

Integral/Perfect formulations

We would like to know when a formulation is “ideal” or “perfect” (or “integral” in the pure integer case). In other words, when does solving the LP relaxation solve the underlying MIP too?

Definition 1. A perfect formulation of a mixed integer set $S \subseteq \mathbb{Z}^n \times \mathbb{R}^d$ is a system $Ax + Gy \leq b$ of linear inequalities such that $\text{conv}(S) = \{(x, y) \in \mathbb{R}^n \times \mathbb{R}^d \mid Ax + Gy \leq b\}$.

Definition 2. A polyhedron $P \subseteq \mathbb{R}^n$ is said to be *integral* if $P = \text{conv}(P \cap \mathbb{Z}^n)$.

How can we show that $Ax + Gy \leq b$ is a perfect formulation?

Theorem 1. Let $P := \{(x, y) \in \mathbb{R}^n \times \mathbb{R}^d \mid Ax + Gy \leq b\}$ be a rational polyhedron and let $S := P \cap (\mathbb{Z}^n \times \mathbb{R}^d)$. The following are equivalent.

- (i) P is a perfect formulation for S , i.e., $P = \text{conv}(S)$;
- (ii) Every nonempty face of P contains a point of S ;
- (iii) Every extreme point (i.e., vertex) of P belongs to S ;
- (iv) For every $(c, h) \in \mathbb{R}^n \times \mathbb{R}^d$ such that $z := \max\{cx + hy \mid (x, y) \in P\}$ is finite, there is a point $(x^*, y^*) \in S$ for which $z = cx^* + hy^*$.

Definition 3 (extreme point). A point x from a polyhedron P is an *extreme point* (a.k.a. *vertex*) of P if $\{x\}$ is a face of P . More generally, a point x belonging to a convex set P is said to be an *extreme point* if it cannot be expressed as a strict convex combination of distinct points from P .

Equivalently, you should not be able to write $x = \lambda x^1 + (1 - \lambda)x^2$ for $\lambda \in (0, 1)$; $x^1 \neq x^2$; $x^1, x^2 \in P$.

Example. Use Theorem 1(iii) to show that P , defined as the set of $(x, y) \in \mathbb{R}^m \times \mathbb{R}$ satisfying:

$$\begin{aligned} 0 &\leq y \leq 1 \\ 0 &\leq x_i \leq y, \quad i = 1, \dots, m. \end{aligned}$$

is a perfect formulation for

$$S := \left\{ (x, y) \in \mathbb{R}^m \times \{0, 1\} \mid \mathbf{0} \leq x \leq \mathbf{1}, \sum_{i=1}^m x_i \leq my \right\}.$$

Example. Use Theorem 1(iv) to show that P , defined as the set of $(x, y) \in \mathbb{R}^m \times \mathbb{R}$ satisfying:

$$\begin{aligned} 0 &\leq y \leq 1 \\ 0 &\leq x_i \leq y, \quad i = 1, \dots, m. \end{aligned}$$

is a perfect formulation for

$$S := \left\{ (x, y) \in \mathbb{R}^m \times \{0, 1\} \mid \mathbf{0} \leq x \leq \mathbf{1}, \sum_{i=1}^m x_i \leq my \right\}.$$

Total Unimodularity

Definition 4. A matrix A is said to be totally unimodular (TUM) if the determinant of every square submatrix belongs to $\{0, +1, -1\}$.

Example. Determine which of the following matrices are totally unimodular and why.

$$\begin{bmatrix} -1 & 2 & 0 \\ 1 & 0 & 1 \\ 0 & 1 & 1 \end{bmatrix} \quad \begin{bmatrix} -1 & 1 & 0 \\ 1 & 1 & 1 \\ 0 & 1 & 1 \end{bmatrix} \quad \begin{bmatrix} -1 & 1 & 0 \\ 1 & 0 & 1 \\ 0 & 1 & 1 \end{bmatrix}$$

Why do we care about TUM matrices?

Theorem 2 (Hoffman and Kruskal, 1956). Let A be an $m \times n$ integral matrix. The polyhedron $P(b) := \{x \mid Ax \leq b, x \geq 0\}$ is integral for every $b \in \mathbb{Z}^m$ if and only if A is totally unimodular.

Notice that TUM refers to the constraint matrix, and not to the formulation as a whole.

Proposition 1. The following are equivalent regarding a matrix A .

1. A is totally unimodular.
2. A^T is totally unimodular.
3. (A, I) is totally unimodular.

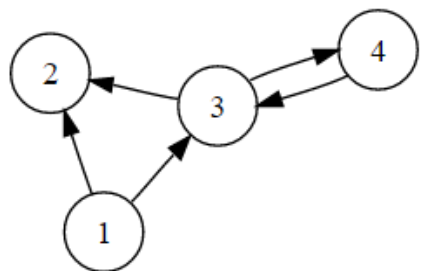
Theorem 3. Let A be an $m \times n$ integral matrix. The polyhedron $Q := \{x \mid c \leq Ax \leq d, l \leq x \leq u\}$ is integral for all integral vectors c, d, l, u if and only if A is totally unimodular.

Applications of TUM

Consider directed graph $G = (V, A)$. The (node-arc) incidence matrix B_G of G has dimension $n \times m$, where $n = |V|$ and $m = |A|$. There is a row for each vertex v and column for each arc $a = (u, w)$. The associated entry B_{va} takes a value as follows:

$$B_{va} = \begin{cases} -1 & \text{if } v = u \\ 1 & \text{if } v = w \\ 0 & \text{o.w.} \end{cases}$$

Example. Give the incidence matrix for this graph.



Theorem 4. *The incidence matrix of a directed graph is totally unimodular.*

By this theorem, many LP relaxations for network problems are integral. For example, consider the following formulation for the minimum cost flow problem.

Remark 1. *The constraint matrix associated with the flow balance constraints above is the graph's incidence matrix. Thus, the above formulation is integral when b and u are integral.*

Using TUM for non-network problems

Exercise. Use total unimodularity to show that P , defined as the set of $(x, y) \in \mathbb{R}^m \times \mathbb{R}$ satisfying:

$$\begin{aligned} 0 &\leq y \leq 1 \\ 0 &\leq x_i \leq y, \quad i = 1, \dots, m. \end{aligned}$$

is integral and is hence a perfect formulation for

$$S := \left\{ (x, y) \in \mathbb{R}^m \times \{0, 1\} \mid \mathbf{0} \leq x \leq \mathbf{1}, \sum_{i=1}^m x_i \leq my \right\}.$$