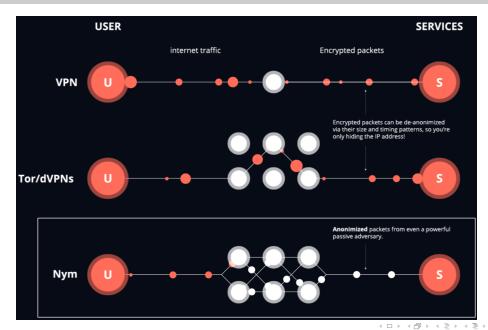
# Sphinx Packets

Decentralized Header Construction

Aurélien Chassagne

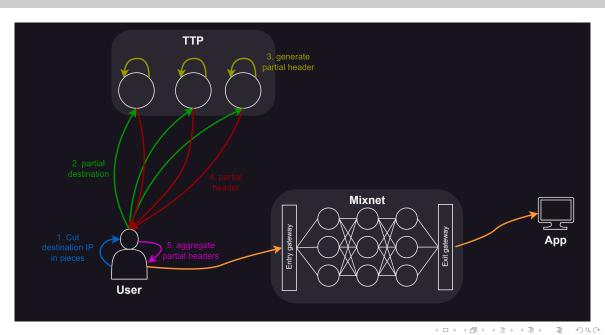
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#### Schema overview



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### Desired properties

#### Generic properties

- Correctness: schema works without adversary
- Compactness: Minimal overhead
- **Efficiency**: Easy and fast to compute (e.g. XOR, hash, exponentiation,...)

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- Depends on the mixnode
  - Forward / reply Undistinguishibility: Cannot distinguish forward from reply packet
  - Replay attack resistant:: Cannot reused previous packet

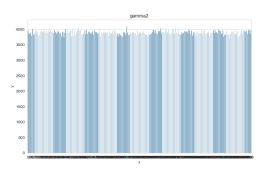
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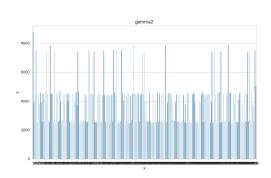
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  - Forward / reply Undistinguishibility: Cannot distinguish forward from reply packet
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- Depends on the header
  - Integrity: Maximum size path
  - Wrap-resistance: Unable to increase the intial path
  - Unlinkability: Cannot link incoming and outgoing packet from a mixnode

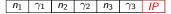
#### Original schema

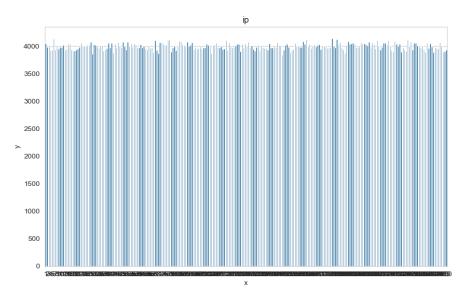


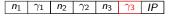
#### My schema

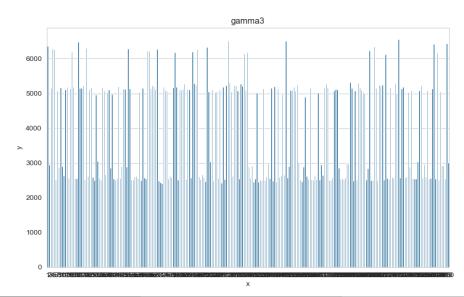


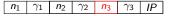
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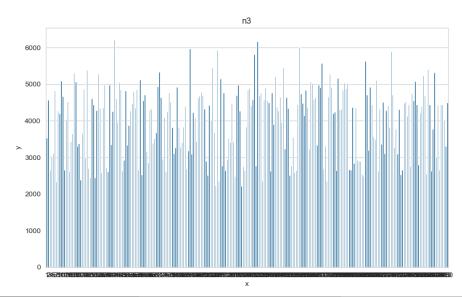


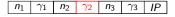


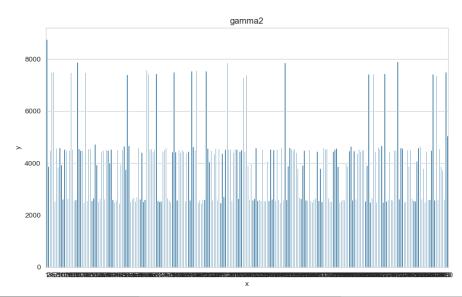


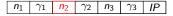


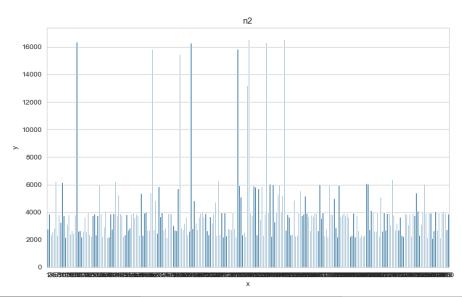


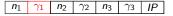


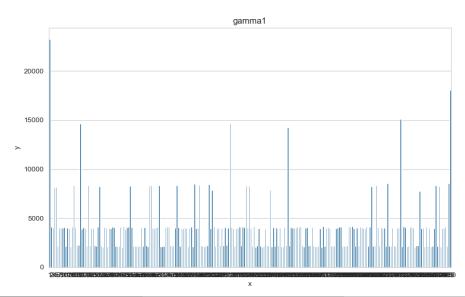




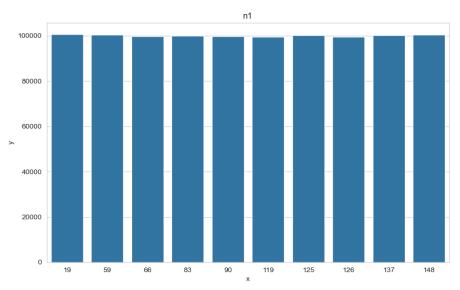






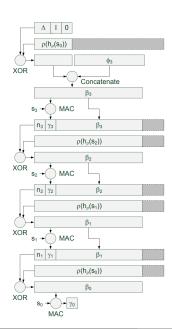


 $n_1$   $\gamma_1$   $n_2$   $\gamma_2$   $n_3$   $\gamma_3$  IP

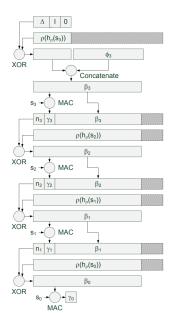


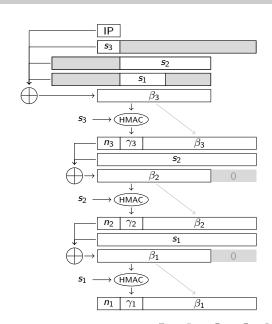
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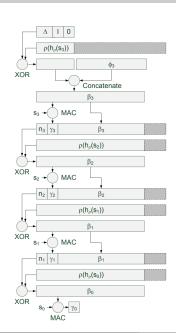
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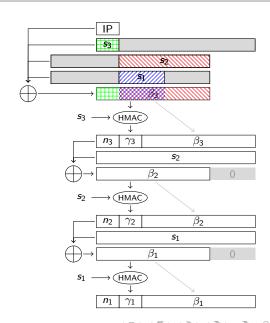


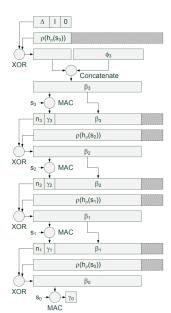
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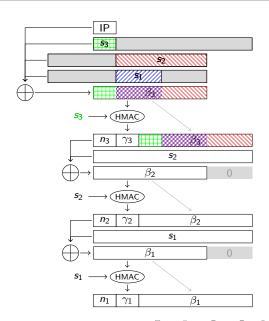


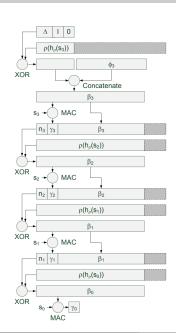


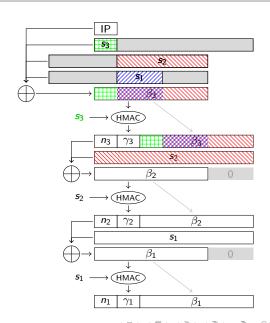


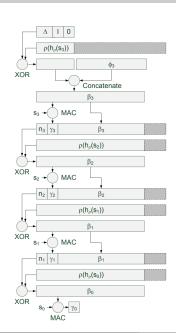


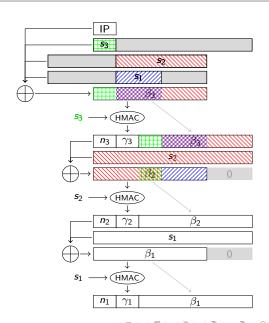


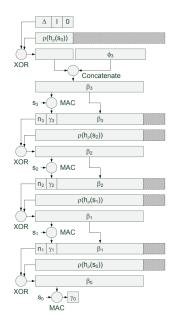


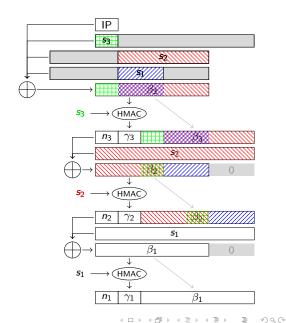


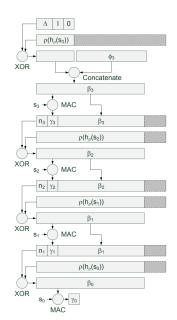


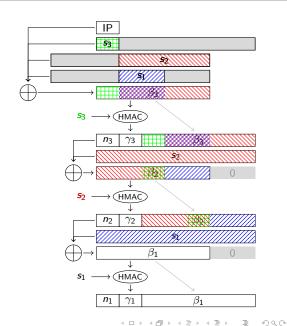


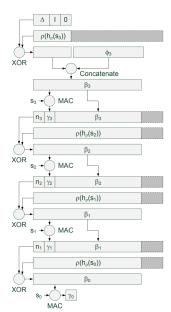


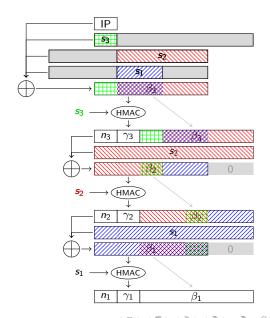


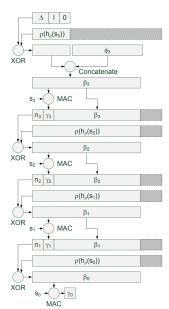


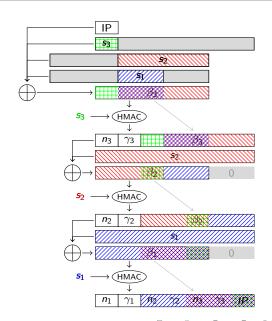


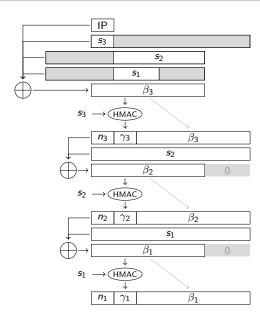




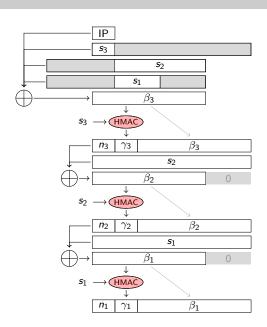






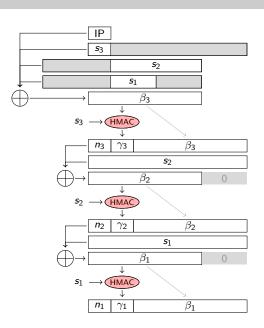


Main problem: Decentralizing a Hash?



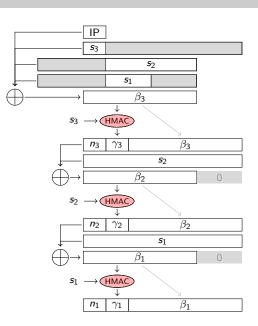
#### Main problem: Decentralizing a Hash?

 Need homomorphic properties to split computation and aggregate results.



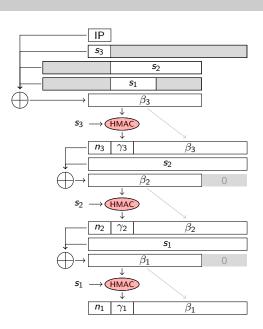
#### Main problem: Decentralizing a Hash?

- Need homomorphic properties to split computation and aggregate results.
- Secure homomorphic hash seems impractical (no promising solutions found).



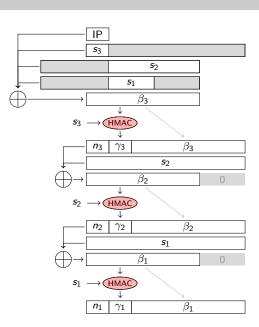
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  - ② ElGamal
  - 3 Paillier



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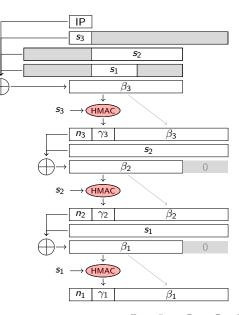
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$$\mathcal{E}(m_1) \cdot \mathcal{E}(m_2) = (g^{m_1} r_1^n) (g^{m_2} r_2^n) \mod n^2$$

$$= g^{m_1 + m_2} (r_1 r_2)^n \mod n^2$$

$$= \mathcal{E}(m_1 + m_2).$$

Problem: Mix of different operations... order matters!



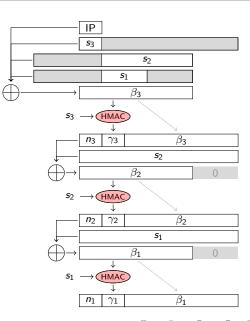
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$$\mathcal{E}(m_1) \cdot \mathcal{E}(m_2) = (g^{r_1}, m_1 \cdot h^{r_1})(g^{r_2}, m_2 \cdot h^{r_2})$$
$$= (g^{r_1 + r_2}, (m_1 \cdot m_2)h^{r_1 + r_2})$$
$$= \mathcal{E}(m_1 \cdot m_2)$$

3 Paillier

Limitation: Increase ciphertext size...



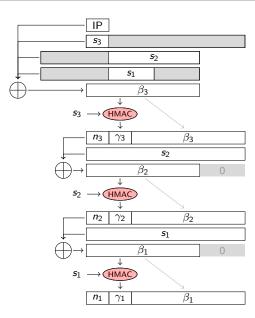
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$$\mathcal{E}(m_1) \cdot \mathcal{E}(m_2) = m_1^e m_2^e \mod n$$
  
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- **ElGamal**
- Paillier

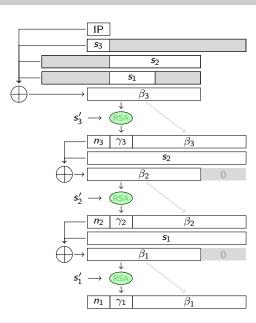


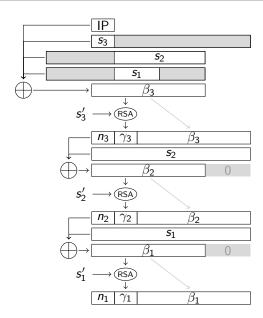
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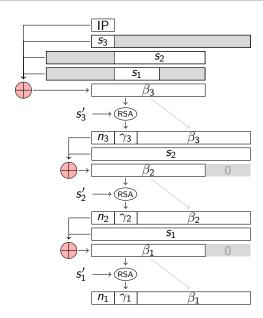
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#### Selected solution: RSA for integrity tag

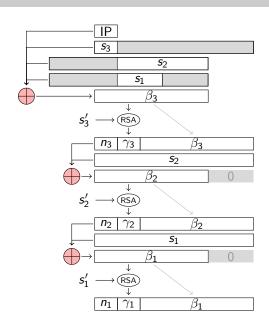
**NB:**  $s_i$  is different for each TTP but RSA required the same e... Thus, create a new shared secret  $s_i'$  common to all TTP







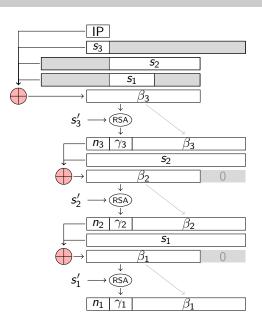
Since we use **RSA** for integrity tag  $\gamma_i$ 



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=  $(m_1 m_2)^e \mod n$   
=  $\mathcal{E}(m_1 \cdot m_2)$ 



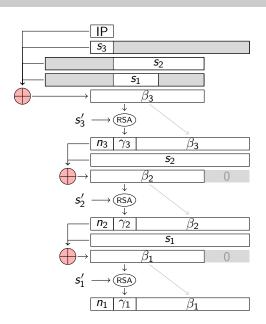
Since we use **RSA** for integrity tag  $\gamma_i$ 

$$\mathcal{E}(m_1) \cdot \mathcal{E}(m_2) = m_1^e m_2^e \mod n$$

$$= (m_1 m_2)^e \mod n$$

$$= \mathcal{E}(m_1 \cdot m_2)$$

Modular multiplication of integrity tags gives integrity tag of headers' modular product.



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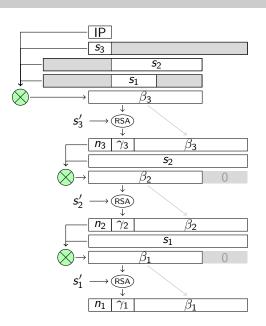
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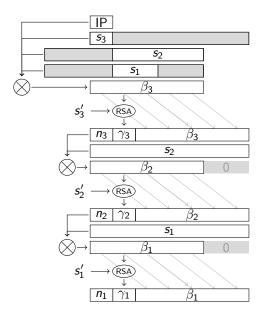
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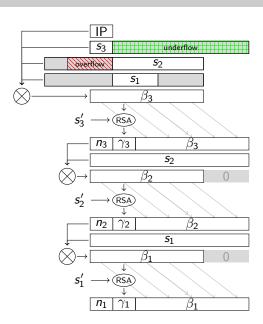
Thus, header elements must be combined via modular multiplication rather than XOR.



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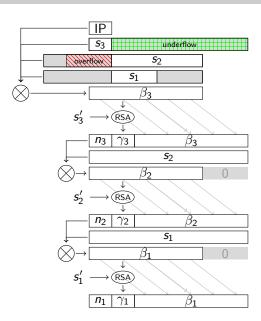


• Current Challenge: Overflow issues



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- Handling these issues is challenging.
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   (further research is required).



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- Current Challenge: Overflow issues
- Handling these issues is challenging.
   It may lead to information loss
   (further research is required).
- Proposed solution: Simplify by dividing data into small chunks and processing each chunk modulo its size.

