

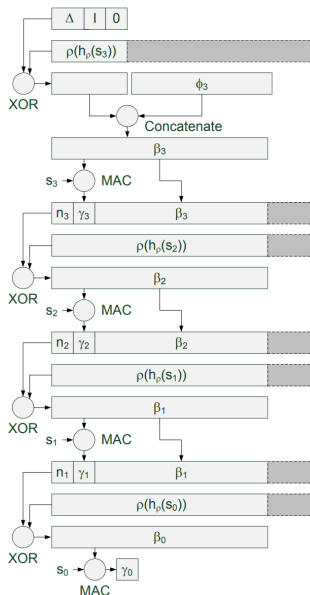
Sphinx Packets

Decentralized Header Construction

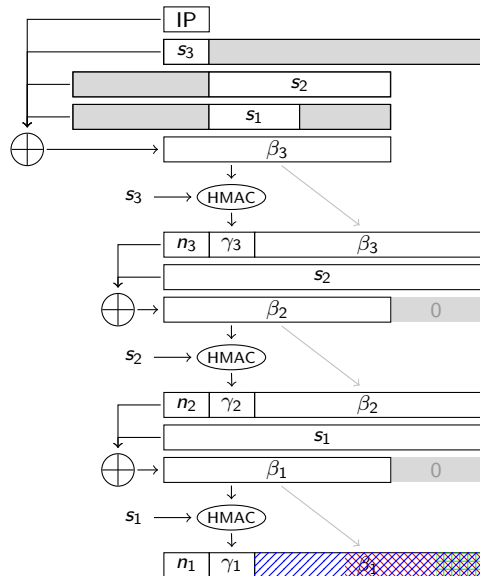
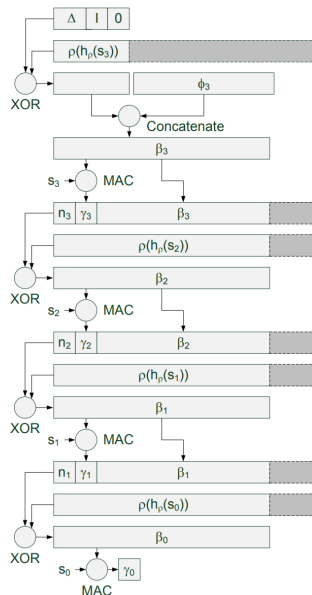
Aurélien Chassagne

February 18, 2025

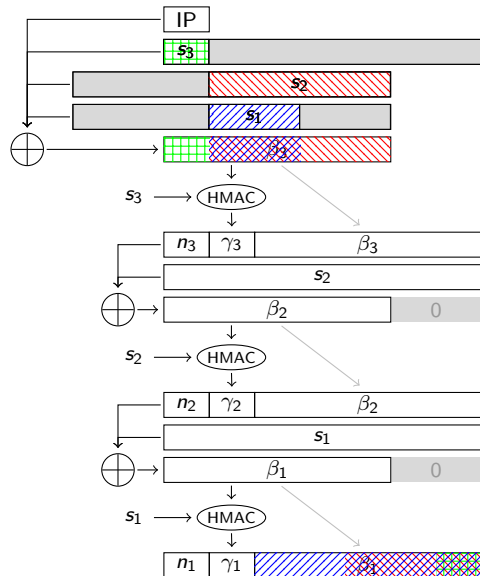
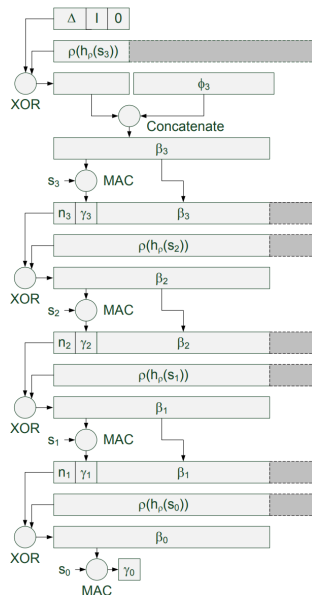
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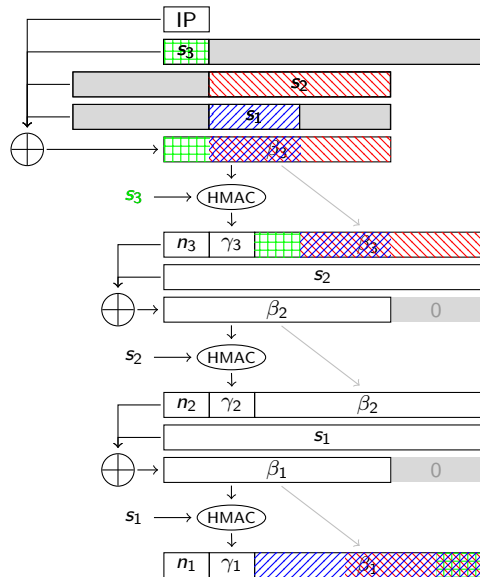
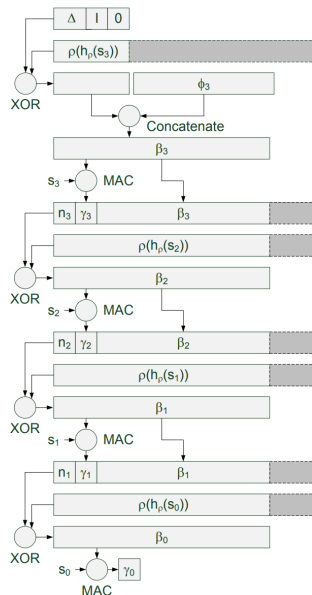
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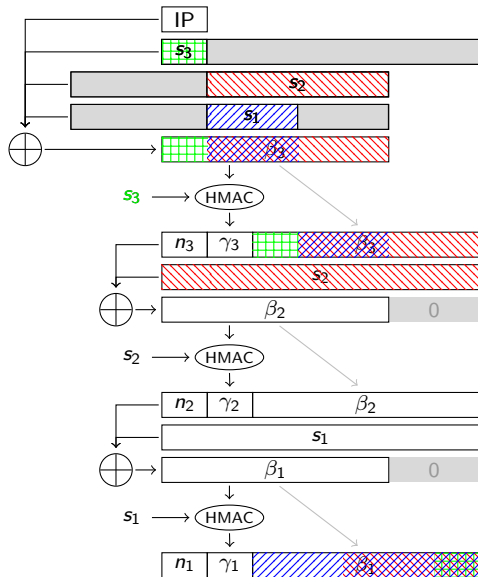
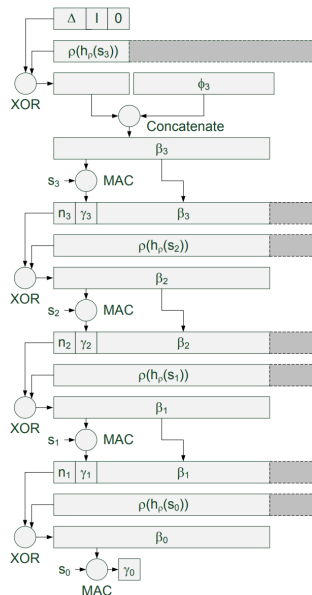
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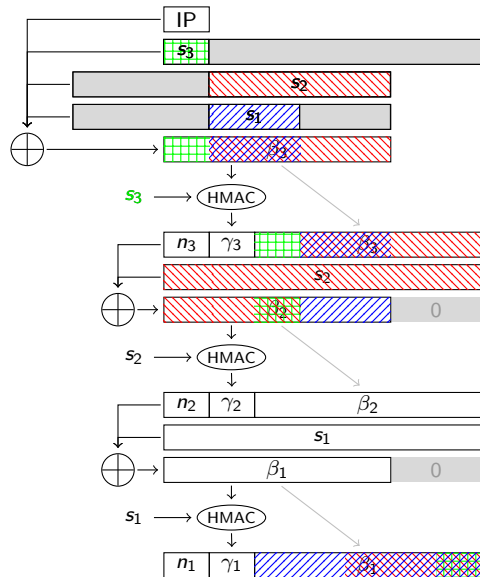
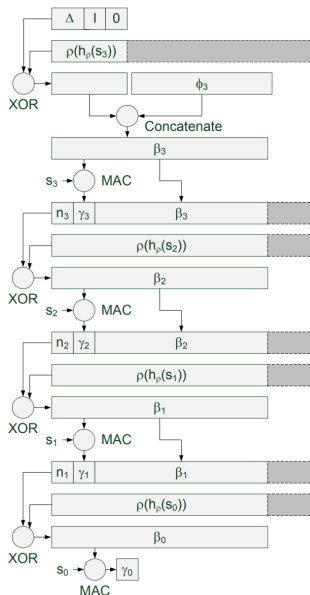
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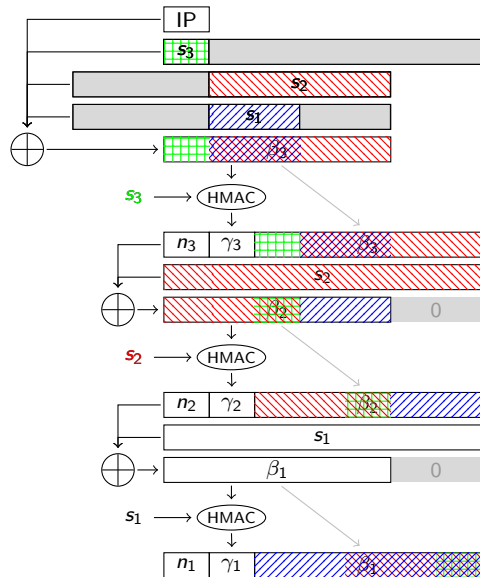
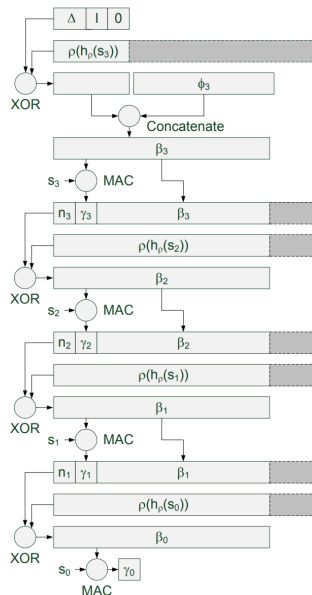
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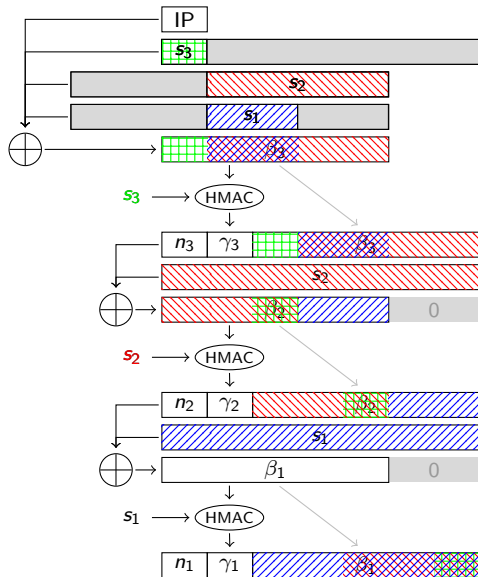
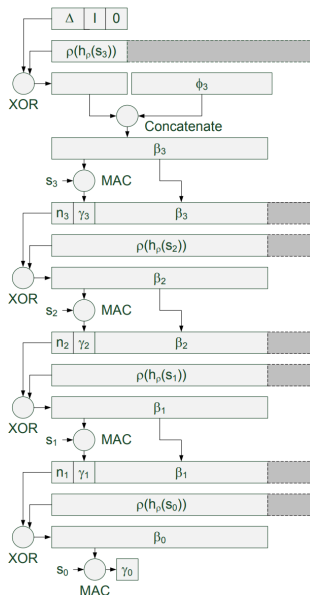
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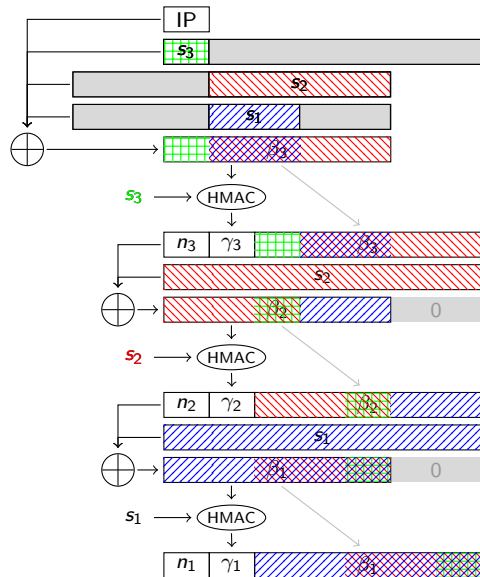
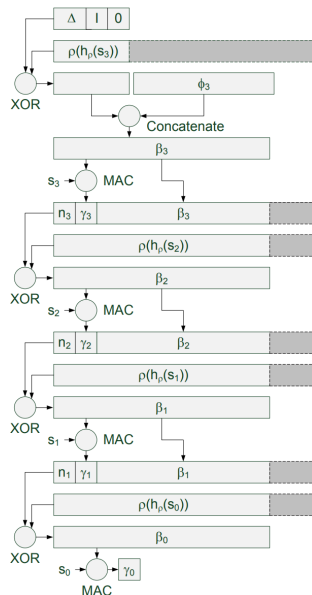
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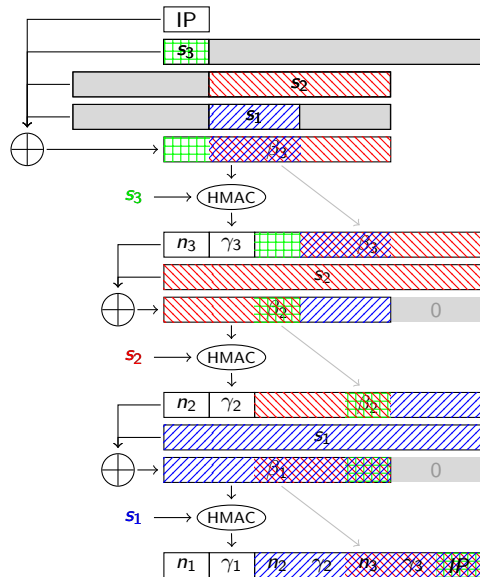
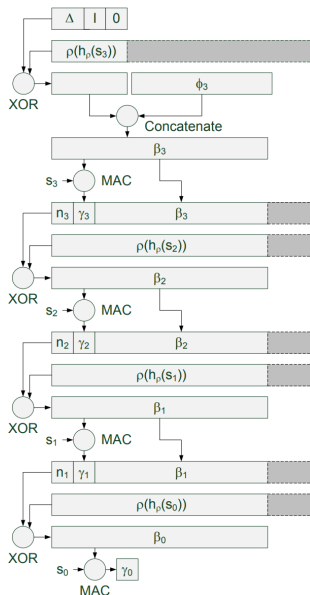
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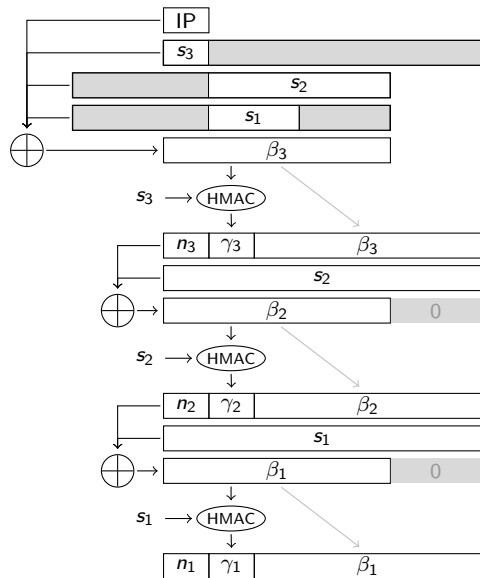
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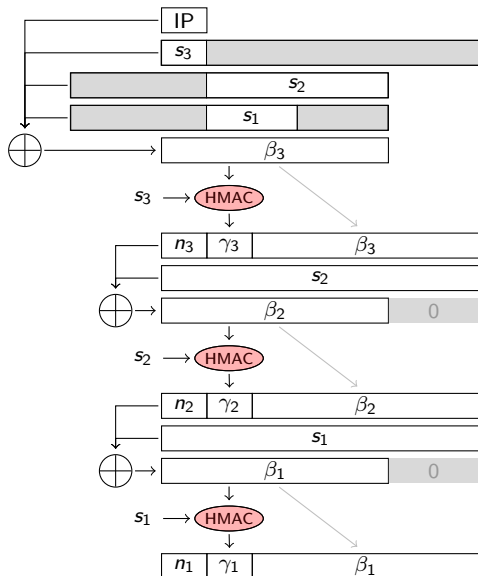


Adapting HMAC



Adapting HMAC

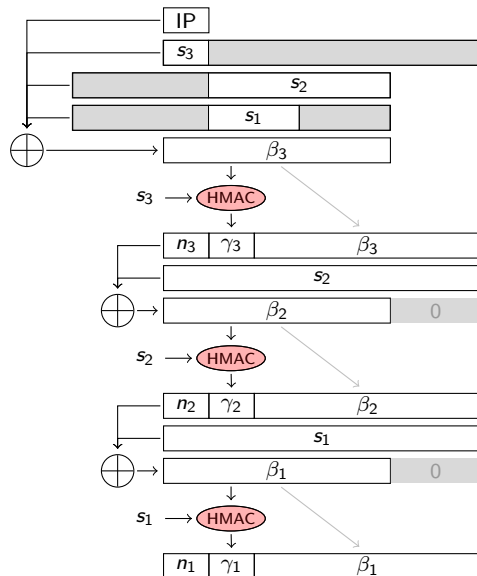
Main problem: Decentralizing a **Hash** ?



Adapting HMAC

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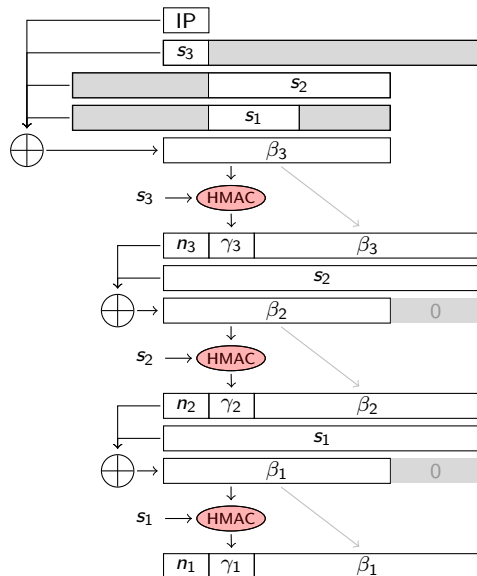
- Need **homomorphic properties** to split computation and aggregate results.



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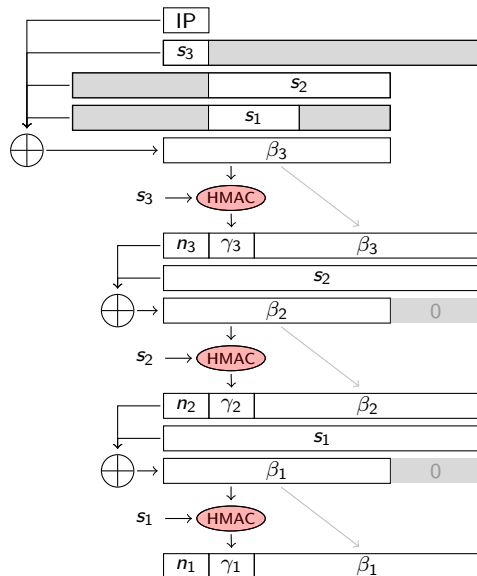
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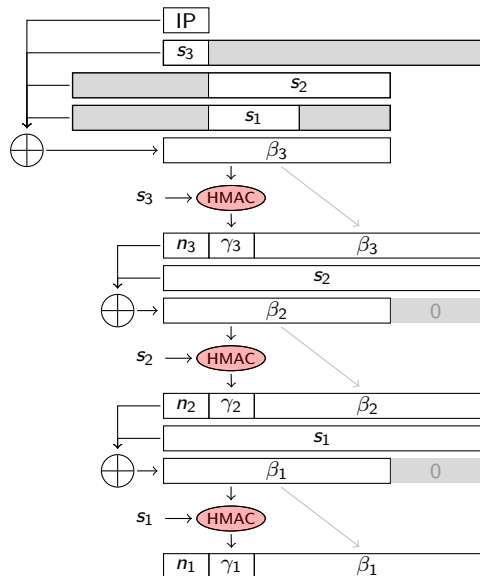


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- 1 RSA
- 2 ElGamal
- 3 Paillier



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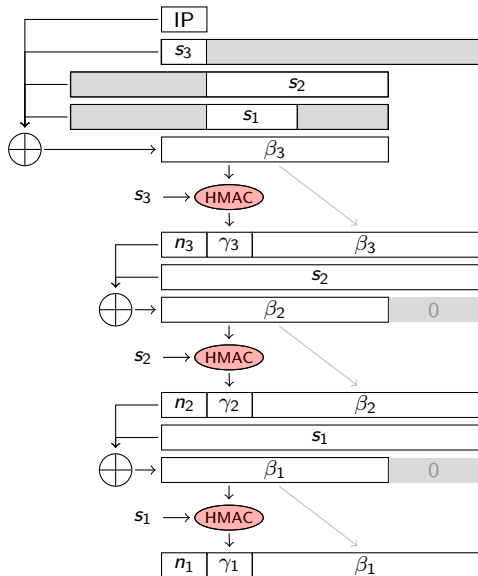
③ Paillier

$$\mathcal{E}(m_1) \cdot \mathcal{E}(m_2) = (g^{m_1} r_1^n)(g^{m_2} r_2^n) \bmod n^2$$

$$= g^{m_1+m_2} (r_1 r_2)^n \bmod n^2$$

$$= \mathcal{E}(m_1 + m_2).$$

Problem: Mix of different operations... order matters!



Adapting HMAC

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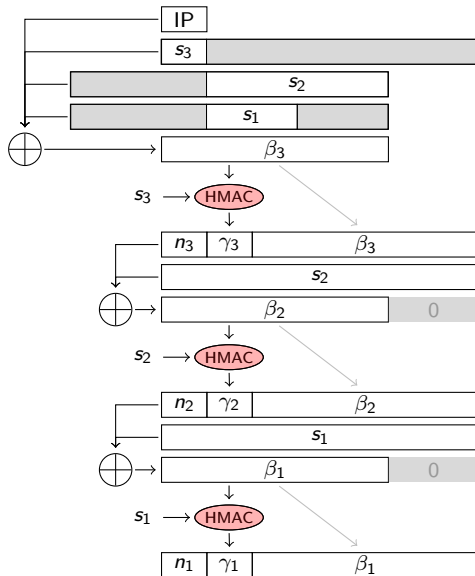
$$\mathcal{E}(m_1) \cdot \mathcal{E}(m_2) = (g^{r_1}, m_1 \cdot h^{r_1})(g^{r_2}, m_2 \cdot h^{r_2})$$

$$= (g^{r_1+r_2}, (m_1 \cdot m_2)h^{r_1+r_2})$$

$$= \mathcal{E}(m_1 \cdot m_2)$$

3 Paillier

Limitation: Increase ciphertext size...



Adapting HMAC

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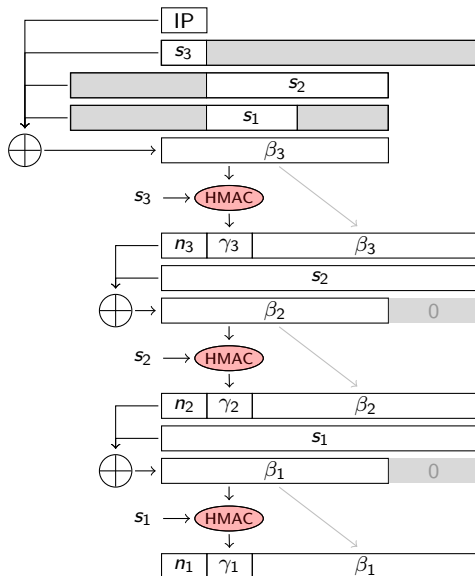
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Adapting HMAC

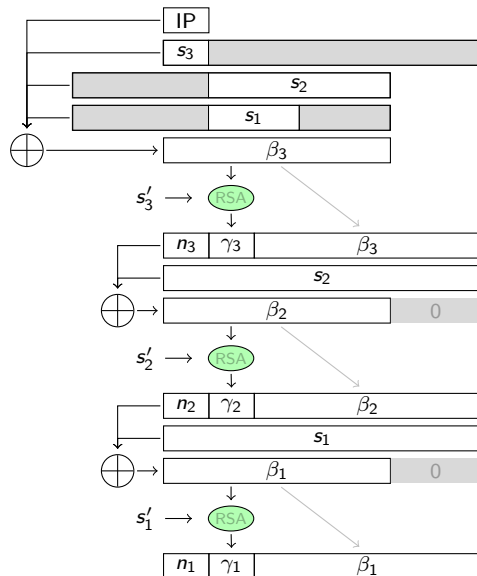
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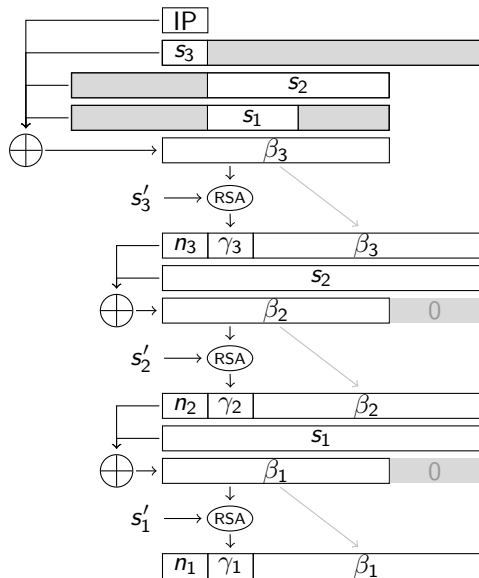
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Selected solution: **RSA** for integrity tag

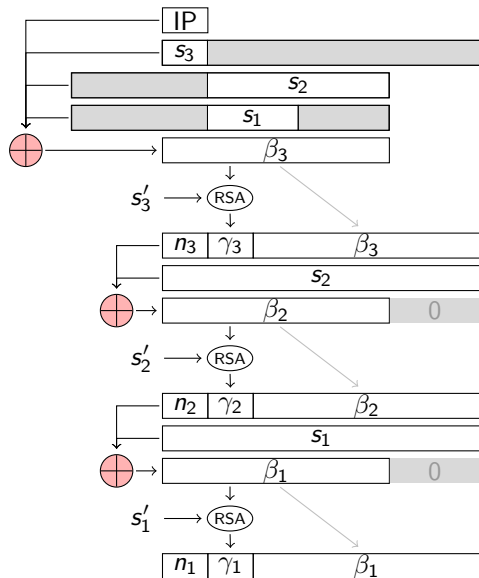
NB: s_i is different for each TTP but RSA required the same e...
Thus, create a new shared secret s'_i common to all TTP



Adapting XOR

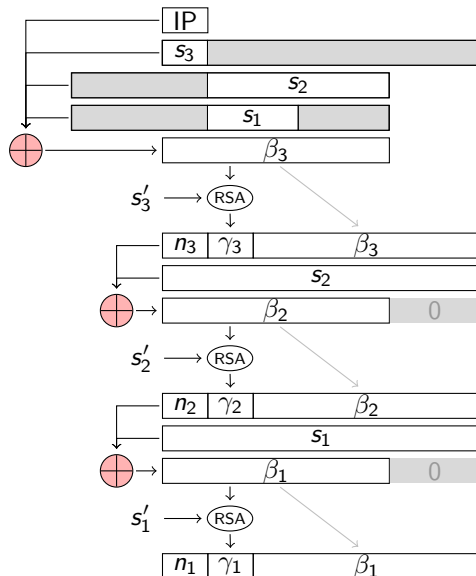


Adapting XOR



Adapting XOR

Since we use **RSA** for integrity tag γ_i



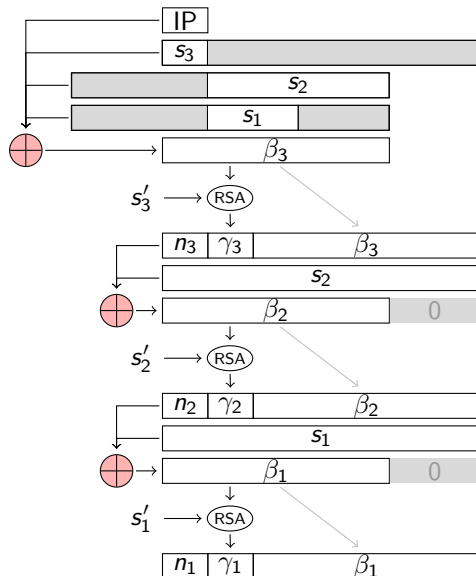
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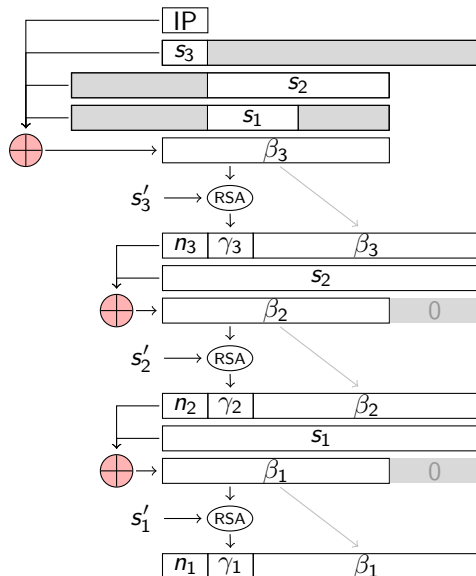


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Modular multiplication of integrity tags gives integrity tag of headers' modular product.



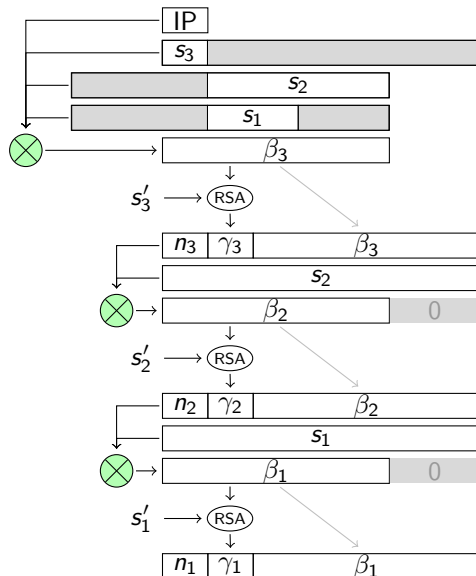
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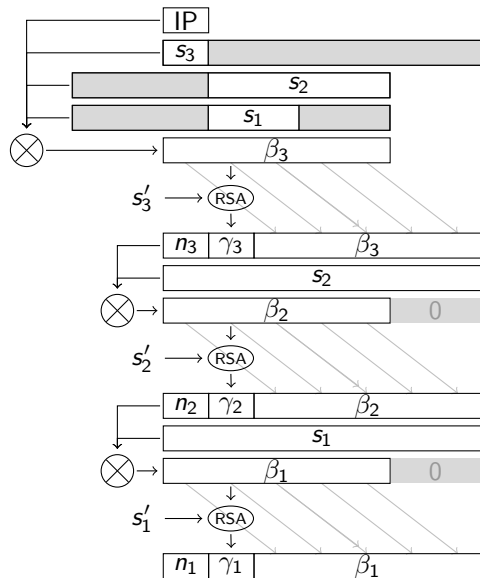
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Modular multiplication of integrity tags gives integrity tag of headers' modular product.

Thus, header elements must be combined via **modular multiplication** rather than **XOR**.

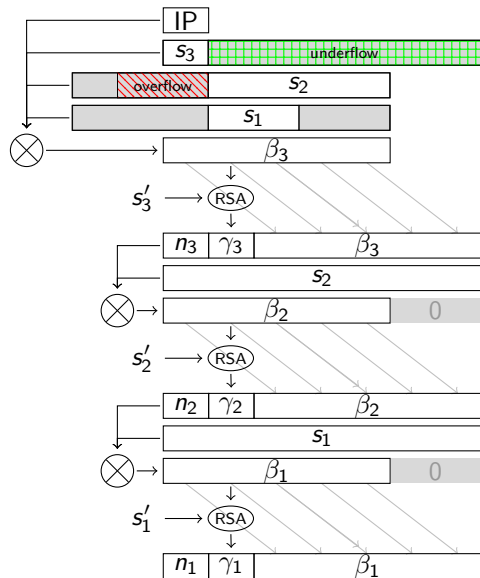


Handle overflows



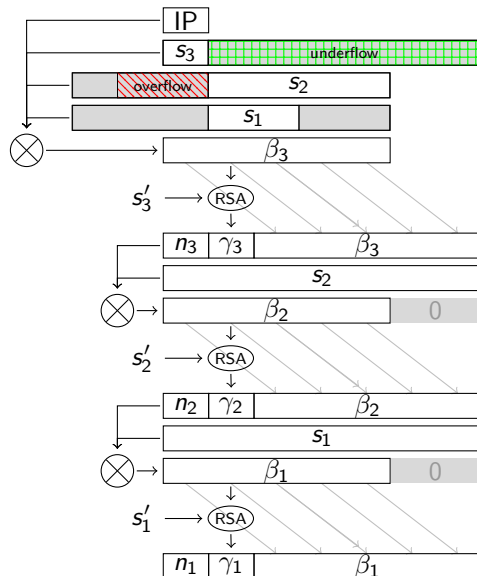
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- Handling these issues is challenging.
It may lead to **information loss** (further research is required).



Handle overflows

- **Current Challenge:** **Overflow** issues
- Handling these issues is challenging.
It may lead to **information loss** (further research is required).
- **Proposed solution:** Simplify by **dividing data into small chunks** and processing each chunk modulo its size.

