# NetPoirot: Taking The Blame Game Out of Data Center Operations

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#### Datacenters can fail ...





A Microsoft Azure encountered major Rer customer who experienced a nine-day service disruption last month is furious over the software giant's lack of clear and transparent communication about the continuously for two

unable

online

Netflix

Reddit

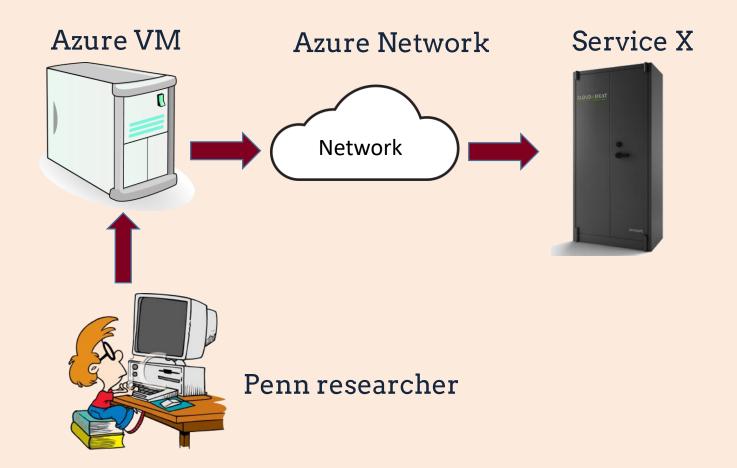
problems with two of his foll VMs running Microsoft's Iowa data center starting April 7.

One of the VMs was down straight days, and was then available The customer, who uses the intermittently for seven exp Azure Virtual Machines days. The other VM was in l Infrastructure-as-a-Service available intermittently for its

## Failures are disruptive

- They can cause significant user downtime
  - Loss of revenue for network providers
  - Lower QoE (Quality of Experience) for the users
- This introduces the need for debugging tools

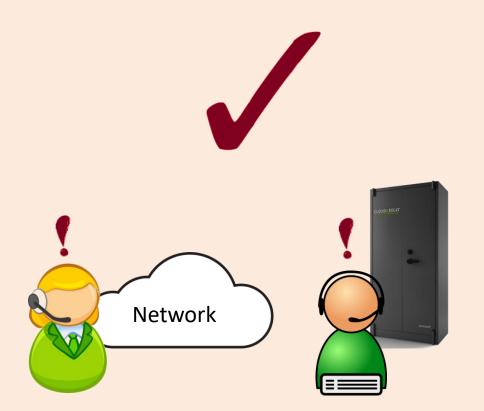
## Why is debugging hard?

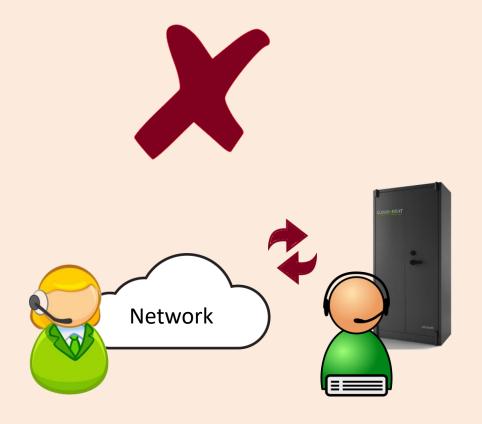


#### In the case of a failure...

Someone accepts responsibility

Each blames the other



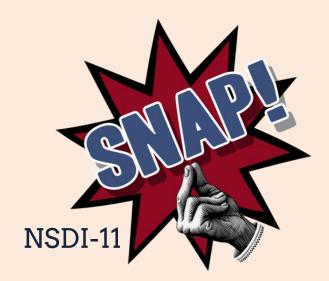


## A real example... Event X

- Azure hypervisors connect to a remote service
- If these connections fail, the VM has uncertain state
  - VM has to reboot
- Did the service fail, or was it the network?

#### Current tools are insufficient



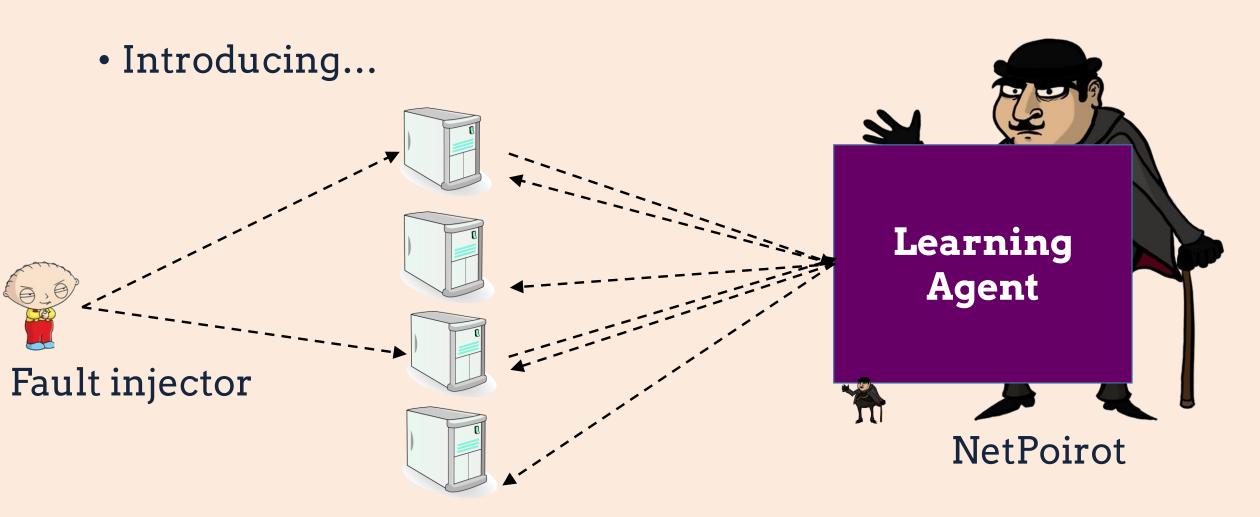






Sherlock SIGCOMM-07 NetMedic SIGCOMM-09

## Can we do better? (Overview)



## The monitoring agent

- Runs on all *clients* in our data center
- Captures and digests Windows TCP events
- Reports digests every 30 seconds
- Examples of metrics captured
  - Number of duplicate Acks
  - Number of timeouts
  - Time spent in zero window probing

## What is the TCP event digest?

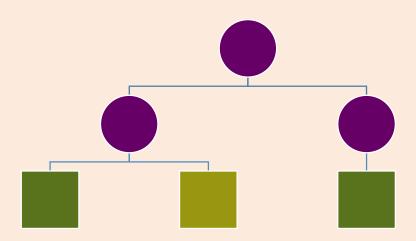
- We aggregate the captured TCP data into epoch digests
- Keep the min, max,  $10^{th}$ ,  $50^{th}$ ,  $95^{th}$  percentile, as well as mean and standard deviation across all connections in an epoch for each metric
  - Helps compare performance across the various connections

## Why do we think this can work?

- TCP observes the entire communication path
  - It goes through the client, the network, and the server
- It "sees" the failure no matter where it happened
- We know how network failures impact TCP
- How does it react to end point failures?
  - Hard to predict based on protocol design

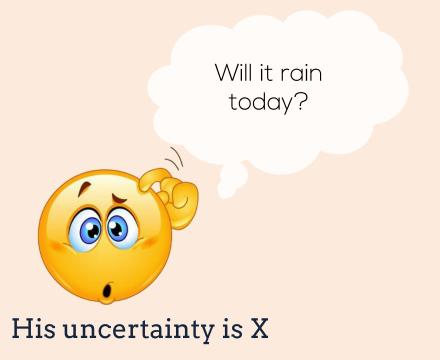
## To distinguish failures...

- We use variants of decision trees
- Other algorithms combine/manipulate features
  - Makes it hard to reason about why they arrive at a decision



#### Decision trees...

• Greedily pick features that maximize information gain

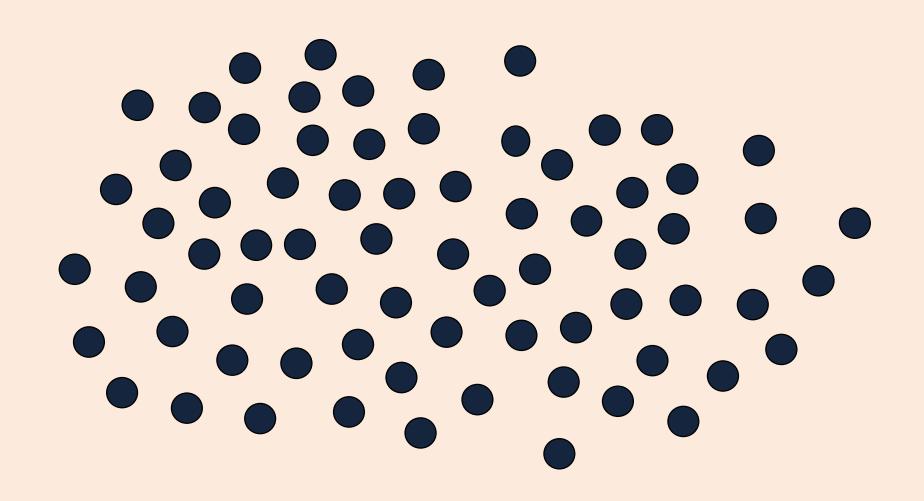


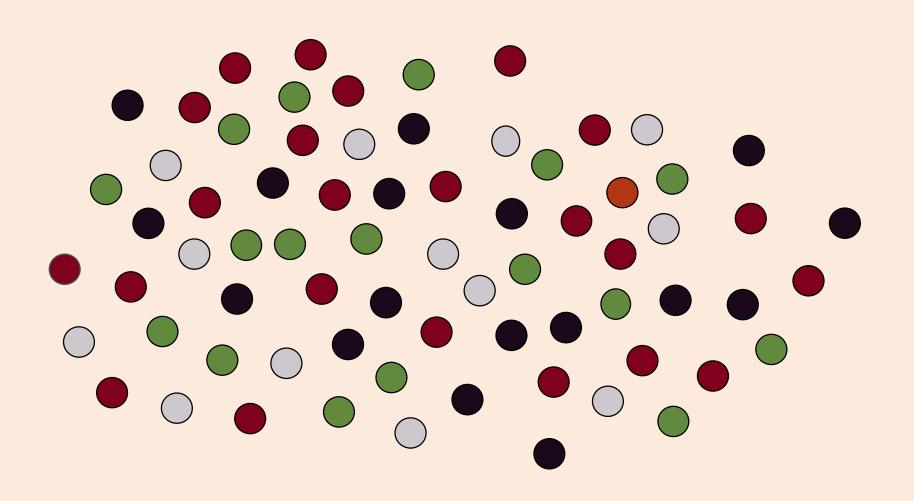


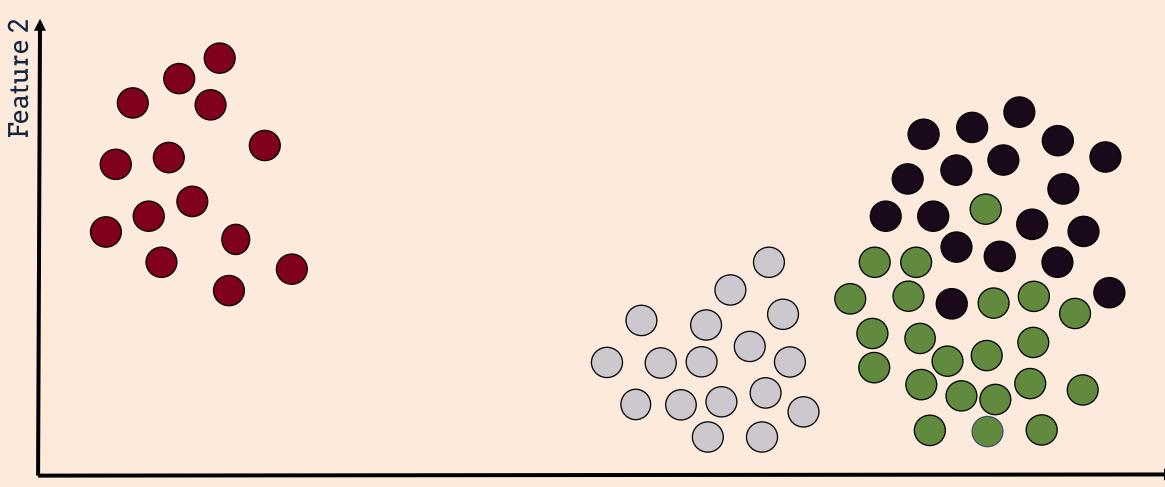
#### Decision trees...

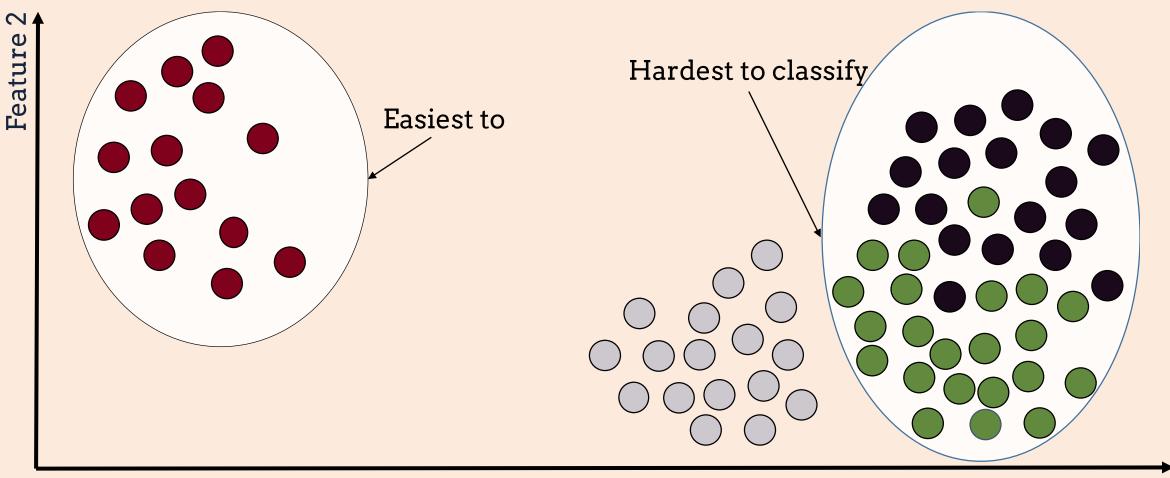
- Greedily pick features that maximize information gain
  - Pick the most "informative" features in each step



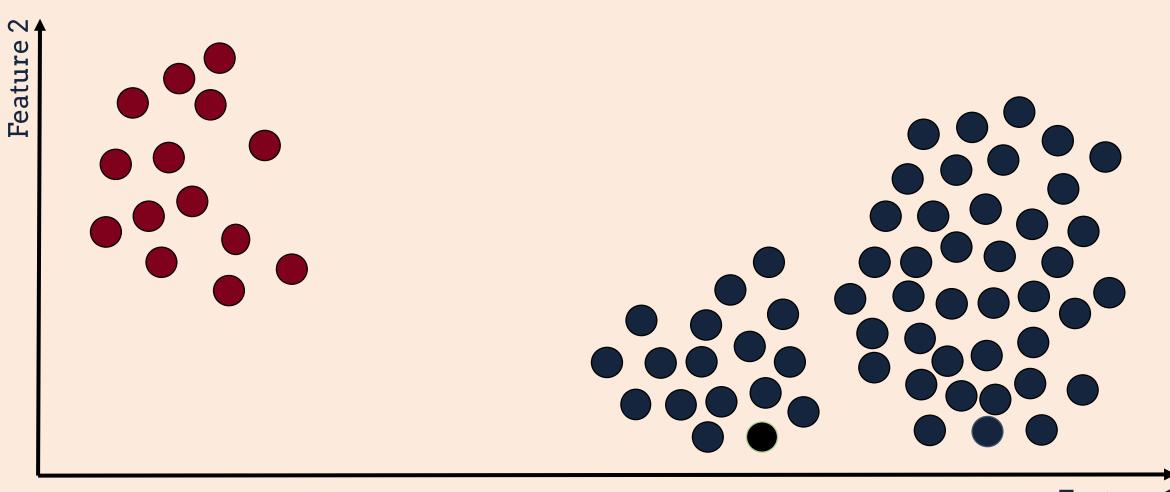




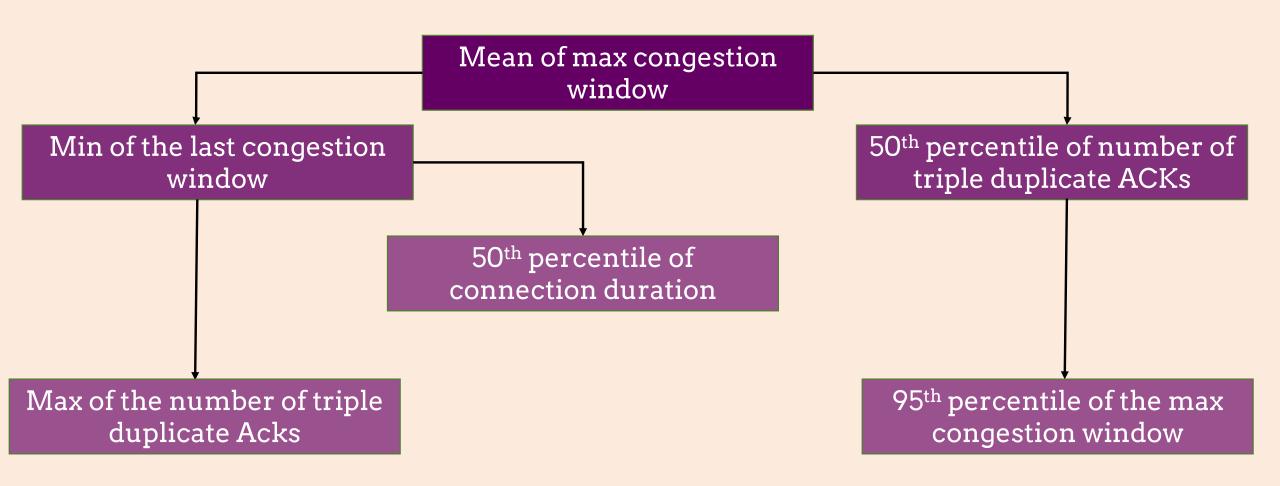




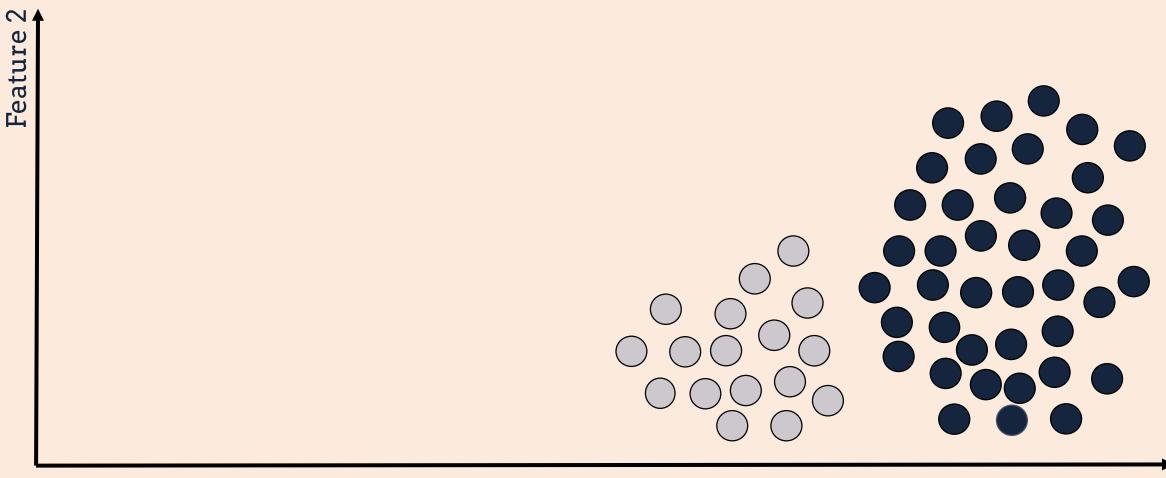
#### What we do to deal with this



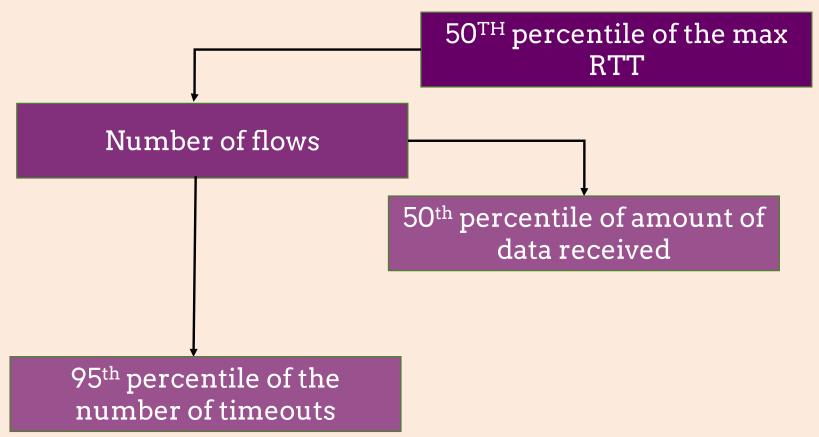
## Upper portion of an example tree...

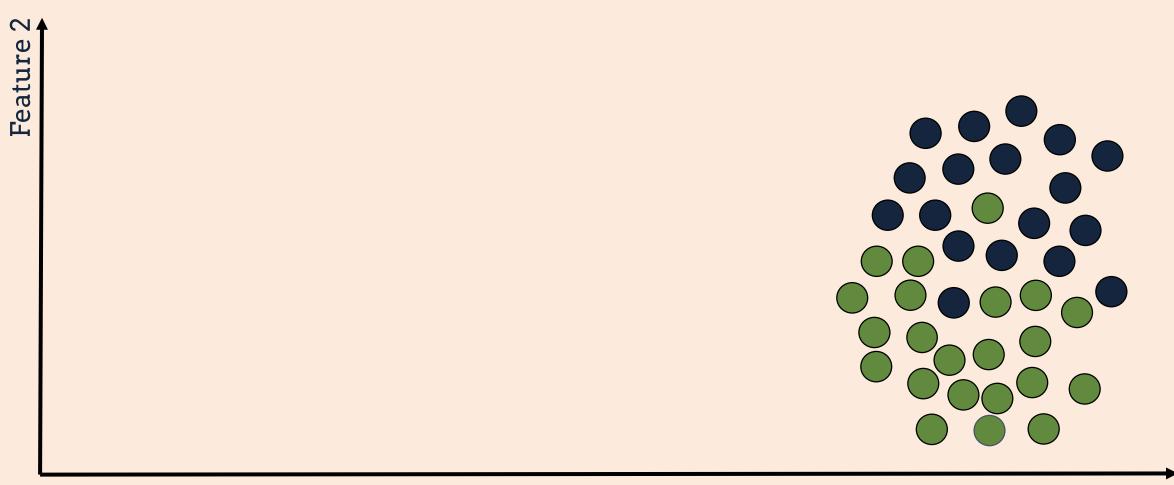


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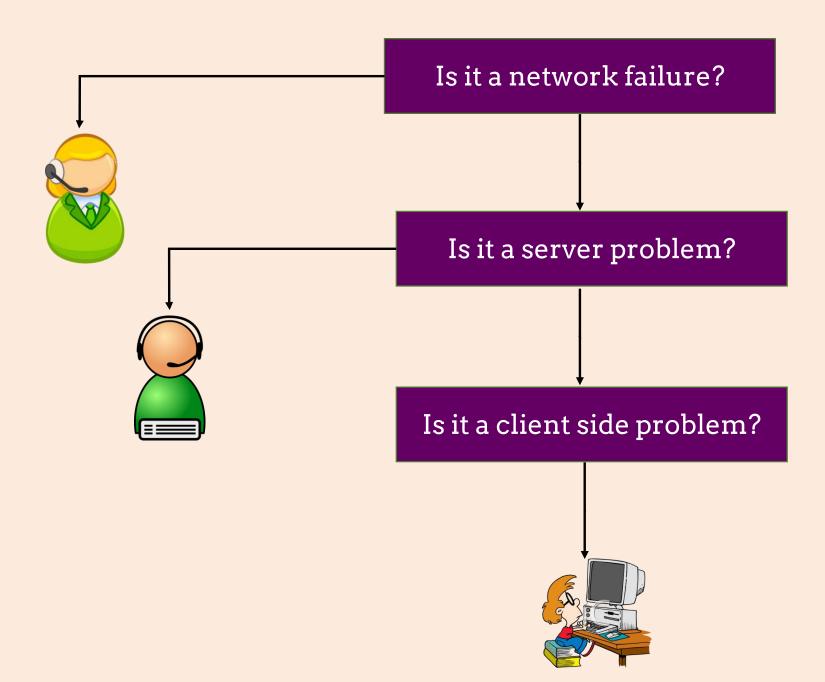
## Upper portion of an example tree...





The upper portion of an example tree...

Mean time spent in zero window probing 95th percentile of the ratio Number of flows of number of bytes posted to received 95th percentile of connection durations Minimum of the number of bytes received Number of flows



#### Other details

 We had to use random forest More stable If thirthroughpent < x: Send Dapen data en the Per application training samecconnections Normalize the data

#### What did we learn from all this?

- TCP sees everything, even at a single end point
  - Allows us to find who was responsible for a failure
- Failures in a group (Client/Server/Network) are similar
  - Makes individual failure classification more challenging
  - Helps NetPoirot be resilient to failures we haven't seen in the past
- The relationship between failures and TCP metrics is non-linear
  - Pearson correlation is low
- Two features suffice to describe each failure as observed by TCP
  - Two largest eigenvalues of the data matrix capture 95% of its variance

#### **Evaluation**

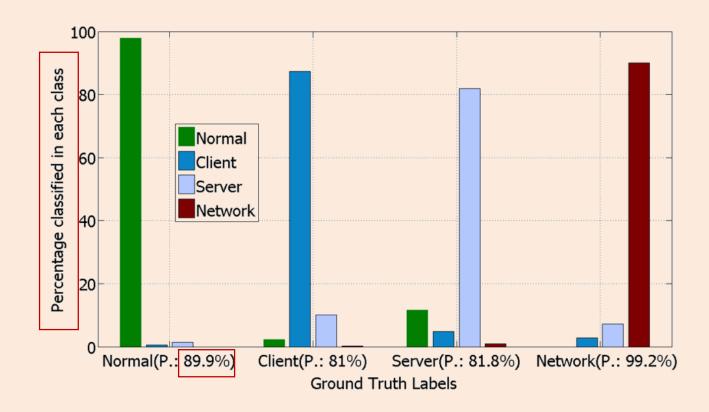
- What is the worst case performance?
  - Applications react to failures
  - Their reactions provide useful information
  - But what if this information is not available?
- What if we did not anticipate a failure type?
  - Dormant failures
  - Unknown failures

## How did we get labeled data?

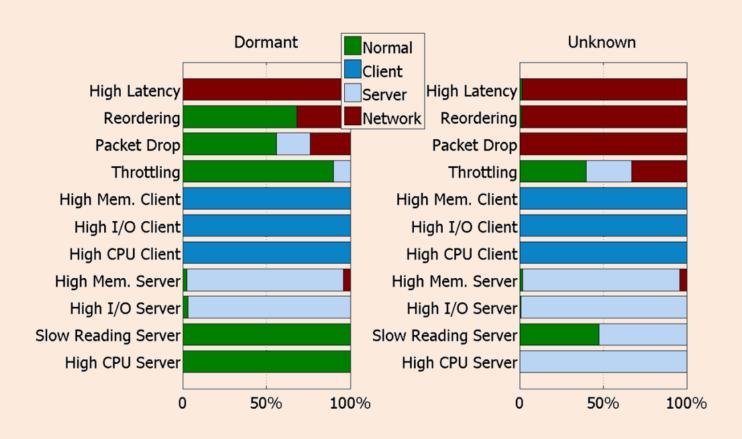
- We inject faults into the communication
- Over 6 months of data
- Examples:
  - High CPU load on the client
  - High I/O load on the server
  - Bandwidth throttling in the network
  - Packet reordering in the network

## Worse case application

• Only TCP statistics are used from the client side machine

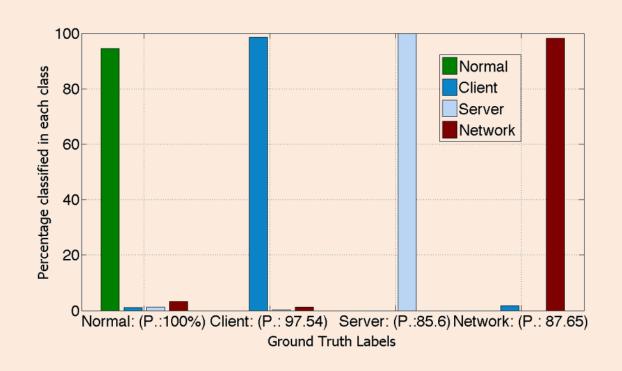


## What if we haven't seen the failure before?



## Performance on real applications

General label	Normal	Client	Networ k
Precisio n	97.78%	99.7%	100%
Recall	99.68%	98.25%	99.37
YouTube			



Event X

## Things we did not talk about

- Identifying the actual type of failure
- Sensitivity to machine location
- Aggregation vs per connection classification
- Sensitivity to failure duration
- Modifications to traditional cross validation required

#### What's next?

- Can we make this application independent?
  - Transfer learning
- Can the end point identify the device causing the failure?
  - Correlate information across clients

#### Conclusion

- TCP's reactions to network and endpoint failures are significantly different
- We can utilize these differences to find the entity that caused the failure



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