

Balanced Brackets

```

import java.util.Stack;
public class BalancedBrackets {

    public static String isBalanced(String brackets) {
        // O(1) time and Space
        if(brackets.length() == 0) {
            return "NO";
        }

        // O(1) time and Space
        Stack<Character> bracketStack = new Stack<>(); //this line is O(1) time and space

        int listSize = 0; // O(1) Time

        // O(N) time where N is the number of characters in the brackets String because each
loop
        // iterations results in a constant number of operations being performed and
        // N loop iterations are performed
        for(int i = 0; i < brackets.length(); i++) {
            char bracket = brackets.charAt(i); // O(1) space (allocates every loop iteration)
            if(bracket == '{' || bracket == '[' || bracket == '(') {
                // O(1) time
                bracketStack.push(bracket);
                listSize++;
            } else if(listSize == 0) {
                return "NO"; // no open brackets indicating it can never be balanced
            } else if (bracket == '}' && bracketStack.peek() == '{') {
                // O(1) time
                bracketStack.pop();
                listSize--;
            } else if (bracket == ']' && bracketStack.peek() == '[') {
                // O(1) time
                bracketStack.pop();
                listSize--;
            } else if (bracket == ')' && bracketStack.peek() == '(') {
                // O(1) time
                bracketStack.pop();
                listSize--;
            } else {
                // O(1) time
                return "NO"; // if a non balanced bracket or non bracket character appears
            }
        }

        // Because we have to push brackets into the stack, which requires an increasing amount
        // of memory, which results in O(M)space complexity where M is the number of open
        // brackets.

        // O(1) time and space
        if (listSize == 0) {
            return "YES";
        } else {
            return "NO";
        }
    }
}

```

Time Complexity: $O(N)$ where N is the number of characters in brackets

Space Complexity: $O(M)$ where M is the number of open brackets in brackets

DNA to RNA

```

import java.util.Stack;

public class DNatoRNA {
    public static String dnaToRna(String dna) {
        // O(1) time and space
        String rna = "";
        char base;

        // O(N^2) time, where N is the number of characters in the dna String, because
        // each string concatenation operation must copy the previous string character by
        // character which is O(N) time and the loop iterates N times.

        // O(N) space, where N is the number of characters in the dna String, because
        // the resulting RNA sequence will have a length of N after the outer loop
        // finishes. However, there are not N Strings being stored in memory, just one String
        // of length N, being continuously rewritten.
        for(int i = 0; i < dna.length(); i++){
            base = dna.charAt(i); // this line is O(1) time and space
            if(base == 'T') {
                rna = rna + "U"; // this line is O(N) and space
            } else {
                rna = rna + base; // this line is O(N) and space
            }
        }

        // O(1) time and space
        return rna;
    }
}

```

Time Complexity: $O(N^2)$ where N is the number of characters in the dna String

Space Complexity: $O(N)$ where N is the number of characters in the dna String

```

import java.util.LinkedList;
import java.util.Queue;
import java.util.HashMap;

public class RnaToProteins {
    public static String rnaToProteins(String rna) {
        // O(1) time and space
        HashMap<String, Character> lookUpTable = createLookUpTable();
        Queue<Character> protQueue = new LinkedList<>();
        Queue<Character> rnaQueue = new LinkedList<>();

        // O(N) time because the loop iterates N times where N is the number of characters
        // in the rna string, and inside the loop O(1) operations occur.
        // O(N) space where N is the number of characters in the rna string. Worst case, if
        // rna.length % 3 == 0, then the size of the rnaQueue will be of size N, and the size
of        // protQueue is N/3.

        for(int i = 0; i < rna.length(); i += 3) {
            // O(1) time because add() is O(1) and the loop always iterates 3 times
            if(i+3 > rna.length()) {
                protQueue.add('.');
                break;
            }
            for(int j = 0; j < 3; j++) {
                // O(1) time because add() is O(1) and the loop always iterates 3 times
                // O(1) space 3 elements are always added to the queue
                rnaQueue.add(rna.charAt(i + j));
            }
            char protein = protLookUpTable(rnaQueue, lookUpTable);
            protQueue.add(protein);
        }

        // O(N^2) time, where N is the number of characters in the rna String, because
        // each string concatenation operation must copy the previous string character by
        // character which is O(N) time. The loop iterates equal to the number of elements in
the        // protQueue, but the size of the protQueue in the worst case would be N/3.

        // O(N) space, where N is the number of characters in the rna String, because
        // the resulting result String will have a length of N after the outer loop
        // finishes. However, there are not N Strings being stored in memory, just one String
        // of length N, being continuously rewritten.
        String protSeq = "";
        int len = protQueue.size();
        for(int i = 0; i < len; i++) {
            protSeq += protQueue.remove();
        }

        return protSeq;
    }

    private static char protLookUpTable(Queue<Character> rnaQueue, HashMap<String, Character>
lookUpTable) {
        // O(1) time and space, 3 loop iterations
        String rnaSeq = "";
        for(int i = 0; i < 3; i++) {
            rnaSeq += rnaQueue.remove();
        }

        // Retrieving from a hashmap has an average time complexity of O(1), but

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        // there are only 64 key-value pairs in the hashmap, so worst case, given 63
collisions,
        // no more than 64 elements must be searched, so O(1) time and space.
        char protein = lookUpTable.get(rnaSeq);
        return protein;
    }
private static HashMap<String, Character> createLookUpTable() {
    // https://en.wikipedia.org/wiki/DNA\_and\_RNA\_codon\_tables
    // O(1) time and space, hashmaps always has 64 elements pushed to it irrelevant of any
input
    HashMap<String, Character> LookUpTable = new HashMap<>();
    LookUpTable.put("UUU", 'F');
    LookUpTable.put("UUC", 'F');
    LookUpTable.put("UUA", 'L');
    LookUpTable.put("UUG", 'L');
    LookUpTable.put("CUU", 'L');
    LookUpTable.put("CUC", 'L');
    LookUpTable.put("CUA", 'L');
    LookUpTable.put("CUG", 'L');
    LookUpTable.put("AUU", 'I');
    LookUpTable.put("AUC", 'I');
    LookUpTable.put("AUA", 'I');
    LookUpTable.put("AUG", 'M');
    LookUpTable.put("GUU", 'V');
    LookUpTable.put("GUC", 'V');
    LookUpTable.put("GUA", 'V');
    LookUpTable.put("GUG", 'V');
    LookUpTable.put("UCU", 'S');
    LookUpTable.put("UCC", 'S');
    LookUpTable.put("UCA", 'S');
    LookUpTable.put("UCG", 'S');
    LookUpTable.put("CCU", 'P');
    LookUpTable.put("CCC", 'P');
    LookUpTable.put("CCA", 'P');
    LookUpTable.put("CCG", 'P');
    LookUpTable.put("ACU", 'T');
    LookUpTable.put("ACC", 'T');
    LookUpTable.put("ACA", 'T');
    LookUpTable.put("ACG", 'T');
    LookUpTable.put("GCU", 'A');
    LookUpTable.put("GCC", 'A');
    LookUpTable.put("GCA", 'A');
    LookUpTable.put("GCG", 'A');
    LookUpTable.put("UAU", 'Y');
    LookUpTable.put("UAC", 'Y');
    LookUpTable.put("UAA", '.');
    LookUpTable.put("UAG", '.');
    LookUpTable.put("CAU", 'H');
    LookUpTable.put("CAC", 'H');
    LookUpTable.put("CAA", 'Q');
    LookUpTable.put("CAG", 'Q');
    LookUpTable.put("AAU", 'N');
    LookUpTable.put("AAC", 'N');
    LookUpTable.put("AAA", 'K');
    LookUpTable.put("AAG", 'K');
    LookUpTable.put("GAU", 'D');
    LookUpTable.put("GAC", 'D');
    LookUpTable.put("GAA", 'E');
    LookUpTable.put("GAG", 'E');
    LookUpTable.put("UGU", 'C');
    LookUpTable.put("UGC", 'C');
    LookUpTable.put("UGA", '.');

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```
LookUpTable.put("UGG", 'W');  
LookUpTable.put("CGU", 'R');  
LookUpTable.put("CGC", 'R');  
LookUpTable.put("CGA", 'R');  
LookUpTable.put("CGG", 'R');  
LookUpTable.put("AGU", 'S');  
LookUpTable.put("AGC", 'S');  
LookUpTable.put("AGA", 'R');  
LookUpTable.put("AGG", 'R');  
LookUpTable.put("GGU", 'G');  
LookUpTable.put("GGC", 'G');  
LookUpTable.put("GGA", 'G');  
LookUpTable.put("GGG", 'G');  
  
return LookUpTable;  
}
```

Time Complexity: $O(N^2)$ where N is the number of characters in the rna String

Space Complexity: $O(N)$ where N is the number of characters in the rna String

Infix to Postfix

```

import java.util.Stack;

public class InfixToPostfix {

    public static String infixToPostFix(String infix) {
        // O(1) time and space
        Stack<Character> operationStack = new Stack<>();
        Stack<Character> postFix = new Stack<>();
        char character;

        // Loop iterates N times where N is the number of characters in the infix String
        for(int i = 0; i < infix.length(); i++) {
            character = infix.charAt(i);
            if(character == '+' || character == '-' || character == '*' || character == '/' ||
character == '(' || character == ')') || character == '^') {
                // Check whether to push operation to result or add it to operation stack
                // Either O(M) or O(P) time which either is greater as defined below, and O(N)
                // time where N is the number of characters in the infix String
                updateStack(operationStack, character, postFix); // O(1)
            } else {
                postFix.push(character);
            }
        }

        // Loop iterates R times where R is the elements in the operationStack after the first
loop
        // Worst case R=N, if the infix string is all the characters + - ^ * / ( )
        // Add any operations that are not in the result already from highest precedence to
lowest
        while(!operationStack.isEmpty()) {
            postFix.push(operationStack.pop());
        }

        // O(N^2) time, where N is the number of characters in the infix String, because
        // each string concatenation operation must copy the previous string character by
        // character which is O(N) time. The loop iterates equal to the number of elements in
the
        // postFix stack, but the size of the postFix stack in the worst case would be N.

        // O(N) space, where N is the number of characters in the infix String, because
        // the resulting result String will have a length of N after the outer loop
        // finishes. However, there are not N Strings being stored in memory, just one String
        // of length N, being continuously rewritten.
        // Turn the stack to a string
        String result = "";
        for(char output : postFix) {
            result += output;
        }
        return result;
    }

    private static int quantifyOperation(char operation) {
        // O(1) time and space
        if(operation == '+' || operation == '-') {
            return 1;
        } else if(operation == '*' || operation == '/') {
            return 2;
        } else if(operation == '^') {
            return 3;
        } else {

```

```

        return 0;
    }
}

private static void updateStack(Stack<Character> operationStack, char operation,
Stack<Character> postFix) {
    if(operation == '('){
        // O(1) time and space
        operationStack.push(operation);
        return;
    } else if(operation == ')') {
        // O(M) time where M is the number of characters between the open and closed
        // parentheses
        while(!operationStack.isEmpty() && operationStack.peek() != '(') {
            postFix.push(operationStack.pop());
        }

        // O(1) time and space
        if(operationStack.peek() == '(') {
            operationStack.pop();
        }
        return;
    }
    // O(1) time and space
    if(operationStack.isEmpty()) {
        operationStack.push(operation);
    } // O(1) time and space
    } else if(operationStack.peek() == '(') {
        operationStack.push(operation);
    } // O(1) time and space
    } else if(quantifyOperation(operationStack.peek()) < quantifyOperation(operation)) {
        operationStack.push(operation);
    } else {
        // O(P) time where P is the number of operations in the stack that have a greater
        // than or equal to precedence as defined in quantifyOperation() and are not
        // succeeded with an open parenthesis.
        char currentChar = operation;
        while(!operationStack.isEmpty() && quantifyOperation(currentChar) <=
operationStack.peek() && currentChar != '(') {
            postFix.push(operationStack.pop());

            currentChar = !operationStack.isEmpty() ? operationStack.peek(): '_';
        }
        operationStack.push(operation);
    }
}

// O(N) space complexity where N is the number of characters in the infix String,
because
// the total amount of elements in the operation stack and the postFix stack
// will never exceed the number of characters in the infix String.
}

```

Time Complexity: $O(N^2)$ where N is the number of characters in the infix String

Space Complexity: $O(N)$ where N is the number of characters in the infix String

