c, or 3e)					
FP Divide Single	div.s	FR	F[fd] = F[fs] / F[ft]		11/10/--/3
FP Divide Double	div.d	FR	{F[fd],F[fd+1]} = {F[fs],F[fs+1]} / {F[ft],F[ft+1]}		11/11/--/3
FP Multiply Single	mul.s	FR	F[fd] = F[fs] * F[ft]		11/10/--/2
FP Multiply Double	mul.d	FR	{F[fd],F[fd+1]} = {F[fs],F[fs+1]} * {F[ft],F[ft+1]}		11/11/--/2
FP Subtract Single	sub.s	FR	F[fd]=F[fs] - F[ft]		11/10/--/1
FP Subtract Double	sub.d	FR	{F[fd],F[fd+1]} = {F[fs],F[fs+1]} - {F[ft],F[ft+1]}		11/11/--/1
Load FP Single	lwc1	I	F[rt]=M[R[rs]+SignExtImm]	(2)	31/--/--/--
Load FP Double	ldc1	I	F[rt]=M[R[rs]+SignExtImm]; F[rt+1]=M[R[rs]+SignExtImm+4]	(2)	35/--/--/--
Move From Hi	mfhi	R	R[rd] = Hi		0 /--/--/10
Move From Lo	mflo	R	R[rd] = Lo		0 /--/--/12
Move From Control	mfc0	R	R[rd] = CR[rs]		10 /0/--/0
Multiply	mult	R	{Hi,Lo} = R[rs] * R[rt]		0/--/--/18
Multiply Unsigned	multu	R	{Hi,Lo} = R[rs] * R[rt]	(6)	0/--/--/19
Shift Right Arith.	sra	R	R[rd] = R[rt] >> shamt		0/--/--/3
Store FP Single	swc1	I	M[R[rs]+SignExtImm] = F[rt]	(2)	39/--/--/--
Store FP Double	sdc1	I	M[R[rs]+SignExtImm] = F[rt]; M[R[rs]+SignExtImm+4] = F[rt+1]	(2)	3d/--/--/--

FLOATING-POINT INSTRUCTION FORMATS



PSEUDOINSTRUCTION SET

NAME	MNEMONIC	OPERATION
Branch Less Than	blt	if(R[rs]<R[rt]) PC = Label
Branch Greater Than	bgt	if(R[rs]>R[rt]) PC = Label
Branch Less Than or Equal	ble	if(R[rs]<=R[rt]) PC = Label
Branch Greater Than or Equal	bge	if(R[rs]>=R[rt]) PC = Label
Load Immediate	li	R[rd] = immediate
Move	move	R[rd] = R[rs]

REGISTER NAME, NUMBER, USE, CALL CONVENTION

NAME	NUMBER	USE	PRESERVED ACROSS A CALL?
$zero	0	The Constant Value 0	N.A.
$at	1	Assembler Temporary	No
$v0-$v1	2-3	Values for Function Results and Expression Evaluation	No
$a0-$a3	4-7	Arguments	No
$t0-$t7	8-15	Temporaries	No
$s0-$s7	16-23	Saved Temporaries	Yes
$t8-$t9	24-25	Temporaries	No
$k0-$k1	26-27	Reserved for OS Kernel	No
$gp	28	Global Pointer	Yes
$sp	29	Stack Pointer	Yes
$fp	30	Frame Pointer	Yes
$ra	31	Return Address	Yes

OPCODES, BASE CONVERSION, ASCII SYMBOLS ③

MIPS opcode (31:26)	(1) MIPS funct (5:0)	(2) MIPS funct (5:0)	Binary	Decimal	Hexadecimal	ASCII Character	Decimal	Hexadecimal	ASCII Character
(1)	sll	add.f	00 0000	0	0	NUL	64	40	@
		sub.f	00 0001	1	1	SOH	65	41	A
j	srl	mul.f	00 0010	2	2	STX	66	42	B
jal	sra	div.f	00 0011	3	3	ETX	67	43	C
beq	sllv	sqrt.f	00 0100	4	4	EOT	68	44	D
one		abs.f	00 0101	5	5	ENQ	69	45	E
blez	srlv	mov.f	00 0110	6	6	ACK	70	46	F
bgtz	srav	neg.f	00 0111	7	7	BEL	71	47	G
addi	jr		00 1000	8	8	BS	72	48	H
addiu	jalr		00 1001	9	9	HT	73	49	I
slti	movz		00 1010	10	a	LF	74	4a	J
sltiu	movn		00 1011	11	b	VT	75	4b	K
andi	syscall	round.w.f	00 1100	12	c	FF	76	4c	L
ori	break	trunc.w.f	00 1101	13	d	CR	77	4d	M
xori		ceil.w.f	00 1110	14	e	SO	78	4e	N
lui	sync	floor.w.f	00 1111	15	f	SI	79	4f	O
	mfhi		01 0000	16	10	DLE	80	50	P
(2)	mthi		01 0001	17	11	DC1	81	51	Q
	mflo	movz.f	01 0010	18	12	DC2	82	52	R
	mtlo	movn.f	01 0011	19	13	DC3	83	53	S
			01 0100	20	14	DC4	84	54	T
			01 0101	21	15	NAK	85	55	U
			01 0110	22	16	SYN	86	56	V
			01 0111	23	17	ETB	87	57	W
	mult		01 1000	24	18	CAN	88	58	X
	multu		01 1001	25	19	EM	89	59	Y
	div		01 1010	26	1a	SUB	90	5a	Z
	divu		01 1011	27	1b	ESC	91	5b	[
			01 1100	28	1c	FS	92	5c	\
			01 1101	29	1d	GS	93	5d	]
			01 1110	30	1e	RS	94	5e	^
			01 1111	31	1f	US	95	5f	_
lb	add	cvt.s.f	10 0000	32	20	Space	96	60	`
lh	addu	cvt.d.f	10 0001	33	21	!	97	61	a
lwl	sub		10 0010	34	22	"	98	62	b
lw	subu		10 0011	35	23	#	99	63	c
lbu	and	cvt.w.f	10 0100	36	24	$	100	64	d
lhu	or		10 0101	37	25	%	101	65	e
lwr	xor		10 0110	38	26	&	102	66	f
	nor		10 0111	39	27	'	103	67	g
sb			10 1000	40	28	(	104	68	h
sh			10 1001	41	29	)	105	69	i
swl	slt		10 1010	42	2a	*	106	6a	j
sw	sltu		10 1011	43	2b	+	107	6b	k
			10 1100	44	2c	,	108	6c	l
			10 1101	45	2d	-	109	6d	m
swr			10 1110	46	2e	.	110	6e	n
cache			10 1111	47	2f	/	111	6f	o
ll	tge	c.f.f	11 0000	48	30	0	112	70	p
lwc1	tgeu	c.un.f	11 0001	49	31	1	113	71	q
lwc2	tlt	c.eq.f	11 0010	50	32	2	114	72	r
pref	tltu	c.ueq.f	11 0011	51	33	3	115	73	s
	teq	c.olt.f	11 0100	52	34	4	116	74	t
ldc1		c.ult.f	11 0101	53	35	5	117	75	u
ldc2	tne	c.ole.f	11 0110	54	36	6	118	76	v
		c.ule.f	11 0111	55	37	7	119	77	w
sc		c.sf.f	11 1000	56	38	8	120	78	x
swc1		c.ngle.f	11 1001	57	39	9	121	79	y
swc2		c.seq.f	11 1010	58	3a	:	122	7a	z
		c.ngl.f	11 1011	59	3b	;	123	7b	{
		c.lt.f	11 1100	60	3c	<	124	7c	|
sdc1		c.nge.f	11 1101	61	3d	=	125	7d	}
sdc2		c.le.f	11 1110	62	3e	>	126	7e	~
		c.ngt.f	11 1111	63	3f	?	127	7f	DEL

(1) opcode(31:26) == 0

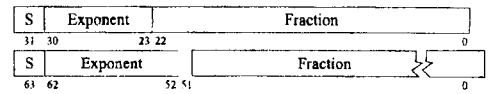
(2) opcode(31:26) == 17_{ten} (11_{hex}); if fmt(25:21)==16_{ten} (10_{hex}) f = s (single); if fmt(25:21)==17_{ten} (11_{hex}) f = d (double)

IEEE 754 FLOATING-POINT STANDARD

$(-1)^S \times (1 + \text{Fraction}) \times 2^{(\text{Exponent} - \text{Bias})}$

where Single Precision Bias = 127, Double Precision Bias = 1023.

IEEE Single Precision and Double Precision Formats:

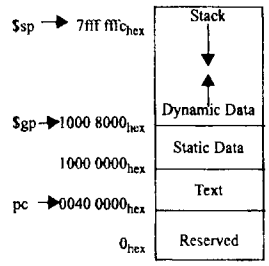


④ IEEE 754 Symbols

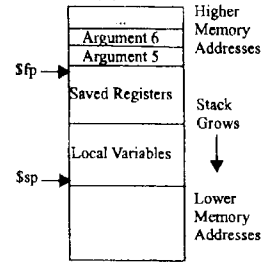
Exponent	Fraction	Object
0	0	± 0
0	≠0	± Denorm
1 to MAX - 1	anything	± Fl. Pt. Num.
MAX	0	±∞
MAX	≠0	NaN

S.P. MAX = 255, D.P. MAX = 2047

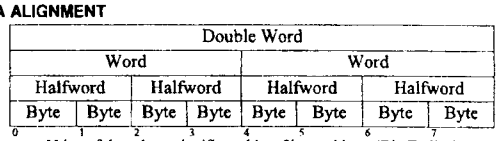
MEMORY ALLOCATION



STACK FRAME

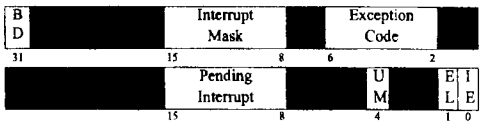


DATA ALIGNMENT



Value of three least significant bits of byte address (Big Endian)

EXCEPTION CONTROL REGISTERS: CAUSE AND STATUS



BD = Branch Delay, UM = User Mode, EL = Exception Level, IE =Interrupt Enable

EXCEPTION CODES

Number	Name	Cause of Exception	Number	Name	Cause of Exception
0	Int	Interrupt (hardware)	9	Bp	Breakpoint Exception
4	AdEL	Address Error Exception (load or instruction fetch)	10	RI	Reserved Instruction Exception
5	AdES	Address Error Exception (store)	11	CpU	Coprocessor Unimplemented
6	IBE	Bus Error on Instruction Fetch	12	Ov	Arithmetic Overflow Exception
7	DBE	Bus Error on Load or Store	13	Tr	Trap
8	Sys	Syscall Exception	15	FPE	Floating Point Exception

SIZE PREFIXES (10^x for Disk, Communication; 2^x for Memory)

SIZE	PRE-FIX	SIZE	PRE-FIX	SIZE	PRE-FIX	SIZE	PRE-FIX
10^3, 2^{10}	Kilo-	10^{15}, 2^{50}	Peta-	10^{-3}	milli-	10^{-15}	femto-
10^6, 2^{20}	Mega-	10^{18}, 2^{60}	Exa-	10^{-6}	micro-	10^{-18}	atto-
10^9, 2^{30}	Giga-	10^{21}, 2^{70}	Zetta-	10^{-9}	nano-	10^{-21}	zepto-
10^{12}, 2^{40}	Tera-	10^{24}, 2^{80}	Yotta-	10^{-12}	pico-	10^{-24}	yocto-

The symbol for each prefix is just its first letter, except μ is used for micro.

MIPS Reference Data Card ("Green Card") 1. Pull along perforation to separate card 2. Fold bottom side (columns 3 and 4) together