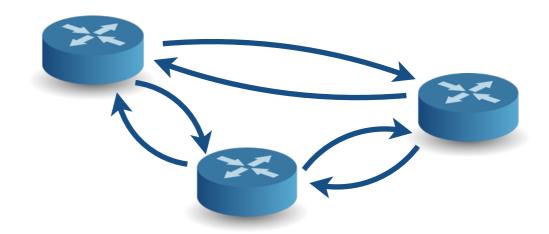
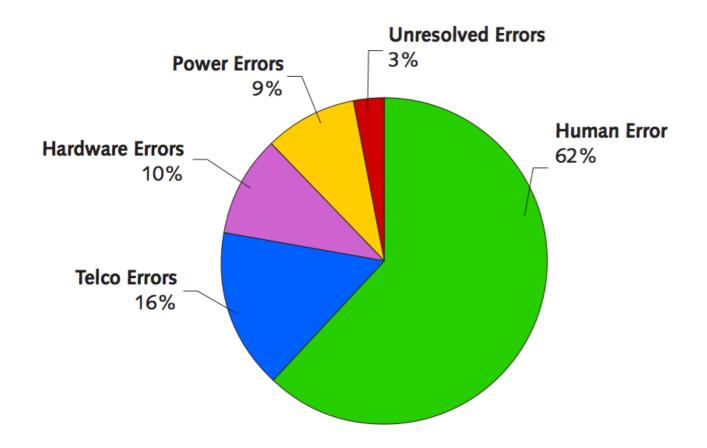
Don't Mind the Gap:

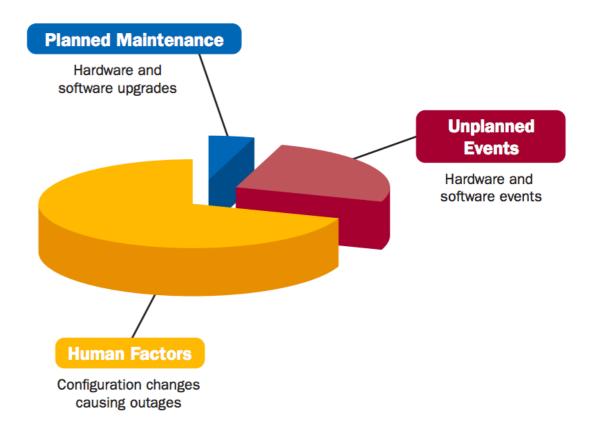
Bridging Network-wide Objectives and Device-level Configurations



Ryan Beckett (Princeton, MSR)
Ratul Mahajan (MSR)
Todd Millstein (UCLA)
Jitu Padhye (MSR)
David Walker (Princeton)







~60% of network downtime is caused by human error

-Yankee group 2002

50-80% of outages from configuration changes

-Juniper 2008





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NETWORKWORLD

YouTube/Pakistan incident: Could something similar whack your site?

Configuring BGP properly is key to avoidance, 'Net registry official says







By Carolyn Duffy Marsan

Network World | Mar 10, 2008 1:00 AM PT

In light of Pakistan Telecom/YouTube incident, Internet registry official explains how you can avoid having your web site victimized by such an attack.

When Pakistan Telecom blocked <u>YouTube</u>'s traffic one Sunday evening in February, the ISP created an international incident that wreaked havoc on the popular video site for more than two hours.

RIPE NCC, the European registry for Internet addresses, has conducted an analysis of what happened during Pakistan Telecom's hijacking of YouTube's traffic and the steps that YouTube took <u>to stop the attack</u>.

We posed some questions to RIPE NCC's Chief Scientist Daniel Karrenberg about the YouTube incident. Here's what he had to say:

How frequently do hijacking incidents like the Pakistan Telecom/YouTube incident happen?

Misconfigurations of iBGP (internal BGP, the protocol used between the routers in the same Autonomous System) happen regularly and are usually the result of an error. One such misconfiguration caused the Pakistan Telecom/YouTube incident. It appears that the Pakistan Telecom/YouTube incident was not an "attack" as some have labeled it, but a configuration error. (See Columnist Johna Till Johnson's take on the topic.)

What is significant about the YouTube incident?



Configuring BGP proper says







By Carolyn Duffy Marsan Network World | Mar 10, 2008 1:00 AM PT

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(https://www.thousandeyes.com) Product Solutions Customers (https://www.thousandeves.com/customers) Resources Login (https://app.thousandeyes.com) (https://www.thousandeyes Sign Up (https://www.thousandeves.com/signup) ← Blog Home (/)

> Time Warner Cable Outage Causes Widespread Routing and DNS Impacts Posted by Pete Anderson (https://blog.thousandeyes.com/author/pete/) on August 28, 2014

By now a lot of you have probably read about the Time Warner Cable (TWC) outage on August 27th, Yesterday morning I was greeted with a slew of alarms, names that wouldn't resolve, websites that wouldn't load and home office employees without any Internet access. It hadn't hit the news yet, but I could sense that a major outage was occurring and quickly opened up the ThousandEyes platform to get a handle on the situation.

Time Warner Outage

The alerts started coming in a little before 930 UTC (5:30 Eastern). I observed several different issues including inaccessible websites, DNS names failing to resolve, BGP reachability issues and agents losing access to the Internet. Companies that peer with Time Warner experienced degraded HTTP availability, affecting critical services such as supply chain portals (Figure 1).

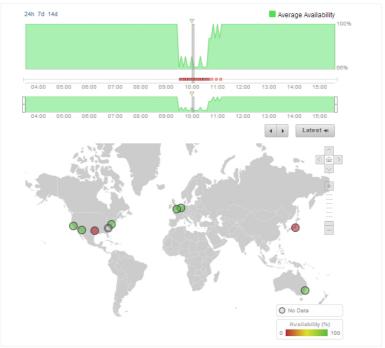


Figure 1: Supply chain portal with limited availability during Time Warner Cable outage

I could see right away that users and networks that connect through Time Warner were unable to reach the supply chain portal, indicating the issue was in the Time Warner network. In this case, I'd expect to see a brief service interruption while all the traffic re-routed through their other upstream ISP AT&T. However, the availability issues continued for the entire duration of the outage. I was surprised to still see issues, so I took a look at the path visualization view to figure out exactly where traffic was getting dropped on the way to this site Normally, two locations (Tokyo and Dallas) transit Road Runner (Time Warner) to reach this supply chain portal, while the rest go through AT&T (Figure 2).



Configuring BGP proper says







By Carolyn Duffy Marsan Network World | Mar 10, 2008 1:00 AM PT

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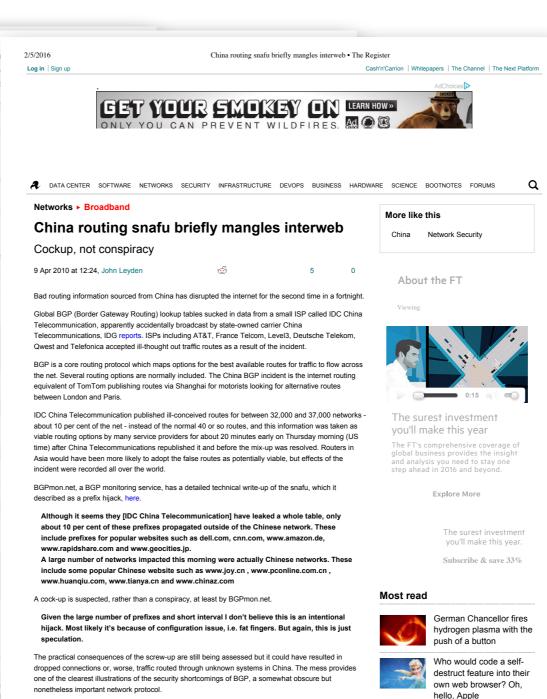
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Who wants a quad-core

4.2GHz, 64GB, 5TB SSD

The China BGP global routing represents a rare but not unprecedented mix-up in global internet traffic

management. For example, just two weeks ago bad routing data resulted in the redirection of Chilean

internet traffic through a DNS (Domain Name System) server in China, as explained in a detailed post-

mortem by internet monitoring firm Renesys here. Bad BGP routing information from Pakistan caused

http://www.theregister.co.uk/2010/04/09/china_bgp_interweb_snafu/



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IDC China Telecommunication published ill-conceived routes about 10 per cent of the net - instead of the normal 40 or so ro viable routing options by many service providers for about 20 r time) after China Telecommunications republished it and befor

Telecommunications, IDG reports, ISPs including AT&T, Fran

Qwest and Telefonica accepted ill-thought out traffic routes as BGP is a core routing protocol which maps options for the best

the net. Several routing options are normally included. The Chi

equivalent of TomTom publishing routes via Shanghai for moto

BGPmon.net. a BGP monitoring service, has a detailed technidescribed as a prefix hijack, here

Asia would have been more likely to adopt the false routes as

incident were recorded all over the world

Although it seems they [IDC China Telecon about 10 per cent of these prefixes propagated outside include prefixes for popular websites such as dell.com, www.rapidshare.com and www.geocities.ip.

A large number of networks impacted this morning wer include some popular Chinese website such as www.joy www.huanqiu.com, www.tianya.cn and www.chinaz.com

A cock-up is suspected, rather than a conspiracy, at least by B

Given the large number of prefixes and short interval I d hijack. Most likely it's because of configuration issue, i.e

The practical consequences of the screw-up are still being ass dropped connections or, worse, traffic routed through unknown one of the clearest illustrations of the security shortcomings of I nonetheless important network protoco

The China BGP global routing represents a rare but not unpred management. For example, just two weeks ago bad routing da internet traffic through a DNS (Domain Name System) server in mortem by internet monitoring firm Renesys here. Bad BGP re

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Internet-Wide Catastrophe - Last Year - Dvn Research | The New Home Of Renesvs



COMMENTS (0)

TODD UNDERWOOD



2/5/2016

3 ₹ m



O DECEMBER 24, 2005

☐ ENGINEERING





One year ago today TTNet in Turkey (AS9121) pretended to be the entire Internet. And unfortunately for the rest of the Internet, many large network providers believed them (or at least believed them in part). As far as anyone knows, it was a mistake, not a malicious act. But the consequences were far from benign: for several hours a large number of « Previous Story

able to reach a large number of Internet sites. re can take a look at what happened, and whether 1 the intervening time.

10rning 2004, TTNet (AS9121) started announcing

...I VIEWS: 3038

OUTAGES

Search

Q

Targeted Internet Traffic Misdirection

Popular Authors Archives

NOVEMBER 19, 2013



Egypt Leaves the

Threat:

CONTENT

IANUARY 27, 2011



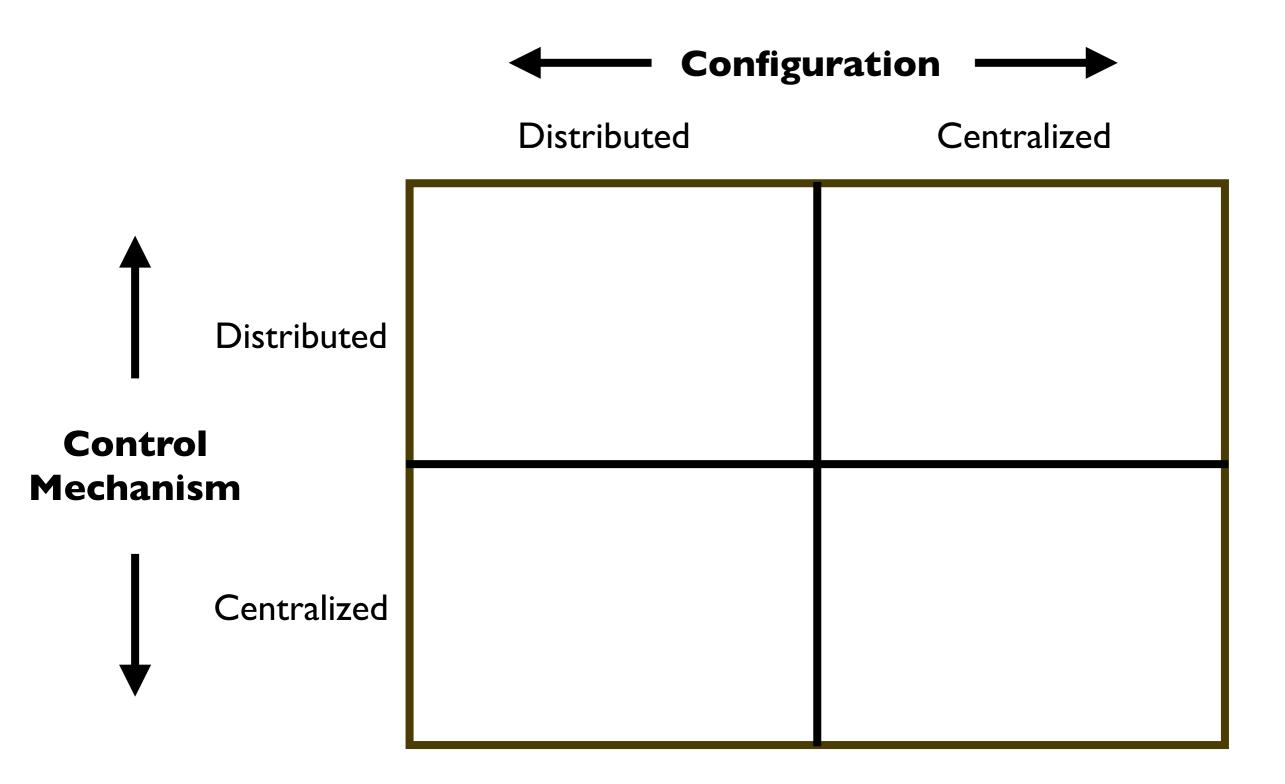
Internet

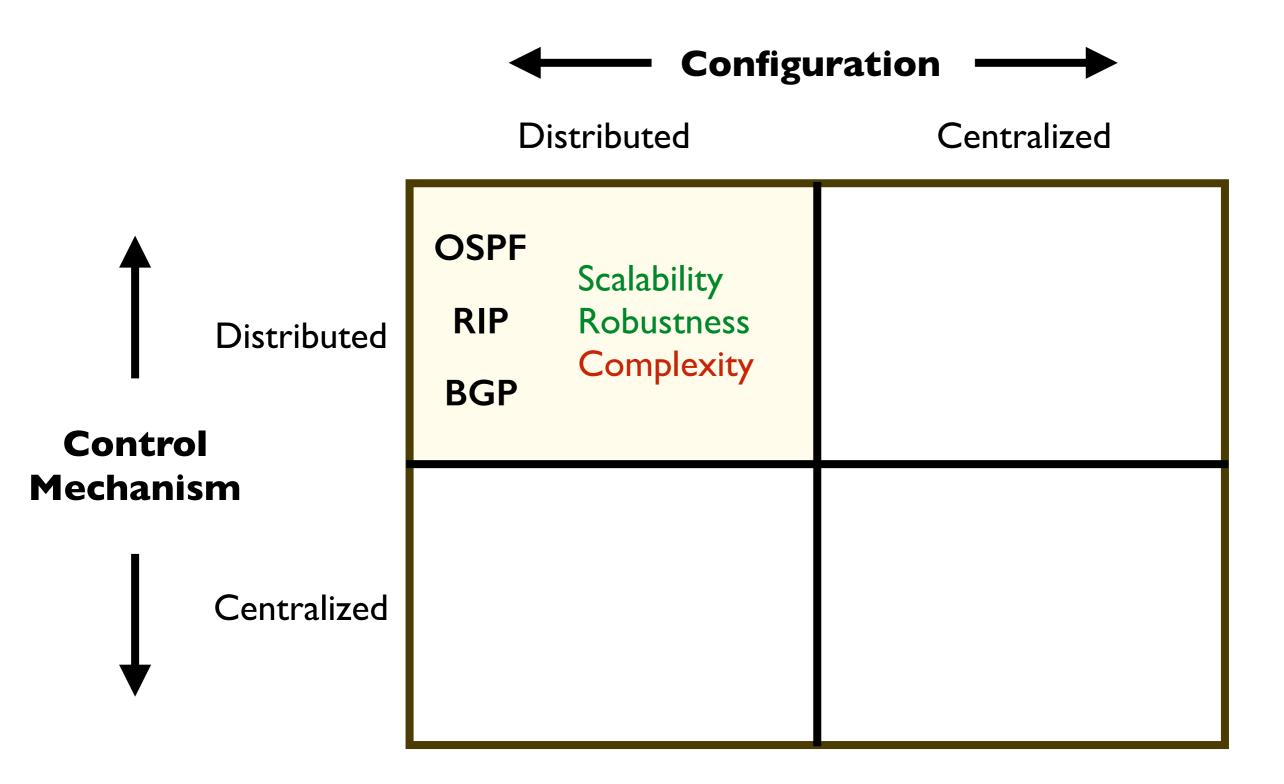
Next Story

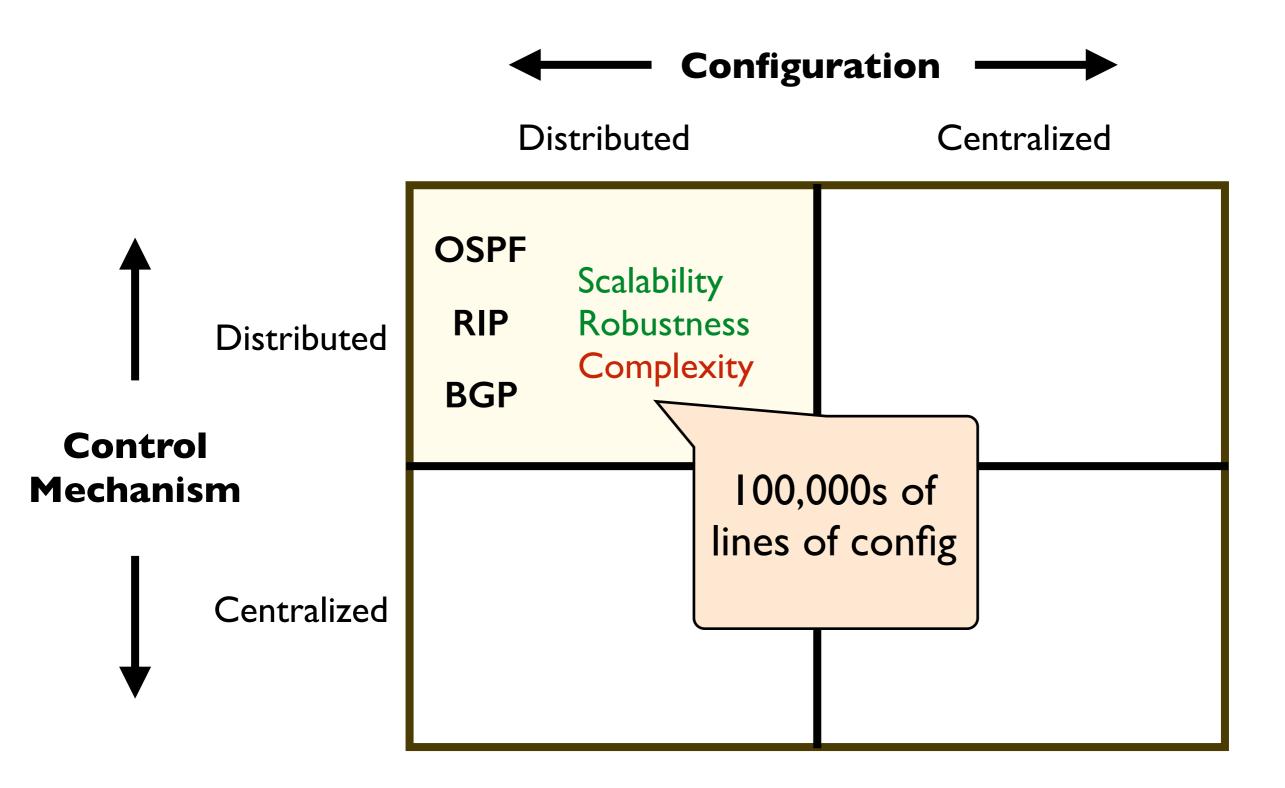
Possible N

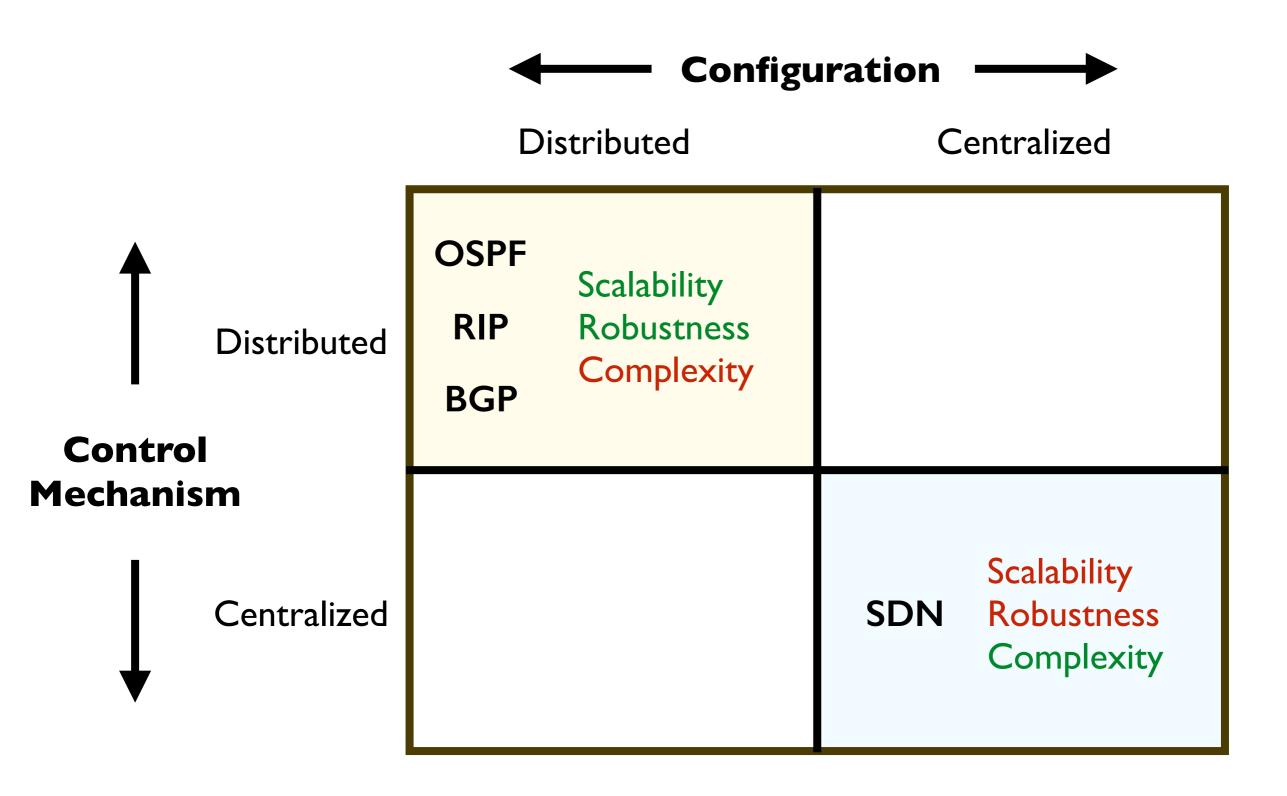
AUGUST 13

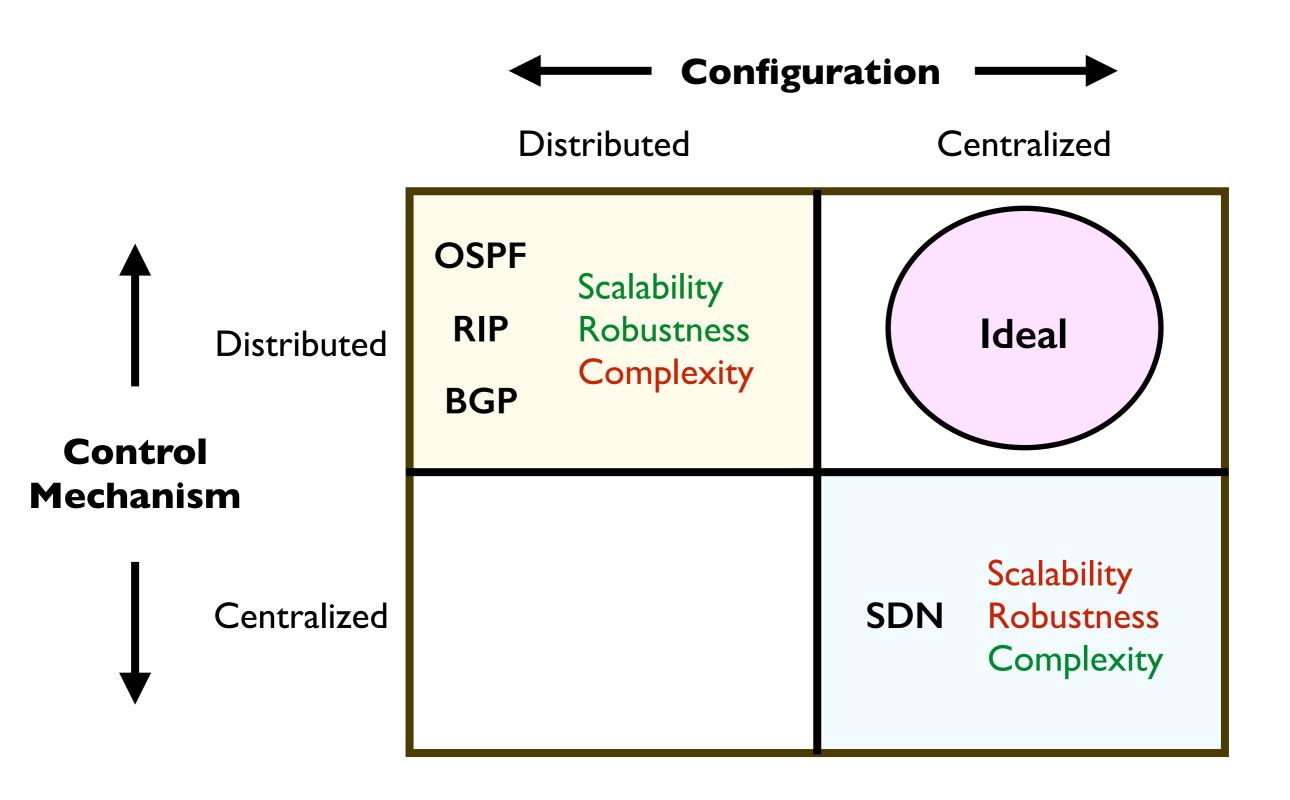
http://research.dyn.com/2005/12/internetwide-nearcatastrophela



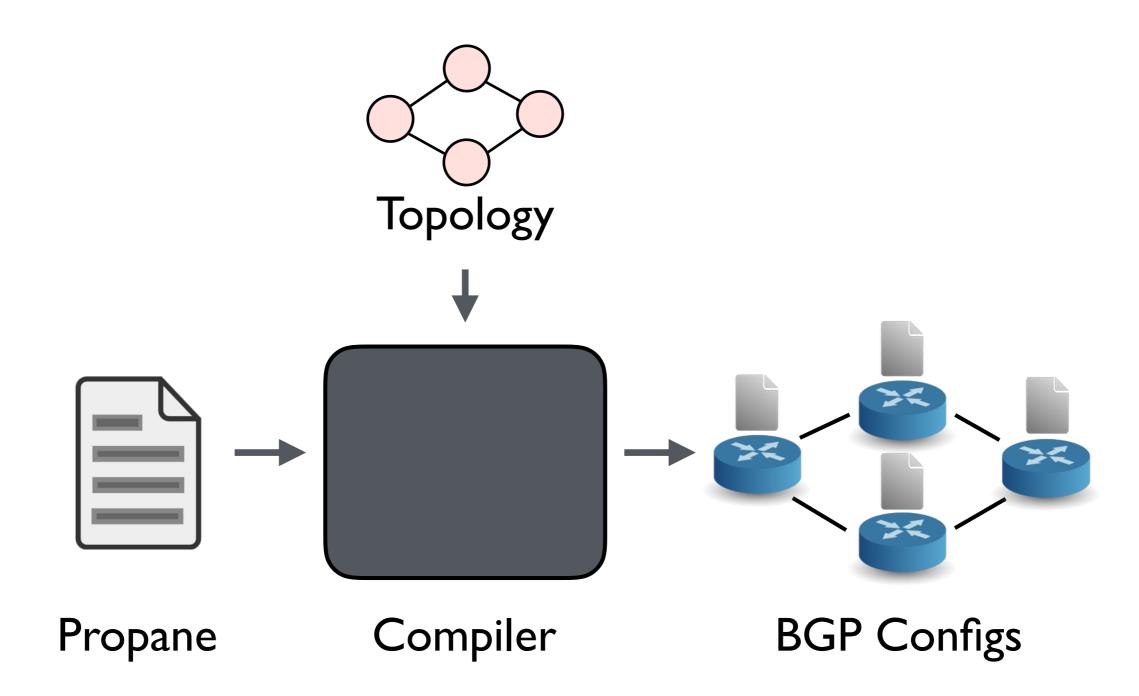








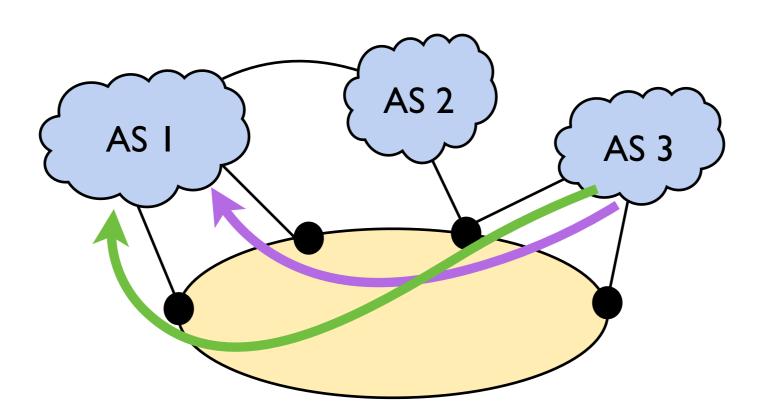
Propane Overview



Propane System

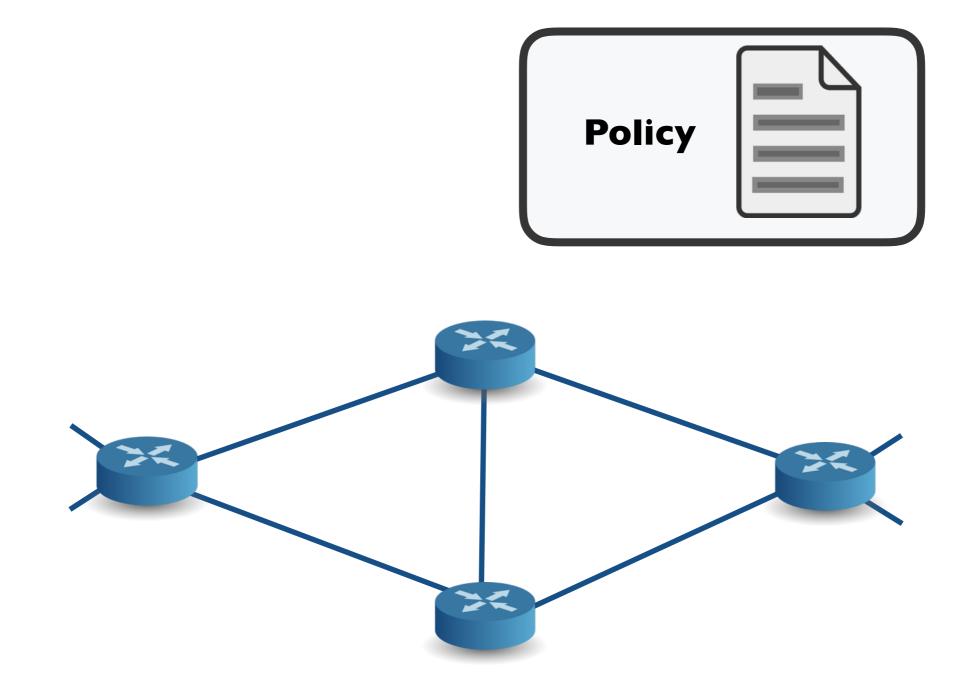
I) Language for expressing network-wide objectives with:

- Path constraints and preferences in case of failures
- Uniform abstractions for intra- and inter-domain routing



Propane System

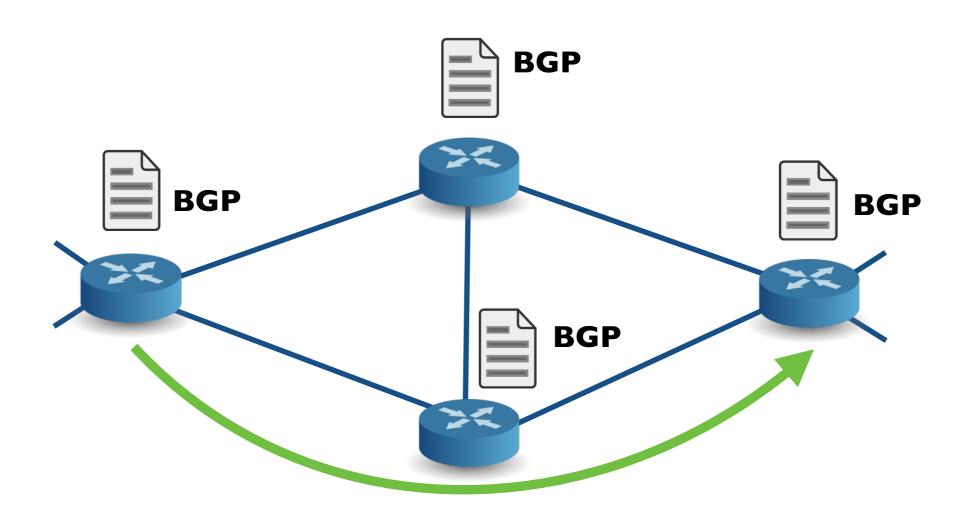
2) Compiler for a purely distributed implementation



Propane System

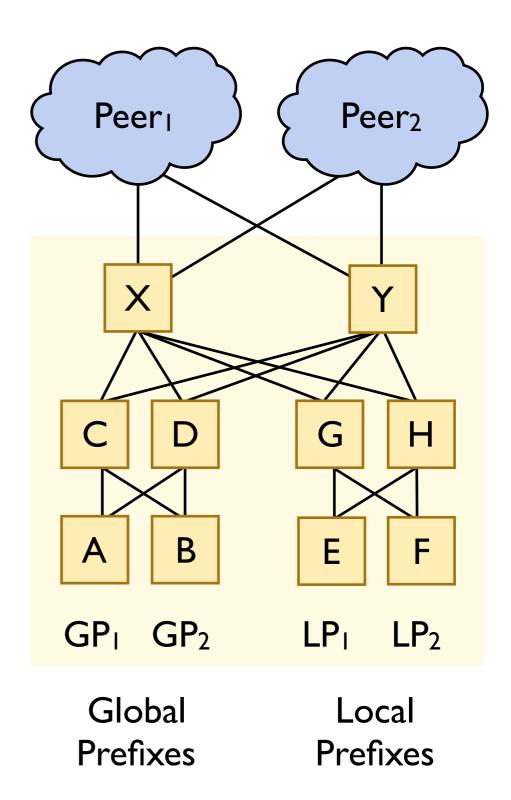
2) Compiler for a purely distributed implementation

- Generate BGP configs for each router
- Compiler guarantees policy-compliance for all failures



Goals

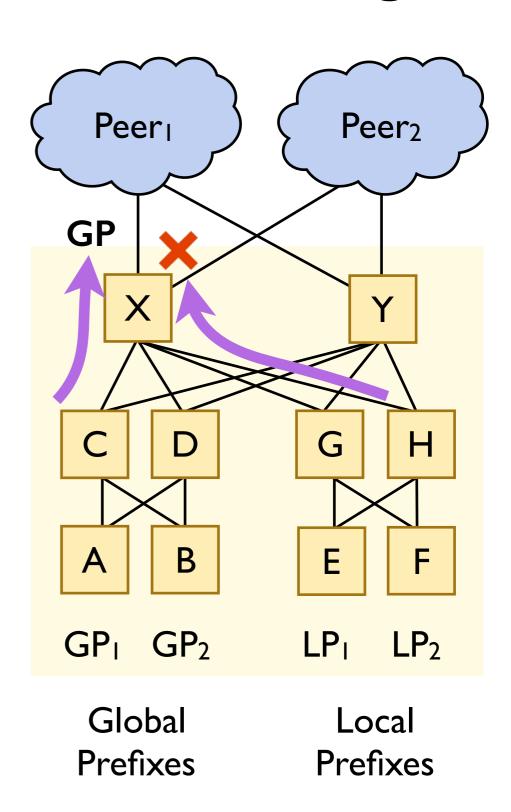
- Local prefixes reachable only internally
- Global prefixes reachable externally
- Aggregate global prefixes as GP
- Prefer leaving through Peer₁ over Peer₂
- Prevent transit traffic between peers



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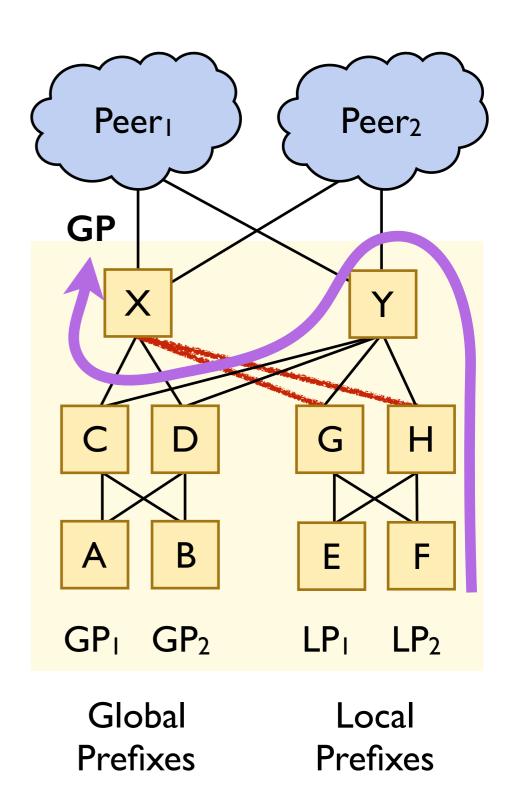
- Don't export from G, H to external
- Aggregate externally as GP



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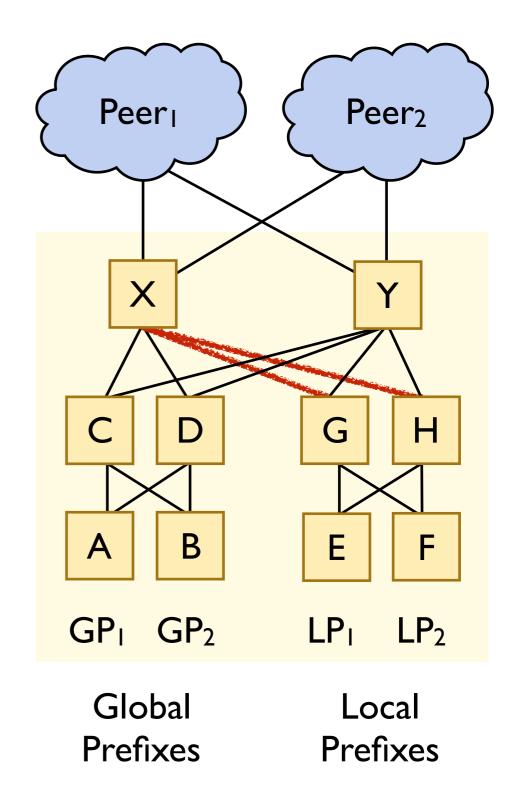
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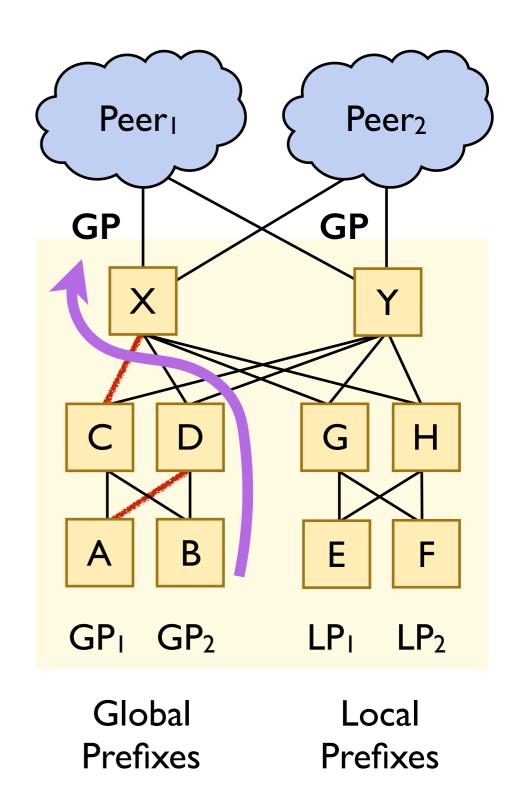
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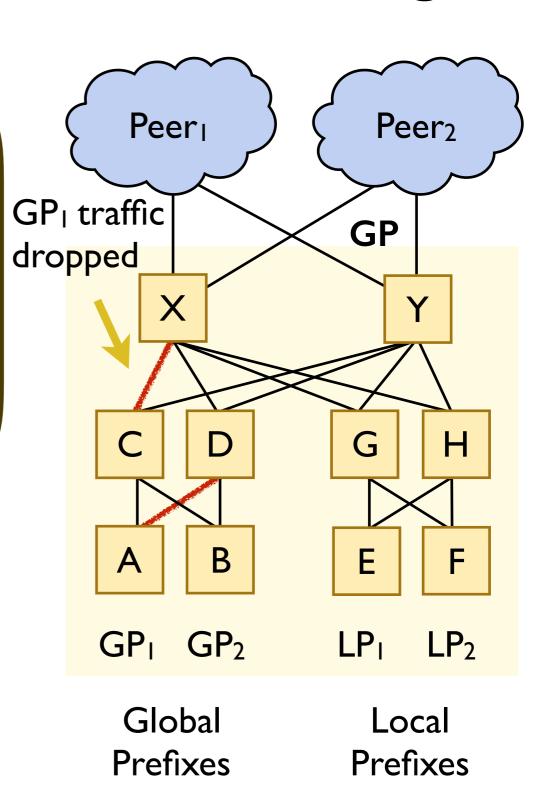
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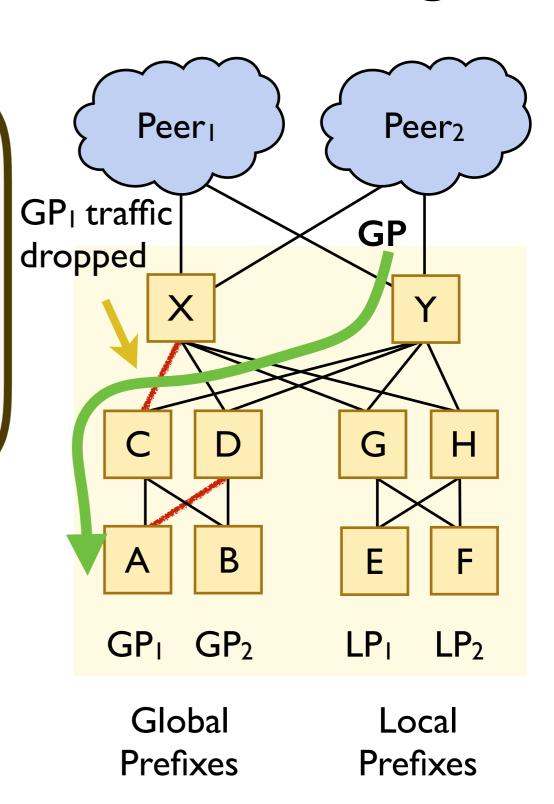


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Configuration Attempt

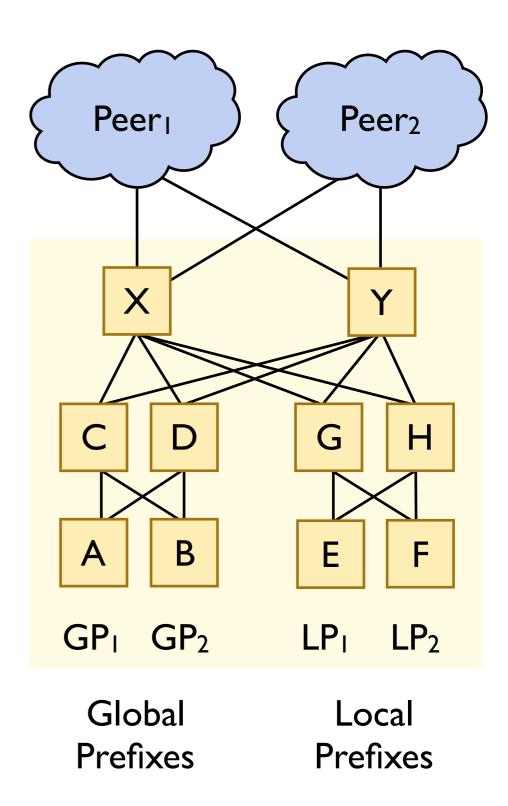
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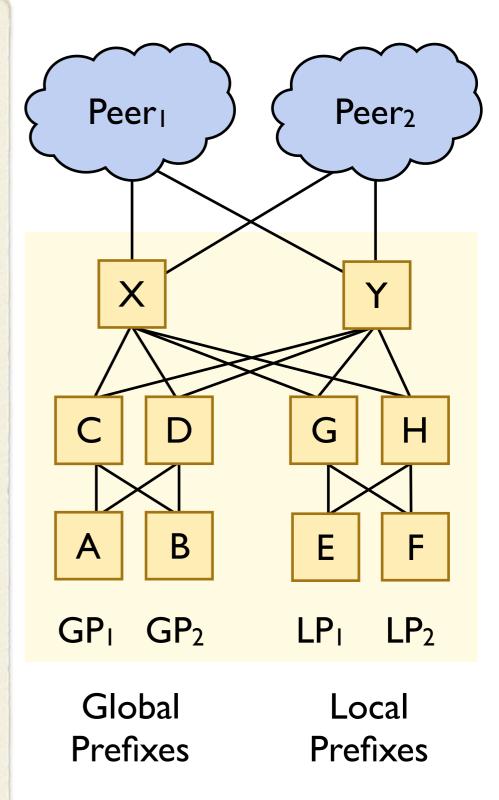
Aggregation-Induced Black Hole!

Goals

- Local prefixes reachable only internally
- Global prefixes reachable externally
- Aggregate global prefixes as GP
- Prefer leaving through Peer₁ over Peer₂
- Prevent transit traffic between peers

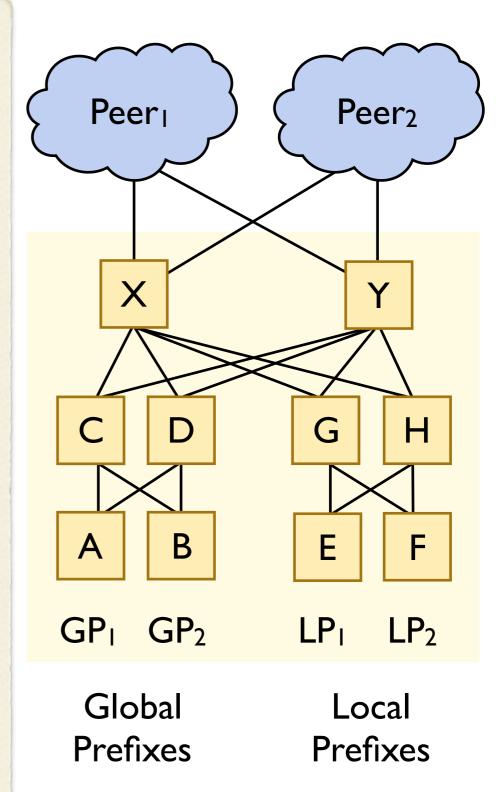


```
define Destination =
  {GP1 => end(A)
   GP2 => end(B)
   LP1 => end(E)
   LP2 => end(F)
   true => exit(Peer1 >> Peer2) }
```

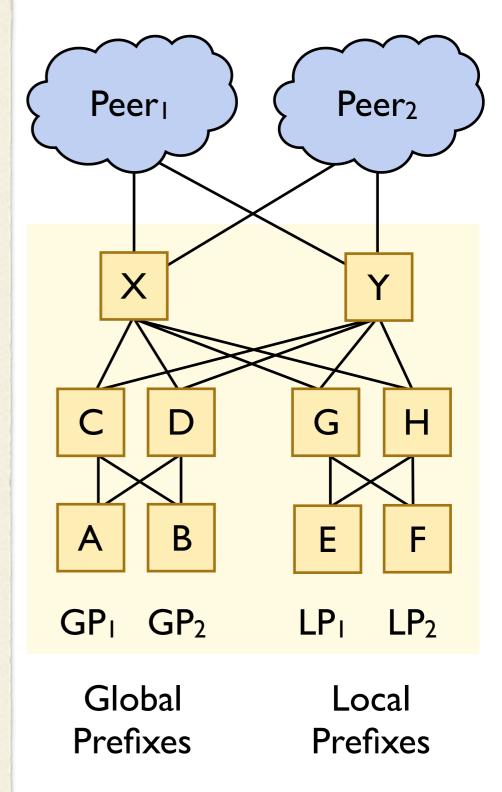


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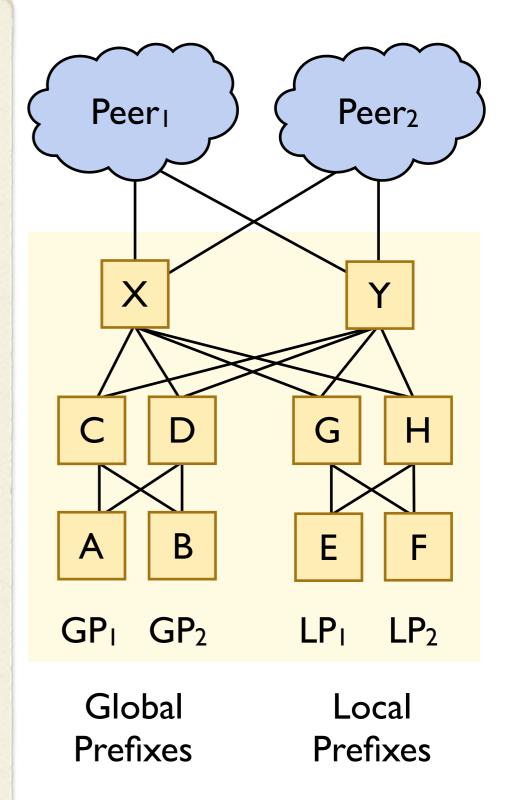
define Locality =
  {LP1 | LP2 => internal}
```



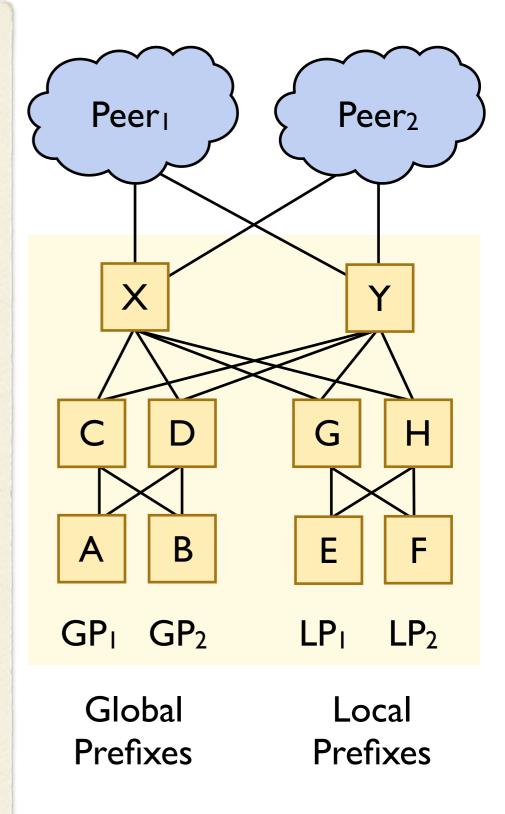
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 \{GP1 => end(A)\}
  GP2 => end(B)
 LP1 => end(E)
 LP2 => end(F)
  true => exit(Peer1 >> Peer2) }
define Locality =
 {LP1 | LP2 => internal}
define transit(X,Y) =
 enter (X | Y) and exit (X | Y)
```



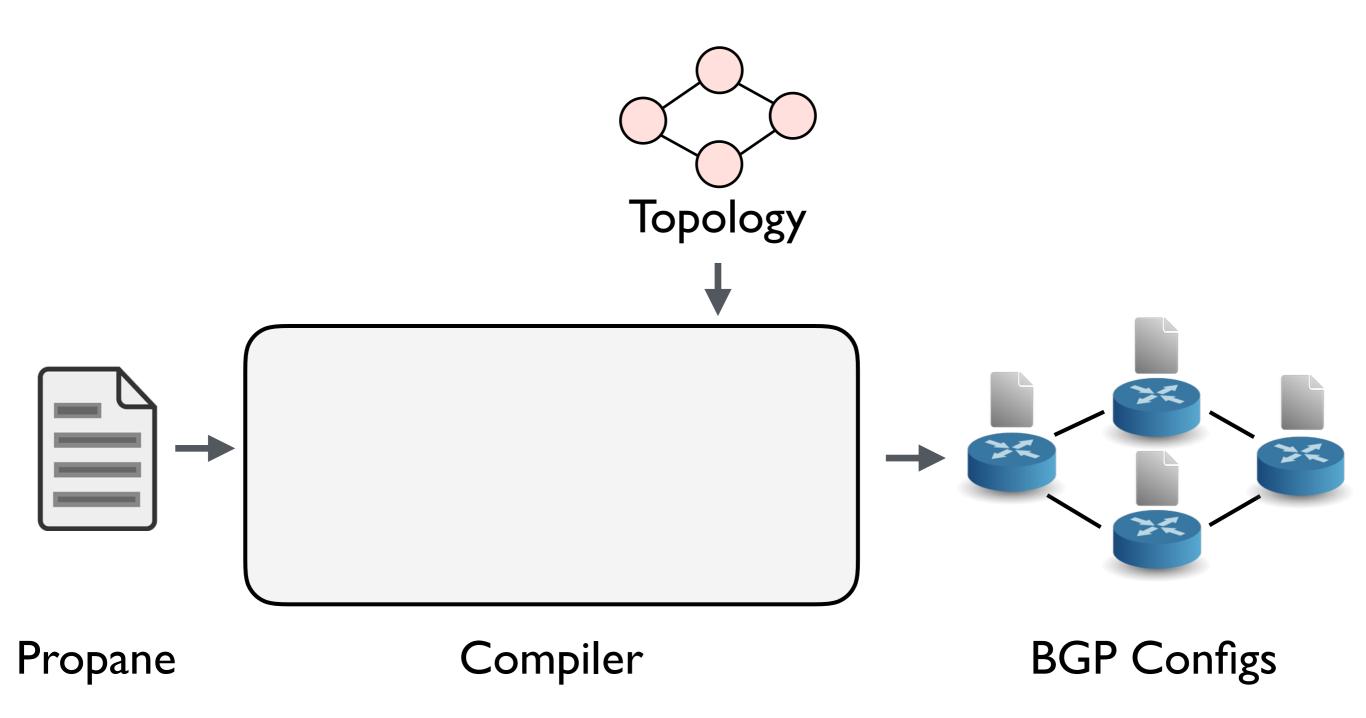
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 LP2 => end(F)
  true => exit(Peer1 >> Peer2)}
define Locality =
 {LP1 | LP2 => internal}
define transit(X,Y) =
 enter (X | Y) and exit (X | Y)
define NoTransit =
 {true => !transit(Peer1, Peer2)}
```



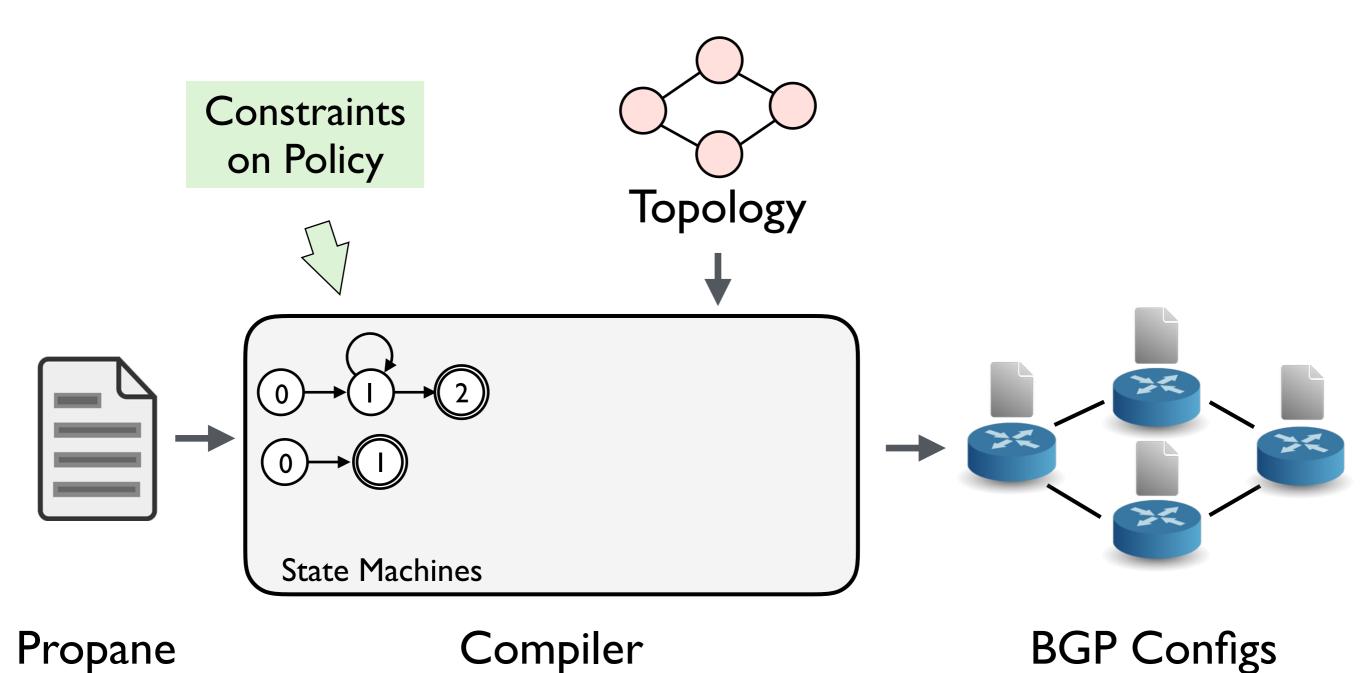
```
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 \{GP1 => end(A)\}
 GP2 => end(B)
 LP1 => end(E)
 LP2 => end(F)
  true => exit(Peer1 >> Peer2)}
define Locality =
 {LP1 | LP2 => internal}
define transit(X,Y) =
 enter (X | Y) and exit (X | Y)
define NoTransit =
 {true => !transit(Peer1, Peer2)}
define Main =
   Destination & Locality &
   NoTransit & agg(GP, in -> out)
```



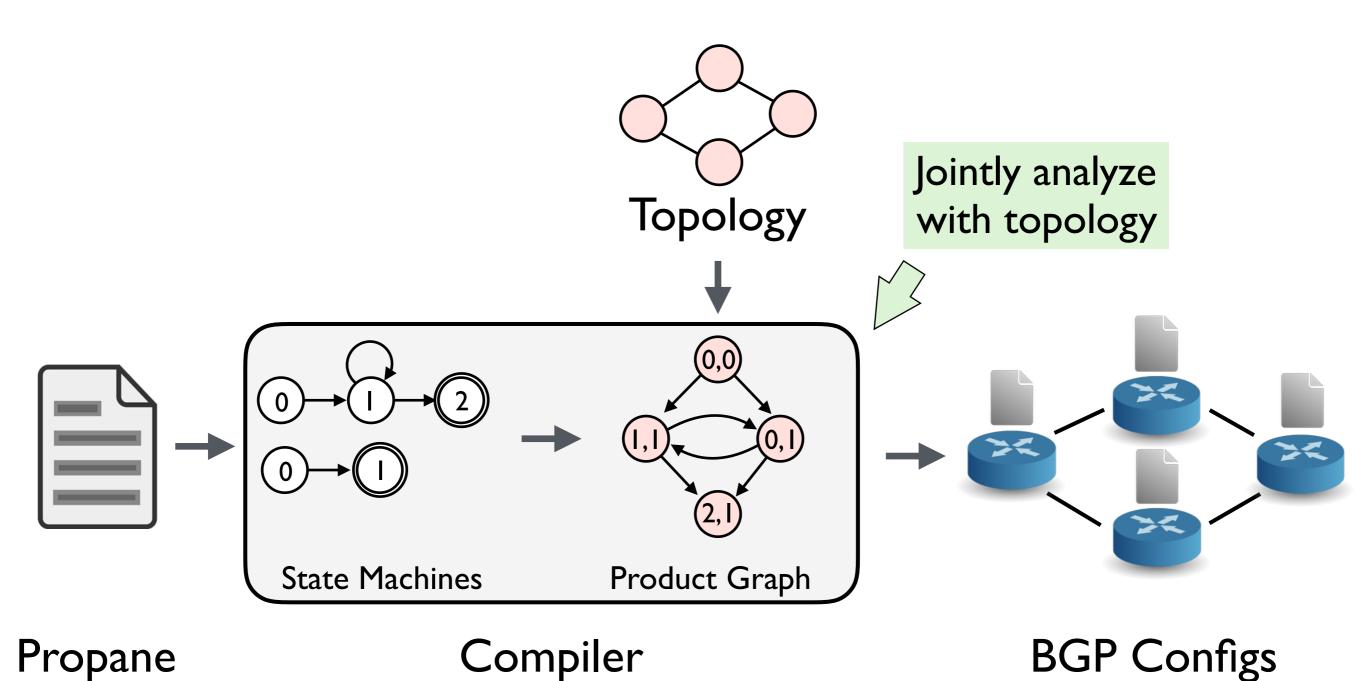
Compilation



Compilation

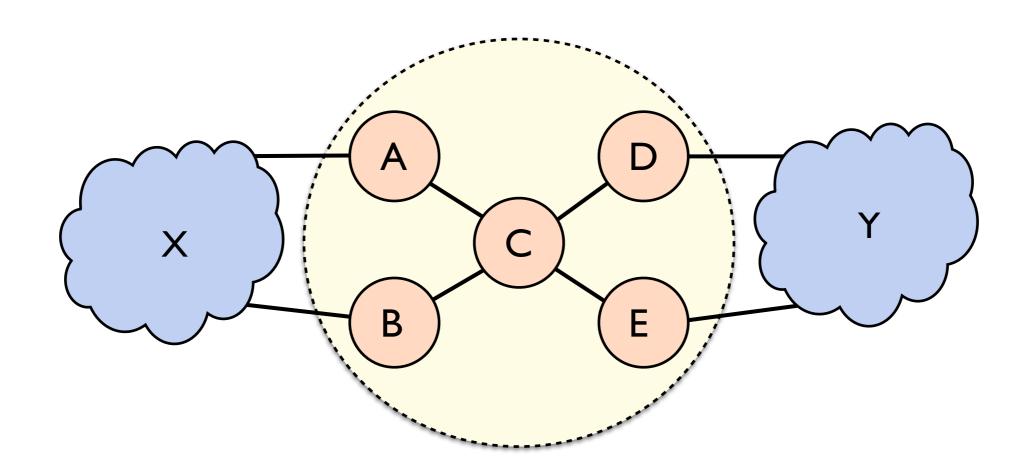


Compilation



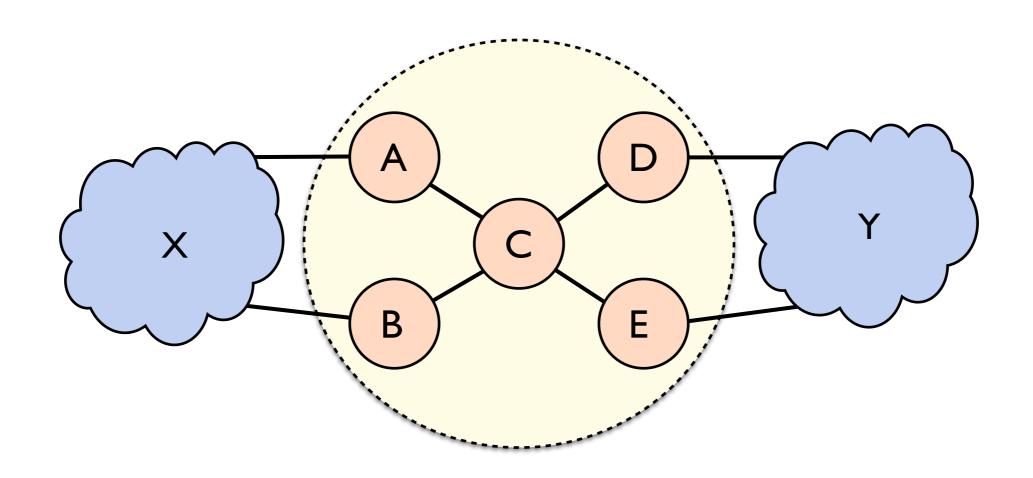
31

Compilation: A simple Example



end(Y) & (path(A,C,D) >> any)

Compilation: A simple Example

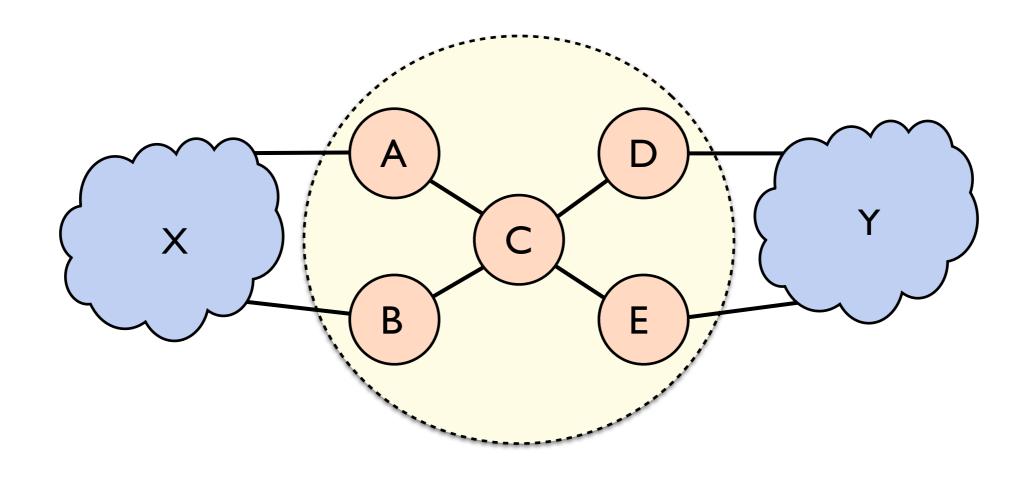


end(Y) & (path(A,C,D) >> any)



 $XACDY >> (\Sigma^*)Y$

Reversed Automata from Policies



Policy:

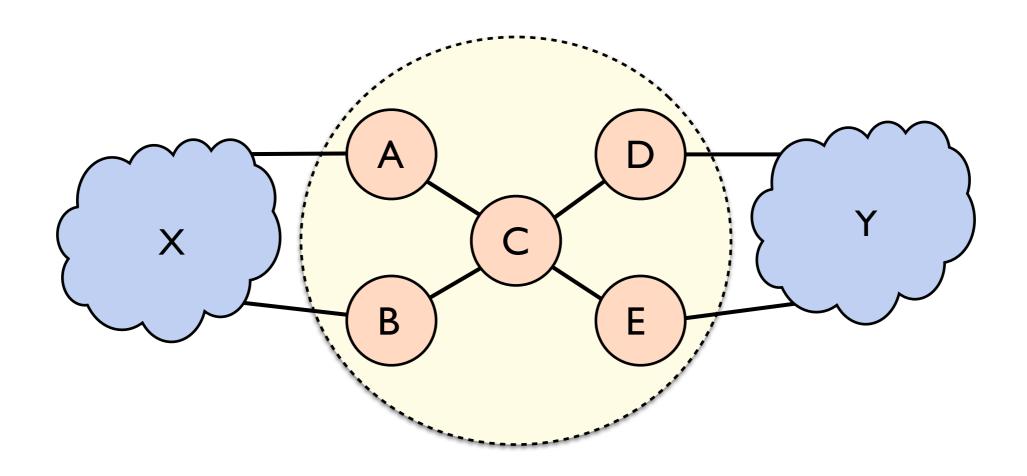
1. XACDY

2. $(\Sigma^*)Y$

More preferred paths

Less preferred paths

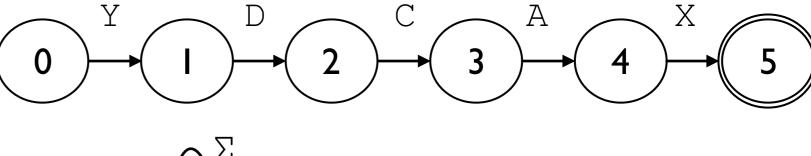
Reversed Automata from Policies

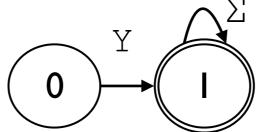


Policy:

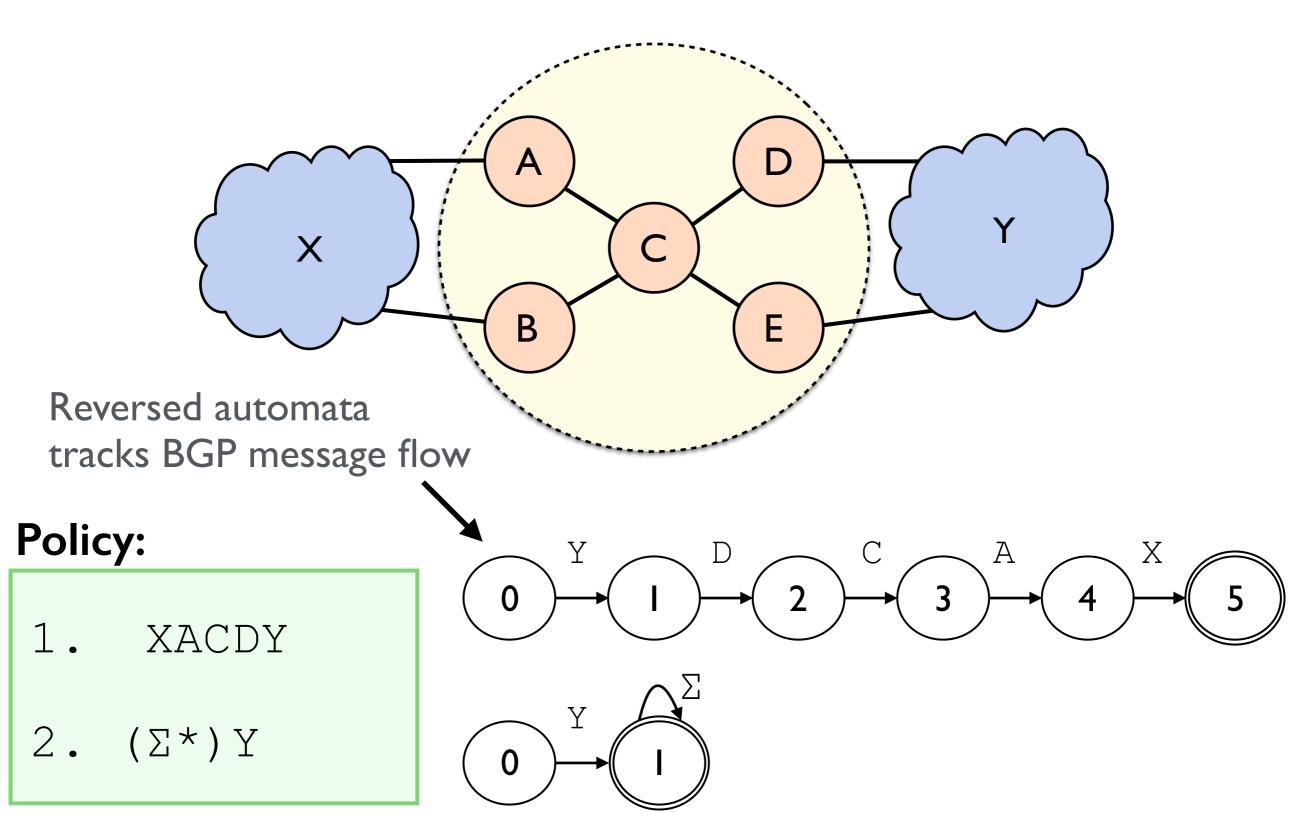
1. XACDY

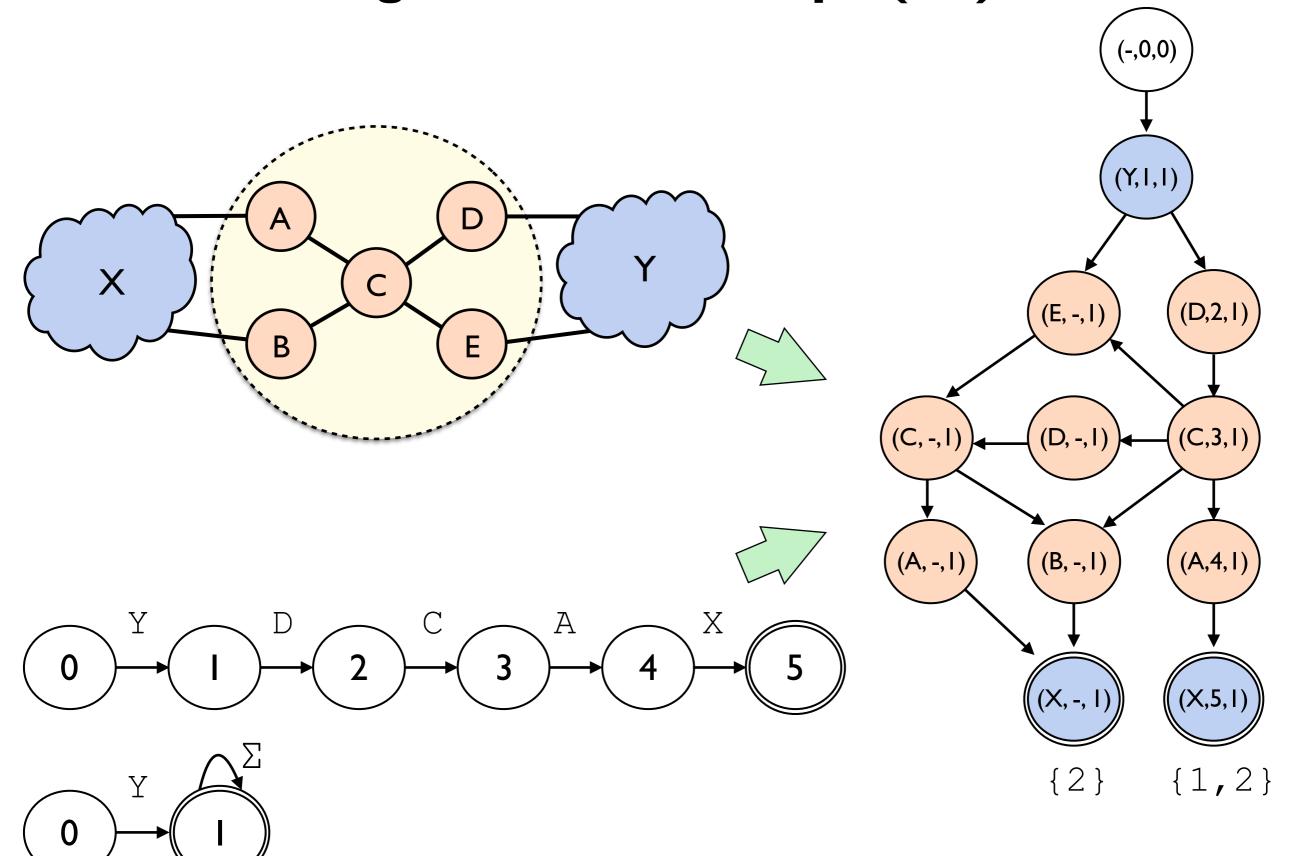
2. $(\Sigma^*)Y$

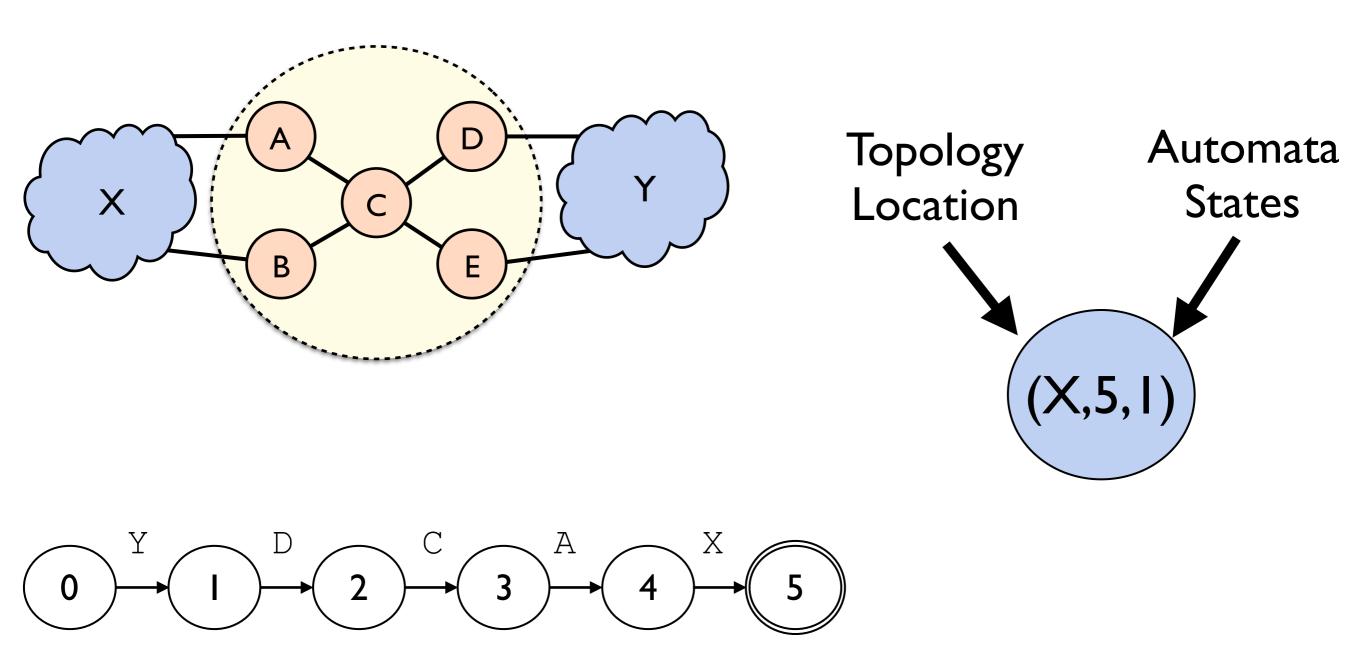


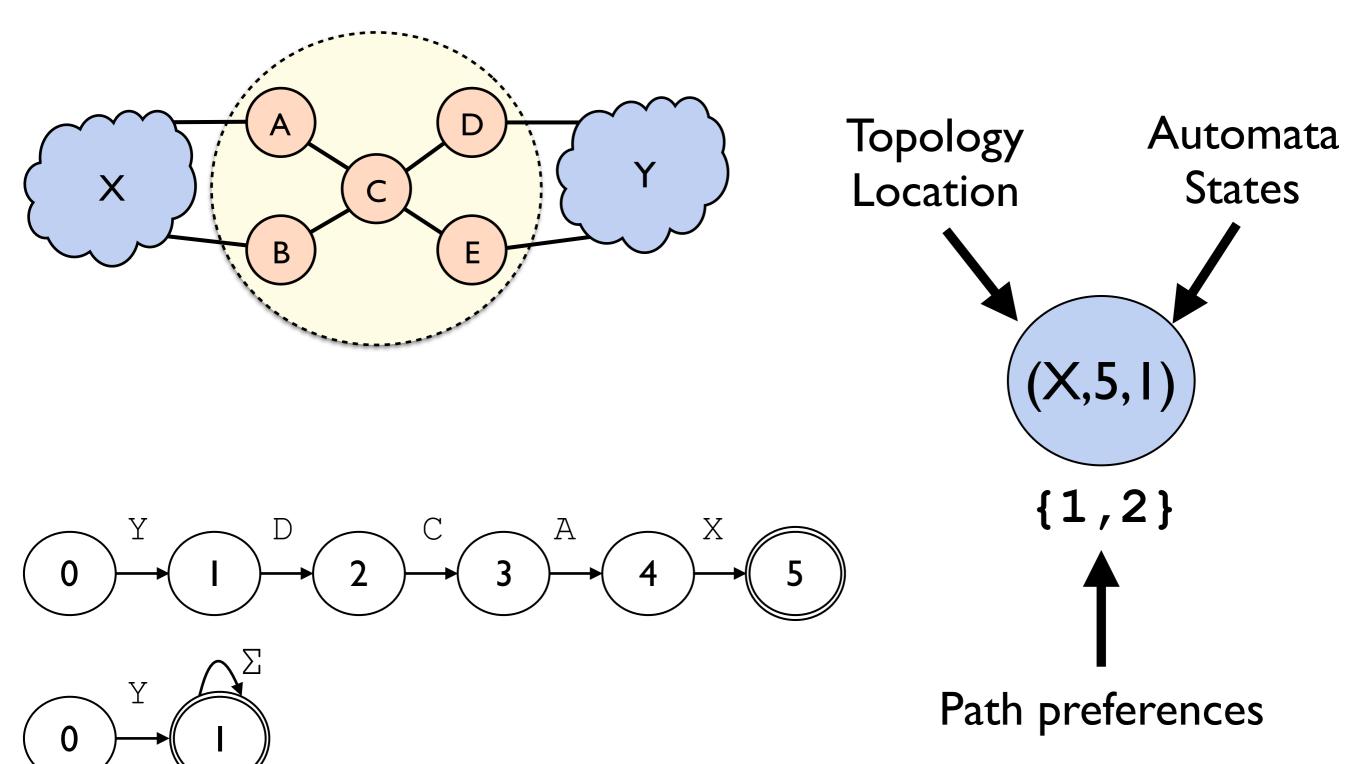


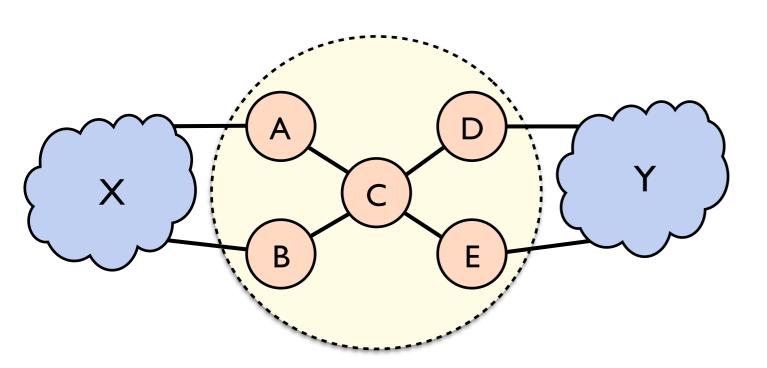
Reversed Automata from Policies

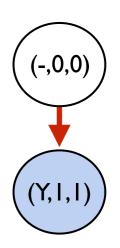


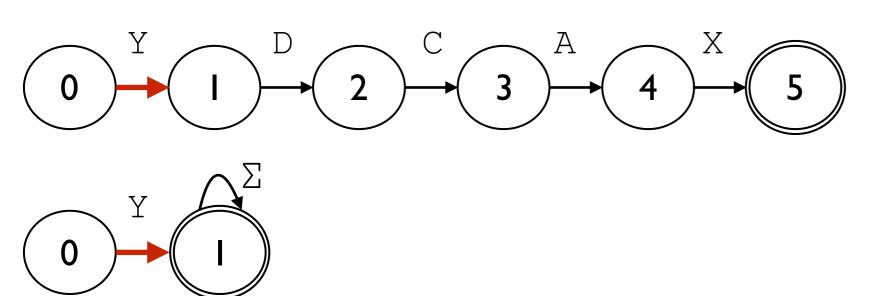


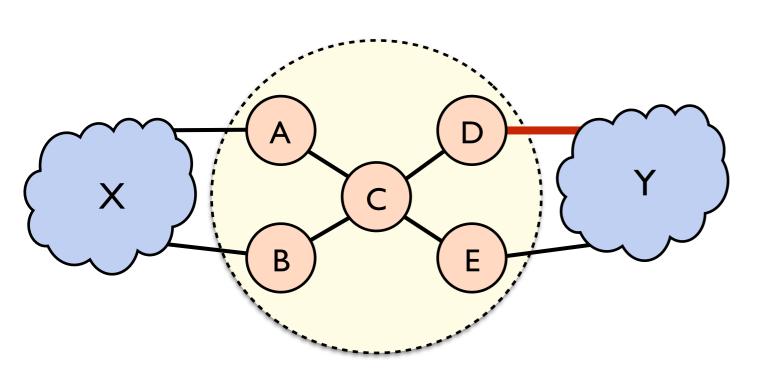


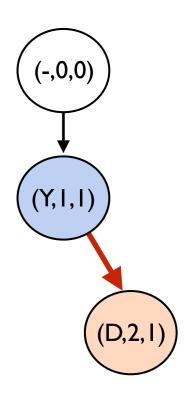


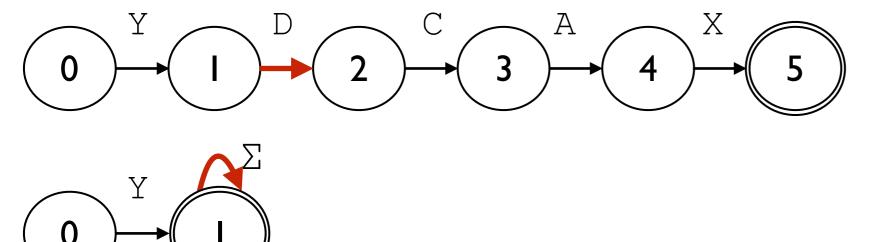


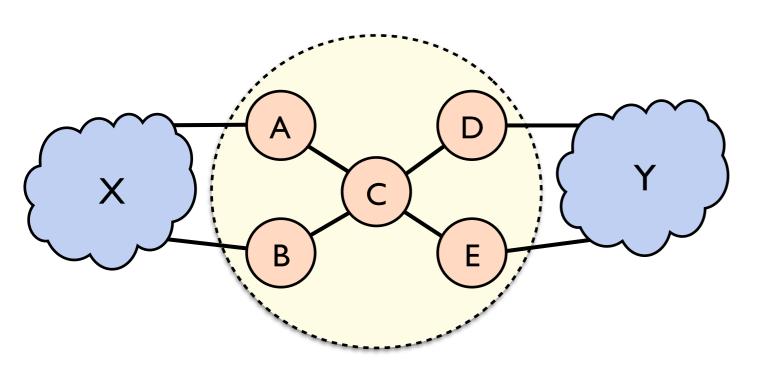




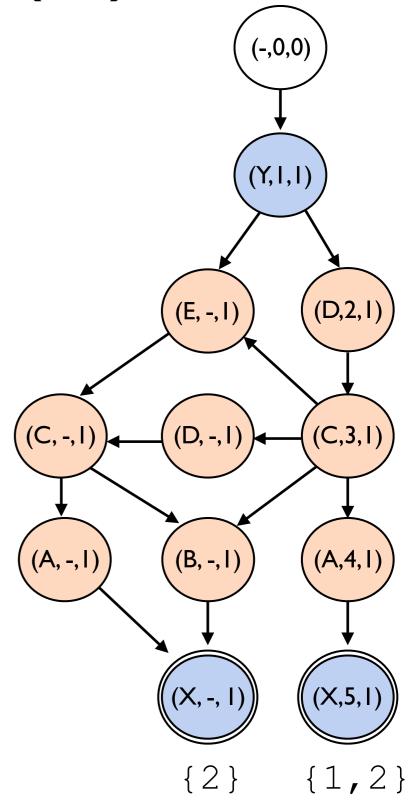


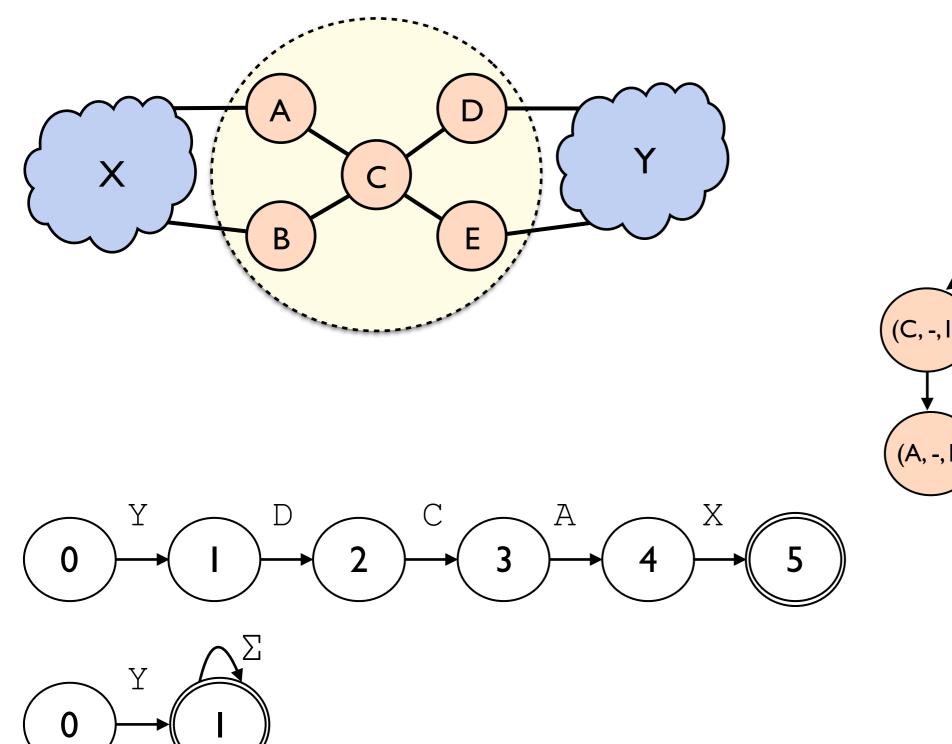


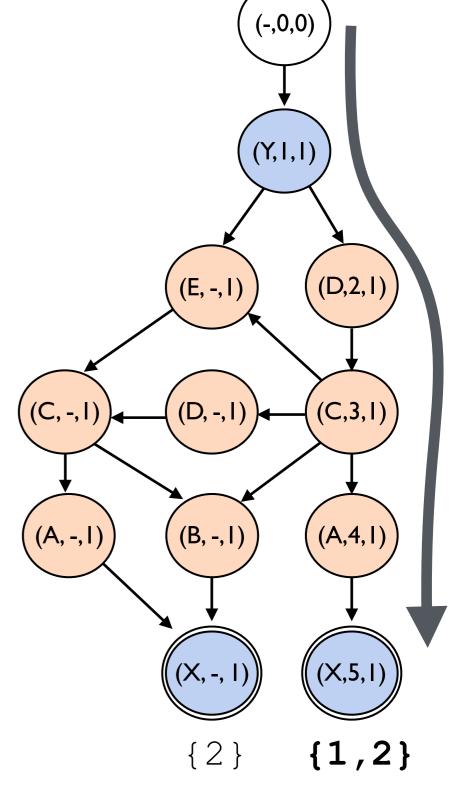




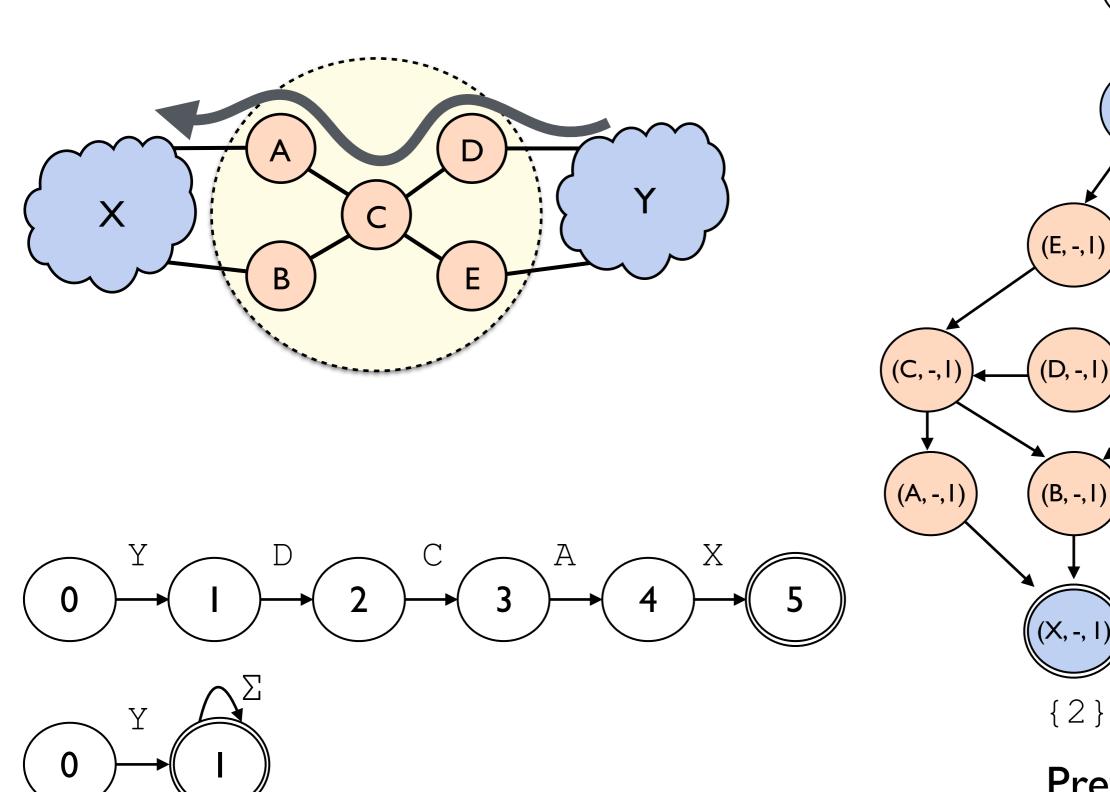
Graph capturing all possible policycompliant paths through the topology

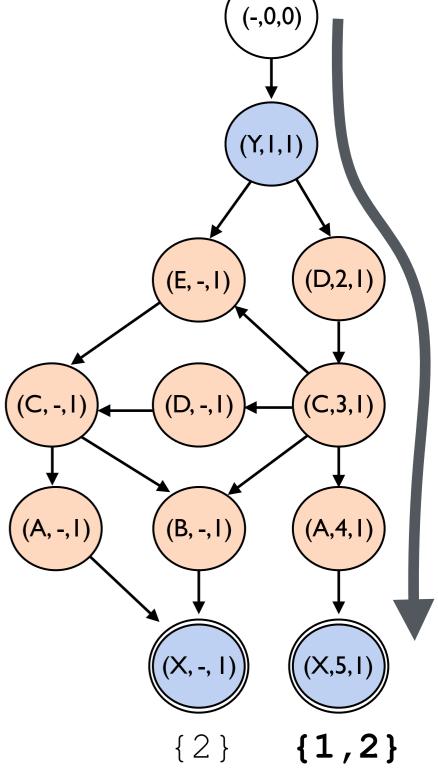




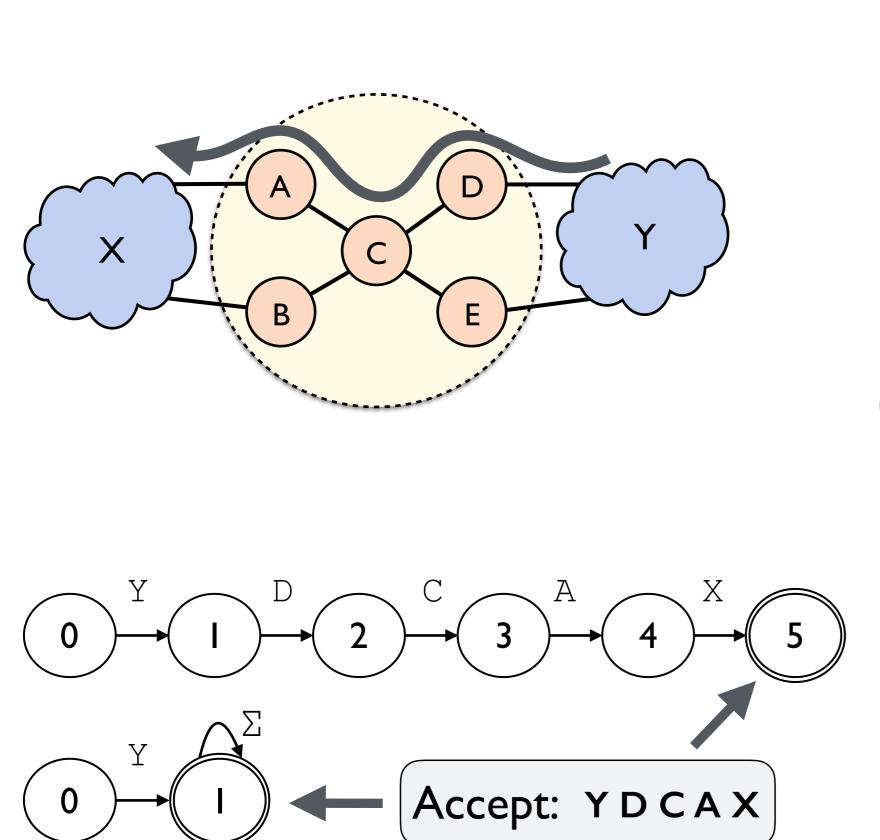


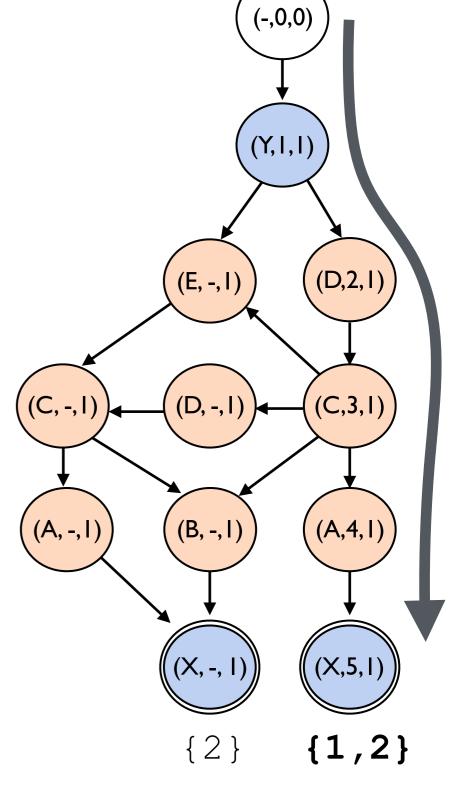
Preferences



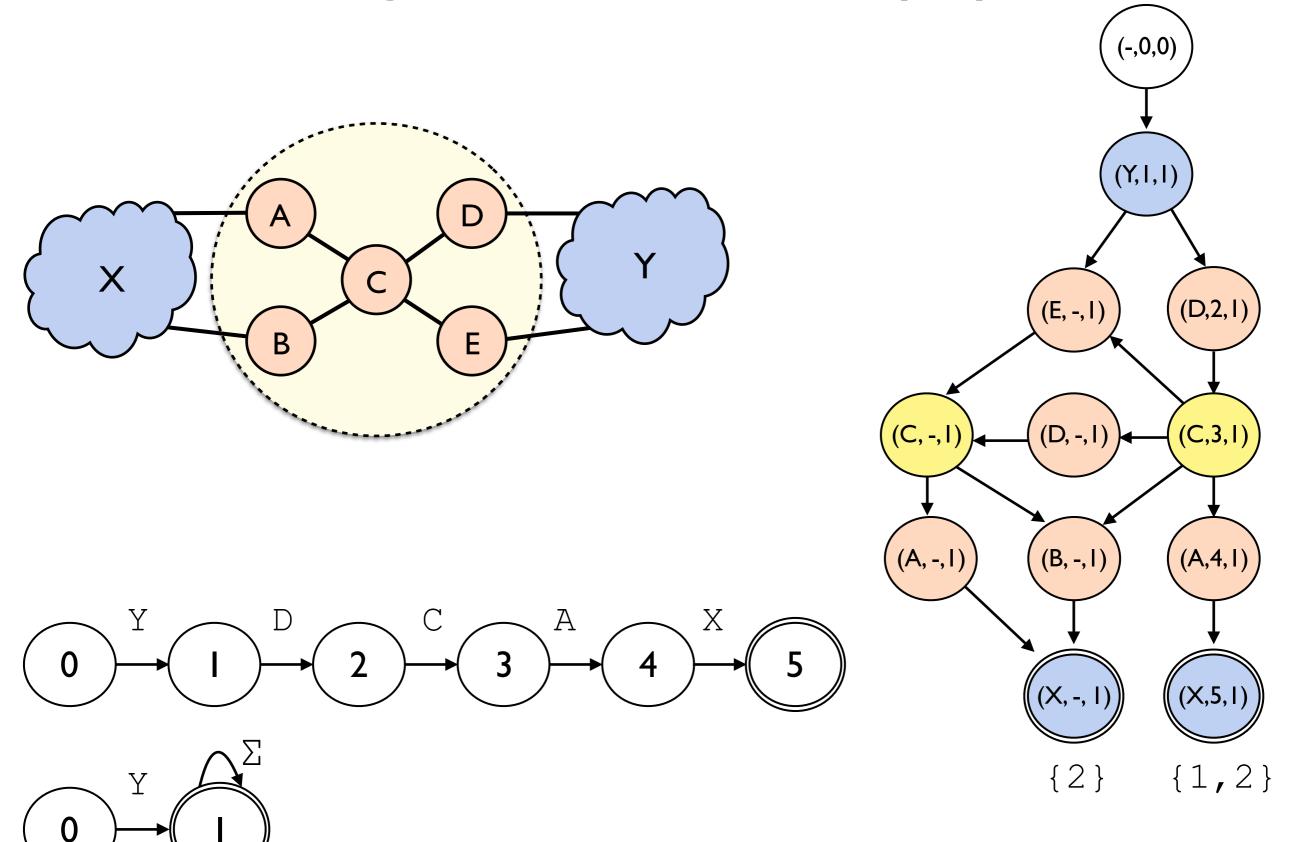


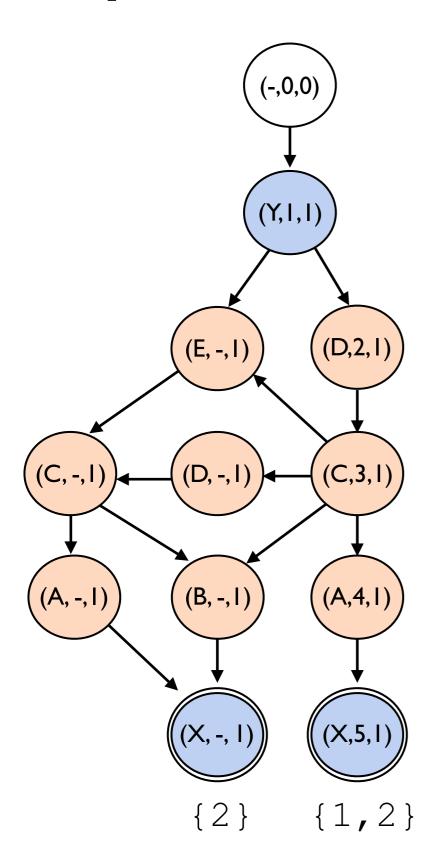
Preferences





Preferences



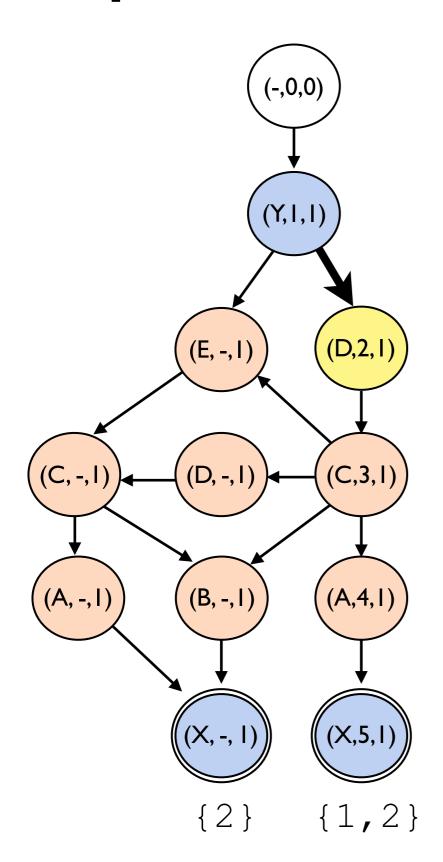


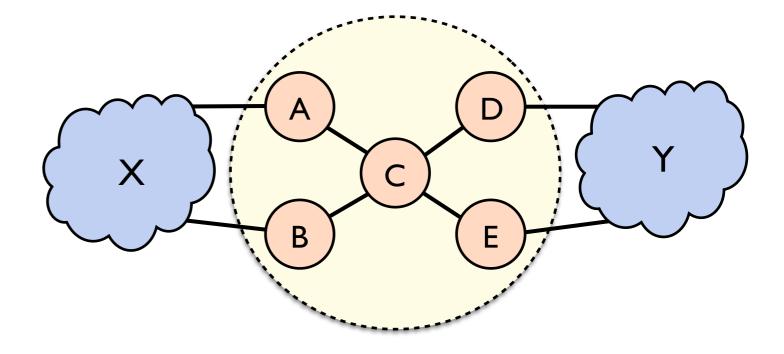
Idea I: Restrict advertisements to edges

- Encode state in a BGP community tag
- Incoming edges import filters
- Outgoing edges export filters

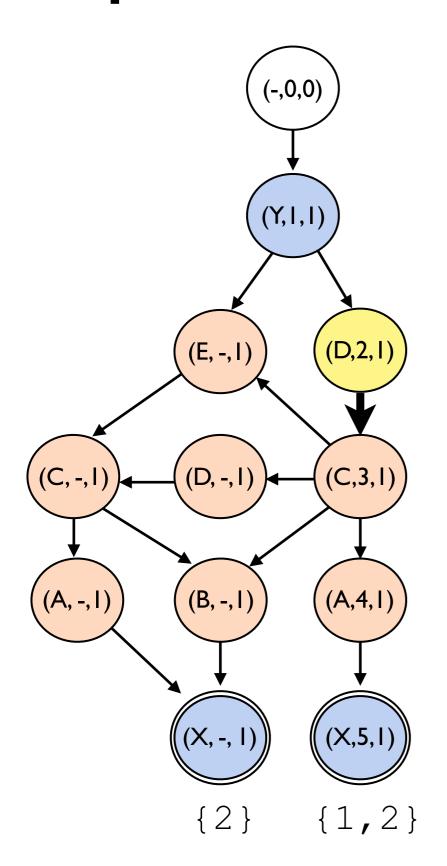


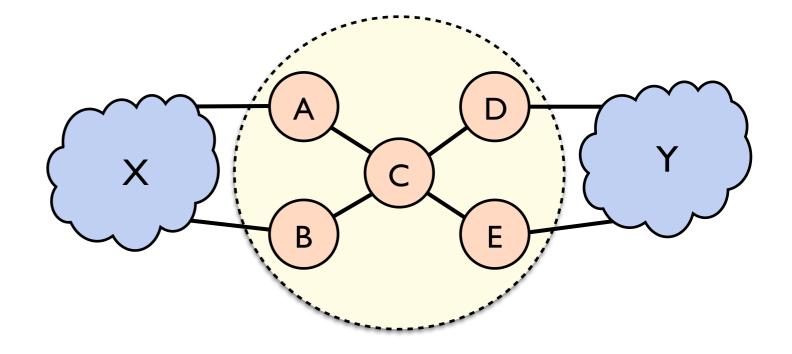
Let BGP find some allowed path dynamically



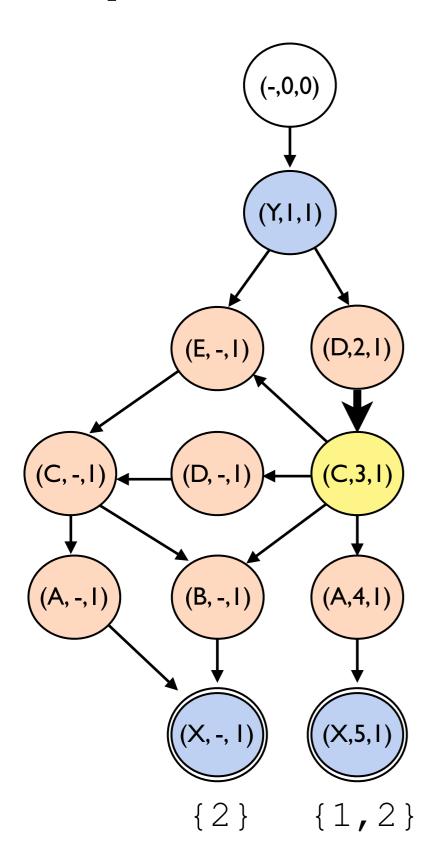


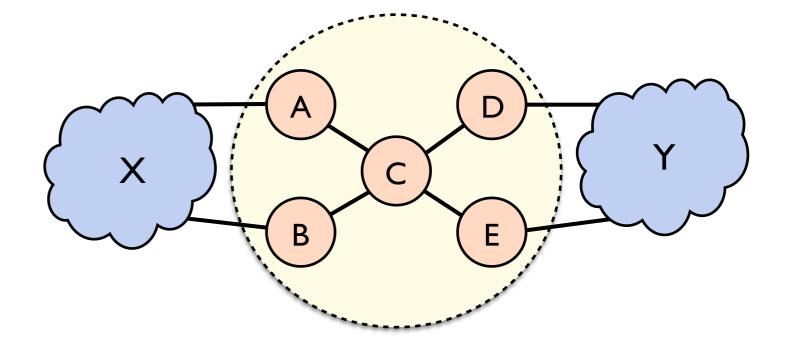
D allows import matching regex(Y)



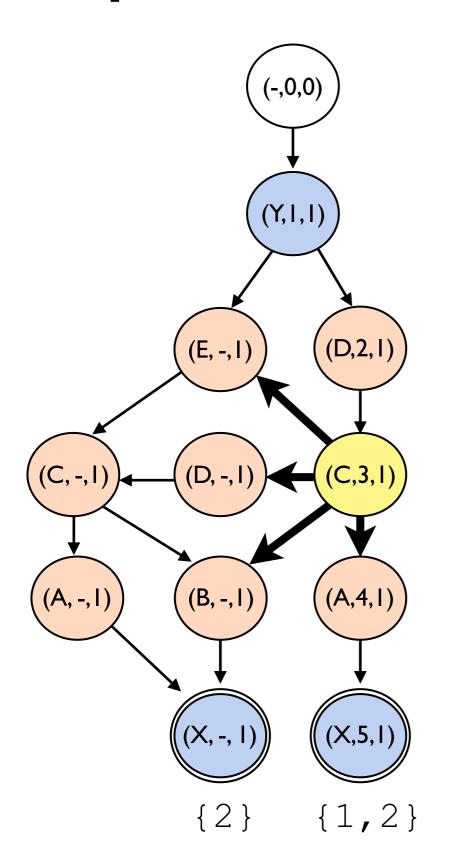


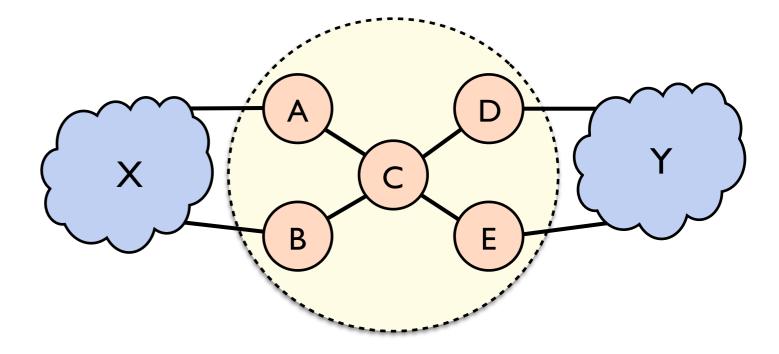
D exports to C with tag (2, I)



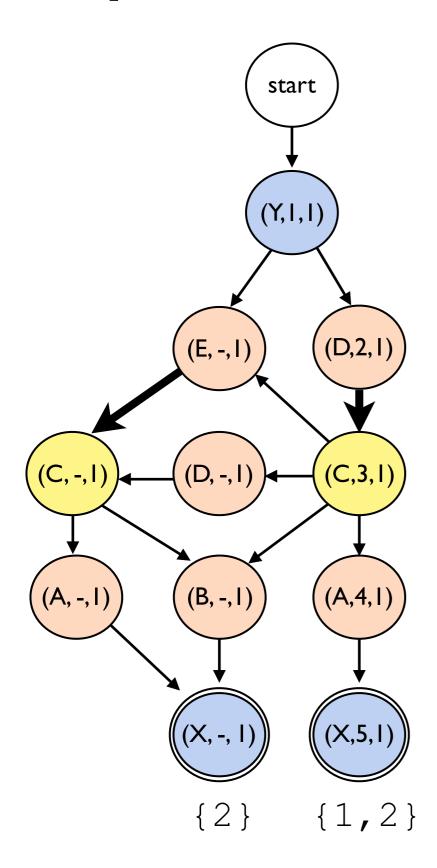


C allows import from D with tag (2,1)





C exports to A,B,D,E with tag (3,1)

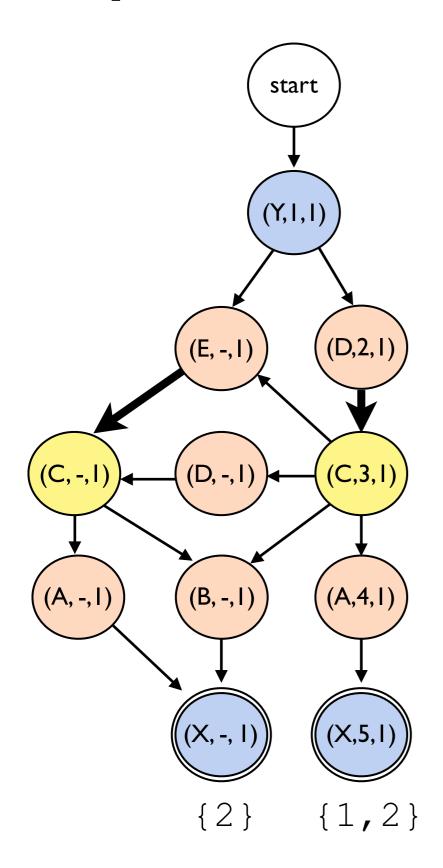


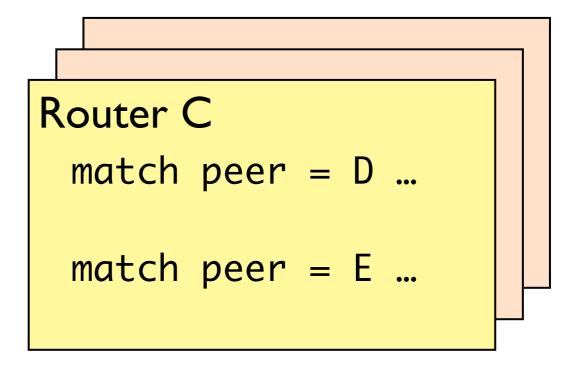
Idea 2: Find preferences

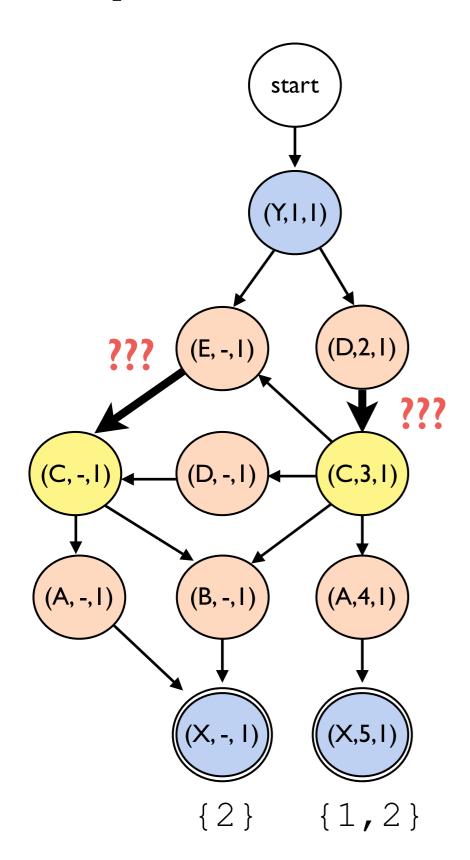
- Direct BGP towards best path
- Under all combinations of failures



Let BGP find the best path dynamically



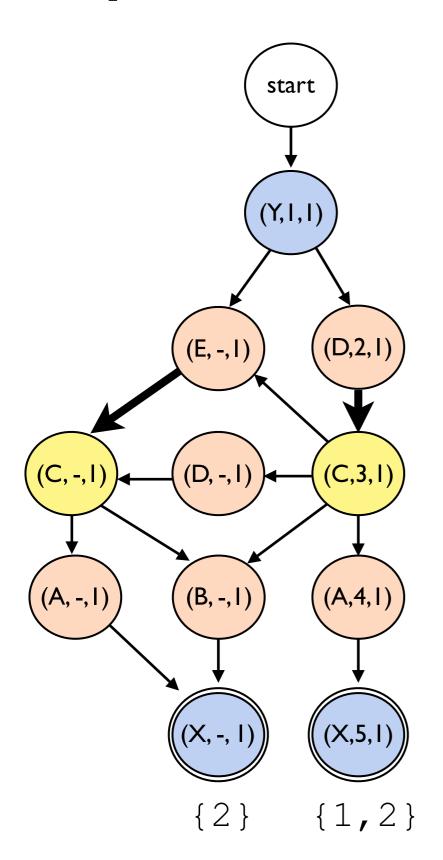




```
Router C

match peer = D ...
local-pref ← ???

match peer = E ...
local-pref ← ???
```



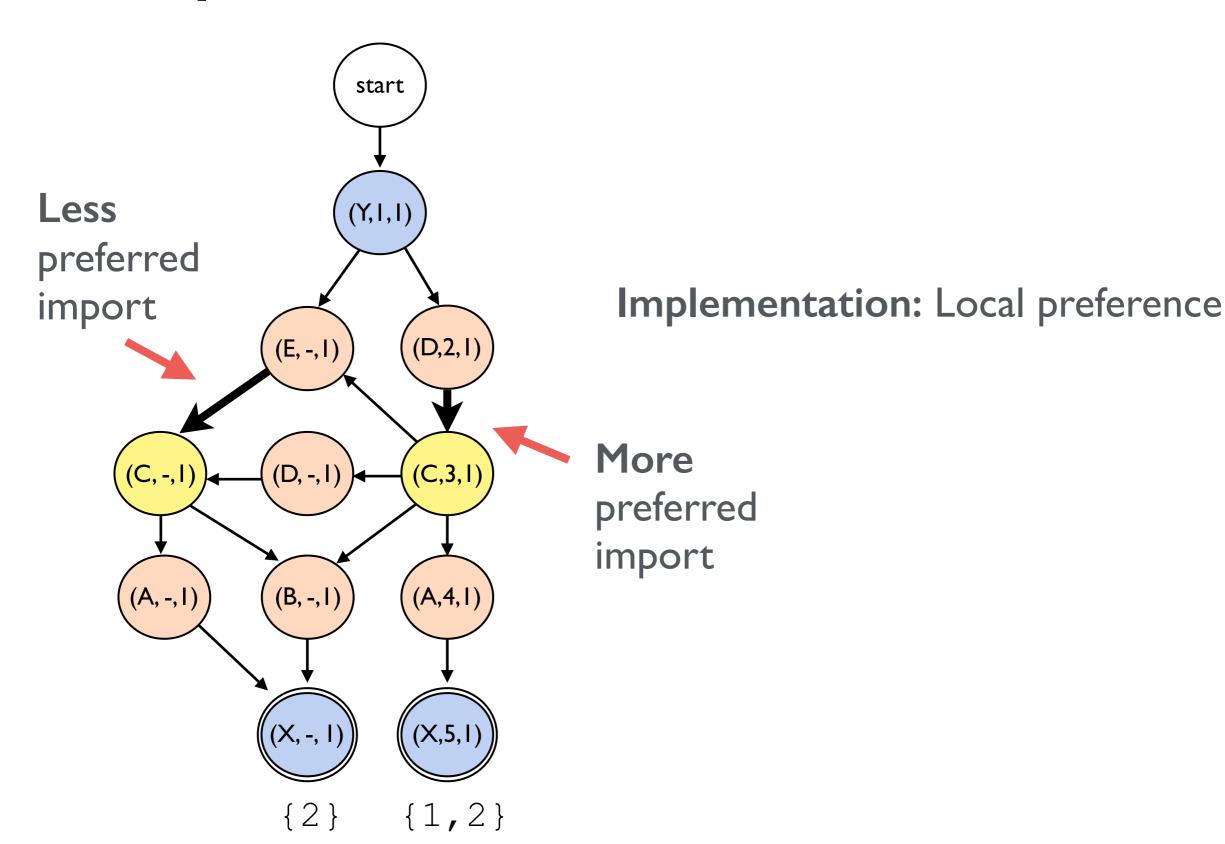
```
Router C

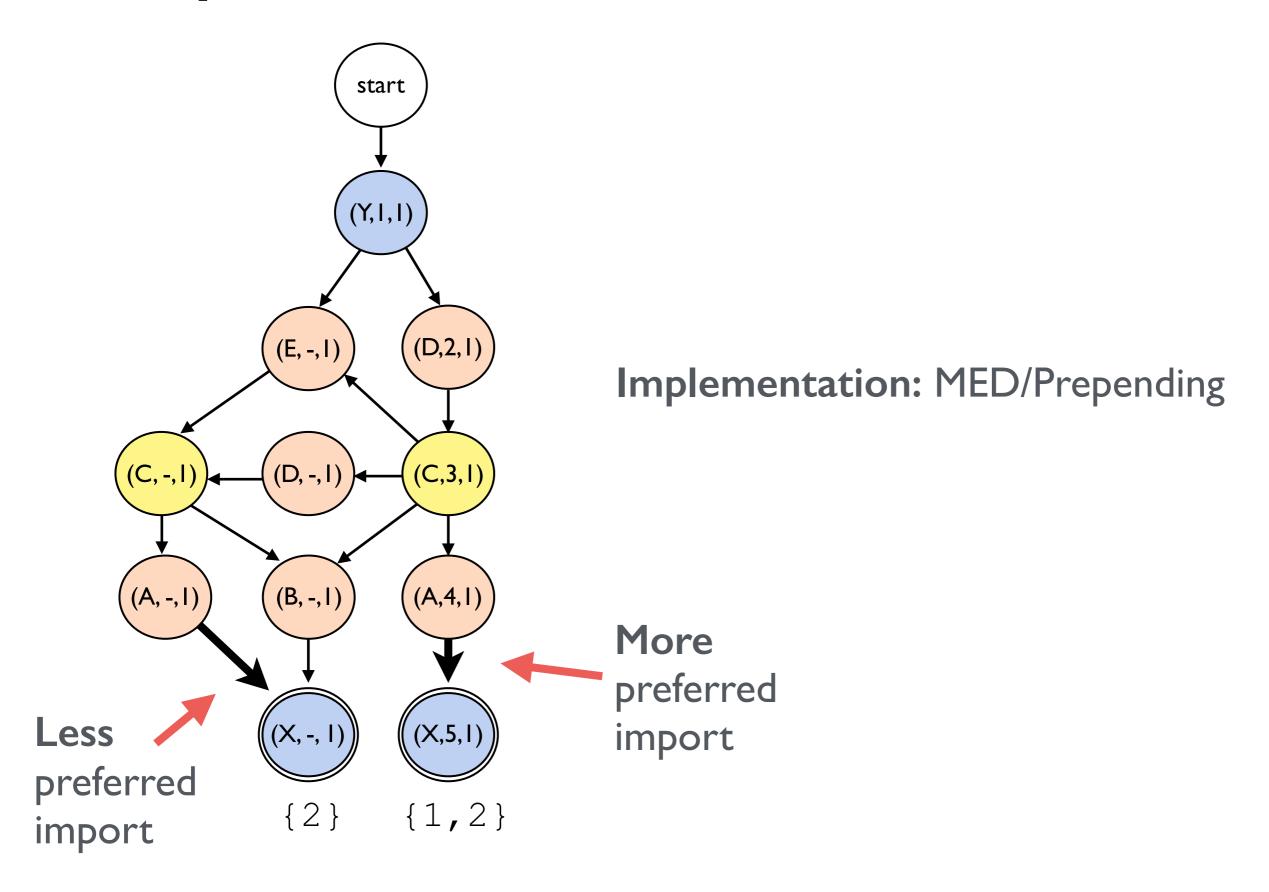
match peer = D ...
local-pref ← ???

match peer = E ...
local-pref ← ???
```

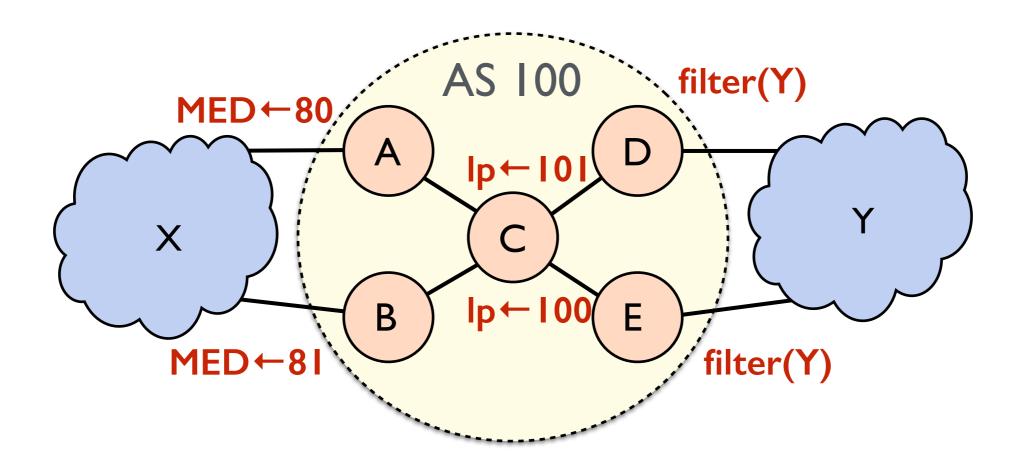
Efficient algorithm to assign preferences that forces BGP to find the best paths for all possible failures

See the paper for details!





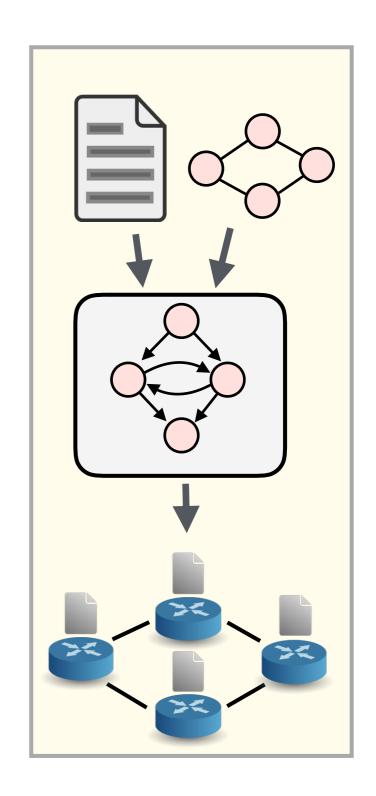
Compilation: A simple Example



end(Y) & (path(A,C,D) >> any)

Implementation

- Written in 7000 lines of F#
- Generates Quagga configurations
- A number of other analyses & features



Benchmarks

- Configurations from a large cloud provider
- Policy described in English documents
- Datacenter and Backbone policies

Policy Size

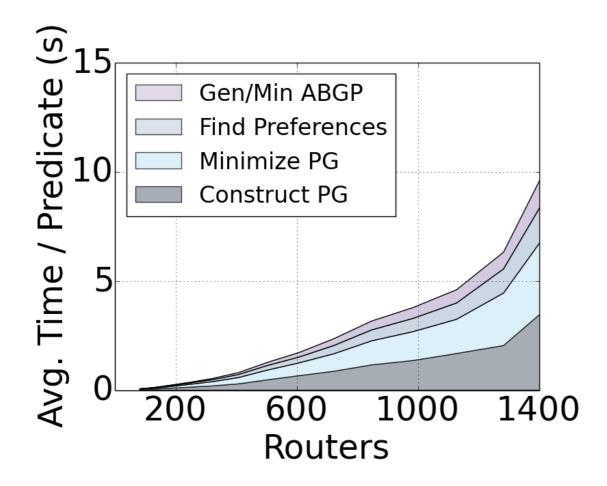
Without prefix/peer definitions

- Datacenter policy: 31 lines of Propane
- Backbone policy: 43 lines of Propane

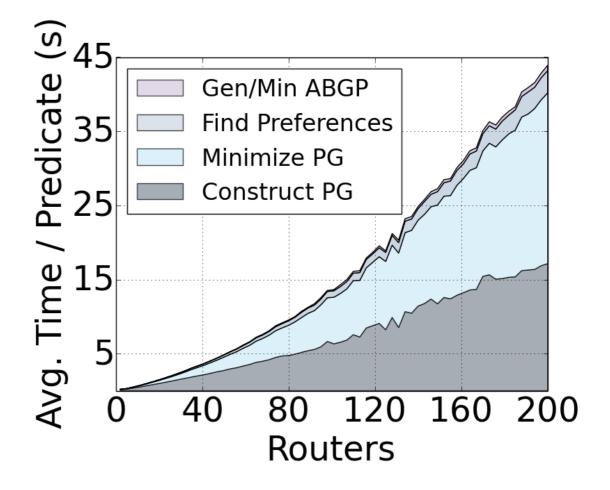
Conventional BGP configurations are 1000s of lines

Compilation Time

- Compile for each prefix equivalence class
- Compile for each equivalence class in parallel
- 8 core, 3.6 GHz Intel Xeon processor



Data center (< 9 min)



Backbone (< 3 min)

Configuration Size

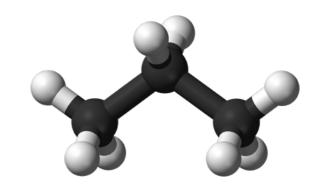
Optimizations

- Avoid using community tags when unambiguous
- Reuse community values across peers
- Merge import/export behaviors across peers

Results

- Optimizations yield 50-100x decrease in config size
- Configurations ~ 1000-10000 lines per router

Propane: Summary



High-level language

- Centralized network programmability
- Constraints specify preferred paths and backups in case of failure
- Uniform abstractions for Inter- and Intra-domain routing
- Core policy in 30-50 lines of Propane vs. 1000s

Compiler

- Distributed implementation via BGP
- Static analysis guarantees policy compliance for all failures
- Scales to reasonably sized network topologies