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TE COMPS

Experiment No.5

Aim:- Write a program to find first and follow sets for the given grammar. Program should accept the grammar from the user and output the first and follow sets for each of the grammar symbols.

Theory:-

$\text{First}(\alpha)$ is a set of terminals that begins strings derived from α . If $\alpha \Rightarrow \epsilon$ then ϵ is also in $\text{First}(\epsilon)$.

In predictive parsing when we have $A \rightarrow \alpha|\beta$, if $\text{First}(\alpha)$ and $\text{First}(\beta)$ are disjoint sets then we can select appropriate A-production by looking at the next input.

$\text{Follow}(A)$, for any nonterminal A, is set of terminals a that can appear immediately after A in some sentential form

If we have $S \Rightarrow \alpha A a \beta$ for some α and β then a is in $\text{Follow}(A)$

If A can be the rightmost symbol in some sentential form, then \$ is in $\text{Follow}(A)$

Finding First:-

First(x) for all grammar symbols X

Apply following rules:

1. If X is terminal, $\text{FIRST}(X) = \{X\}$.
2. If $X \rightarrow \epsilon$ is a production, then add ϵ to $\text{FIRST}(X)$.
3. If X is a non-terminal, and $X \rightarrow Y_1 Y_2 \dots Y_k$ is a production, and ϵ is in all of $\text{FIRST}(Y_1), \dots, \text{FIRST}(Y_k)$, then add ϵ to $\text{FIRST}(X)$.
4. If X is a non-terminal, and $X \rightarrow Y_1 Y_2 \dots Y_k$ is a production, then add a to $\text{FIRST}(X)$ if for some i, a is in $\text{FIRST}(Y_i)$, and ϵ is in all of $\text{FIRST}(Y_1), \dots, \text{FIRST}(Y_{i-1})$.

Applying rules 1 and 2 is obvious. Applying rules 3 and 4 for $\text{FIRST}(Y_1 Y_2 \dots Y_k)$ can be done as follows:

Add all the non- ϵ symbols of $\text{FIRST}(Y_1)$ to $\text{FIRST}(Y_1 Y_2 \dots Y_k)$. If $\epsilon \in \text{FIRST}(Y_1)$, add all the non- ϵ symbols of $\text{FIRST}(Y_2)$. If $\epsilon \in \text{FIRST}(Y_1)$ and $\epsilon \in \text{FIRST}(Y_2)$, add all the non- ϵ symbols of $\text{FIRST}(Y_3)$, and so on. Finally, add ϵ to $\text{FIRST}(Y_1 Y_2 \dots Y_k)$ if $\epsilon \in \text{FIRST}(Y_i)$, for all $1 \leq i \leq k$.

Example:

Consider the following grammar.

$E \rightarrow E + T \mid T$

$T \rightarrow T * F \mid F$

$F \rightarrow (E) \mid \text{id}$

Grammar after removing left recursion:

$E \rightarrow TX$

$X \rightarrow +TX \mid \epsilon$

$T \rightarrow FY$

$Y \rightarrow *FY \mid \epsilon$

$F \rightarrow (E) \mid \text{id}$

For the above grammar, following the above rules, the FIRST sets could be computed as follows:

$\text{FIRST}(E) = \text{FIRST}(T) = \text{FIRST}(F) = \{ (, \text{id} \}$

$\text{FIRST}(X) = \{ +, \epsilon \}$

$\text{FIRST}(Y) = \{ *, \epsilon \}$

Finding Follow:-

To compute Follow(A) for all nonterminals A, apply following rules until nothing can be added to any follow set:

1. Place \$ in Follow(S) where S is the start symbol
2. If there is a production $A \rightarrow \alpha B \beta$ then everything in First(β) except ϵ is in Follow(B).
3. If there is a production $A \rightarrow B$ or a production $A \rightarrow \alpha B \beta$ where First(β) contains ϵ , then everything in Follow(A) is in Follow(B)

Example:-

$E \rightarrow TX$

$X \rightarrow +TX \mid \epsilon$

$T \rightarrow FY$

$Y \rightarrow *FY \mid \epsilon$

$F \rightarrow (E) \mid \text{id}$

$\text{FOLLOW}(E) = \{ \$,) \}$

$\text{FOLLOW}(X) = \{ \$,) \}$

$\text{FOLLOW}(T) = \{ +, \$,) \}$

$\text{FOLLOW}(Y) = \{ +, \$,) \}$

$\text{FOLLOW}(F) = \{ *, +, \$,) \}$

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Code:-

```
n=int(input("Enter Number of productions "))
Sym=[]
V=[]
T=[]
P=[]
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for i in range(n):
    prod=input("Enter production of the form V->(V U T) ")
    v=prod[0]
    if v not in V:
        V.append(v)
        P.append([])
    if(v not in Sym):
        Sym.append(v)
    ind=V.index(v)
    for j in range(3,len(prod)):
        if(prod[j] not in Sym and prod[j]!='|' and prod[j]!='ε'):
            Sym.append(prod[j])
    lst=prod[3:].split('|')
    for z in lst:
        P[ind].append(z)
for j in Sym:
    if j not in V:
        T.append(j)
First=[]
for i in range(len(Sym)):
    First.append([])
Sym=[]
Sym.extend(T)
Sym.extend(V)
eps=[]
t=len(T)
for i in range(t):
    First[i].append(Sym[i])
for j in range(t,len(Sym)):
    for x in P[j-t]:
        if(x[0] in T):
            First[j].append(x[0])

    if('ε' in P[j-t]):
        First[j].append('ε')
        eps.append(V[j-t])
change=1

```

```

while (change!=0) :
    change=0
    for j in range(t,len(Sym)):
        if(Sym[j] not in eps):
            for p in P[j-t]:
                flag=0
                for k in p:
                    if(k not in eps):
                        flag=1
                        break
                if(flag==0):
                    First[j].append('ε')
                    eps.append(V[j-t])
                    change+=1
                    break
change=1
while (change!=0) :
    change=0
    for j in range(t,len(Sym)):
        for p in P[j-t]:
            if(p!='ε'):
                k=0
                idx=Sym.index(p[k])
                for z in First[idx]:
                    if(z not in First[j] and z!='ε'):
                        First[j].append(z)
                        change+=1
                while('ε' in First[idx] and k<len(p)):
                    k+=1
                    idx=Sym.index(p[k])
                    for z in First[idx]:
                        if (z not in First[j] and z!='ε'):
                            First[j].append(z)
                            change+=1
print("Symbol    First-Pos")
for i in range(len(V)):
    print(V[i], "    ",First[t+i])

```

```

Follow=[]
for i in range(len(V)):
    Follow.append([])
Follow[0].append('$')
for i in P:
    for j in i:
        for k in range(len(j)-1):
            var=j[k]
            if(var in V):
                idx=V.index(var)
                l=k+1
                while(l<len(j)):
                    if(j[l] in T):
                        Follow[idx].append(j[l])
                        break
                    else:
                        ind=Sym.index(j[l])
                        for x in First[ind]:
                            if(x not in Follow[idx] and x!='ε'):
                                Follow[idx].append(x)
                            if(j[l] not in eps):
                                break
                        l+=1
change=1
while(change!=0):
    change=0
    for i in range(len(P)):
        for j in P[i]:
            l=len(j)-1
            if(j[l] in V):
                ind=Sym.index(j[l])-t
                for x in Follow[i]:
                    if(x not in Follow[ind] and x!='ε'):
                        Follow[ind].append(x)
                        change+=1
                l-=1
                while(l>=0):

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        if(j[l+1] in eps):
            ind=Sym.index(j[l])-t
            for x in Follow[i]:
                if(x not in Follow[ind] and x!='ε'):
                    Follow[ind].append(x)
                    change+=1

            l-=1
        else:
            break
print("Symbol      Follow-Pos")
for i in range(len(Follow)):
    print(V[i], "      ", Follow[i])

```

Output:-

```

C:\Users\aumkar\Downloads>python exp5.py
Enter Number of productions 6
Enter production of the form V->(V U T) S->ABCDE
Enter production of the form V->(V U T) A->a|ε
Enter production of the form V->(V U T) B->b|ε
Enter production of the form V->(V U T) C->c
Enter production of the form V->(V U T) D->d|ε
Enter production of the form V->(V U T) E->e|ε
Symbol      First-Pos
S           ['a', 'b', 'c']
A           ['a', 'ε']
B           ['b', 'ε']
C           ['c']
D           ['d', 'ε']
E           ['e', 'ε']
Symbol      Follow-Pos
S           ['$']
A           ['b', 'c']
B           ['c']
C           ['d', 'e', '$']
D           ['e', '$']
E           ['$']

```

```

C:\Users\aumkar\Downloads>python exp5.py
Enter Number of productions 5
Enter production of the form V->(V U T) E->TG
Enter production of the form V->(V U T) G->+TG|ε
Enter production of the form V->(V U T) T->FH
Enter production of the form V->(V U T) H->*FH|ε
Enter production of the form V->(V U T) F->i|(E)
Symbol    First-Pos
E          ['i', '(']
G          ['+', 'ε']
T          ['i', '(']
H          ['*', 'ε']
F          ['i', '(']
Symbol    Follow-Pos
E          ['$ ', ')']
G          ['$ ', ')']
T          ['+', '$ ', ')']
H          ['+', '$ ', ')']
F          ['*', '+', '$ ', ')']

```

Conclusion:-

Thus the given program can determine the first and follow pos for all non-terminals in a given grammar.