

# Let it Rot

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			. 1	- 1	•	1

```
vector<int>ans; // reverse ans in the end
     int vis[sz];
   | int dfs(int x) {
     vector<int>t:
       while (V[x].size()) {
           auto [to,id]=V[x].back();
           V[x].pop back();
          if (!vis[abs(id)])

    vis[abs(id)]=1,t.push_back(dfs(to)),ans.push_back(id);

      rep(i,1,(int)t.size()-1) if (t[i]!=x) ans.clear();
     return t.size()?t[0]:x;
   | }
   | int n,m;
    pii e[sz];
    int deg[sz], vv[sz];
18
     void clr() {
     | rep(i,1,n) V[i].clear(),deg[i]=vv[i]=0;
     | rep(i,1,m) vis[i]=0;
        ans.clear();
       n=m=<mark>0</mark>;
   | }
   void addedge(int x,int y) {
     | chkmax(n,x),chkmax(n,y); ++m;
        e[m]={x,y};
27
        if (directed) {
           V[x].push_back({y,m});
29
        ++deg[x],--deg[y],vv[x]=vv[y]=1;
     ∣ else {
           V[x].push_back({y,m});
           V[y].push_back({x,-m});
          ++deg[x],++deg[y],vv[x]=vv[y]=1;
37
    using vi=vector<int>;
    pair<vi,vi> work() {
   | | if (!m) return clr(),pair<vi,vi>{{1},{}};
      int S=1;
       rep(i,1,n) if (vv[i]) S=i;
        rep(i,1,n) if (deg[i]>0&&deg[i]%2=1) S=i;
        if ((int)ans.size()!=m) return clr(),pair<vi,vi>();
        reverse(ans.begin(),ans.end());
     vi ver,edge=ans;
  | | if (directed) {
          ver={e[ans[0]].fir};
        | for (auto t:ans) ver.push back(e[t].sec);
50
     | }
51
      ∣ else {
           ver={ans[0]>0?e[ans[0]].fir:e[-ans[0]].sec};
           for (auto t:ans) ver.push back(t>0?e[t].sec:e[-t].fir);
```

```
1.2 二分图匹配 | 最小边覆盖
 |A| // 匈牙利,左到右单向边,O(M|match|)
std::vector<int> edge[N];
3 | bool dfs(int x, std::vector<int> & vis, std::vector<int> & match) {
for(int y : edge[x]) if(!vis[y])
s | | if(vis[y] = 1, !match[y] || dfs(match[y], vis, match))
  | return 0;
std::vector<int> match(int nl, int nr) {
| | std::vector<int> vis(nr + 1), match(nr + 1), ret(nl + 1);
| | | | memset(vis.data(), 0, vis.size() << 2);
return ret[0] = 0, ret;
15 }
16 // 最小边覆盖
| 17 | std::pair<std::vector<int>, std::vector<int>> minedgecover(int nl, int nr) {
| std::vector<int> vis(nr + 1), match(nr + 1), ret(nl + 1);
|_{19}| | for(int i = 1;i <= nl;++i) if(dfs(i, vis, match))
20 | | memset(vis.data(), 0, vis.size() << 2);</pre>
|_{22}| | ret[0] = 0:
| for(int i = 1;i <= nl;++i) if(!ret[i]) dfs(i, vis, match);
| std::vector<int> le, ri;
| 25 | | for(int i = 1;i <= nl;++i) if(ret[i] && !vis[ret[i]]) le.push_back(i);
    for(int i = 1;i <= nr;++i) if(vis[i]) ri.push_back(i);</pre>
    return std::make_pair(le, ri);
|_{29} // 匈牙利,左到右单向边,bitset,O(n^2/w|match|)
| using set = std::bitset<N>;
set edge[N]:
bool dfs(int x, set & unvis, std::vector<int> & match) {
s3 | | for(set z = edge[x];;) {
  | | z &= unvis;
  int y = z._Find_first();
  | if(y = N) return 0;
  if(unvis.reset(y), !match[y] || dfs(match[y], unvis, match))
  | | return match[y] = x, 1;
39 | }
std::vector<int> match(int nl, int nr) {
| set unvis; unvis.set();
| std::vector<int> match(nr + 1), ret(nl + 1);
  | for(int i = 1;i <= nl;++i)
  | | if(dfs(i, unvis, match))
45
  for(int i = 1;i <= nr;++i) ret[match[i]] = i;</pre>
```

```
| return ret[0] = 0, ret;
49 }
_{50} // HK, 左到右单向边, O(M\sqrt{|match|})
std::vector<int> e[N];
std::vector<int> matchl, matchr, a, p;
std::vector<int> match(int nl, int nr) {
     matchl.assign(nl + 1, 0), matchr.assign(nr + 1, 0);
    for(;;) {
55
       a.assign(nl + 1, 0), p.assign(nl + 1, 0);
        static std::queue<int> Q;
        for(int i = 1;i <= nl;++i)</pre>
         | if(!matchl[i]) a[i] = p[i] = i, Q.push(i);
       int succ = 0;
     | for(;Q.size();) {
           int x = Q.front(); Q.pop();
           if(matchl[a[x]]) continue;
           for(int y : e[x]) {
           | if(!matchr[y]) {
               \mid for(succ = 1;y;x = p[x])
                 | matchr[y] = x, std::swap(matchl[x], y);
              | break;
         | | if(!p[matchr[y]])
           | | Q.push(y = matchr[y]), p[y] = x, a[y] = a[x];
     if(!succ) break;
75
   | return matchl;
77 } // matchl 是左边每个点匹配的右边点编号
78 std::pair<std::vector<int>, std::vector<int>> minedgecover(int nl, int nr) {
    match(nl, nr);
     std::vector<int> l, r;
   | for(int i = 1;i <= nl;++i) if(!a[i]) l.push_back(i);
     for(int i = 1;i <= nr;++i) if(a[matchr[i]]) r.push back(i);</pre>
  | return {l, r};
83
84 }
```

## 1.3 网络最大流 | dinic

```
// S 编号最小, T 最大, 或者改一下清空
struct Dinic {
    | struct T {
    | int to, nxt, v;
    | } e[N << 3];
    | int h[N], head[N], num = 1;
    | void link(int x, int y, int v) {
    | | e[++num] = {y, h[x], v}, h[x] = num;
    | | e[++num] = {x, h[y], 0}, h[y] = num; // !!!
    | }
    | int dis[N];
    | bool bfs(int s, int t) {
    | | std::queue<int> Q;
```

```
| for(int i = s;i <= t;++i) dis[i] = -1, head[i] = h[i]; //如果编号不是
          → [S,T], 只要改这里
  | | for(Q.push(t), dis[t] = 0;!Q.empty();) {
16 | | int x = Q.front(); Q.pop();
        for(int i = h[x]; i; i = e[i].nxt) if(e[i ^ 1].v && dis[e[i].to] < 0)
           | dis[e[i].to] = dis[x] + 1, Q.push(e[i].to);
20 | | }
| | | | return dis[s] >= 0;
| | int dfs(int s, int t, int lim) {
  | if(s = t || !lim) return lim;
  | | int ans = 0, mn;
  | | for(int & i = head[s];i;i = e[i].nxt) {
  | \cdot | if (dis[e[i].to] + 1 = dis[s] & (mn = dfs(e[i].to, t, std::min(lim,
            \rightarrow e[i].v)))) {
           | e[i].v -= mn, e[i ^ 1].v += mn;
  30 | | | | if(!lim) break;
31 | | | }
32 | | }
33 | return ans;
  | int flow(int s, int t) {
  | | int ans = 0;
  | for(;bfs(s, t);) ans += dfs(s, t, 1e9);
38 | return ans;
39 | }
40 } G;
```

## 1.4 最小费用流

```
ı // S 编号最小, T 最大, 或者改一下清空
2 namespace mcmf {
  using pr = std::pair<ll, int>;
  \mid const int N = 10005, M = 1e6 + 10;
    struct edge {
  | | int to, nxt, v, f;
  | } e[M << 1];
    int h[N], num = 1;
     void link(int x, int y, int v, int f) {
  | e[++num] = \{y, h[x], v, f\}, h[x] = num;
  | | e[++num] = \{x, h[y], 0, -f\}, h[y] = num;
12 | }
14 | int vis[N], fr[N];
| 15 | | bool spfa(int s, int t) {
  | | std::queue<int> Q;
| | | | std::fill(d + s, d + t + 1, 1e18); // CHECK
| | | for(d[s] = 0, Q.push(s);!Q.empty();) {
  | | int x = Q.front(); Q.pop(); vis[x] = 0;
19
        | for(int i = h[x];i;i = e[i].nxt)
           | if(e[i].v && d[e[i].to] > d[x] + e[i].f) {
```

```
d[e[i].to] = d[x] + e[i].f;
              | fr[e[i].to] = i;
       | | | if(!vis[e[i].to]) vis[e[i].to] = 1, Q.push(e[i].to);
     | return d[t] < 1e17;
    bool dijkstra(int s, int t) { // 正常题目不需要 dijk
     | std::priority_queue<pr, std::vector<pr>, std::greater<pr>> Q;
     | for(int i = s; i <= t; ++i) dis[i] = d[i], d[i] = 1e18, vis[i] = fr[i] =
         → 0; // CHECK
       for(Q.emplace(d[s] = 0, s);!Q.empty();) {
          int x = Q.top().second; Q.pop();
          if(vis[x]) continue;
          vis[x] = 1;
          for(int i = h[x];i;i = e[i].nxt) {
          | const ll v = e[i].f + dis[x] - dis[e[i].to];
             if(e[i].v \& d[e[i].to] > d[x] + v) {
             | fr[e[i].to] = i;
             \mid Q.emplace(d[e[i].to] = d[x] + v, e[i].to);
       for(int i = s;i <= t;++i) d[i] += dis[i]; // CHECK</pre>
       return d[t] < 1e17;
    std::pair<ll, ll> EK(int s, int t) {
     | spfa(s, t); // 如果初始有负权且要 dijk
  | for(;dijkstra(s, t);) { // 正常可以用 spfa
          ll fl = 1e18;
        for(int i = fr[t];i;i = fr[e[i ^ 1].to]) fl = std::min<ll>(e[i].v,
        | for(int i = fr[t];i;i = fr[e[i ^ 1].to]) e[i].v -= fl, e[i ^ 1].v +=
            ن fl;
        | f += fl, c += fl * d[t];
     return std::make pair(f, c);
57
58 }
59 // in flow problems with lower bounds (or with negative cycles), flow the
    60 // after the first round, revert the auxiliary edges
 1.5 二分图最大权匹配 | KM
```

```
namespace KM {
    | int nl,nr;
    | ll e[sz][sz];
    | ll lw[sz],rw[sz];
    | int lpr[sz],rpr[sz];
    | int vis[sz],fa[sz];
    | ll mnw[sz];
    | void work(int x) {
    | int xx=x;
}
```

```
rep(i,1,nr) vis[i]=0,mnw[i]=1e18;
        while (233) {
         | rep(i,1,nr) if (!vis[i]\delta\delta chkmin(mnw[i],lw[x]+rw[i]-e[x][i]))
              \hookrightarrow fa[i]=x;
           ll mn=1e18; int y=-1;
13
           rep(i,1,nr) if (!vis[i]&&chkmin(mn,mnw[i])) y=i;
           lw[xx]-=mn; rep(i,1,nr) if (vis[i]) rw[i]+=mn,lw[rpr[i]]-=mn; else

    mnw[i]-=mn;
   | | if (rpr[y]) x=rpr[y],vis[y]=1; else { while (y)

    rpr[y]=fa[y],swap(y,lpr[fa[y]]); return; }

17
18 | }
   void init(int nl,int nr) {
   | | assert(nl<=nr);</pre>
        KM::nl=nl,KM::nr=nr;
        rep(i,1,nl) lw[i]=-1e18;
  | | rep(i,1,nl) rep(j,1,nr) e[i][j]=0; // or -1e18
24
25
   void clr() {
        rep(i,1,nl) lpr[i]=lw[i]=0;
        rep(i,1,nr) rpr[i]=rw[i]=vis[i]=fa[i]=mnw[i]=0;
     | rep(i,1,nl) rep(j,1,nr) e[i][j]=0;
29
     void addedge(int x,int y,ll w){chkmax(e[x][y],w),chkmax(lw[x],w);}
    ll work() {
        rep(i,1,nl) work(i);
        ll tot=0;
        rep(i,1,nl) tot+=e[i][lpr[i]];
35
        return tot;
  | }
36
```

## 1.6 一般图最大匹配 | 带花树

```
namespace blossom {
  vector<int>V[sz];
   | int f[sz];
     int n,match[sz];
     int getfa(int x){return f[x]=x?x:f[x]=getfa(f[x]);}
     void link(int x,int y){V[x].push_back(y),V[y].push_back(x);}
     int pre[sz].mk[sz];
     int vis[sz],T;
     queue<int>q;
     int LCA(int x,int y) {
     | T++;
     \mid for (;; x=pre[match[x]], swap(x,y))
  | | if (vis[x=getfa(x)]=T) return x;
14 | | else vis[x]=x?T:0;
15
| | | void flower(int x,int y,int z) {
| | | | | while (getfa(x)!=z) {
| | | | | pre[x]=y; y=match[x]; f[x]=f[y]=z; x=pre[y];
        | if (mk[y]=2) q.push(y),mk[y]=1;
```

```
void aug(int s){
      for (int i=1;i<=n;i++) pre[i]=mk[i]=vis[i]=0,f[i]=i;</pre>
       mk[s]=1; q.push(s);
       while (q.size()) {
          int x=q.front(); q.pop();
          for (auto v:V[x]) {
             int y=v,z;
             if (mk[y]=2) continue;
             if (mk[y]=1) z=LCA(x,y),flower(x,y,z),flower(y,x,z);
             else if (!match[y]) {
              for (pre[y]=x;y;) x=pre[y],match[y]=x,swap(y,match[x]);
              return;
           else pre[y]=x,mk[y]=2,q.push(match[y]),mk[match[y]]=1;
  | }
  | int work() {
     | rep(i,1,n) if (!match[i]) aug(i);
     | int res=0;
     | rep(i,1,n) res+=match[i]>i;
     | return res;
45
```

#### 1.7 最小树形图

抄罗大的,返回值是边的集合,如果没有最小树形图会返回空集(注意 n=1),可以修改建图。

```
namespace DMST {
  | struct edge {
  | | int u, v, id; ll w;
    | bool operator < (const edge & y) const {
       return w < y.w;
  int ls[M], rs[M], size[M], cc; ll tag[M];
    int fs[N], fw[N], rt[N];
    void put(int x, ll v) {
     | if(x) val[x].w += v, tag[x] += v;
  void pushdown(int x) {
       put(ls[x], tag[x]);
       put(rs[x], tag[x]);
       tag[x] = 0;
17
  int merge(int x, int y) {
     | if(!x || !y) return x | y;
       if(val[y] < val[x]) std::swap(x, y);</pre>
       pushdown(x), rs[x] = merge(rs[x], y);
       if(size[rs[x]] > size[ls[x]]) {
          std::swap(ls[x], rs[x]);
```

```
size[x] += size[y];
  | | return x;
26
27 | }
  | void ins(int & x, const edge & z) {
28
     | val[++cc] = z, size[cc] = 1;
     \mid x = merge(x, cc);
   | void pop(int \delta x) { x = merge(ls[x], rs[x]); }
     edge top(int x) { return val[x]; }
     int find(int x, int * anc) {
     return anc[x] = x ? x : anc[x] = find(anc[x], anc);
     void link(int u, int v, int w, int id) {
     | ins(rt[v], {u, v, id, w});
     int pa[N * 2], tval[N * 2], up[N * 2], end_edge[M], cmt, baned[M];
     std::vector<int> solve(int r) {
   | | std::queue<int> roots;
   | | for(int i = 1;i <= n;++i) {
   | | | fs[i] = fw[i] = i, tval[i] = ++ cmt;
        if(i != r) roots.push(i);
45
46
       std::vector<edge> H;
       std::vector<int> ret;
   | | for(;!roots.empty();) {
        int k = roots.front(); roots.pop();
          if(!rt[k]) return ret;
           edge e = top(rt[k]); pop(rt[k]);
           int i = e.u, j = e.v;
           if(find(i, fs) = k) roots.push(k);
           else {
              H.push_back(e); end_edge[e.id] = tval[k];
              if(find(i, fw) != find(j, fw)) {
               | fw[find(j, fw)] = i;
                ent[k] = e;
            pa[tval[k]] = ++ cmt, up[tval[k]] = e.id;
                put(rt[k], -e.w);
                for(;(e = ent[find(e.u, fs)]).u;) {
                   int p = find(e.v, fs);
                   pa[tval[p]] = cmt;
                   up[tval[p]] = e.id;
                   put(rt[p], -e.w);
                   rt[k] = merge(rt[k], rt[p]);
                   fs[p] = k;
                tval[k] = cmt;
                roots.push(k);
  74
  | | }
75
     | reverse(H.begin(), H.end());
       for(edge i : H) if(!baned[i.id]) {
```

## 1.8 缩点 | kasaraju

时间复杂度  $O(\frac{n^2}{N})$ ,可以对于边修改不多的图快速计算。

```
using set = std::bitset<N>;
2 // re 是反向边,需要连好
set e[N], re[N], vis;
std::vector<int> sta;
void dfs0(int x, set * e) {
  vis.reset(x);
     for(;;) {
     | int go = (e[x] & vis)._Find_first();
  | if(go = N) break;
     | dfs0(go, e);
     sta.push_back(x);
std::vector<std::vector<int>> solve() {
     for(int i = 1;i <= n;++i) if(vis.test(i)) dfs0(i, e);</pre>
    vis.set();
    auto s = sta;
    std::vector<std::vector<int>> ret;
    for(int i = n - 1; i >= 0; --i) if(vis.test(s[i])) {
     | sta.clear(), dfs0(s[i], re), ret.push_back(sta);
     }
22
    return ret;
23
```

## 1.9 缩点 | Tarjan

```
int dfn[sz],low[sz],cc;
stack<int>S; int in[sz];
int bel[sz],T;
void dfs(int x) {
    | dfn[x]=low[x]=++cc; S.push(x),in[x]=1;
    | for (auto v:V[x]) {
    | if (!dfn[v]) dfs(v,x),chkmin(low[x],low[v]);
    | else if (in[v]) chkmin(low[x],dfn[v]);
    | }
    | if (dfn[x]=low[x]) {
    | int y; ++T;
    | do y=S.top(),S.pop(),in[y]=0,bel[y]=T; while (y!=x);
    | }
}
```

## 1.10 缩点 | 点双

## 1.11 缩点 | 边双

#### 1.12 仙人掌

```
vector<int>V2[sz], V[sz]; // V2: cactus edges; V: reconstructed tree edges
  int m; // set to n before dfs
  void dfs(int x,int f) {
     static int mark[sz],fa[sz],vis[sz],dep[sz];
  fa[x]=f; vis[x]=1;dep[x]=dep[f]+1;
  | for (auto v:V2[x]) if (v!=f) {
     | if (!vis[v]) dfs(v,x);
       else if (dep[v]<dep[x]) {</pre>
           ++m;
           V[v].push_back(m);
          for (int y=x;y!=v;y=fa[y]) V[m].push_back(y),mark[y]=1;
12
13
  | if (!mark[x]) {
       V[fa[x]].push_back(m),V[m].push_back(x);
17
```

int semi[sz];

```
1.13 2-Sat

rep(i,1,n) if (bel[i<<1]=bel[i<<1|1]) return puts("IMPOSSIBLE"),0;
puts("POSSIBLE");
rep(i,1,n) printf("%d ",bel[i<<1]>bel[i<<1|1]);
1.14 支配树
```

```
1.14 支配树
namespace BuildTree {
  int idom[sz]:
     vector<int>V[sz],ANS[sz]; // ANS: final tree
     int deg[sz];
    int fa[sz][25],dep[sz];
    int lca(int x,int y) {
     if (dep[x]<dep[y]) swap(x,y);</pre>
        drep(i,20,0)
          if (fa[x][i]&&dep[fa[x][i]]>=dep[y])
           | x=fa[x][i]:
        if (x=y) return x;
       drep(i,20,0)
          if (fa[x][i]!=fa[y][i])
         | | x=fa[x][i],y=fa[y][i];
       return fa[x][0];
     void work() {
       queue<int>q;q.push(1);
     while (!q.empty()) {
          int x=q.front();q.pop();
           ANS[idom[x]].push\_back(x);fa[x][0]=idom[x];dep[x]=dep[idom[x]]+1;
           rep(i,1,20) fa[x][i]=fa[fa[x][i-1]][i-1];
           for (int v:V[x]) {
              --deg[v];if (!deg[v]) q.push(v);
             if (!idom[v]) idom[v]=x;
              else idom[v]=lca(idom[v],x);
29
30 }
namespace BuildDAG {
    vector<int>V[sz],rV[sz];
  int dfn[sz],id[sz],anc[sz],cnt;
     void dfs(int x) {
     | id[dfn[x]=++cnt]=x;
     for (int v:V[x]) if (!dfn[v])
          BuildTree::V[x].push_back(v), BuildTree::deg[v]++, anc[v] = x,
             \hookrightarrow dfs(v):
  | }
   int fa[sz].mn[sz];
  int find(int x) {
     | if (x=fa[x]) return x;
        int tmp=fa[x];fa[x]=find(fa[x]);
       chkmin(mn[x],mn[tmp]);
     return fa[x];
```

```
47 | | void work() {
        dfs(1);
        rep(i,1,n) fa[i]=i,mn[i]=1e9,semi[i]=i;
        drep(w,n,2) {
        int x=id[w];int cur=1e9;
           if (w>cnt) continue;
           for (int v:rV[x]) {
            | if (!dfn[v]) continue;
           if (dfn[v]<dfn[x]) chkmin(cur,dfn[v]);</pre>
           | else find(v),chkmin(cur,mn[v]);
           semi[x]=id[cur];mn[x]=cur;fa[x]=anc[x];
58
           BuildTree::V[semi[x]].push back(x);BuildTree::deg[x]++;
  1 }
61
     void link(int x,int y){V[x].push_back(y),rV[y].push_back(x);}
```

#### 1.15 三/四元环

```
static int id[sz].rnk[sz];
  rep(i,1,n) id[i]=i;
 | sort(id+1,id+n+1,[](int x,int y){return pii{deg[x],x}<pii{deg[y],y};});
 | rep(i,1,n) rnk[id[i]]=i;
  rep(i,1,n) for (auto v:V[i]) if (rnk[v]>rnk[i]) V2[i].push_back(v);
6 int ans3=0; // 3-cycle
7 rep(i,1,n) {
s | static int vis[sz];
     for (auto v:V2[i]) vis[v]=1;
10 | for (auto v1:V2[i]) for (auto v2:V2[v1]) if (vis[v2]) ++ans3; // (i,v1,v2)
  | for (auto v:V2[i]) vis[v]=0;
13 | ll ans4=0: // 4-cvcle
14 rep(i,1,n) {
15 | static int vis[sz];
16 | for (auto v1:V[i]) for (auto v2:V2[v1]) if (rnk[v2]>rnk[i])

    ans4+=vis[v2],vis[v2]++;
     for (auto v1:V[i]) for (auto v2:V2[v1]) vis[v2]=0;
```

#### 1.16 双极定向

```
vector<int>G[sz];
namespace bipolar_orientation {
    int dfn[sz],low[sz],cc,p[sz],inv[sz],topo[sz];
    | bool flg,sgn[sz];
    | void dfs(int x,int fa,int s,int t) {
    | | dfn[x]=low[x]=++cc; inv[cc]=x,p[x]=fa;
    | | if (x=s) dfs(t,x,s,t);
    | | for (int y:G[x]) {
    | | | if (x=s&by=t) continue;
    | | | | dfs(y,x,s,t);
    | | | | dfs(y,x,s,t);
    | | | | | chkmin(low[x],low[y]);
    | | | | if (x=s|low[y]>=dfn[x]) flg=1;
```

```
else if (dfn[y]<dfn[x]&&y!=fa) chkmin(low[x],dfn[y]);
int check(int s,int t,int n) { // return topo
 | if (n=1) return topo[1]=1,1;
 | if (s=t) return 0:
   | cc=flg=0; dfs(s,s,s,t);
    if (flg) return 0;
    sgn[s]=0;
    static int pre[sz], suf[sz];
    suf[0]=s,pre[s]=0,suf[s]=t;
    pre[t]=s,suf[t]=n+1,pre[n+1]=t;
    rep(i,3,n) {
     | int v=inv[i];
     | if (!sgn[inv[low[v]]]) {
     | | int P=pre[p[v]];
       | pre[v]=P,suf[v]=p[v];
       suf[P]=pre[p[v]]=v;
     ∣ else {
     | | int S=suf[p[v]];
      pre[v]=p[v],suf[v]=S;
       | suf[p[v]]=pre[S]=v;
       sgn[p[v]]=!sgn[inv[low[v]]];
  for (int x=s,cnt=0;x!=n+1;x=suf[x]) topo[++cnt]=x;
  return 1;
void clr(int n) {
  | rep(i,1,n) dfn[i]=low[i]=p[i]=inv[i]=topo[i]=sgn[i]=0,G[i].clear();
```

## 1.17 Tree And Graph

#### 1.17.1 树的计数 Prufer序列

树和其prufer编码——对应, 一颗 n 个点的树, 其prufer编码长度为 n-2, 且度数为  $d_i$  的点在prufer 编码中出现  $d_i-1$  次.

由树得到序列: 总共需要 n-2 步,第 i 步在当前的树中寻找具有最小标号的叶子节点,将与其相连的点的标号设为Prufer序列的第 i 个元素  $p_i$ ,并将此叶子节点从树中删除,直到最后得到一个长度为 n-2 的Prufer 序列和一个只有两个节点的树。

由序列得到树: 先将所有点的度赋初值为 1, 然后加上它的编号在Prufer序列中出现的次数, 得到每个点的度; 执行 n-2 步, 第 i 步选取具有最小标号的度为 1 的点 u 与 v =  $p_i$  相连, 得到树中的一条边, 并将 u 和 v 的度减一. 最后再把剩下的两个度为 1 的点连边, 加入到树中.

相关结论: n 个点完全图,每个点度数依次为  $d_1,d_2,...,d_n$ ,这样生成树的棵树为:  $\frac{(n-2)!}{(d_1-1)!(d_2-1)!...(d_n-1)!}$ . 左边有  $n_1$  个点,右边有  $n_2$  个点的完全二分图的生成树棵树为  $n_1^{n_2-1} \times n_2^{n_1-1}$ . m 个连通块,每个连通块有  $c_i$  个点,把他们全部连通的生成树方案数:  $(\sum c_i)^{m-2} \prod c_i$ 

#### 1.17.2 有根树的计数

首先, 令  $S_{n,j}=\sum_{1\leq j\leq n/j}$ ; 于是 n+1 个结点的有根树的总数为  $a_{n+1}=\frac{\sum_{j=1}^n j a_j S_{n-j}}{n}$ . 注:  $a_1=1,a_2=1,a_3=2,a_4=4,a_5=9,a_6=20,a_9=286,a_{11}=1842$ .

#### 1.17.3 无根树的计数

n 是奇数时, 有  $a_n - \sum_{i=1}^{n/2} a_i a_{n-i}$  种不同的无根树.

n 时偶数时,有  $a_n - \sum_{i=1}^{n/2} a_i a_{n-i} + \frac{1}{2} a_{n/2} (a_{n/2} + 1)$  种不同的无根树

#### 1.17.4 生成树计数 Kirchhoff's Matrix-Tree Thoerem

Kirchhoff Matrix T = Deg - A, Deg 是度数对角阵, A 是邻接矩阵. 无向图度数矩阵是每个点度数; 有向图度数矩阵是每个点入度.

邻接矩阵 A[u][v] 表示  $u \to v$  边个数, 重边按照边数计算, 自环不计入度数.

无向图生成树计数: c = |K| 的任意1个 n1 阶主子式 |

有向图外向树计数: c = | 去掉根所在的那阶得到的主子式 |

#### 1.17.5 有向图欧拉回路计数 BEST Thoerem

$$\operatorname{ec}(G) = t_{w}(G) \prod_{v \in V} (\operatorname{deg}(v) - 1)!$$

其中  $\deg$  为入度 (欧拉图中等于出度),  $t_w(G)$  为以 w 为根的外向树的个数. 相关计算参考生成树计数. 欧拉连通图中任意两点外向树个数相同:  $\mathbf{t}_v(G) = \mathbf{t}_w(G)$ .

以 1 结尾的欧拉路径计数就是把 deg 视为出度,把 deg(1) 的贡献改为 deg(1)!.

#### 1.17.6 Tutte Matrix

Tutte matrix A of a graph G = (V, E):

$$A_{ij} = \begin{cases} x_{ij} & \text{if } (i,j) \in E \text{ and } i < j \\ -x_{ij} & \text{if } (i,j) \in E \text{ and } i > j \\ 0 & \text{otherwise} \end{cases}$$

where  $x_{ij}$  are indeterminates. The determinant of this skew-symmetric matrix is then a polynomial (in the variables  $x_{ij}$ , i < j): this coincides with the square of the pfaffian of the matrix A and is non-zero (as a polynomial) if and only if a perfect matching exists.

#### 1.17.7 Edmonds Matrix

Edmonds matrix A of a balanced (|U| = |V|) bipartite graph G = (U, V, E):

$$A_{ij} = \begin{cases} x_{ij} & (u_i, v_j) \in E \\ 0 & (u_i, v_j) \notin E \end{cases}$$

where the  $x_{ij}$  are indeterminates. G 有完美匹配当且仅当关于  $x_{ij}$  的多项式  $det(A_{ij})$  不恒为 0. 完美匹配的个数等于多项式中单项式的个数.

#### 1.1/8 拟阵交

```
// max size, minimum weight
namespace MatroidIntersection {
   int K;
   ll W[sz]; // weight
   int in[sz]; // ans
   namespace Check { // implementation needed
```

```
// recommend writing checker here
          void init() {}
          // return {-1} if no cycle; return cycle otherwise.
          vector<int> cycleA(int x) {}
          vector<int> cycleB(int x) {}
          // not necessary
          void check() {init();}
      bool work() { // try augment
          using pli=pair<ll,int>;
          static vector<int> V[sz];
          static pli dis[sz]:
18
          static int fr[sz];
          Check::init();
          rep(i,1,K) V[i].clear();
          vector<int>A,B;
          rep(i,1,K) if (!in[i]) {
23
              auto cyca=Check::cycleA(i);
              if (cyca.size()=1u&&cyca[0]=-1) A.push_back(i);
25
              else for (auto y:cyca) V[y].push_back(i);
              auto cycb=Check::cycleB(i);
              if (cycb.size()=1u&&cycb[0]=-1) B.push_back(i);
              else for (auto y:cycb) V[i].push_back(y);
          rep(i,1,K) dis[i]={ll(1e18),K+1},fr[i]=0;
          priority_queue<pair<pli,int>, vector<pair<pli,int>>,
32

    greater<pair<pli,int>>>q;
          for (auto x:A) dis[x]=\{W[x],0\},q.push(\{dis[x],x\});
33
          while (!q.empty()) {
              auto [ww,x]=q.top(); q.pop();
              if (dis[x]!=ww) continue;
              for (auto v:V[x])
                  if (chkmin(dis[v],{dis[x].fir+W[v],dis[x].sec+1}))
                      q.push(\{dis[v],v\}),fr[v]=x;
          pli mn={ll(1e18),K+1}; int mnp=-1;
          for (auto x:B) if (chkmin(mn,dis[x])) mnp=x;
          if (mnp=-1) return 0;
          for (int x=mnp;x;x=fr[x]) in[x]^=1;
          Check::check();
          return 1;
     void clr() {
          rep(i,1,K) in[i]=0;
50
```

```
| | std::swap(x, y), std::swap(a, b);
| | a -= x / y * b, x %= y;
| | return {x, a};
| // q = x/a (mod p), x \le A, |a| 取到最小值
```

#### 2.2 扩展欧几里得

```
// result : -b < x < b AND -a < y <= a when a,b != 0
void exgcd(ll a, ll b, ll & x, ll & y) {
  | if(!b) return x = 1, y = 0, void();
  | exgcd(b, a % b, y, x), y -= a / b * x;
}
```

#### 2.3 万能欧几里得

```
」 // 万欧
2 // 前提 : r < q, r >= q 先提几个 U 出来再用
₃ / // 使用: Y * q <= X * p + r, 斜率 p/q, U表示向上, R表示到达一个顶点, 先一些 U 再一个 R
4 template<class T>
5 T power(T a, ll k) {
。 // 有效率需求可以改为半群乘法
  | if(!k) return T();
8 | T res = a;
9 | for(--k;k;) {
  \mid if(k & 1) res = res + a;
| | | | | if(k >>= 1) a = a + a;
12 | }
13 | return res;
15 template<class T>
16 T solve(ll p, ll q, ll r, ll l, T U, T R) {
| | | return solve(p % q, q, r, l, U, power(U, p / q) + R);
| if (!m) return power(R, l);
| ll cnt = l - ((__int128)q * m - r - 1) / p;
|z| | return power(R, (q - r - 1) / p) + U + solve(q, p, (q - r - 1) % p, m - 1,
      \hookrightarrow R, U) + power(R, cnt);
```

#### 2.4 直线下点数|欧几里得

 $n < 2^{32}, 1 \le m < 2^{32}$ 

$$result = \sum_{i=0}^{n-1} \lfloor \frac{ai+b}{m} \rfloor \pmod{2^{63}}$$

## 2 数论

#### 2.1 取模还原分数

```
u64 floor_sum(u64 n, u64 m, u64 a, u64 b) {
    | u64 ans = 0;
    | for(;;) {
    | | if(a >= m) ans += n * (n - 1) / 2 * (a / m), a %= m;
    | | if(b >= m) ans += n * (b / m), b %= m;
    | | u64 ymax = a * n + b; // use u128 if it's big
```

#### 2.5 Stern-Brocot Tree 二分

```
using cp = std::complex<ll>;
cp fracBS(ll n, ll m, auto f) {
  | bool dir = 1, A = 1, B = 1;
    cp lo(0, 1), hi(1, 1); // hi can be (1, 0), f(hi) must be true
  | if (f(lo)) return lo;
  | while(A || B) {
     | ll adv = 0, s = 1;
     | for (int x = 0;s;(s *= 2) >>= x) {
          adv += s;
          cp mid = lo * adv + hi;
          if (mid.real() > n || mid.imag() > m || dir = !f(mid)) {
          | adv -= s, x = 2;
       hi += lo * adv, dir = !dir;
       swap(lo, hi);
       A = B, B = adv;
  | return dir ? hi : lo;
20 } // 返回值是最小的使得 f 为真的
   // 另外一个是最大的使得 f 为假的
```

#### 2.6 扩展中国剩余定理

```
ll exCRT(ll a1, ll p1, ll a2, ll p2) {
    | ll a, b, gcd = std::gcd(p1, p2);
    | if((a1 - a2) % gcd)
    | | return -1;
    | exgcd(p1, p2, a, b);
    | ll k = i128((a2 - a1) % p2 + p2) * (a + p2) % p2;
    | return p1 / gcd * k + a1;
}
```

#### 2.7 Miller-Rabin

```
using f64 = long double;
ll p;
f64 invp;
void setmod(ll x) {
    | p = x, invp = (f64) 1 / x;
}
ll mul(ll a, ll b) {
    | ll z = a * invp * b + 0.5;
    | ll res = a * b - z * p;
    | return res + (res >> 63 & p);
}
```

```
|_{12}| ll pow(ll a, ll x, ll res = 1) {
|_{13}| | for(;x;x >>= 1, a = mul(a, a))
|_{14}| | | if(x & 1) res = mul(res, a);
15 | return res;
16 }
| bool checkprime(ll p) {
|_{18}| | if(p = 1) return 0;
   | setmod(p);
   \mid ll d = __builtin_ctzll(p - 1), s = (p - 1) >> d;
   | for(ll a : {2, 3, 5, 7, 11, 13, 82, 373}) {
   | | if(a \% p = 0)
   | | | continue;
23
   \mid  \mid ll x = pow(a, s), y;
   | | for(int i = 0; i < d; ++i, x = y) {
   | | | if(y = 1 \delta \delta x != 1 \delta \delta x != p - 1)
   | | | return 0;
   | | if(x != 1) return 0;
32
   return 1;
33 }
```

## 2.8 Pollard-rho

```
ı|ll rho(ll n) {
2 | if(!(n & 1)) return 2;
  static std::mt19937 64 gen((size t)"hehezhou");
  | ll x = 0, y = 0, prod = 1;
  | auto f = [8](ll o) { return mul(o, o) + 1; };
  | for(int t = 30, z = 0;t % 64 || std::gcd(prod, n) = 1;++t) {
  | if (x = y) x = ++ z, y = f(x);
  if(ll q = mul(prod, x + n - y)) prod = q;
|x| = f(x), y = f(f(y));
  | return std::gcd(prod, n);
std::vector<ll> factor(ll x) {
std::vector<ll> res;
|a| | auto f = [\delta](auto f, ll x) {
|x| | if(x = 1) return;
if(checkprime(x)) return res.push_back(x);
| | | | | f(f, y), f(f, x / y); 
  | f(f, x), sort(res.begin(), res.end());
  | return res;
23
```

#### 3 Math

## 3.1 拉格朗日反演

$$G(F(x)) = H(x) \Rightarrow [x^n]G(x) = \frac{1}{n}[u^{n-1}]H'(u)(\frac{u}{F(u)})^n$$

$$G(F(x)) = x \Rightarrow [x^n]H(G(x)) = \frac{1}{n}[u^{n-1}]H'(u)(\frac{u}{F(u)})^n$$

$$G(F(x)) = x \Rightarrow [x^n]G^k(x) = \frac{k}{n}[u^{n-k}](\frac{u}{F(u)})^n$$

#### 3.2 分拆数 五边形数

$$\prod_{i\geq 1} (1-x^i) = \sum_{k=-\infty}^{\infty} (-1)^k x^{\frac{k(3k-1)}{2}}$$

#### 3.3 Fast Fourier Transform

```
using db = double:
using cp = std::complex<db>;
3 // cp::real, cp::imag, std::conj, std::arg
const db pi = std::acos(-1);
5 int rev[N], lim;
6 cp wn[N];
void init(int len) {
  | lim = 2 << std::__lg(len - 1);
  | for(static int i = 1;i < lim;i += i) {
     | for(int j = 0; j < i; ++j) {
        | wn[i + j] = std::polar(1., db(j) / i * pi);
12
     for(int i = 1;i < lim;++i) {</pre>
     | rev[i] = rev[i >> 1] >> 1 | (i % 2u * lim / 2);
18 void DFT(cp * a) {
     for(int i = 0;i < lim;++i) {</pre>
      if(rev[i] < i) std::swap(a[rev[i]], a[i]);</pre>
     for(int i = 1;i < lim;i += i) {</pre>
     | for(int j = 0; j < lim; j += i + i) {
         | for(int k = 0; k < i; ++k) {
         | cp x = a[i + j + k] * wn[i + k];
           | a[i + j + k] = a[k + j] - x;
           | a[k + j] += x;
32 void IDFT(cp * a) {
   | DFT(a), std::reverse(a + 1, a + lim);
   | for(int i = 0;i < lim;++i)
     | a[i] /= lim;
35
```

#### 3.4 Number Theoretic Transform

```
int rev[N], wn[N], lim, invlim;
  int pow(int a, int b, int ans = 1) {
  | for(;b;b >>= 1, a = (u64) a * a % mod) if(b & 1)
    \mid ans = (u64) ans * a % mod;
    return ans;
  void init(int len) {
  | lim = 2 << std::__lg(len - 1);
  | invlim = mod - (mod - 1) / lim;
    for(static int i = 1;i < lim;i += i) {</pre>
       wn[i] = 1;
       const int w = pow(3, mod / i / 2);
     | for(int j = 1; j < i; ++ j) {
       | wn[i + j] = (u64) wn[i + j - 1] * w % mod;
  | for(int i = 1;i < lim;++i) {
    | rev[i] = rev[i >> 1] >> 1 | (i % 2u * lim / 2);
|<sub>21</sub>|void DFT(int * a) {
| for(int i = 0;i < lim;++i) t[i] = a[rev[i]];
  | for(int i = 1;i < lim;i += i) {
    | for(int k = i & (1 << 19); k--;)
     | | if(t[k] >= mod * 9ull) t[k] -= mod * 9ull;
  | | for(int j = 0; j < lim; j += i + i) {
  | | | for(int k = 0; k < i; ++k) {
  | | | | t[i + j + k] = t[k + j] + (mod - x), t[k + j] += x;
  | for(int i = 0;i < lim;++i) a[i] = t[i] % mod;
36 void IDFT(int * a) {
37 | DFT(a), std::reverse(a + 1, a + lim);
  | for(int i = 0;i < lim;++i)
    | a[i] = (u64) a[i] * invlim % mod;
```

## 3.5 Generating function

```
void cpy(int * a, int * b, int n) {
    | if(a != b) memcpy(a, b, n << 2);
    | memset(a + n, 0, (lim - n) << 2);
}

void inv(int * a, int * b, int n) { // b = inv(a) mod x^n
    | if(n = 1) return void(*b = pow(*a, mod - 2));

static int c[N], d[N];
| int m = (n + 1) / 2;
| inv(a, b, m);
| init(n + m), cpy(c, b, m), cpy(d, a, n);</pre>
```

```
DFT(c), DFT(d):
     for(int i = 0; i < \lim; ++i) c[i] = (u64) c[i] * c[i] % mod * d[i] % mod;
     IDFT(c);
13
     for(int i = m; i < n; ++i) b[i] = norm(mod - c[i]);
15
_{16} | void log(int * a, int * b, int n) { // b = log(a) (mod x^n)
     static int c[N], d[N];
     inv(a, c, n), init(n + n);
     for(int i = 1; i < n; ++i) d[i - 1] = (u64) a[i] * i % mod;
     cpy(d, d, n - 1), cpy(c, c, n);
     DFT(c), DFT(d);
     for(int i = 0; i < \lim_{i \to +i} c[i] = (u64) c[i] * d[i] % mod;
     IDFT(c), *b = 0;
     for(int i = 1; i < n; ++i) b[i] = pow(i, mod - 2, c[i - 1]);
24
25 }
```

#### 3.6 全在线卷积

```
struct oc {
     std::vector<int> f, g, res;
     std::vector<std::vector<int>> fa, fb;
    int n, p;
     oc(int n) : f(n), g(n), res(n), n(n), p(0) { }
     void push(int v0, int v1) {
     | f[p] = v0;
       res[p] = (res[p] + (u64) f[0] * v1 + (u64) g[0] * v0) % mod;
       g[p++] = v1;
       static int A[N], B[N];
       int lb = p & -p;
       init(lb * 2);
       memset(A, 0, lim << 2);
       memset(B, 0, lim << 2);
       for(int i = 0; i < lb; ++i) A[i] = g[p - lb + i], B[i] = f[p - lb + i];
15
       DFT(A), DFT(B);
       if(lb = p) {
          fa.emplace back(A, A + lim);
          fb.emplace back(B, B + lim);
          for(int i = 0; i < \lim_{i \to +i} A[i] = (u64) A[i] * B[i] % mod;
          auto & C = fb[std:: lg(lim)], & D = fa[std:: lg(lim)];
          for(int i = 0; i < \lim_{t \to 0} (u64) A[i] * C[i * 2] + (u64)
             \hookrightarrow B[i] * D[i * 2]) \% mod:
     | IDFT(A);
     | for(int j = p; j 
          → lb]) % mod;
27 | }
28 };
29 struct Exp : oc {
    std::vector<int> res;
     Exp(int n) : oc(n), res(n) { }
  void push(int v) {
     | if(!res[0]) return void(res[0] = 1);
     \mid oc::push(res[p], v * u64(p + 1) % mod);
```

```
| res[p] = (u64) oc :: res[p - 1] * inv[p] % mod;
  | }
37 };
38 | struct Ln : oc {
   | std::vector<int> res; int fi;
   | Ln(int n) : oc(n), res(n), fi(0) {}
     void push(int v) {
   | | if(!fi) return void(fi = 1);
   | | oc::push(res[p] * (u64) p % mod, v);
   | res[p] = ((u64) v * p + mod - oc::res[p - 1]) % mod * inv[p] % mod;
46 };
47 struct Inv : oc {
  | std::vector<int> res; int fi;
     Inv(int n) : oc(n), res(n), fi(0) {}
     void push(int v) {
     | res[p] = fi ? (oc::res[p] + (u64) v * res[0]) % mod * (mod - res[0]) %
          \hookrightarrow mod : pow(fi = v, mod - 2);
        oc::push(res[p], v);
53
54 };
```

## 3.7 Berlekamp Massey

```
vector<int> berlekamp_massey(const vector<int> &a) {
vector<int> v, last; // v is the answer, 0-based
  \mid int k = -1, delta = 0;
  | for (int i = 0; i < (int)a.size(); i++) {
     | int tmp = 0;
       for (int j = 0; j < (int)v.size(); j++)</pre>
        | tmp = (tmp + (long long)a[i - j - 1] * v[j]) % p;
       if (a[i] = tmp) continue;
       if (k < 0) {
10
       | k = i; delta = (a[i] - tmp + p) % p;
       v = vector<int>(i + 1); continue; }
       vector<int> u = v;
       int val = (long long)(a[i] - tmp + p) *
        | qpow(delta, p - 2) % p;
      | if (v.size() < last.size() + i - k)
        v.resize(last.size() + i - k);
     | (v[i - k - 1] += val) \% = p;
     | for (int j = 0; j < (int)last.size(); j++) {
       | v[i - k + j] = (v[i - k + j] -
          | (long long)val * last[j]) % p;
        | if (v[i - k + j] < 0) v[i - k + j] += p; }
| if ((int)u.size() - i < (int)last.size() - k) {
  | if (delta < 0) delta += p; } }
  \mid for (auto \delta x : v) x = (p - x) % p;
  | v.insert(v.begin(), 1); //一般是需要最小递推式的, 处理一下
  | return v; }
  // \forall i, \sum_{j=0}^{m} a_{i-j} v_j = 0
```

#### 3.8 线性规划 | 单纯形法

```
using db = long double;
2 const db eps = 1e-16;
int sgn(db x) { return x < -eps ? -1 : x > eps; }
4 namespace LP {
    const int N = 21, M = 21;
     int n, m; // n : 变量个数, m : 约束个数
  | db a[M + N][N], x[N + M];
  | // 约束: 对于 1 <= i <= m : a[i][0] + \sum j x[j] * a[i][j] >= 0
   | // x[j] >= 0
    // 最大化 \sum_j x[j] * a[0][j]
     int id[N + M];
     void pivot(int p, int o) {
       std::swap(id[p], id[o + n]);
       db w = -a[o][p];
       for(int i = 0; i <= n; ++i) a[o][i] /= w;
       a[o][p] = -1 / w;
       for(int i = 0;i <= m;++i) if(sgn(a[i][p]) && i != o) {</pre>
           db w = a[i][p]; a[i][p] = 0;
           for(int j = 0; j <= n; ++ j) a[i][j] += w * a[o][j];
     db solve() { // nan : 无解, inf : 无界, 否则返回最大值
     | for(int i = 1; i <= n + m; ++i) id[i] = i;
        for(;;) {
           int p = 0, min = 1;
           for(int i = 1;i <= m;++i) {
           | if(a[i][0] < a[min][0]) min = i;
           if(a[min][0] >= -eps) break;
           for(int i = 1; i <= n; ++i) if(a[min][i] > eps && id[i] > id[p]) {
           | p = i;
           if(!p) return nan("");
           pivot(p, min);
        for(;;) {
           int p = 1;
           for(int i = 1; i \le n; ++i) if(a[0][i] > a[0][p]) p = i;
           if(a[0][p] < eps) break;
           db min = INFINITY; int o = 0;
           for(int i = 1;i <= m;++i) if(a[i][p] < -eps) {</pre>
            | db w = -a[i][0] / a[i][p]; int d = sgn(w - min);
           | if(d < 0 || !d && id[i] > id[o]) o = i, min = w;
           if(!o) return INFINITY;
           pivot(p, o);
     for(int i = 1;i <= m;++i) x[id[i + n]] = a[i][0];</pre>
       return a[0][0];
50
```

#### 3.9 Simpson 积分

$$\int_{a}^{b} f(x)dx \approx \frac{b-a}{3n} (f(x_{0}) + 4\sum_{i=1}^{n/2} f(x_{2i-1}) + 2\sum_{i=1}^{n/2-1} f(x_{2i}) + f(x_{n}))$$

$$\approx \frac{3(b-a)}{8n} (f(x_{0}) + 3\sum_{i=1}^{n/3} (f(x_{3i-1}) + f(x_{3i-2})) + 2\sum_{i=1}^{n/3-1} f(x_{3i}) + f(x_{n}))$$

#### 3.10 黄金三分

## 4 字符串

## 4.1 后缀自动机 | SAM

需要两倍点数量。

```
int ch[N][26], lk[N], len[N], nd = 1, las = 1;
void extend(int c, int k) {
  \mid int x = ++ nd, p = las; las = x;
  \mid len[x] = len[p] + 1;
  | for(; p && !ch[p][c]; p = lk[p]) ch[p][c] = x;
    if(!p) return lk[x] = 1, void();
    int q = ch[p][c];
    if(len[q] = len[p] + 1)
    | return lk[x] = q, void();
    int cl = ++ nd;
    len[cl] = len[p] + 1;
    memcpy(ch[cl], ch[q], 104);
    lk[cl] = lk[q], lk[q] = lk[x] = cl;
  for(; p \&\& ch[p][c] = q; p = lk[p]) ch[p][c] = cl;
| static int bin[N];
    memset(bin, 0, sizeof (int) * (n + 1));
  | for(int i = 1; i <= nd; i++) ++ bin[len[i]];
  | for(int i = 1; i <= n; i++) bin[i] += bin[i - 1];
    for(int i = nd; i; i--) A[bin[len[i]]--] = i;
```

#### 4.2 基本子串字典

```
for(int i = 2; i <= T[0].nd; i++) {</pre>
   + int x = T[0].A[i];
   | int R = T[0].r[x], L = R - T[0].len[x] + 1;
   | int y = T[1].fnd(T[1].ed[L], R - L + 1);
   | if(T[1].len[y] = R - L + 1) {
   | | ++ cnt; T[0].tag[x] = cnt; T[1].tag[y] = cnt;
     | rt[0][cnt] = x, rt[1][cnt] = y;
_{10} | for(int o = 0; o < 2; o++)
in for(int i = T[o].nd; i > 1; i--) {
   | int x = T[o].A[i];
     if(T[o].tag[x]) continue;
     for(int k = 0; k < 26; k++)
      if(T[o].ch[x][k]) T[o].tag[x] = T[o].tag[T[o].ch[x][k]];
15
_{17} | for(int o = 0; o < 2; o++)
_{18} for(int i = 2; i <= T[o].nd; i++) {
| | |  int x = T[o].A[i];
20 | | vec[o][T[o].tag[x]].pb(x);
|z_1| // vec[0] : from left to right, node id of the columm , vec[1] : from down
_{22} // U : T[0].r[rt] - T[0].len[rt] + 1, D = U + vec[1][t].size() - 1, L = R -
     \rightarrow vec[0][t].size() + 1, R = T[0].r[rt]
int x = T[0].fnd(T[0].ed[r], r - l + 1);
_{24} int blk = T[0].tag[x];
_{25} // distance to the right , 0 - base
int rp = T[0].r[rt[0][blk]] - T[0].r[x];
_{27} // distance to the up // the upper - right point is (T[0].r[rt] -
     \hookrightarrow T[0].len[rt] + 1, T[0].r[rt])
int lp = T[0].r[x] - (r - l) - (T[0].r[rt[0][blk]] - T[0].len[rt[0][blk]] +
    → 1);
```

#### 4.3 DAG 剖分

```
void build() {
  | for(int i = 2; i <= nd; i++)
     | e[lk[i]].pb(i);
    static int q[N], d[N];
  | for(int i = 1; i <= nd; i++)
     | for(int j = 0; j < 26; j++)|
        | if(ch[i][j]) ++ d[ch[i][j]];
  | int hd = 1, tl = 0;
  | q[++ tl] = 1;
    while(hd <= tl) {</pre>
  \mid | int x = q[hd ++];
  | | for(int i = 0; i < 26; i++)
         | if(ch[x][i]) {
13
         \mid int v = ch[x][i];
         | if((--d[v]) = 0) q[++tl] = v;
  | static ll f[N], h[N];
```

```
19 | for(int i = tl, x; i; i--) {
  | | f[x = q[i]] ++;
  | | for(int j = 0; j < 26; j++)|
  | | | if(ch[x][j]) f[x] += f[ch[x][j]];
23
  | for(int i = 1, x; i <= tl; i++) {
  | | h[x = q[i]] ++;
   | | for(int j = 0; j < 26; j++)
     | | if(ch[x][j]) h[ch[x][j]] += h[x];
28
  | static int nx[N], fr[N];
  | for(int i = 1; i <= nd; i++) {
     | for(int j = 0; j < 26; j++)|
        | if(ch[i][j] && f[ch[i][j]] > f[nx[i]]) nx[i] = ch[i][j];
   | | for(int j = 0; j < 26; j++)|
     | | if(ch[i][j] && h[i] > h[fr[ch[i][j]]]) fr[ch[i][j]] = i;
  fr[0] = nx[0] = 0;
   | static bool vis[N];
  | for(int i = 1; i <= nd; i++) {
     | if(fr[nx[i]] = i) son[i] = nx[i], vis[son[i]] = 1;
```

#### 4.4 exKMP

```
| static int lcp[N]:
1 int mx=1, pt=1; lcp[1]=n;
for(int i=2; i<=n; i++){</pre>
4 | if(i<=mx) lcp[i]=min(lcp[i-pt+1],mx-i+1);</pre>
    while(i+lcp[i]<=n && S[i+lcp[i]]=S[1+lcp[i]]) ++lcp[i];</pre>
    if(i+lcp[i]-1>mx) pt=i, mx=i+lcp[i]-1;
```

## 4.5 log 个最小后缀

```
for(int i = 1; i <= n; i ++) {</pre>
                      St.pb(i); vector<int> nw;
                       for(auto t : St) {
                       | bool ok = true:
                                   while(!nw.empty()){
                                      int x = nw.back();
                                                if(S[i] > S[i-t+x]) ok = false;
                                       if(S[i] >= S[i-t+x]) break; nw.pop_back();
             | if(ok \&\& (nw.empty() || (i - t + 1 <= t - nw.back()))) nw.pb(t);
             | } St = nw;
|_{13}| for (int x : St){
14 | bool FLAG = true;
| while(nx.size()){
|_{17}| | if(S[x + lcp] > S[y + lcp]){ FLAG = false; break; } nx.pop back();
|x| + |x| + |x| = x - x + 1 \le x - x + 1 
 <sub>19</sub>|} // in segmentree, work(L, ans, rpos), work(R, ans, rpos), then return ans
```

```
4.6 SA
 _{1} | char s[N]; int m, rk[N * 2], sa[N], tmp[N * 2], h[N], y[N];
void Sort(){
   | for(int i=1;i<=m;i++) c[i] = 0;
     for(int i=1;i<=n;i++) c[rk[i]]++;</pre>
     for(int i=1;i<=m;i++) c[i] += c[i-1];
     for(int i=n;i>=1;i--) sa[c[rk[y[i]]]--] = y[i];
s void get sa(){
     for(int i=1;i<=n;i++) rk[i] = s[i], y[i] = i; Sort();</pre>
     for(int k=1;k<=n;k<<=1){</pre>
     | int ret = 0;
       | for(int i=n-k+1;i<=n;i++) y[++ret] = i;
        for(int i=1;i<=n;i++) if(sa[i] > k) y[++ret] = sa[i] - k;
        for(int i = 1; i <= n; i++) swap(rk[i], tmp[i]);</pre>
        rk[sa[1]] = 1; int num = 1;
        for(int i=2;i<=n;i++){</pre>
            if(tmp[sa[i]] = tmp[sa[i-1]] \delta tmp[sa[i]+k] = tmp[sa[i-1]+k])
             | rk[sa[i]] = num;
         | else rk[sa[i]] = ++num;
      | } m = num;
22
23 }
24 void get_h(){
   | int k = 0;
     for(int i=1;i<=n;i++){</pre>
      | if(rk[i]=1) continue;
      | int j = sa[rk[i]-1]; if(k) k--;
        while(i+k \le n \delta \delta j+k \le n \delta \delta s[i+k] = s[j+k]) k++;
     | h[rk[i]] = k;
31
 4.7 PAM
```

```
namespace pam {
  int ch[N][26], len[N], lk[N], rp, las, nd, top[N], d[N];
  | void init() { rp = 0, las = nd = 1, len[1] = -1, lk[0] = 1; }
  | // remember to set S[0] = *
    int jmp(int x) { while(S[rp - len[x] - 1] != S[rp]) x = lk[x]; return x; }
    void ins(int c) {
    | ++ rp; int p = jmp(las);
       if(!ch[p][c]) {
          int x = ++ nd;
          len[x] = len[p] + 2;
          lk[x] = ch[jmp(lk[p])][c];
          ch[p][c] = x;
          if(len[x] - len[lk[x]] = d[lk[x]])
           \mid top[x] = top[lk[x]], d[x] = d[lk[x]];
          else {
             top[x] = x;
             d[x] = len[x] - len[lk[x]];
```

#### 4.8 AC 自动机

#### 4.9 Manacher

```
S[1] = '%';
for(int i = 1; i <= len; i++){
    | S[i << 1] = '&';
    | S[i << 1|1] = s[i];
}
len = len << 1 | 1;
S[++len] = '&';
S[++len] = '$;
int mx = 0, id = 0, ans = 0;
for(int i = 1; i <= len; i++){
    | if(mx > i) p[i] = min(p[id * 2 - i], mx - i);
    | else p[i] = 1;
    | while(S[i - p[i]] = S[i + p[i]]) ++p[i];
    | if(i + p[i] > mx) id = i, mx = i + p[i];
    | ans = max(ans, p[i] - 1);
}
```

## 4.10 Lyndon/最小表示法

```
vector <int> duval(vector <int> S) {
  | int i = 0, j, k, s = S.size(); vector <int> ans;
   while(i < s) {</pre>
  | | j = i, k = i + 1;
    | while(j < s && k < s && S[j] <= S[k]) {
       | if(S[j] = S[k]) ++ j;
    | | else j = i; ++ k;
     | } while (i \le j) \{ ans.pb(i + k - j - 1); i += k - j; \}
  vector <int> min_rep(vector <int> S) {
  | int k = 0, i = 0, j = 1, n = S.size();
    while (k < n && i < n && j < n) {
    | if (S[(i + k) \% n] = S[(j + k) \% n]) k ++;
     ∣ else {
15
         S[(i + k) \% n] > S[(j + k) \% n] ? i = i + k + 1 : j = j + k + 1;
        | if (i = j) i ++; k = 0;
18
    } i = min(i, j);
    rotate(S.begin(), S.begin() + i, S.end()); return S;
```

#### 4.11 Runs

```
1 // need lcp and lcs
bool cmp(int x, int y){
| | int l = lcp(x, y);
4 | if(x + l > n) return true;
5 | if(y + l > n) return false;
   \mid return S[x + l] < S[y + l];
s set <pi> ex;
void ins(int l, int r) {
| | | | int p = r - 1;
   | int l1 = lcp(l, r);
   | int l2 = lcs(l - 1, r - 1);
   | int L = l - l2, R = r + l1 - 1;
   | if(R - L + 1 >= 2 * p) {
   | | auto iter = ex.lower_bound(pi(L, R));
   \mid if(iter != ex.end() && *iter = pi(L, R)) return;
     | ex.emplace_hint(iter, pi(L, R));
     | runs.pb((run){L, R, p});
19
20 }
21 void Run(int o){
     static int s[N];
     int top = 0; s[++top] = n + 1;
     for(int i = n; i; i--){
   | | while(top > 1 && cmp(i, s[top]) = o) --top;
     | ins(i, s[top]), s[++top] = i;
27 | }
28 }
```

## 5 数据结构

#### 5.1 区间加区间求和树状数组

```
// 后缀加,前缀求和
struct BIT {
    | ll a[N], b[N];
    | void add(ll p, int v) {
    | | for(int i = p;i < N;i += i & -i)
    | | | a[i] += v, b[i] += p * v;
    | }
    | ll qry(ll p) {
    | | ll res = 0;
    | | for(int i = p;i;i &= i - 1) res += (p + 1) * a[i] - b[i];
    | return res;
    | }
    | void add(int l, int r, int v) { add(l, v), add(r + 1, -v); }
    | | ll qry(int l, int r) { return qry(r) - qry(l - 1); }
} bit;
```

#### 5.2 zkw 线段树

```
struct seg {
   | ll o[1 << 20]; int L;
     void upt(int x) {
   | o[x] = o[x << 1] + o[x << 1 | 1];
  void init(int n, int * w) {
   | L = 2 << std:: lg(n + 1);
s | | for(int i = 1; i <= n; ++i) o[i + L] = w[i];</pre>
9 | | for(int i = L;i >= 1;--i) upt(i);
10 | }
| | | void upt(int p, int v) {
| | | | | for(o[p += L] += v;p >>= 1;upt(p));
  | ll gry(int l, int r) {
15 | | l += L - 1, r += L + 1;
| 16 | | | 11 ans = 0;
17 | | for(; l ^ r ^ 1; l >>= 1, r >>= 1) {
  | | | if((l & 1) = 0) ans += o[l ^ 1];
        | if((r \& 1) = 1) ans += o[r ^ 1];
   | | return ans;
  | // if there is no I
\mid if(l = r) return o[l + L];
   | | | ll le = o[l + L], ri = o[r + L];
|<sub>27</sub>| | | l += L, r += L;
  | | for(; l ^ r ^ 1; l >>= 1, r >>= 1) {
|_{29}| + |_{1} \text{ if}((l & 1) = 0) \text{ le = le + o[l }^1];
|_{30}| | | if((r & 1) = 1) ri = o[r ^ 1] + ri:
31 | | }
     | return le + ri;
```

```
34 } sgt;
 5.3 Link Cut Tree
int son[N][2], fa[N], rev[N];
int get(int x, int p = 1) { return son[fa[x]][p] = x; }
void update(int x) { }
int is_root(int x) { return !(get(x) || get(x, 0)); }
void rotate(int x) {
  | int y = fa[x], z = fa[y], b = get(x);
  if(!is_root(y)) son[z][get(y)] = x;
  | son[y][b] = son[x][!b], son[x][!b] = y;
  fa[son[y][b]] = y, fa[y] = x, fa[x] = z;
  | update(y);
void put(int x) {
     if(x) rev[x] = 1, std::swap(son[x][0], son[x][1]);
void down(int x) {
    if(rev[x]) {
     | put(son[x][0]);
     | put(son[x][1]);
     | rev[x] = 0;
void pushdown(int x) {
     if(!is_root(x)) pushdown(fa[x]);
     down(x);
void splay(int x) {
     for(pushdown(x);!is_root(x);rotate(x)) if(!is_root(fa[x]))
     | rotate(get(x) ^{\circ} get(fa[x]) ? x : fa[x]);
     update(x):
void access(int x) {
     for(int t = 0;x;son[x][1] = t, t = x, x = fa[x])
     \mid splay(x);
void makeroot(int x) {
    access(x), splay(x), put(x);
```

## 5.4 FHQ Treap

```
int root.cc:
struct hh{int w,pri,ch[2],size;}tr[sz];
#define ls(x) tr[x].ch[0]
#define rs(x) tr[x].ch[1]
5 | void pushup(int x){tr[x].size=1+tr[ls(x)].size+tr[rs(x)].size;}
6 int newnode(int w) {
7 ++CC;
  | tr[cc].w=w,tr[cc].pri=rnd(1,int(1e9)),tr[cc].size=1;
  | return cc;
int merge(int x,int y) {
```

```
| | | | |  if (!x||!y) return x+y;
   if (tr[x].pri<tr[y].pri) return rs(x)=merge(rs(x),y),pushup(x),x;</pre>
   return ls(y)=merge(x,ls(y)),pushup(y),y;
void split(int x,int w,int &a,int &b) {
| if (!x) return a=b=0,void();
| if (tr[x].w<=w) a=x,split(rs(x),w,rs(x),b);
else b=x,split(ls(x),w,a,ls(x));
   | pushup(x);
```

#### 5.5 pbds tree

```
#include <bits/extc++.h>
using namespace __gnu_pbds;
s template<class T> // insert, erase, join, order_of_key, find_by_order(return

    iterator), order is 0-index

using Tree = tree<T, null_type, std::less<T>, rb_tree_tag,

    tree_order_statistics_node_update>;
```

## geometry

#### 6.1 向量

```
using db = long double:
2 const db eps = 1e-10;
 _{3} db sgn(db x) { return x < -eps ? -1 : x > eps; }
db eq(db x, db y) { return !sgn(x - y); }
s struct p2 {
6 | db x, y;
    db norm() const { return x * x + y * y; }
    db abs() const { return std::sqrt(x * x + y * y); }
    db arg() const { return atan2(y, x); }
10 };
|| db arg(p2 x, p2 y) |
| | | db a = y.arg() - x.arg();
  \mid if(a > pi) a -= pi * 2;
| if(a < -pi) a += pi * 2;
  | return a;
p2 r90(p2 x) { return {-x.y, x.x}; }
19 p2 operator - (p2 x, p2 y) { return {x.x - y.x, x.y - y.y}; }
p2 operator / (p2 x, db y) { return {x.x / y, x.y / y}; }
|p|_{21} p2 operator * (p2 x, db y) { return {x.x * y, x.y * y}; }
|_{23} | db operator * (p2 x, p2 y) { return x.x * y.y - x.y * y.x; }
| 25 | int half(p2 x){return x.y < 0 || (x.y = 0 && x.x <= 0); }
int half(p2 x){return x.y < -eps || (std::fabs(x.y) < eps \delta\delta x.x < eps);}
|z| bool cmp(p2 a, p2 b) { return half(a) = half(b) ? a * b > 0 : half(b); }
bool cmp_eq(p2 A, p2 B) { return half(A) = half(B) & eq(A \star B, 0); }
|29|// 判断 A, B, C 三个向量是否是逆时针顺序
30 // 如果是, 返回 1
|₃₁|// 如果 (A, B), (C, B) 同方向共线,返回 -1
```

#### 6.2 直线半平面

```
struct line : p2 {
2 | db z;
  | // a * x + b * y + c (= or >) 0
| | | line() = default:
   | line(db a, db b, db c) : p2\{a, b\}, z(c) \{\}
   | line(p2 a, p2 b) : p2(r90(b - a)), z(a * b) { } // 左侧 > 0
    db operator ()(p2 a) const { return a % p2(*this) + z; }
    line perp() const { return {y, -x, 0}; } // 垂直
    line para(p2 o) { return {x, y, z - (*this)(o)}; } // 平行
10 };
p2 operator & (line x, line y) {
p2{p2{x.z, x.y}} * p2{y.z, y.y}, p2{x.x, x.z} * p2{y.x, y.z} /
       \hookrightarrow -(p2(x) * p2(y));
  | // 注意此处精度误差较大,以及 res.y 需要较高精度
<sup>15</sup>|p2 proj(p2 x, line l){return x - p2(l) * (l(x) / l.norm());}//投影
16 | p2 refl(p2 x, line l){return x - p2(l) * (l(x) / l.norm()) * 2;}//对称
n/db dist(line l, p2 x={0, 0}){return l(x) / l.abs();}//有向点到线距离
18 | bool is_para(line x, line y){return eq(p2(x) * p2(y), 0);}//判断线平行
polis_perp(line x, line y){return eq(p2(x) % p2(y), 0);}//判断线垂直
20 | bool online(p2 x, line l) { return eq(l(x), 0); } // 判断点在线上
int ccw(p2 a, p2 b, p2 c) {
  \mid int sign = sgn((b - a) * (c - a));
  \mid if(sign = 0) {
   | if(sgn((b - a) \% (c - a)) = -1) return 2;
   \mid if((c - a).norm() > (b - a).norm() + eps) return -2;
26
  | return sign;
27
29 db det(line a, line b, line c) {
     p2 A = a, B = b, C = c;
   | return c.z * (A * B) + a.z * (B * C) + b.z * (C * A);
33 db check(line a, line b, line c) { // sgn same as c(a & b), 0 if error
```

```
| return sgn(det(a, b, c)) * sgn(p2(a) * p2(b)); | bool paraS(line a, line b) { // 射线同向 | return is_para(a, b) && p2(b) > 0; | | |
```

#### 6.3 半平面交

```
std::vector<p2> HPI(std::vector<line> vs) {
| auto cmp = [](line a, line b) {
  if(paraS(a, b)) return dist(a) < dist(b);</pre>
4 | return ::cmp(p2(a), p2(b));
5 | };
sort(vs.begin(), vs.end(), cmp);
  int ah = 0, at = 0, n = size(vs);
     std::vector<line> deq(n + 1);
     std::vector<p2> ans(n);
     deg[0] = vs[0]:
  | for(int i = 1; i <= n; ++ i) {
  | | line o = i < n ? vs[i] : deq[ah];
     if(paraS(vs[i - 1], o)) continue;
  | | for(;ah < at && check(deq[at - 1], deq[at], o) < 0;) -- at;//maybe <=
  \mid if(i != n) for(;ah < at && check(deq[ah], deq[ah + 1], o) < 0;) ++ ah;
  | | if(!is para(o, deg[at])) {
17 | | | ans[at] = o & deq[at];
          deq[++at] = o;
19
20 | }
  | if(at - ah <= 2) return {};
    return {ans.begin() + ah, ans.begin() + at};
```

#### 6.4 线段

```
1 struct seg {
 2 | p2 x, y;
     seg() {}
     seg(const p2 \& A, const p2 \& B) : x(A), y(B) {}
     bool onseg(const p2 & o) const {
     | return (o - x) % (o - y) < eps &\text{$\text{s}} std::fabs((o - x) * (o - y)) < eps;
 9 db dist(const seg & o, const p2 & x) {
| if((o.x - o.y)) \% (x - o.y) <= eps) return (x - o.y).abs();
| if((o.y - o.x) % (x - o.x) <= eps) return (x - o.x).abs();
   | return fabs((0.x - x) * (0.y - x) / (0.x - 0.y).abs());
13 | }
14 bool is isc(const seg & x, const seg & y) {
| 16 | | ccw(x.x, x.y, y.x) * ccw(x.x, x.y, y.y) <= 0 &&
     | ccw(y.x, y.y, x.x) * ccw(y.x, y.y, x.y) <= 0;
| db dist(const seg & x, const seg & y) {
|z_0| | if(is isc(x, y)) return 0;
|z_1| | return std::min({dist(y, x.x), dist(y, x.y), dist(x, y.x), dist(x, y.y)});
```

```
| }
  6.5 多边形
                                                                                                                                        | return 1;
 using polygon = std::vector<p2>;
2 // counter-clockwise
                                                                                                                                     6.7 图形求交
db area(const polygon & x) {
_{4} | db res = 0:
                                                                                                                                    struct circle : p2 { db r; };
    for(int i = 2;i < (int) x.size();++i) {</pre>
                                                                                                                                    2 std::vector<p2> operator δ (circle o, line l) {
        | res += (x[i - 1] - x[0]) * (x[i] - x[0]);
                                                                                                                                    | p2 v = 1, Rv = r90(v); db L = 1.abs();
                                                                                                                                    | db d = l(p2(0)) / L, x = 0.r * 0.r - d * d;
     | return res / 2;
                                                                                                                                       | if(x < -eps) return {};
                                                                                                                                       \mid x = std::sqrt(x * sgn(x));
10 bool is convex(const polygon & x, bool strict = 1) {
                                                                                                                                       | p2 z = p2(o) - v * (d / L), p = Rv * (x / L);
    | // warning, maybe wrong
                                                                                                                                       \mid return \{z + p, z - p\};
    const db z = strict ? eps : -eps;
                                                                                                                                    。 } // l 如果是构造函数给出,那么返回交点按射线顺序
    for(int i = 2;i < (int) x.size() + 2;++i) {</pre>
                                                                                                                                    10 std::vector<p2> operator & (circle o, seg s) {
        if((x[(i-1) % x.size()] - x[i-2]) * (x[i % x.size()] - x[i-2]) < (x[i % 
                                                                                                                                       | std::vector<p2> b;
                                                                                                                                       | for(p2 x : (o & s.to_l()))
15 | }
                                                                                                                                       | | if(s.onseg(x)) b.push_back(x);
    return 1;
                                                                                                                                       | return b;
17 | }
ɪs int contain(const std::vector<p2> & a, p2 o) { // 简单多边形包含判定
                                                                                                                                   16 std::vector<p2> operator δ (circle o0, circle o1) {
        bool in = 0:
                                                                                                                                   |p| | p2 tmp = (p2(o1) - p2(o0)) * 2.;
        for(int i = 0;i < (int) a.size();++i) {</pre>
                                                                                                                                   18 | return 00 & line(tmp.x, tmp.y, o1.r * o1.r - o0.r * o0.r + o0.norm() -
        | p2 x = a[i] - o, y = a[(i + 1) \% a.size()] - o;
                                                                                                                                               → o1.norm());
        if(x.y > y.y) std::swap(x, y);
         | if(x.y \le eps \& y.y > eps \& x * y < -eps) in ^= 1;
                                                                                                                                   20 std::vector<p2> tang(circle o, p2 x) {
         if(std::fabs(x * y) < eps & x % y < eps) return 2; // 在线段上,看情况改
                                                                                                                                       | db d = (x - p2(o)).abs();
25
                                                                                                                                   | | if(d <= o.r + eps) return {};
        return in;
26
                                                                                                                                        | return o & circle{x, sqrt(d * d - o.r * o.r)};
                                                                                                                                   25 // 三角形 (0, a, b) 和圆 o 的交的有向面积 * 2
  6.6 线段 in 多边形
                                                                                                                                   db intersect(circle o, p2 a, p2 b) {
bool contains(p2 x, p2 y, const std::vector<p2> & a) {
                                                                                                                                   |a| = a - p2(0), b = b - p2(0); o.x = o.y = 0;
using pr = std::pair<double, int>;
                                                                                                                                       | int va = a.abs() <= o.r + eps;
std::vector<pr> e = {pr(-inf, 0), pr(inf, 0)};
                                                                                                                                       | int vb = b.abs() <= o.r + eps;
    | p2 t = y - x;
                                                                                                                                        | if(va && vb) return a * b;
     | auto f = [&](p2 a, p2 b, p2 c, p2 d) {
                                                                                                                                           auto v = o \delta seg\{a, b\}; // 注意这里, 有必要改一下 onseg, 去掉平行判定
        | return (b - a).abs() * ((c - a) * (d - c)) / ((b - a) * (d - c));
                                                                                                                                           if(v.empty()) return arg(a, b) * o.r * o.r;
     | };
                                                                                                                                           db sum = 0;
     | for(int i = 0;i < (int) a.size();++i) {
                                                                                                                                           sum += va ? a * v[0] : arg(a, v[0]) * o.r * o.r;
        | p2 u = a[i], v = a[(i + 1) \% a.size()];
                                                                                                                                       \mid sum += vb ? v.back() * b : arg(v.back(), b) * o.r * o.r;
     \mid int a = sgn(t * (u - x));
                                                                                                                                       | if(v.size() > 1) sum += v[0] * v[1];
        \mid int b = sgn(t * (v - x));
                                                                                                                                           return sum;
        \mid if(a != b) e.emplace_back(f(x, y, u, v), b - a);
                                                                                                                                   39 // 有向弓形面积 * 2, arg 不能改
sort(e.begin(), e.end());
                                                                                                                                   40 db csegS(circle o, p2 a, p2 b) {
    | int sum = 0; db R = t.abs();
                                                                                                                                   |a_1| + a = a - p2(o);
     | for(int i = 0;i + 1 < (int) e.size();++i) {
                                                                                                                                       | b = b - p2(o);
17 | | sum += e[i].second:
                                                                                                                                       | db d = b.arg() - a.arg();
        | if(sum = 0 && std::max(e[i].first, 0.) + eps < std::min(e[i +
                                                                                                                                       | if(d < 0) d += pi * 2;
                \hookrightarrow 1].first, R)) {
                                                                                                                                       | return d * o.r * o.r - a * b;
              | return 0;
                                                                                                                                   46 }
```

# 6.8 凸包

结果为逆时针。

```
_{1} db cross(p2 x, p2 y, p2 z) { return (y.x - x.x) * (z.y - x.y) - (y.y - x.y) *
    \hookrightarrow (z.x - x.x); }
std::vector<p2> gethull(std::vector<p2> o) {
  | rgs::sort(o, [](p2 x, p2 y) { return eq(x.x, y.x) ? x.y < y.y : x.x < y.x;
  | o.erase(unique(o.begin(), o.end(), [](p2 x, p2 y) {
  \mid return eq(x.x, y.x) && eq(x.y, y.y);
  | }), o.end());
7 | std::vector<p2> s;
  | for(int i = 0;i < (int) o.size();++i) {
  | | for(;s.size() >= 2 && cross(s.rbegin()[1], s.back(), o[i]) <= eps;)
     | s.push_back(o[i]);
12
   for(int i = o.size() - 2, t = s.size(); i >= 0; --i) {
     | for(;s.size() > t && cross(s.rbegin()[1], s.back(), o[i]) <= eps;)
        | s.pop_back();
     | s.push_back(o[i]);
  | if(s.size() > 1) s.pop back();
  | return s;
20 } // 把两个 eps 改成 -eps 可求出所有在凸包上的点
int findmin(std::vector<p2> & a, auto cmp) {
  | int l = 0, r = a.size() - 1, d = 1;
  | if(cmp(a.back(), a[0])) std::swap(l, r), d = -1;
  | for(;(r - l) * d > 1;) {
  | | int mid = (l + r) >> 1;
     if(cmp(a[mid], a[mid - d]) && cmp(a[mid], a[l])) {
     \mid r = mid;
30
  | }
32 | return l;
33 } // cmp is less, and a.size()>0 plz
int contains(std::vector<p2> & a, p2 x) {
     auto it = lower bound(a.begin() + 2, a.end(), x, [\delta](p2 x, p2 y) {
     | return cross(a[0], x, y) > 0;
  | });
  | ll c0 = cross(it[-1], *it, x), c1 = cross(a[0], a[1], x);
```

```
| if(it != a.end() && c0 >= 0 && c1 >= 0) {
| return c0 > 0 && c1 > 0 && cross(a.back(), a[0], x) > 0 ? IN : ON;
| else {
| return 0;
| }
| // a.size()>2 plz
```

#### 6.9 上凸壳

结果显然为顺时针。

```
std::vector<p2> gethull(std::vector<p2> o) {
  | sort(o.begin(), o.end(), [](p2 x, p2 y) {
     | if(x.x = y.x) 
        | return x.y > y.y; // gt \Rightarrow lt
     | | return x.x < y.x;
  | });
   | std::vector<p2> stack;
  \mid for(p2 x : o) {
  \mid if(stack.size() && stack.back().x = x.x) {
        continue;
13
14 | | for(;stack.size() >= 2 && cross(stack.rbegin()[1], stack.back(), x) >=
          \hookrightarrow 0;) { // gt \Rightarrow lt
15 | | | stack.pop_back();
16 | | }
| | | stack.push_back(x);
18 | }
19 | return stack;
```

#### 6.10 最小圆覆盖

```
struct circle : p2 { db r; };
     circle incircle(p2 a, p2 b, p2 c) {
     | db A = (b - c).abs(), B = (c - a).abs(), C = (a - b).abs();
                 | return \{(a * A + b * B + c * C) / (A + B + C), fabs((b - a) * (c - a)) / (a + B + C)\} 
                                            \hookrightarrow (A + B + C)};
     5 } // 三点确定内心,不是最小圆覆盖内容
     6 circle circumcenter(p2 a, p2 b, p2 c) {
     |p| + |p| 
     |p| = |p| 
    9 | return {0, (a - 0).abs()};
|10|}// 三点确定外心
| circle cir(p2 a, p2 b) { // 根据直径生成圆
| return {(a + b) / 2, (a - b).abs() / 2};
| s | circle mincircle(std::vector<p2> a) { // 最小圆覆盖, 需要 shuffle
| | circle o = {a[0], 0};
| 17 | | int n = a.size();
| | | | for(int i = 1;i < n;++i) {
| if(in(o, a[i])) continue;
                                  | o = cir(a[0], a[i]);
```

```
| for(int j = 1; j < i; ++ j) {
          if(in(o, a[j])) continue;
          o = cir(a[j], a[i]);
          for(int k = 0; k < j; ++k) {
          if(in(o, a[k])) continue;
          | o = circumcenter(a[i], a[j], a[k]);
  return o;
6.11 最近点对
db mindist(std::vector<p2> a) {
  | db ans = 1e18:
  | sort(a.begin(), a.end(), [](p2 x, p2 y) { return x.x < y.x; });
  \mid ans = (a[0] - a[1]).abs();
  auto solve = [8](auto s, int l, int r) {
```

## | | if(l + 1 = r) return; $\mid$ int mid = $(l + r) \gg 1$ ;

 $\mid$  db mx = a[mid].x; | | s(s, l, mid), s(s, mid, r); static std::vector<p2> b; b.clear(); inplace\_merge(a.begin() + l, a.begin() + mid, a.begin() + r, [](p2 x,  $\hookrightarrow$  p2 y) { return x.y < y.y; }); | for(int i = l;i < r;++i) if(fabs(a[i].x - mx) <= ans) b.push\_back(a[i]);</pre>

| | for(int j = i - 1; j >= 0 && b[i].y <= b[j].y + ans;--j) ans = $\hookrightarrow$  std::min(ans, (b[i] - b[j]).abs()); | }; solve(solve, 0, a.size()); return ans;

for(int i = 1;i < (int) b.size();++i)</pre>

## 6.12 凸包直径

18

```
db convex_diameter(std::vector<p2> & o) {
 \mid int n = size(o);
  \mid db max = 0:
  | for(int i = 0, j = 0; i < n; ++i) {
  | for(;j + 1 < n \& (o[j] - o[i]).abs() < (o[j + 1] - o[i]).abs();) ++ j;
    | max = std::max(max, (o[j] - o[i]).abs());
  | return max;
9 // 凸包直径
```

## 6.13 切凸包

```
std::vector<p2> cut(const std::vector<p2> & o, line l) {
std::vector<p2> res;
3 | int n = size(o);
4 | for(int i = 0;i < n;++i) {
  | p2 a = o[i], b = o[(i + 1) \% n];
     | if(sgn(l(a)) >= 0) res.push_back(a); // 注意 sgn 精度
```

```
| if(sgn(l(a)) * sgn(l(b)) < 0) res.push_back(line(a, b) & l);
9 | if(res.size() <= 2) return {};</pre>
10 return res:
11 } // 切凸包
```

#### 6.14 V图

```
| std::vector<line> cut(const std::vector<line> & o, line l) {
1 | std::vector<line> res:
   \mid int n = size(o):
     for(int i = 0;i < n;++i) {</pre>
   | | line a = o[i], b = o[(i + 1) \% n], c = o[(i + 2) \% n];
        int va = check(a, b, l), vb = check(b, c, l);
      | if(va > 0 || vb > 0 || (va = 0 && vb = 0)) {
        | res.push back(b);
   | | if(va >= 0 && vb < 0) {
   | | res.push_back(l);
12
13 | }
   if(res.size() <= 2) return {};</pre>
15 | return res;
16 } // 切凸包
| 17 | line bisector(p2 a, p2 b) { return line(a.x - b.x, a.y - b.y, (b.norm() -
    \rightarrow a.norm()) / 2); }
std::vector<std::vector<line>> voronoi(std::vector<p2> p) {
19 | int n = p.size();
     auto b = p; shuffle(b.begin(), b.end(), gen);
   | const db V = 1e5; // 边框大小, 重要
   std::vector<std::vector<line>> a(n, {
   | | \{ V, 0, V * V \}, \{ 0, V, V * V \},
   | \{-V, 0, V * V\}, \{0, -V, V * V\},
24
25 | });
   | for(int i = 0;i < n;++i) {
   | | for(p2 x : b) if((x - p[i]).abs() > eps) {
   \mid \cdot \mid \cdot \mid a[i] = cut(a[i], bisector(p[i], x));
30
   | }
   | return a;
```

## 6.15 Delaunay 三角剖分

```
using i128 = __int128;
using Q = struct Quad*;
  p2 arb(LLONG_MAX, LLONG_MAX);
4 struct Quad {
    Q rot, o; p2 p = arb; bool mark;
    p2& F() { return r() -> p; }
7 | | Q& r() { return rot->rot; }
s | Q prev() { return rot->o->rot; }
9 | Q next() { return r()->prev(); }
10 } *H;
```

```
|| | | | return (b - a) * (c - a);
13 | }
14 bool circ(p2 p, p2 a, p2 b, p2 c) { // p 是否在 a, b, c 外接圆中
   | i128 p2 = p.norm(), A = a.norm() - p2, B = b.norm() - p2, C = c.norm() -

→ p2;

   | a = a - p, b = b - p, c = c - p;
     return (a * b) * C + (b * c) * A + (c * a) * B > 0;
19 Q link(p2 orig, p2 dest) {
   | Q r = H ? H : new Quad{new Quad{new Quad{new Quad{0}}}};
   | H = r \rightarrow 0; r \rightarrow r() \rightarrow r() = r;
   + for(int i = 0;i < 4;++i)
   | r = r -> rot, r -> p = arb, r -> o = i & 1 ? r : r -> r();
     r \rightarrow p = orig, r \rightarrow F() = dest;
   return r;
25
26 }
void splice(Q a, Q b) {
     std::swap(a -> o -> rot -> o, b -> o -> rot -> o);
     std::swap(a \rightarrow o, b \rightarrow o);
30
31 Q conn(Q a, Q b) {
   \mid Q q = link(a -> F(), b -> p);
     splice(q, a -> next());
   \mid splice(q -> r(), b);
   | return q;
std::pair<Q, Q> rec(const std::vector<p2> & s) {
   | int N = size(s);
   \mid if(N \le 3) {
   | Q a = link(s[0], s[1]), b = link(s[1], s.back());
   | if(N = 2) return {a, a -> r()};
   \mid | splice(a -> r(), b);
   | | | ll side = cross(s[0], s[1], s[2]);
   \mid \ \mid \ Q \ c = side ? conn(b, a) : 0;
     | return {side < 0 ? c->r() : a, side < 0 ? c : b -> r() };
   | }
_{47} | #define H(e) e -> F(), e -> p
#define valid(e) (cross(e->F(), H(base)) > 0)
   \mid int half = N / 2;
   | auto [ra, A] = rec({s.begin(), s.end() - half});
   | auto [B, rb] = rec({s.end() - half, s.end()});
   \mid while((cross(B -> p, H(A)) < 0 && (A = A -> next())) \mid
   | | (cross(A \rightarrow p, H(B)) > 0 \& (B = B \rightarrow r() \rightarrow o)));
   \mid Q base = conn(B -> r(), A);
   if(A \rightarrow p = ra \rightarrow p) ra = base \rightarrow r();
   \mid if(B -> p = rb -> p) rb = base;
#define DEL(e, init, dir) Q e = init -> dir; if(valid(e)) \
   | for(;circ(e -> dir -> F(), H(base), e -> F());) { \
   | Qt = e \rightarrow dir; \setminus
      | splice(e, e -> prev()); \
   \mid splice(e -> r(), e -> r() -> prev()); \
      | e -> o = H, H = e, e = t; \setminus
```

```
|64| | for(;;) {
                            DEL(LC, base \rightarrow r(), o);
                            DEL(RC, base, prev());
                           if(!valid(LC) && !valid(RC)) break;
                            if(!valid(LC) || (valid(RC) && circ(H(RC), H(LC))))
          \mid \cdot \mid base = conn(RC, base -> r());
          | | else
 r_1 | | base = conn(base -> r(), LC -> r());
72
 73 | return {ra, rb};
 rs | std::vector<p2> triangulate(std::vector<p2> a) {
                 sort(a.begin(), a.end()); // unique
           if((int)size(a) < 2) return {};</pre>
          | Q e = rec(a).first;
          | std::vector<Q> q = {e};
|so| | while(cross(e -> o -> F(), e -> F(), e -> p) < 0) e = e -> o;
|s_1| #define ADD { Q c = e; do { c -> mark = 1; a.push back(c -> p); \
|q| = |q| + |q| 
83 | ADD; a.clear();
|s_4| | for(int qi = 0;qi < (int) size(q);) if(!(e = q[qi++]) -> mark) ADD;
85 | return a;
 ss|} // 返回若干逆时针三角形 \{t[0][0], t[0][1], t[0][2], t[1][0], \dots\}
```

## 7 geometry3d

## 7.1 向量

```
struct p3 {
2 | db x, y, z;
3 | db norm() const { return x * x + y * y + z * z; }
4 | db abs() const { return std::sqrt(norm()); }
5 };
_{6} | p3 operator + (p3 x, p3 y){ return {x.x + y.x, x.y + y.y, x.z + y.z}; }
7 | p3 operator - (p3 x, p3 y){ return {x.x - y.x, x.y - y.y, x.z - y.z}; }
s|p3 operator * (p3 x, db y) { return {x.x * y, x.y * y, x.z * y}; }
9 p3 operator / (p3 x, db y) { return {x.x / y, x.y / y, x.z / y}; }
10 p3 operator * (p3 x, p3 y) { // 三维叉积需要更高的精度
| | | return {
| | | | | | x.y * y.z - x.z * y.y,
14 | X.X * Y.Y - X.Y * Y.X
15 | };
| 17 | db operator % (p3 x, p3 y) {    return x.x * y.x + x.y * y.y + x.z * y.z; }
|_{18}| p3 perpvec(p3 x) {
| return fabs(x.x) > fabs(x.z) ? p3{ x.y, -x.x, 0 } : p3{0, -x.z, x.y};
|20|} // 找到一个与给定向量垂直的向量
|<sub>21</sub>|db area(p3 a, p3 b, p3 c) {            return ((b - a) * (c - a)).abs();        } // 三角形面积两
| 22 | db volume(p3 d, p3 a, p3 b, p3 c) { // 四面体有向体积六倍
|_{23}| | return (d - a) % ((b - a) * (c - a));
```

```
7.2 平面

struct plane {
    | p3 n; db d; // n dot x = d
    | plane() {}
    | plane(p3 a, p3 b, p3 c) : n((c - a) * (b - a)) { d = n % a; }
    | db side(p3 x) const { return n % x - d; }
    | db dist(p3 w) const { return side(w) / n.abs(); }
    | p3 proj(p3 w) const { return w - n * (side(w) / n.abs()); }
}

7.3 直线
```

```
struct line3 {
    p3 d. o: // kd + o
  line3() {}
   | line3(p3 p, p3 q) : d(q - p), o(p) {}
    line3(plane p1, plane p2): d(p1.n * p2.n) { // 平面交出直线
     | o = (p2.n * p1.d - p1.n * p2.d) * d / d.norm();
    db dist(p3 p) const { return (d * (p - o)).abs() / d.abs(); }
     p3 proj(p3 p) const { return o + d * (d % (p - o)) / d.norm(); }//投影
     p3 relf(p3 p) const { return proj(p) * 2 - p; } // 对称
     p3 operator δ (const plane δ p) const { // 线与平面交
     | return o - d * p.side(o) / (p.n % d);
   }
13
14 };
15 db dist(line3 l1, line3 l2) {
     p3 n = l1.d * l2.d:
    if(n.abs() < eps) return l1.dist(l2.o);</pre>
     return abs((l2.o - l1.o) % n) / n.abs();
p3 closestOnL1(line3 l1, line3 l2) {
     p3 n2 = l2.d * (l1.d * l2.d);
     return l1.0 + l1.d * ((l2.0 - l1.0) % n2) / (l1.d % n2);
23 }
24 bool ispara(plane p1, plane p2){return(p1.n * p2.n).abs() < eps;}//判断是否平行
zs|bool ispara(line3 p1, line3 p2){return(p1.d * p2.d).abs() < eps;}//判断是否平行
26 bool isperp(plane p1, plane p2){return fabs(p1.n % p2.n) < eps;}//判断是否垂直
pool isperp(line3 p1, line3 p2){return fabs(p1.d % p2.d) < eps;}//判断是否垂直
28 line3 perpthrough(plane p, p3 o){return line3(o, o + p.n);}//过平面一点做垂线
```

## 7.4 凸包

#### 8 Misc

## 8.1 Pragma

```
#pragma GCC optimize("Ofast")
#pragma GCC optimize("unroll-loops")
#pragma GCC target("sse,sse2,sse3,ssse4,popcnt,abm,mmx,avx,avx2")
#pragma pack(1) // default=8
```

## 8.2 Barrett

```
| struct DIV {
| u64 x;
| void init(u64 v) { x = -1ull / v + 1; }
| }; // 带误差版本 x = -1ull/v;
| // ret=ans while x*(y-1)<2^64, ans-1<=ret<=ans while x<2^64
| u64 operator / (const u64 & x, const DIV & y) {
| return (u128) x * y.x >> 64;
| }
```

#### 8.3 LCS

```
int lim;
struct bitset {
    static const int B = 63;
    u64 a[N / B + 1];
    void set(int p) { a[p / B] |= 1ull << (p % B); }
    bool test(int p) { return a[p / B] >> (p % B) & 1; }
    void run(const bitset & o) {
        | u64 c = 1;
        | for(int i = 0;i < lim;++i) {
        | u64 x = a[i], y = x | o.a[i];
        | | x += x + c + (~y & (1ull << 63) - 1);
        | | a[i] = x & y, c = x >> 63;
        | | }
        | dp;
}
```

## 8.4 日期公式

#### 8.5 Xorshift

```
u64 xorshift(u64 x) { x ^= x << 13; x ^= x >> 7; x ^= x << 17; return x; }
u32 xorshift(u32 x) { x ^= x << 13; x ^= x >> 17; x ^= x << 5; return x; }
```

## 9 配置

#### 9.1 vimrc

```
set si ci ts=4 sw=4 nu cino=j1 backup undofile
syntax on
map<F9> <ESC>:!make %<<CR>
map<F10> <ESC>:!./%<<CR>
map<F4> <ESC>:!gdb %<<CR>
```

#### 9.2 bashrc

```
export CXXFLAGS='-g -Wall -fsanitize=address,undefined -Dzqj -std=gnu++20'
mk() { g++ -02 -Dzqj -std=gnu++20 $1.cpp -o $1; }
ulimit -s 1048576
ulimit -v 1048576
```

#### 9.3 对拍

需要 chmod +x

#### 9.4 编译参数

-D\_GLIBCXX\_DEBUG : STL debug mode -fsanitize=address : 内存错误检查 -fsanitize=undefined : UB 检查

#### 9.5 随机素数

979345007 986854057502126921 935359631 949054338673679153 931936021 989518940305146613 984974633 972090414870546877 984858209 956380060632801307

#### 9.6 常数表

n	lo	g <sub>10</sub> n	n!	C(n, n/2)	) LCM	$(1 \dots n)$	$P_n$
2	0.3010	2999	2		2	2	2
3	0.4771	7712125 6		:	3	6	
4	4 0.60205999		24	(	5	12	
5	5 0.69897000		120	10	)	60	
6			720	20	-	60 420	
7			5040	3:	-		
8	0.9030	8998	40320	70	-	840 2520	
9	0.9542	24251	362880	120	5		
10		1	3628800	25		2520	42
11	1.0413		39916800	462	2	27720	56
12	12 1.0791812 15 1.1760912		179001600	92		27720 360360	
15			1.31e12	643	5		
20	1.3010		2.43e18	18475	-	232792560	
25	1.39794001		1.55e25	520030		71144400	1958 5604
30	1.47712125		2.65e32	155117520	-	1.444e14	
$P_n$	3733840		20422650	966467 <sub>60</sub>	1905	190569292 <sub>100</sub>	
n	≤	10	100	1e3	1e4	1e5	1e6
max	$\omega(n)$	2	3	4	5	6	7
max	d(n)	4	12	32	64	128	240
π	(n)	4	25	168	1229	9592	78498
n	≤	1e7	1e8	1e9	1e10	1e11	1e12
max	$\omega(n)$	8	8	9	10	10	11
max	d(n)	448	768	1344	2304	4032	6720
$\pi$	(n)	664579	5761455	5.08e7	4.55e8	4.12e9	3.7e10
n	≤	1e13	1e14	1e15	1e16	1e17	1e18
max	$\omega(n)$	12	12	13	13	14	15
max	d(n)	10752	17280	26880	41472	64512	103680
$\pi(n)$ Prime number theorem: $\pi(x) \sim x$							)

## 10 注意事项

#### 10.1 测试项目

pbds tree, float128, int128, long double submit 命令, printfile, MLE ?= RE, pragma, axv2, python,

## 10.2 bugs

看数据范围(多测总和), 变量 shadow, 清空, long long, 数组大小, 模数, MLE?, 对拍记得看输出在不在变, 输出格式, inf 开小, 答案初值, STL 重构导致引用失效, 极端情况 (n=1)

## 11 tables

## 11.1 导数积分

$$\begin{array}{lll} \left(\frac{u}{v}\right)' = \frac{u'v - uo'}{v^2} & (\arctan x)' = \frac{1}{1+x^2} & (\arcsin x)' = \frac{1}{\sqrt{1+x^2}} \\ (ax)' = (\ln a)a^x & (\arccos x)' = -\frac{1}{1+x^2} & (\arccos x)' = \frac{1}{\sqrt{1-x^2}} \\ (\cot x)' = \sec^2 x & (\arccos x)' = -\frac{1}{x\sqrt{1-x^2}} & (\arctan x)' = \frac{1}{\sqrt{x^2-1}} \\ (\sec x)' = \tan x \sec x & (\arccos x)' = \frac{1}{x\sqrt{1-x^2}} & (\arctan x)' = \frac{1}{1-x^2} \\ (\arcsin x)' = \frac{1}{\sqrt{1-x^2}} & (\tanh x)' = \sech^2 x & (\arccos x)' = -\frac{1}{|x|\sqrt{1+x^2}} \\ (\arccos x)' = -\frac{1}{\sqrt{1-x^2}} & (\operatorname{sech} x)' = -\operatorname{sech} x \tanh x \\ (\operatorname{csch} x)' = -\operatorname{sech} x \coth x \\ \end{array}$$

## $ax^2 + bx + c(a > 0)$

1. 
$$\int \frac{\mathrm{d}x}{ax^2 + bx + c} = \begin{cases} \frac{2}{\sqrt{4ac - b^2}} \arctan \frac{2ax + b}{\sqrt{4ac - b^2}} + C & (b^2 < 4ac) \\ \frac{1}{\sqrt{b^2 - 4ac}} \ln \left| \frac{2ax + b - \sqrt{b^2 - 4ac}}{2ax + b + \sqrt{b^2 - 4ac}} \right| + C & (b^2 > 4ac) \end{cases}$$

2. 
$$\int \frac{x}{ax^2+bx+c} dx = \frac{1}{2a} \ln|ax^2+bx+c| - \frac{b}{2a} \int \frac{dx}{ax^2+bx+c}$$

## $\sqrt{\pm ax^2 + bx + c}(a > 0)$

1. 
$$\int \frac{dx}{\sqrt{ax^2 + bx + c}} = \frac{1}{\sqrt{a}} \ln |2ax + b + 2\sqrt{a}\sqrt{ax^2 + bx + c}| + C$$

2. 
$$\int \sqrt{ax^2 + bx + c} dx = \frac{2ax+b}{4a} \sqrt{ax^2 + bx + c} + \frac{4ac-b^2}{8\sqrt{a^3}} \ln|2ax + b| + 2\sqrt{a} \sqrt{ax^2 + bx + c}| + C$$

3. 
$$\int \frac{x}{\sqrt{ax^2 + bx + c}} dx = \frac{1}{a} \sqrt{ax^2 + bx + c} - \frac{b}{2\sqrt{a^3}} \ln|2ax + b| + 2\sqrt{a} \sqrt{ax^2 + bx + c}| + C$$

4. 
$$\int \frac{dx}{\sqrt{c_+ h x_- a x^2}} = -\frac{1}{\sqrt{a}} \arcsin \frac{2ax - b}{\sqrt{h^2 + 4ac}} + C$$

5. 
$$\int \sqrt{c + bx - ax^2} dx = \frac{2ax - b}{4a} \sqrt{c + bx - ax^2} + \frac{b^2 + 4ac}{8\sqrt{a^3}} \arcsin \frac{2ax - b}{\sqrt{b^2 + 4ac}} + C$$

6. 
$$\int \frac{x}{\sqrt{c+bx-ax^2}} dx = -\frac{1}{a} \sqrt{c+bx-ax^2} + \frac{b}{2\sqrt{a^3}} \arcsin \frac{2ax-b}{\sqrt{b^2+4ac}} + C$$

## $\sqrt{\pm \frac{x-a}{x-b}}$ 或 $\sqrt{(x-a)(x-b)}$

1. 
$$\int \frac{\mathrm{d}x}{\sqrt{(x-a)(b-x)}} = 2\arcsin\sqrt{\frac{x-a}{b-x}} + C \ (a < b)$$

2. 
$$\int \sqrt{(x-a)(b-x)} dx = \frac{2x-a-b}{4} \sqrt{(x-a)(b-x)} + \frac{(b-a)^2}{4} \arcsin \sqrt{\frac{x-a}{b-x}} + C, (a < b)$$

#### 三角函数的积分

1. 
$$\int \tan x \, \mathrm{d}x = -\ln|\cos x| + C$$

2. 
$$\int \cot x dx = \ln|\sin x| + C$$

3. 
$$\int \sec x dx = \ln \left| \tan \left( \frac{\pi}{4} + \frac{x}{2} \right) \right| + C = \ln \left| \sec x + \tan x \right| + C$$

4. 
$$\int \csc x dx = \ln \left| \tan \frac{x}{2} \right| + C = \ln \left| \csc x - \cot x \right| + C$$

$$5. \int \sec^2 x dx = \tan x + C$$

6. 
$$\int \csc^2 x \, \mathrm{d}x = -\cot x + C$$

7. 
$$\int \sec x \tan x dx = \sec x + C$$

8. 
$$\int \csc x \cot x dx = -\csc x + C$$

9. 
$$\int \sin^2 x dx = \frac{x}{2} - \frac{1}{4} \sin 2x + C$$

10. 
$$\int \cos^2 x dx = \frac{x}{2} + \frac{1}{4} \sin 2x + C$$

11. 
$$\int \sin^n x dx = -\frac{1}{n} \sin^{n-1} x \cos x + \frac{n-1}{n} \int \sin^{n-2} x dx$$

12. 
$$\int \cos^n x dx = \frac{1}{n} \cos^{n-1} x \sin x + \frac{n-1}{n} \int \cos^{n-2} x dx$$

13. 
$$\int \frac{dx}{\sin^n x} = -\frac{1}{n-1} \frac{\cos x}{\sin^{n-1} x} + \frac{n-2}{n-1} \int \frac{dx}{\sin^{n-2} x}$$

14. 
$$\int \frac{dx}{\cos^n x} = \frac{1}{n-1} \frac{\sin x}{\cos^{n-1} x} + \frac{n-2}{n-1} \int \frac{dx}{\cos^{n-2} x}$$

15.

$$\int \cos^m x \sin^n x dx$$

$$= \frac{1}{m+n} \cos^{m-1} x \sin^{n+1} x + \frac{m-1}{m+n} \int \cos^{m-2} x \sin^n x dx$$

$$= -\frac{1}{m+n} \cos^{m+1} x \sin^{n-1} x + \frac{n-1}{m+1} \int \cos^m x \sin^{n-2} x dx$$

16. 
$$\int \frac{\mathrm{d}x}{a+b\sin x} = \begin{cases} \frac{2}{\sqrt{a^2 - b^2}} \arctan \frac{a\tan \frac{x}{2} + b}{\sqrt{a^2 - b^2}} + C & (a^2 > b^2) \\ \frac{1}{\sqrt{b^2 - a^2}} \ln \left| \frac{a\tan \frac{x}{2} + b - \sqrt{b^2 - a^2}}{a\tan \frac{x}{2} + b + \sqrt{b^2 - a^2}} \right| + C & (a^2 < b^2) \end{cases}$$

17. 
$$\int \frac{dx}{a+b\cos x} = \begin{cases} \frac{2}{a+b} \sqrt{\frac{a+b}{a-b}} \arctan\left(\sqrt{\frac{a-b}{a+b}} \tan \frac{x}{2}\right) + C & (a^2 > b^2) \\ \frac{1}{a+b} \sqrt{\frac{a+b}{a-b}} \ln\left|\frac{\tan \frac{x}{2} + \sqrt{\frac{a+b}{b-a}}}{\tan \frac{x}{2} - \sqrt{\frac{a+b}{b-a}}}\right| + C & (a^2 < b^2) \end{cases}$$

18. 
$$\int \frac{\mathrm{d}x}{a^2 \cos^2 x + b^2 \sin^2 x} = \frac{1}{ab} \arctan\left(\frac{b}{a} \tan x\right) + C$$

19. 
$$\int \frac{dx}{a^2 \cos^2 x - b^2 \sin^2 x} = \frac{1}{2ab} \ln \left| \frac{b \tan x + a}{b \tan x - a} \right| + C$$

20. 
$$\int x \sin ax dx = \frac{1}{a^2} \sin ax - \frac{1}{a} x \cos ax + C$$

21. 
$$\int x^2 \sin ax dx = -\frac{1}{a}x^2 \cos ax + \frac{2}{a^2}x \sin ax + \frac{2}{a^3} \cos ax + C$$

22. 
$$\int x \cos ax dx = \frac{1}{a^2} \cos ax + \frac{1}{a} x \sin ax + C$$

23. 
$$\int x^2 \cos ax dx = \frac{1}{a}x^2 \sin ax + \frac{2}{a^2}x \cos ax - \frac{2}{a^3} \sin ax + C$$

## 反三角函数的积分 (其中 a > 0 )

1. 
$$\int \arcsin \frac{x}{a} dx = x \arcsin \frac{x}{a} + \sqrt{a^2 - x^2} + C$$

2. 
$$\int x \arcsin \frac{x}{a} dx = (\frac{x^2}{2} - \frac{a^2}{4}) \arcsin \frac{x}{a} + \frac{x}{4} \sqrt{x^2 - x^2} + C$$

3. 
$$\int x^2 \arcsin \frac{x}{a} dx = \frac{x^3}{3} \arcsin \frac{x}{a} + \frac{1}{9}(x^2 + 2a^2)\sqrt{a^2 - x^2} + C$$

4. 
$$\int \arccos \frac{x}{a} dx = x \arccos \frac{x}{a} - \sqrt{a^2 - x^2} + C$$

5. 
$$\int x \arccos \frac{x}{a} dx = \left(\frac{x^2}{2} - \frac{a^2}{4}\right) \arccos \frac{x}{a} - \frac{x}{4} \sqrt{a^2 - x^2} + C$$

6. 
$$\int x^2 \arccos \frac{x}{a} dx = \frac{x^3}{3} \arccos \frac{x}{a} - \frac{1}{9} (x^2 + 2a^2) \sqrt{a^2 - x^2} + C$$

7. 
$$\int \arctan \frac{x}{a} dx = x \arctan \frac{x}{a} - \frac{a}{2} \ln(a^2 + x^2) + C$$

8. 
$$\int x \arctan \frac{x}{a} dx = \frac{1}{2} (a^2 + x^2) \arctan \frac{x}{a} - \frac{a}{2} x + C$$

9. 
$$\int x^2 \arctan \frac{x}{a} dx = \frac{x^3}{3} \arctan \frac{x}{a} - \frac{a}{6}x^2 + \frac{a^3}{6} \ln(a^2 + x^2) + C$$

#### 指数函数的积分

1. 
$$\int a^x dx = \frac{1}{\ln a} a^x + C$$

2. 
$$\int e^{ax} dx = \frac{1}{a} a^{ax} + C$$

3. 
$$\int xe^{ax} dx = \frac{1}{a^2} (ax - 1)a^{ax} + C$$

4. 
$$\int x^n e^{ax} dx = \frac{1}{a} x^n e^{ax} - \frac{n}{a} \int x^{n-1} e^{ax} dx$$

5. 
$$\int x a^x dx = \frac{x}{\ln a} a^x - \frac{1}{(\ln a)^2} a^x + C$$

6. 
$$\int x^n a^x dx = \frac{1}{\ln a} x^n a^x - \frac{n}{\ln a} \int x^{n-1} a^x dx$$

7. 
$$\int e^{ax} \sin bx dx = \frac{1}{a^2 + b^2} e^{ax} (a \sin bx - b \cos bx) + C$$

8. 
$$\int e^{ax} \cos bx dx = \frac{1}{a^2 + b^2} e^{ax} (b \sin bx + a \cos bx) + C$$

9. 
$$\int e^{ax} \sin^n bx dx = \frac{1}{a^2 + b^2 n^2} e^{ax} \sin^{n-1} bx (a \sin bx - nb \cos bx) + \frac{n(n-1)b^2}{a^2 + b^2 n^2} \int e^{ax} \sin^{n-2} bx dx$$

10. 
$$\int e^{ax} \cos^n bx dx = \frac{1}{a^2 + b^2 n^2} e^{ax} \cos^{n-1} bx (a \cos bx + nb \sin bx) + \frac{n(n-1)b^2}{a^2 + b^2 n^2} \int e^{ax} \cos^{n-2} bx dx$$

## 对数函数的积分

$$1. \int \ln x dx = x \ln x - x + C$$

2. 
$$\int \frac{dx}{x \ln x} = \ln \left| \ln x \right| + C$$

3. 
$$\int x^n \ln x dx = \frac{1}{n+1} x^{n+1} (\ln x - \frac{1}{n+1}) + C$$

4. 
$$\int (\ln x)^n dx = x(\ln x)^n - n \int (\ln x)^{n-1} dx$$

5. 
$$\int x^m (\ln x)^n dx = \frac{1}{m+1} x^{m+1} (\ln x)^n - \frac{n}{m+1} \int x^m (\ln x)^{n-1} dx$$

## STL 积分/求和 (need std::)

1. 
$$\int_0^1 t^{x-1} (1-t)^{y-1} dt = beta(x, y) = \frac{\Gamma(x)\Gamma(y)}{\Gamma(x+y)}$$

2. 
$$\int_0^{+\infty} t^{num-1}e^{-t}dt = tgamma(num) = e^{lgamma(num)} = \Gamma(num)$$

3. 
$$\int_0^{phi} \frac{d\theta}{\sqrt{1-k^2\sin^2\theta}} = ellint_1(k, phi)$$

4. 
$$\int_0^{phi} \sqrt{1 - k^2 \sin^2 \theta} d\theta = ellint_2(k, phi)$$

5. 
$$\int_{num}^{+\infty} \frac{e^{-t}}{t} dt = -expint(-num)$$

6. 
$$\sum_{n=1}^{+\infty} n^{-num} = riemann\_zeta(num)$$

7. 
$$\frac{2}{\sqrt{\pi}} \int_0^{arg} e^{-t^2} dt = erf(arg)$$