Computer Engg

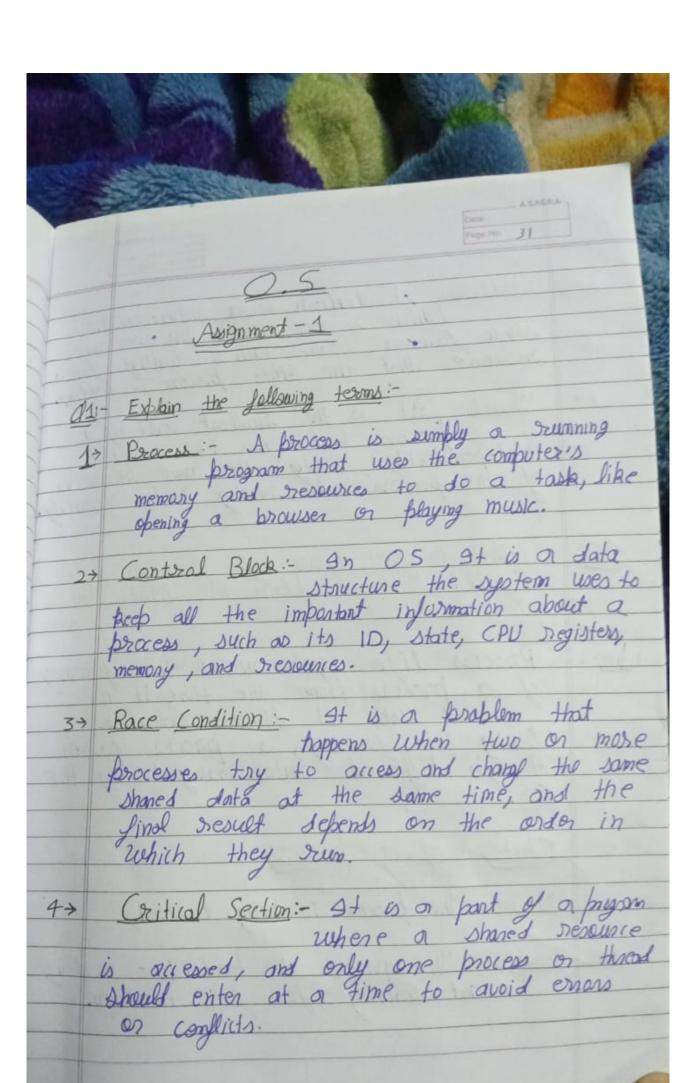
OS-Assignment-1

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Roll No: 02

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1>	14-08-25	CPUC- Assignment-1	1-30	
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		CPUC-Assignment -2		



	Date Page No. 3.2
5>	Deodlock: A deadlock is a situation in uniting for a stuck foreser because each is uniting for a resource that the other process is holding.
6>	11 - 1 coni + 101 a
02:	What do you mean by process life . Cycle & also explain the states of brocess.
Ans:-	Process life cycle means the fourney of a process from the time it is created until it finishes execution. During this journey, a process proses through different states managed by the OS.
241	States of a process:
	New: - The process is loaded into memory and writing for CPU to run.

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	Defection 33
3>	Running: - The process is being executed by the CPU.
47	Waiting / Blocked: - The process is paused, woiting for some event to
53	Terminated / Exit: - The process has finished its execution & is removed from memory.
Q3:-	Write short note on process synchonnization
Ans	At is a technique used in OS to ensure that multiple processes can work together safely when they share resources like memory, files, or data.
· No.	Without synchonization, processes may interje with each other, cousing problems such as sace conditions on in consistent results.
	Yo handle, the critical section problem is exclued using synchrozimation tools like semphones, muter locks and monitors, which make sure only one process use shared resource of a time.
1334	shared resource of a time.

	Date Page No. 5-4	
	34	
04>	Explain PCB in detail.	7
Am:-	Process Control Block [PCB]:-	
Vent la	9+ is a special data structure in the	8
	05 that stones all the light on identily	
of the	cond or necond file for each process,	
1	cond on necond file for each process, so the OS can keep track of it while scheduling and executing.	
	Contents of PCB:	
24	Process ID [PID]: - Unique no to identify each	
the state	process.	
27	Process state: - Current State of the process	
33	CPU Registers: Values of CPU registers: when the process is poused.	
4-3	Program Counter - Albert of the next	
	Program Counter: - Addresses of the next instruction to execute.	1
5-7	Memory information! - Base and limit	1
3/	Memory information! - Base and limit register, page tables, etc.	+
6-7	(PII Scheduling Antonmotion - Prionity, Achedulia	9
0	CPU Scheduling Anjonmotion - Priority, scheduling queure pointers, etc.	+

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7	
	Done: Page No. 35
7-7	INO Anjo: - Jist of I/O devices allocated and files opened.
	Accounting Information: - CPU usage, execution time, process number, etc.
•	Amportonie of PCB:
	The OS uses the PCB to switch between processes: 9t ensures that when a process serumes, it stants from where it left off.
Cali	without PCB, the OS weldsnit be able to monge multiple processes at once.
(15:	necessary conditions for deadlock Explain four deadlock is avoided.
Aw: -	A Seodlocka is a situation in am. OS where 2 or more processes one waiting for resources in such a way that none of them can continue, and they semain stuck forever.
,	Necessary Conditions for Sendlock:
	Deadlack occurs only if all four of these conditions hold at the same time:

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		ASABKA
		Paga No. 36
	1->	Mutual Exclusion: At least one resource
		must be held in a non-
	The state of the s	shareable mode.
3	2.9	Half and Wait: A process is holding at least one resource and
9	-	waiting for mone.
4	3→	No Preemption: - A resource convot be
		a process; it must be released valuntarily.
	4->	Circular libit: - A set of processes are waiting for each other ina
	Para la	cincular chain
	#>	How deadlock is avoided:
À	1>	Avoid Circular Writ: - 3mpose a fixed order of resource allocation so
		ciscular worting cannot occur.
	- Toping T	Avoid Hold and wait > Require processes to
	27	neguest als resources
9	. March.	Avoid Hold and wait :> Require processes to request all resources of selections which they hold before requesting new ones.
5		100 20 1-2 1-10
	3-7	Precaption: Albu resources to be taken
	1	given to another of needed.
-		
98	4	

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		#
		Disks: Page No. 3 7
		#
4>	Bankers Algorithm - A	resource offocation
	granting a resource of system in a safe else usit.	request will leave the
	system in a safe	State; if yes gont,
	The state of the s	TO THE PARTY OF TH
06:-	Differentiate blu moi	no lithic & Micro Rennol
January.		
ns:-	Monolithic Kernel.	MicroKernel 6
13	A single large Kernel.	43 A minimal Kernel
	Juhan & all OC APAULIE	Tichene only essential
	nun in trend mode.	and male and in
2->	Longe and complex	Tennel mode, while
		wer space.
34	Faster	h alarad
43	Difficult to extend on	23 Small and light-wight
id at	modify	3 > Slower.
c 2	ex: - Jinux, UNIX, MS+DOS	
3 /	et. sina, villajos	4-3 E00182 TO WILLIA OUT
		odd new fortunes.
		5> ex: - Minix, Mar OS X,
		QNX.

Communication? Explain its types. Interprocess Communication [IRC] in an OS to exchange data, share information, and co-andinate their activities Since processes run independently, IPC provides a way for them to communicate and stay synchronizes. Types of IPC ... Shared Memory: -A block of memory is shared between processes. Processes read/write data directly in this shored space. very fast but requires syn chronization to

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**	Date Page 76: 39
23	Menage Paning -
	No shared memory is required. Slawer than shared memory but simpler and ager. Process exchange information by rending
a .	ond oreceiving messages through the os. ex:- Client Server Communication.
08	Process AT BT Pl 2 8 P2 0 7 P3 1 9
	Calculate its TAT, WT LRT
Am:-	CT [Completion Time]: Jime when a process finishes.
	TAT [Jum-anaind Time]: CT-AT WITEWAITING time]: TAT-BT RT [Response time]: Jirst CPU start time - AT
1	RT CRESPONSE TIMES. 9150.

		STARBARE SALES
		Date: A Shoka
		Page No. 40
A >	Time quantum = 4	1000
		. 07
1 >	$t = 0-4 \rightarrow P2$ [arrive	ed at U.
- 11	· P2:7 > 3 left	ach the
2.1	1-12:03	when the same of the
27	t=4-8 > P3.	alian. In
	• P3: 9 > 5 left	The Williams
3>	t=8-12 > P1	
	· P1:8 > 4 left	
		at limited 7
43	0	1 31 JIMONES 2
	· CT [P2] =15	
	· Aucur	
53	t=15-19 > P3	
	· P3:5 > 1 left	
,		
63	t = 19-23 > P1 Eximishes]	
1	· CT[PI] = 23	
7	t=23-24 > P2 t 1, id 2	
	t=23-24 > P3[finishes] · (TEP3] = 27	
		-Sanat A a
\$->	Chart	D.fm.dinus
· Landing	0 4 8 12 15 10	1 22 22
	1 P2 P3 P1 P2 P3	PI 031
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