

PSET 0

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1. Q1

Problem 1:

I understand I read the course website.

2. Q2

Problem 2:

To begin we know that the expectation is

$$\mathbb{E}[x] = \frac{1}{2} \cdot 2 + \frac{1}{2} \cdot \frac{1}{2} \cdot 3 + \frac{1}{2} \cdot \frac{1}{2} \cdot \mathbb{E}[\text{die}]$$

and since we know that $\mathbb{E}[\text{die}] = 3.5$ then

$$\mathbb{E}[x] = \frac{1}{2} \cdot 2 + \frac{1}{2} \cdot \frac{1}{2} \cdot 3 + \frac{1}{2} \cdot \frac{1}{2} \cdot 3.5$$

and evaluating this gives us

$$\boxed{\mathbb{E}[x] = 2.625}$$

Problem 3:

(1)

$$\mathbb{E}_{x \in \chi}[|f(x)|] \leq \max_{x \in \chi} f(x)$$

Not true by counterexample: Define $\chi := 0, 1$ and $P(0) = \frac{1}{2}$ and $P(1) = \frac{1}{2}$ with $f(x)$ defined as $0 \mapsto -5, 1 \mapsto -4$. Then

$$\mathbb{E}_{x \sim P}[|f(x)|] = \frac{1}{2}|-5| + \frac{1}{2}|-4| = \frac{1}{2}5 + \frac{1}{2}4 = 4.5 \not\leq \max_{x \in \chi} f(x) = \max\{-4, -5\} = -4$$

(2)

$$\mathbb{E}_{x \sim P}[f(x)] \leq \max_{x \in \chi} f(x)$$

Suppose $\forall x \in \chi, f(x) = c, c \in \mathbb{R}$. Then $\max_{x \in \chi} f(x) = c$ and $\mathbb{E}_{x \sim P}[f(x)] = \sum_{i=1}^{|X|} P(X[i])c = c$, where $|X|$ is the cardinality of the finite set χ and $X[i]$ is the i th number in an arbitrary ordering of the finite set X . Hence $\mathbb{E}_{x \sim P}[f(x)] = \max_{x \in \chi} f(x)$ which implies if this condition is true then this inequality holds.

Now suppose this condition is false. Then $\exists x_1, x_2 \in \chi$ s.t. $f(x_1) \neq f(x_2)$. Now wlog assume $f(x_1) < f(x_2)$ and $f(x_2) = \max_{x \in \chi} f(x)$. Then

$$\mathbb{E}_{x \sim P}[f(x)] = \sum_{i=1}^{|X|} P(X[i])f(X[i]) \leq P(x_1)f(x_1) + (1 - P(x_1)) \max_{x \in \chi} f(x) < \max_{x \in \chi} f(x)$$

Hence this is true in all cases

(3)

$$\mathbb{E}_{x \sim P}[f(x)] \leq \max_{x \in \chi} |f(x)|$$

Now since (b) implies

$$\mathbb{E}_{x \sim P}[f(x)] \leq \max_{x \in \chi} f(x)$$

And since $\forall x \in \chi, f(x) \leq |f(x)|$ implies $\max_{x \in \chi} f(x) \leq \max_{x \in \chi} |f(x)| \implies$

$$\mathbb{E}_{x \sim P}[f(x)] \leq \max_{x \in \chi} |f(x)|$$

Problem 4:

Wlog suppose $\max_{x \in \chi} f_1(x) = c_1, \max_{x \in \chi} f_2(x) = c_2$ and $c_1 \leq c_2$. Then at all $x_i \in \chi$ s.t. $f_2(x_i) = c_2$, $f_1(x_i) \leq c_1$. Suppose $\exists f_1(x_i)$ s.t. $f_1(x_i) = c_1$. Then $\max_{x_i} |f_2(x_i) - f_1(x_i)|$ is at least equal to $|c_2 - c_1| = |\max_{x \in \chi} f_2(x) - \max_{x \in \chi} f_1(x)|$ if not greater which satisfies the inequality. Suppose $\nexists f_1(x_i)$ s.t. $f_1(x_i) = c_1$. Then let $c_3 = \max_{x \in \chi \text{ s.t. } f_2(x)=c_2} f_1(x) \implies c_3 < c_1 \implies$ that

$$\max_{x_i} |f_2(x_i) - f_1(x_i)| = |c_2 - c_3| \geq |c_2 - c_1| = |\max_{x \in \chi} f_2(x) - \max_{x \in \chi} f_1(x)|$$

Problem 5:

Given that χ finite implies

$$\mathbb{E}_{x \sim P}[x] := \sum x_i P(x_i)$$

Implies that by nature of expectation linearity

$$\begin{aligned} |\mathbb{E}_{x \sim P}[f(x)] - \mathbb{E}_{x \sim P}[g(x)]| &= \left| \sum f(x_i)P(x_i) - \sum g(x_i)P(x_i) \right| = \\ \left| \sum (f(x_i) - g(x_i))P(x_i) \right| &= \mathbb{E}_{x \sim P}[f(x) - g(x)] \leq \mathbb{E}_{x \sim P}[|f(x) - g(x)|] \end{aligned}$$

3. Q3

Problem 6:

- True: $\text{Nullity}(A) = 0 \implies \text{Rank}(A) = n - 0 = n$ by Rank Nullity Theorem. Full rank implies invertibility given square matrix.
- False: If such a vector exists implies $\text{Nullity}(A) \neq 0 \implies \text{Rank}(A) < n \implies$ Not invertible.
- True: Range of $A = \mathbb{R}^n \implies$ Full rank which implies invertibility.

Problem 7:

- $\|Bv\|_\infty = \max_{1 \leq i \leq n} |r_i^T v|$: True: Each element in Bv is created by $r_i^T v$
- $\|Bv\|_\infty = \max_{1 \leq i \leq n} |c_i^T v|$: False: Each element in Bv is created by $r_i^T v$ not $c_i^T v$
- $\|Bv\|_\infty = \max_{1 \leq i \leq n} |r_i^T c_i|$: False: This definition does not have any notion of v in it.
- $\|Bv\|_\infty = \max_{1 \leq i \leq n} |e_i^T Bv|$: True: $e_i^T Bv$ is equivalent to plucking out the i th entry in the new vector Bv .

Problem 8:

$$\|v\|_\infty = \max_{1 \leq i \leq n} |v_i| \implies$$

Suppose

$$\forall i, v_i = c \in \mathbb{R} \implies \text{that the } j\text{th row of } |Bv| = \left| \sum_{i=0}^n B_{ji}c \right| = |c \sum_{i=0}^n B_{ji}| \leq |c\gamma| < |c| \implies$$

$$\|Bv\|_\infty \leq |\gamma c| < |c| = \|v\|_\infty$$

Now suppose the condition is false. Then $\exists v_a, v_b \in v$ s.t. wlog $v_a < v_b$ and $\max |v_i| = |v_b|$. Then the j th row of $|Bv| \leq |v_a B_{ja} + v_b(\gamma - B_{ja})| \leq |\gamma v_b| \leq v_b = \max_{1 \leq i \leq n} |v_i|$ which implies the statment is true.

Thus it is true for both cases.

Problem 9:

Now since we know that the L -infinity norm satisfies the reverse triangle inequality implies that

$$|||Iv||_{\infty} - ||Bv||_{\infty}| \leq ||Iv - Bv||_{\infty} = ||(I - B)v||_{\infty}$$

Now since we know that $||Iv||_{\infty} = ||v||_{\infty} > ||Bv||_{\infty}$ implies that $|||Iv||_{\infty} - ||Bv||_{\infty}| > 0 \implies ||(I - B)v||_{\infty} > 0$



Problem 10:

Problem 8 implies that

$$\|Bv\|_\infty \leq |\gamma c| = |\gamma| |c| = \gamma \|v\|_\infty$$

Therefore since $\|B\|_\infty \leq \gamma \|v\|_\infty \forall v \in \mathbb{R}^n$ than since $Bv = \lambda v \implies \|Bv\|_\infty = \|\lambda v\|_\infty$ implies that since $0 < \gamma < 1$ than all eigenvalues $\lambda \leq \gamma < 1$. This implies that the spectral radius of B is less than 1 which implies that the neumann series for operator B , $\sum_{i=0}^{\infty} B^k$ converges which implies $I - B$ is always invertible

4. Q4

Problem 11:

We know that

$$P = \begin{bmatrix} 0.2 & 0.2 & 0.2 \\ 0.6 & 0.2 & 0.2 \\ 0.3 & 0.6 & 0.1 \end{bmatrix}, \quad p_1 = \begin{bmatrix} 0.2 \\ 0.3 \\ 0.5 \end{bmatrix} \implies$$
$$p_2 = [0.2 \quad 0.3 \quad 0.5] \begin{bmatrix} 0.2 & 0.2 & 0.2 \\ 0.6 & 0.2 & 0.2 \\ 0.3 & 0.6 & 0.1 \end{bmatrix} = [0.37 \quad 0.40 \quad 0.23]$$

Problem 12:Knowing $p_2 \implies$

$$p_3 = p_2 P = \begin{bmatrix} 0.383 & 0.292 & 0.325 \end{bmatrix}$$

Problem 13:

Given

$$p_3 = p_2 P = p_1 P^2$$

Problem 14:

Given $k \geq 2$, $p_k = p_1 P^{k-1}$. We can calculate this efficiently with diagonalization given known constants for P