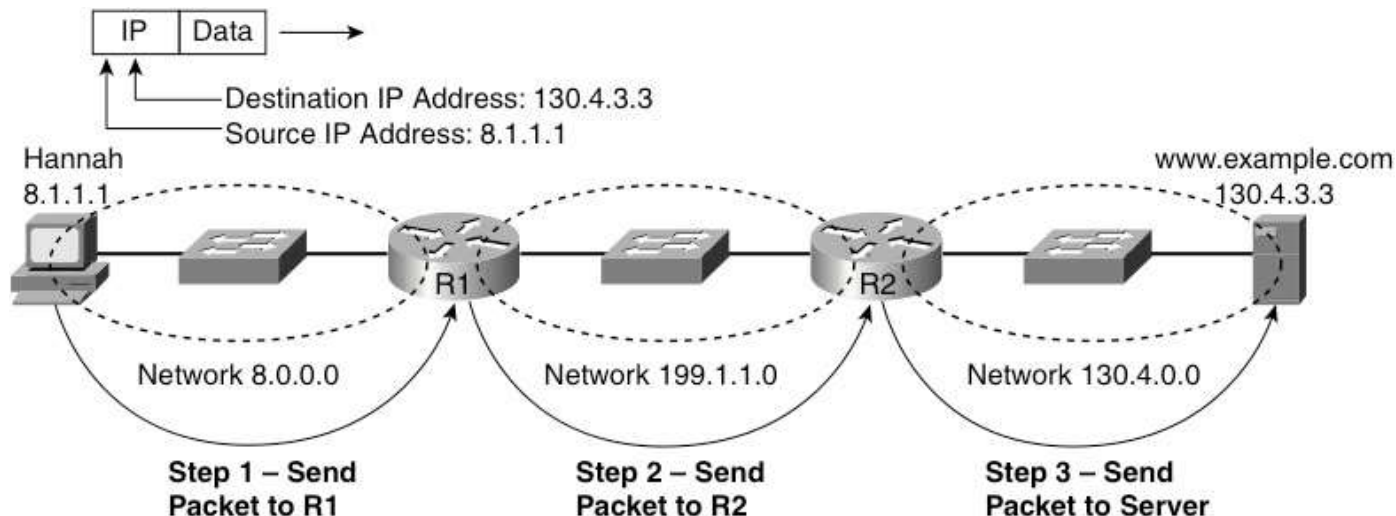


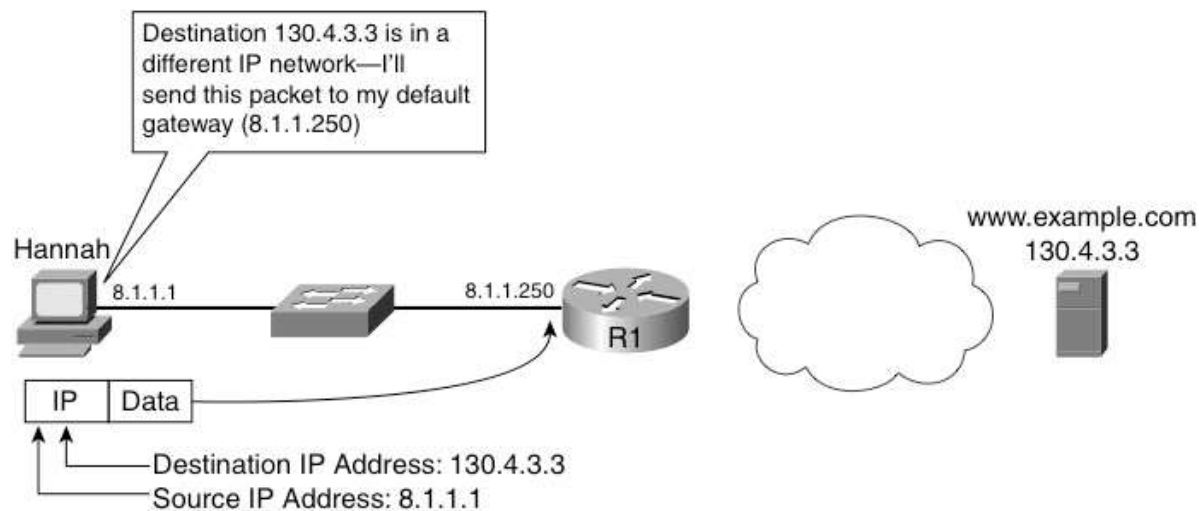
Simplicity in routing

- ◆ Routers sit at the intersections within an internet
 - Forward packets from one network to another
 - ◆ Some of those packets originated on one of the connected networks
 - ◆ Many are simply “passing through”
- ◆ Routers may have a huge number of packets to forward
 - Efficiency is key
 - Routing rules are kept as simple as possible



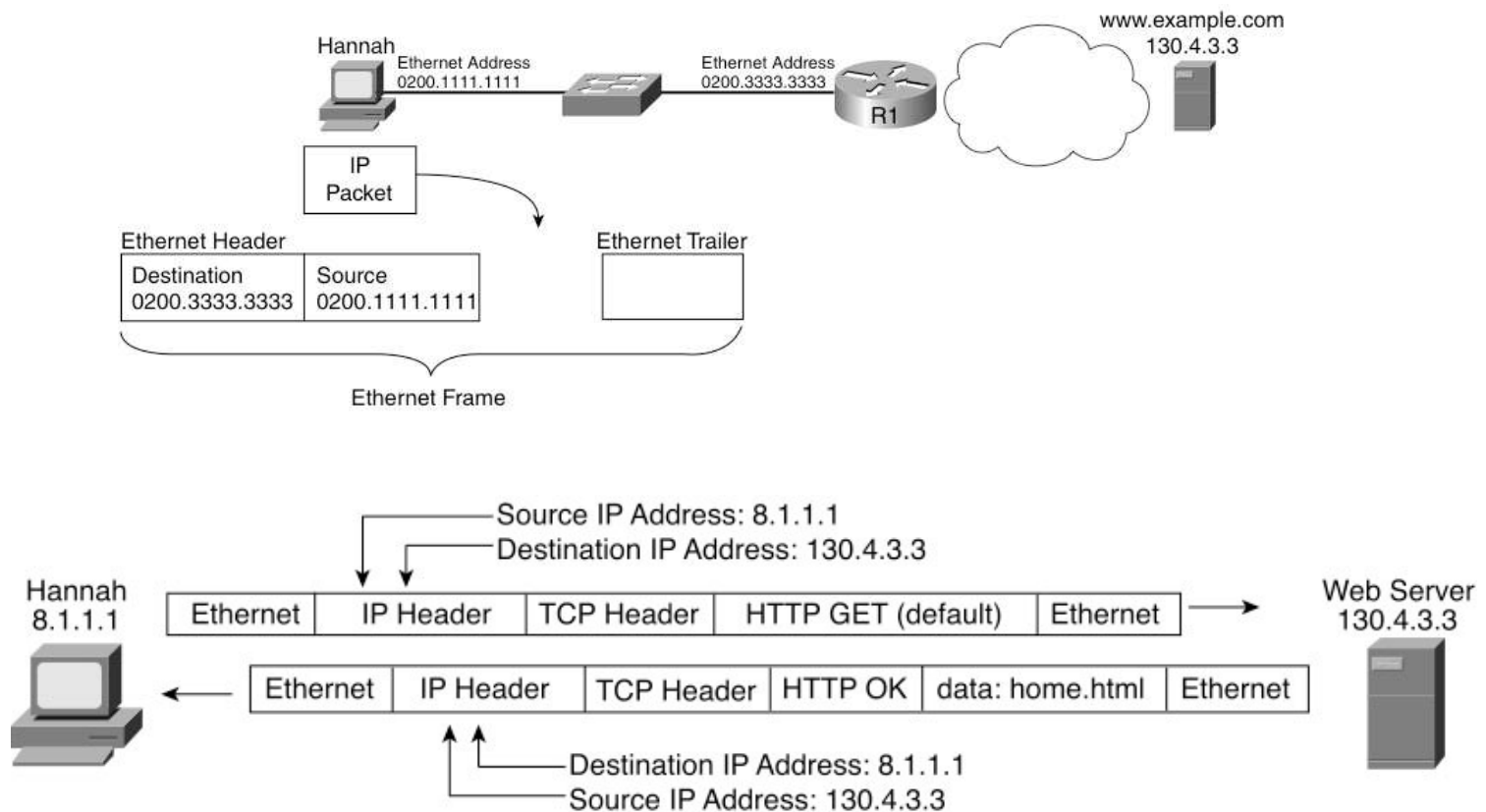
Default routers

- ◆ Every host is configured with the IP address of its network's default router (or gateway)
 - Manually or via Dynamic Host Configuration Protocol (DHCP)
 - Whenever a packet is addressed to a destination on a different network, it is sent to the default router
 - ◆ That router examines the packet and decides how to route it
 - If the recipient is on one of its connected networks, it simply delivers it to that host
 - Otherwise, it forwards the packet to another router on one of its connected networks



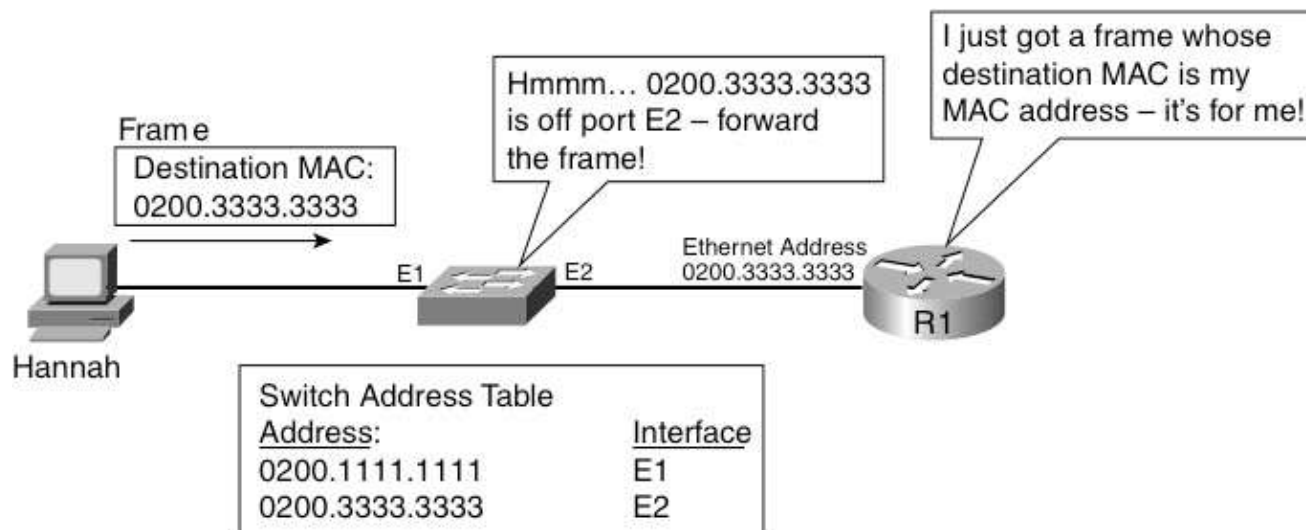
Encapsulating packets in frames

- ◆ Only valid Ethernet frames can travel on an Ethernet network
 - IP packets must therefore be encapsulated in Ethernet frames
 - ◆ Same is true of other network types as well



Addressing the default router

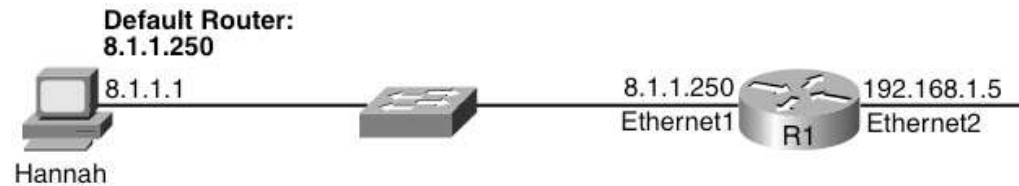
- ◆ IP headers only hold the IP address of the packet's ultimate destination
 - There's no field to hold the IP address of the default router
 - But since the packet must be encapsulated in an Ethernet frame, that frame can be addressed to the default router's MAC address
 - ◆ But hosts are only configured with default router's IP address
 - ◆ Somehow that IP address must be converted to a MAC address



Address Resolution Protocol (ARP)

- ◆ Used by sender to learn MAC address of default router
 - Sender sends an ARP request message to the Ethernet broadcast address
 - ◆ Request simply asks any host with included IP address to respond with its MAC address
 - Owner of IP address in request sends an ARP response back to the sender's unicast address
 - ◆ Response includes responder's MAC address
 - Sender stores answer in ARP cache for subsequent use

ARP illustrated

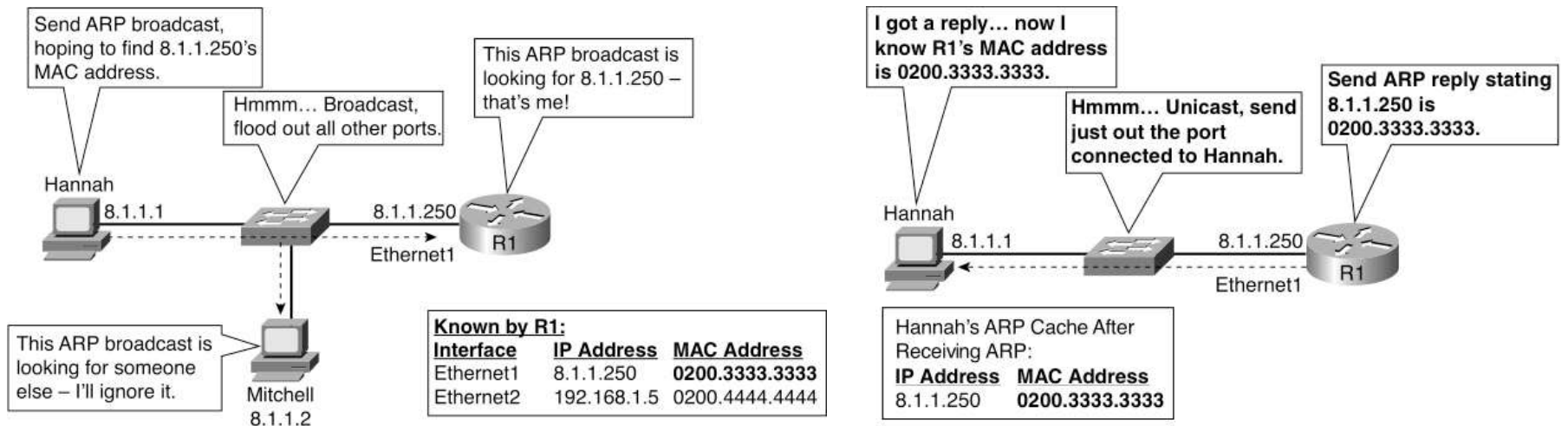


Known by Hannah:

| Interface | IP Address | MAC Address |
|----------------|------------|----------------|
| My NIC: | 8.1.1.1 | 0200.1111.1111 |
| Default Router | 8.1.1.250 | ? |

Known by R1:

| Interface | IP Address | MAC Address |
|-----------|------------|----------------|
| Ethernet1 | 8.1.1.250 | 0200.3333.3333 |

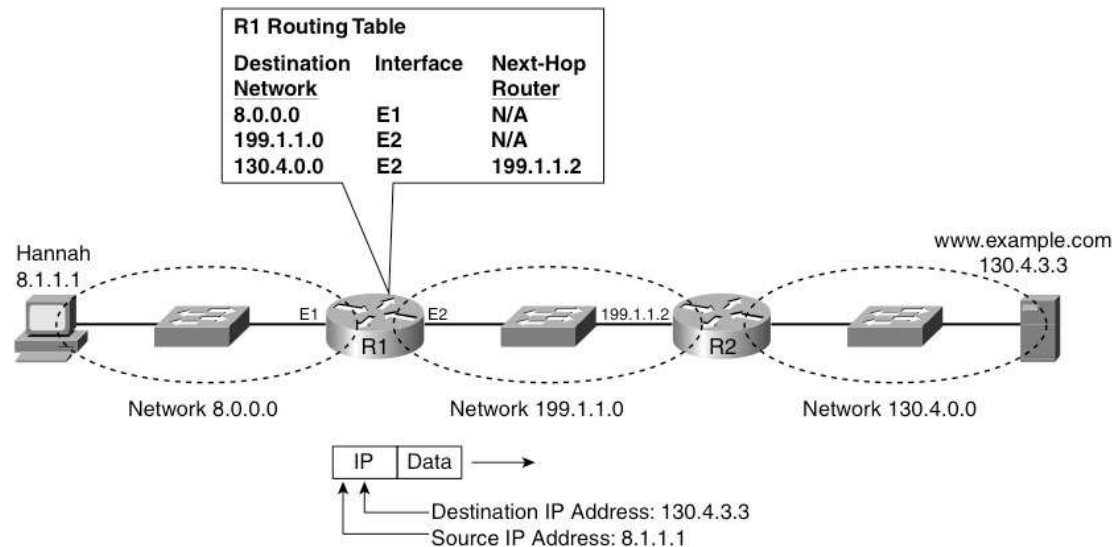


Intelligent intersections

- ◆ Routers are like intelligent intersections
 - Packets are entirely passive
 - Routers make all the decisions and move packets closer to their destinations
- ◆ The decision making is called routing logic
 - Router checks each incoming frame's FCS and discards erroneous frames
 - For good frames, router decapsulates IP packet from frame
 - Router examines packet and decides where to forward it
 - Having made a forwarding decision, router encapsulates packet in a new frame and sends it

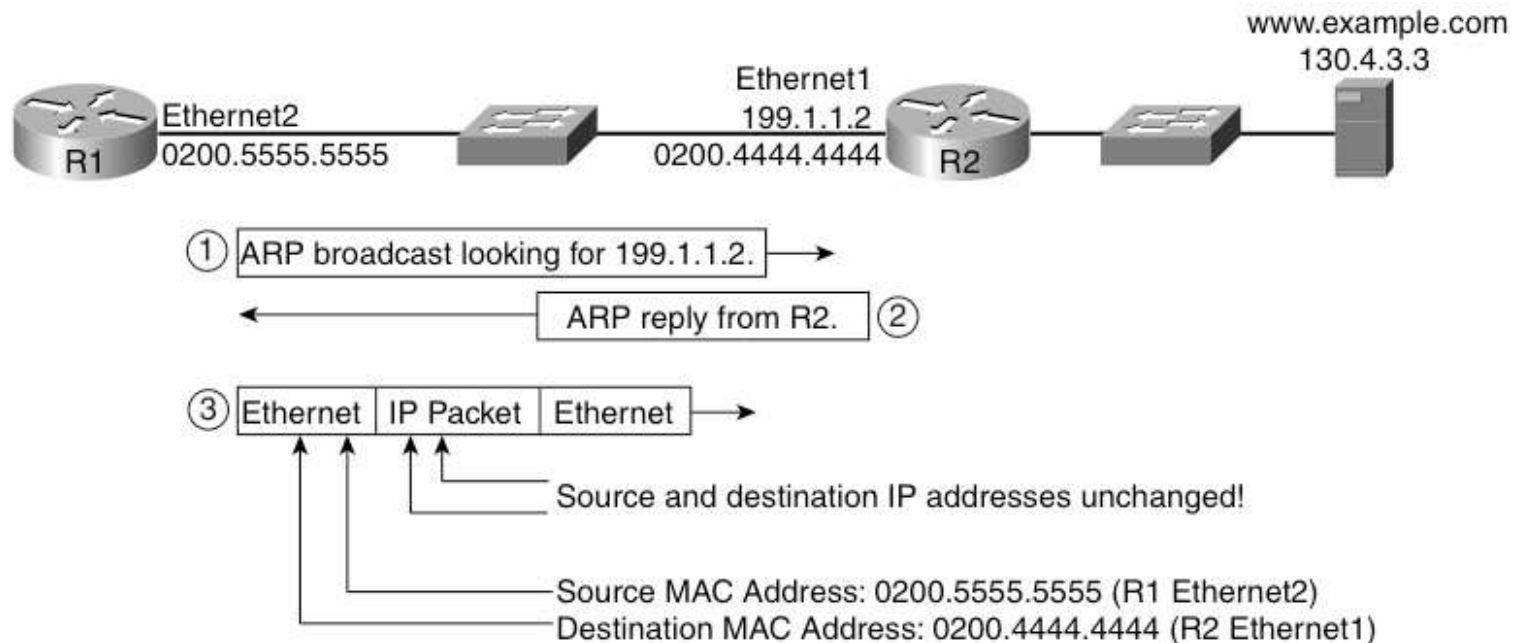
Routing tables

- ◆ Each router builds its own routing table
 - Matches network and subnet numbers to forwarding instructions
 - ◆ Forwarding instructions include outgoing interface and (if necessary) address of next hop router
 - The next hop router is simply the one that gets the packet closer to its destination



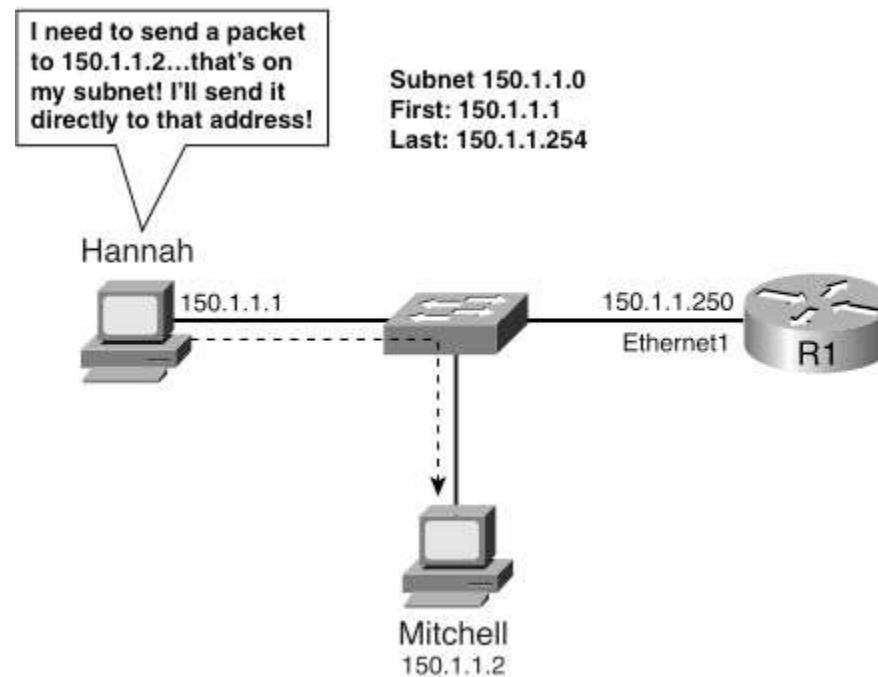
ARP in routing

- ◆ The routing table identifies next hop routers by IP address
 - If that IP address is not already in the router's ARP cache, the router broadcasts an ARP request to find the corresponding MAC address
 - Uses that MAC address on outgoing interface in routing table to forward packet to next hop router



The last hop

- ◆ In some cases, the destination will be on the same subnet as the source
- ◆ If not, the packet eventually gets forwarded to a router on the same subnet as the destination
- ◆ In both cases, the packet needs to be encapsulated in a frame addressed to the MAC address of the destination host
 - If necessary, an ARP broadcast is used to find that MAC address



Starting a routing table

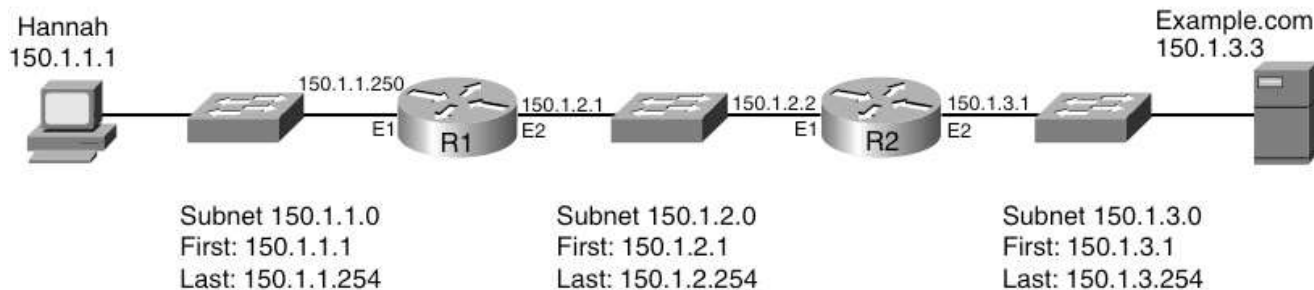
- ◆ Routing is only as good as the info in the routing tables
 - Each router must build its own routing table
- ◆ Start locally
 - Router is configured with info about its interfaces
 - ◆ Has IP address and subnet mask for each interface
 - ◆ Can use that info to determine subnet number for each interface
 - So router knows how to route packets for each of those subnets
 - And that no next hop router is necessary to reach them
 - Adds a route for each of these subnets to its routing table

R1 Routing Table

| Destination Subnets | Interface | Next-Hop Router |
|---------------------|-----------|-----------------|
| 150.1.1.0 | E1 | N/A |
| 150.1.2.0 | E2 | N/A |

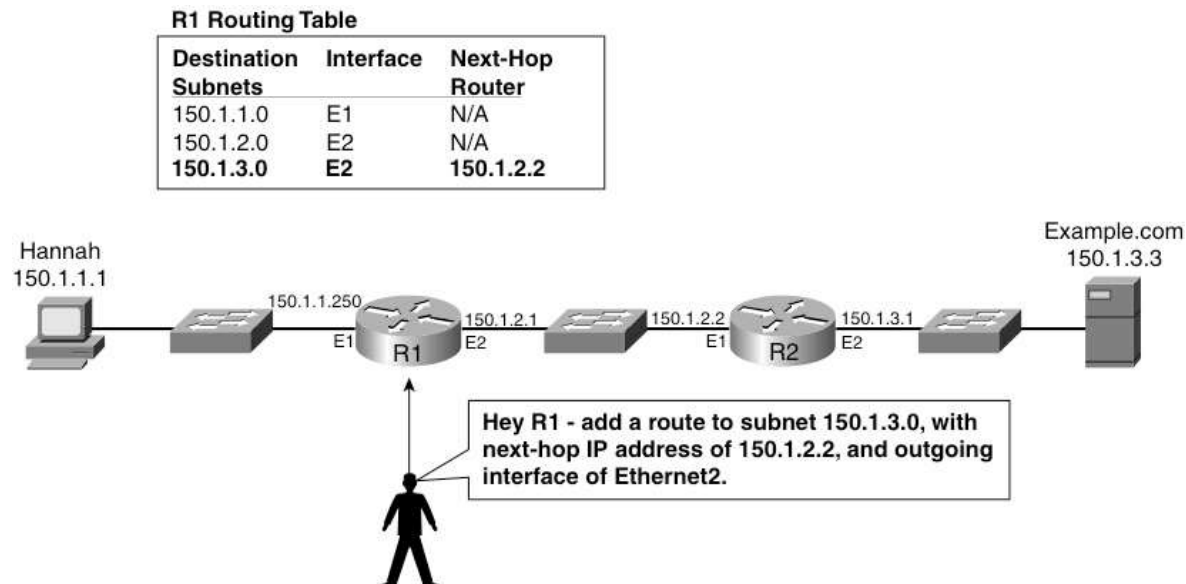
R2 Routing Table

| Destination Subnets | Interface | Next-Hop Router |
|---------------------|-----------|-----------------|
| 150.1.2.0 | E1 | N/A |
| 150.1.3.0 | E2 | N/A |



Adding static routes

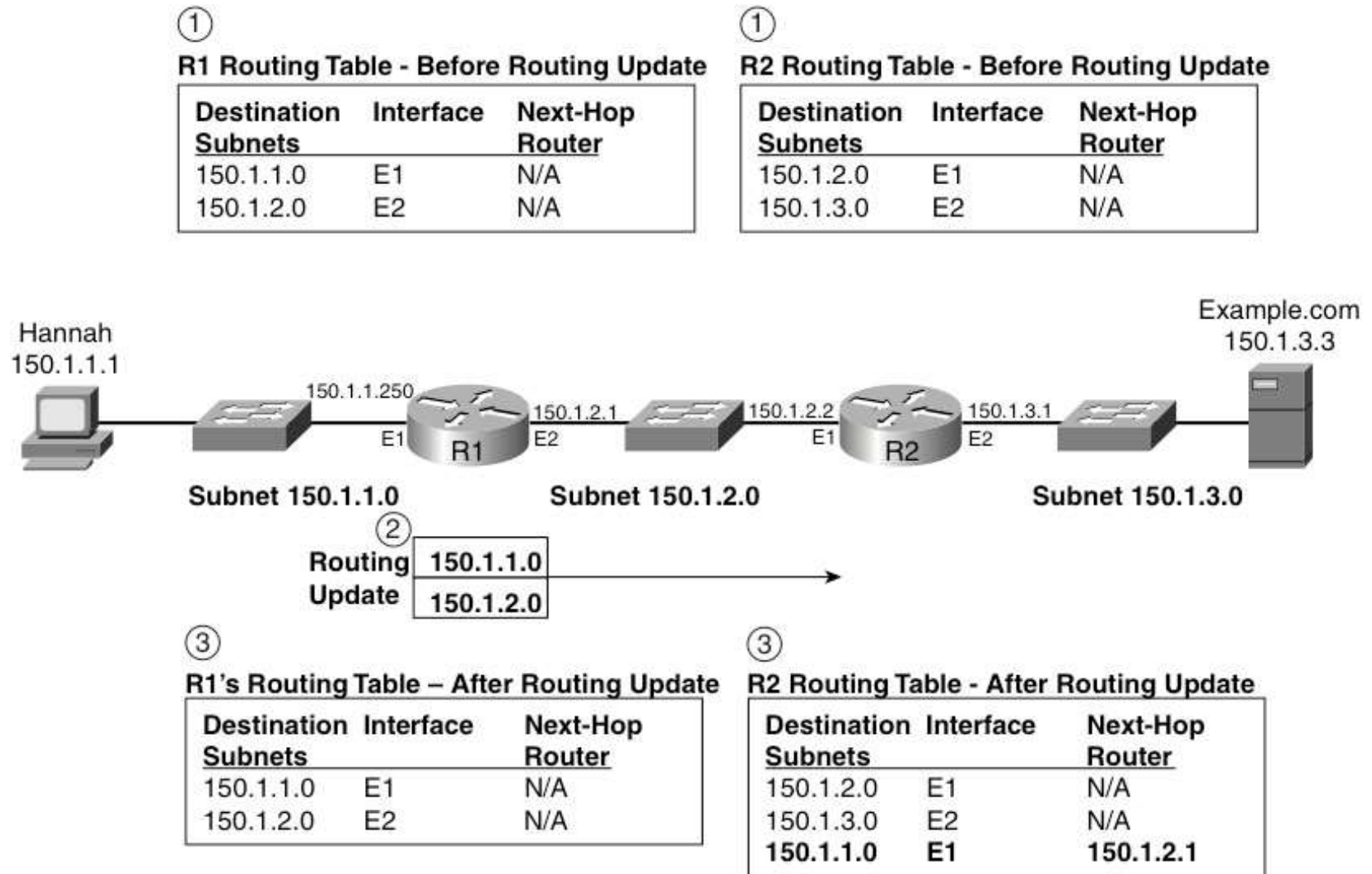
- ◆ Possible to add static routes to a routing table
 - Added to the routing table manually by an administrator
 - Include subnet number, interface, and next hop router
 - ◆ Tell router how to route packets to that subnet
 - Functional, but labor intensive to maintain, so seldom used in practice



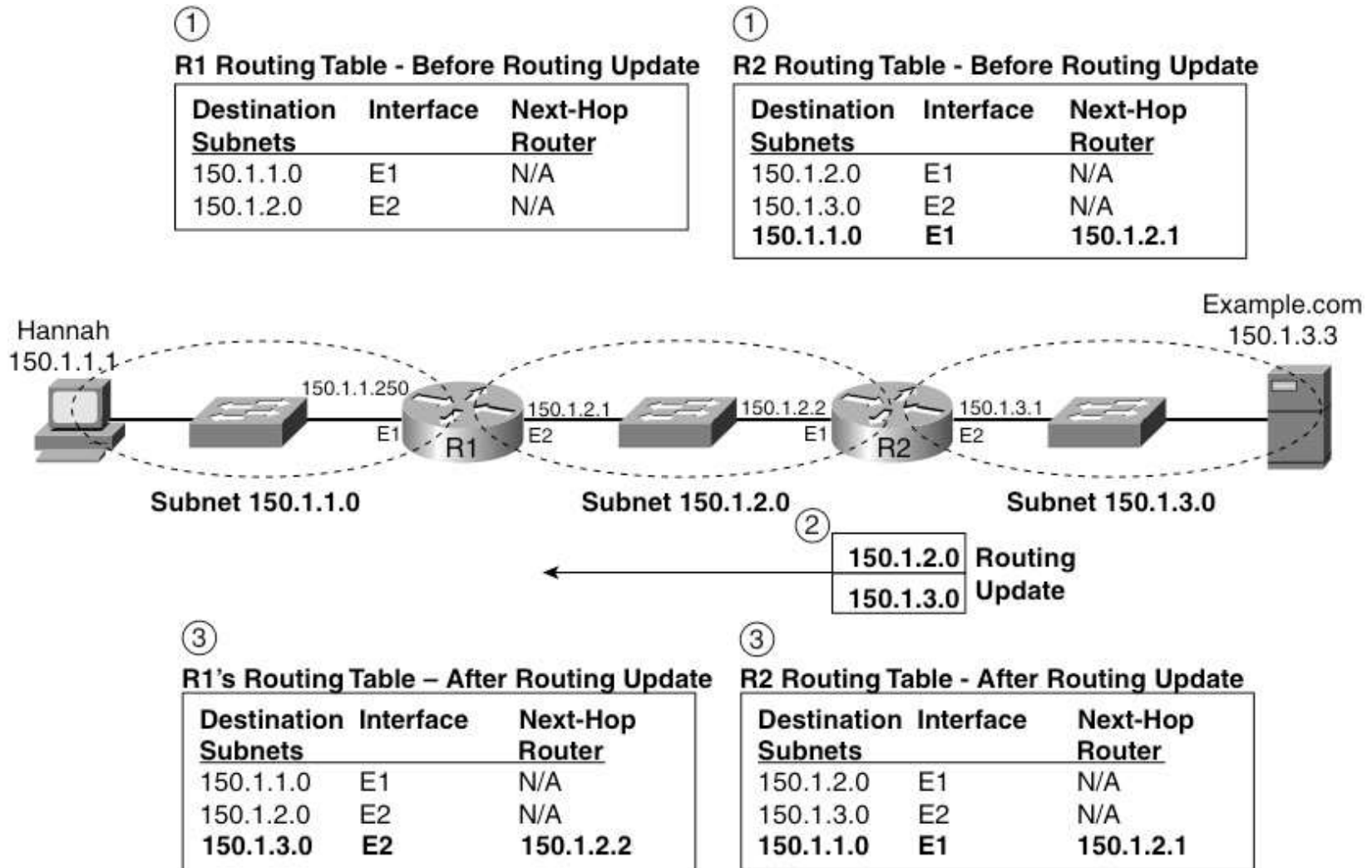
Learning new routes

- ◆ Most routers use a routing protocol to exchange routing info with other routers
 - Routing updates are simply a list of all the subnets in a router's routing table
 - Receiving router can add those subnets to its routing table, listing incoming interface and sending router as next hop router
- ◆ Each router can tell other routers about the routes it knows and learn about new routes from others
 - Allows routers to build and maintain their routing tables autonomously

Learning new routes illustrated – part 1

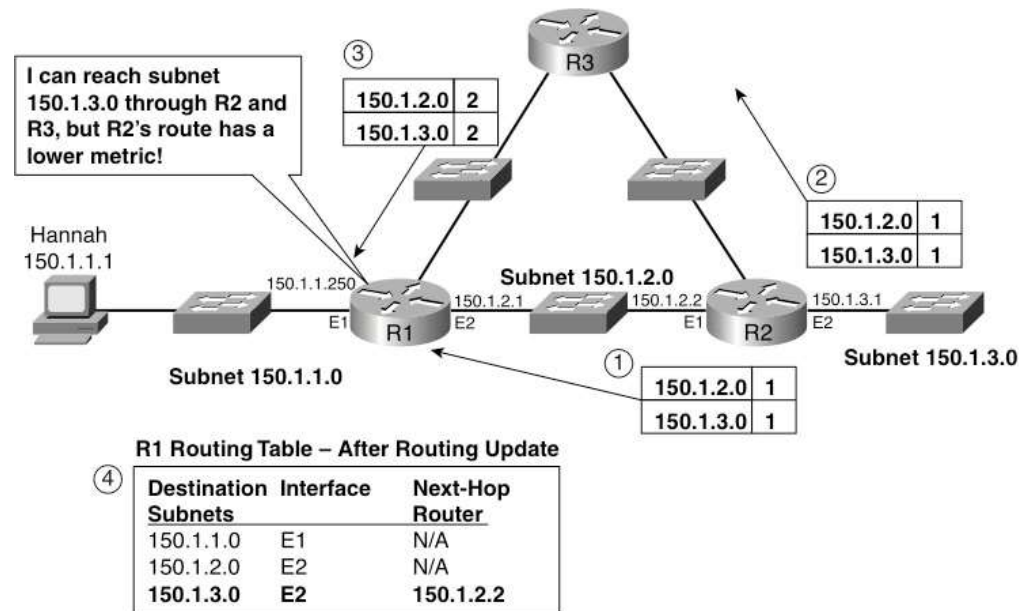


Learning new routes illustrated – part 2



Finding the best route

- ◆ Most internets contain redundant connections
 - When choices exist, routers need to know which is best
 - Routing protocols provide a numeric metric with each route included in a routing update
 - ◆ Lower values indicate better routes
 - ◆ Router typically enters all routes in routing table
 - Uses route with lowest metric if available, alternate route otherwise



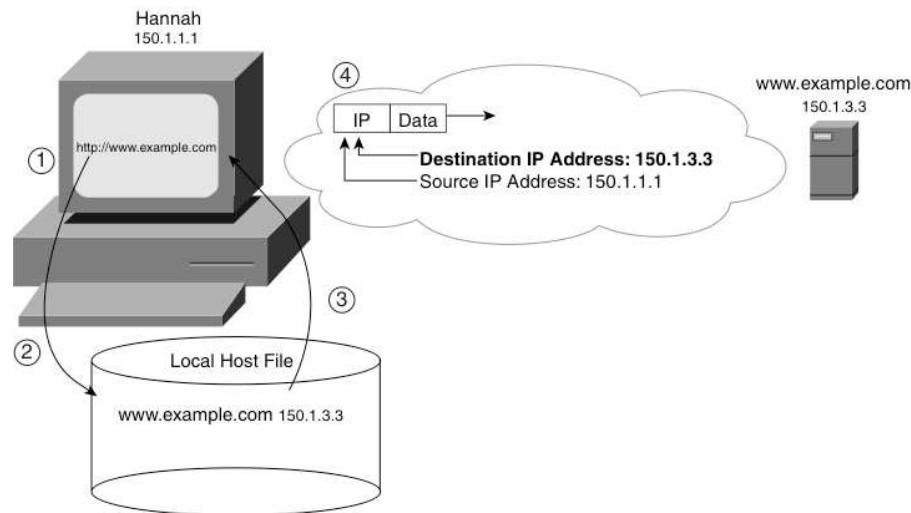
Routing protocols

| Routing Protocol | Public/Proprietary | Interior/Exterior |
|--|--------------------|-------------------|
| Routing Information Protocol (RIP) | Public | Interior |
| Interior Gateway Routing Protocol (IGRP) | Proprietary | Interior |
| Open Shortest Path First (OSPF) | Public | Interior |
| Enhanced IGRP (EIGRP) | Proprietary | Interior |
| Border Gateway Protocol (BGP) | Public | Exterior |

- ◆ Interior routing protocols are primarily intended for use within a single company
- ◆ Exterior routing protocols are primarily intended for use between different companies or between ISPs and their customers
- ◆ IGRP and EIGRP belong to Cisco Systems
- ◆ BGP is used on the Internet

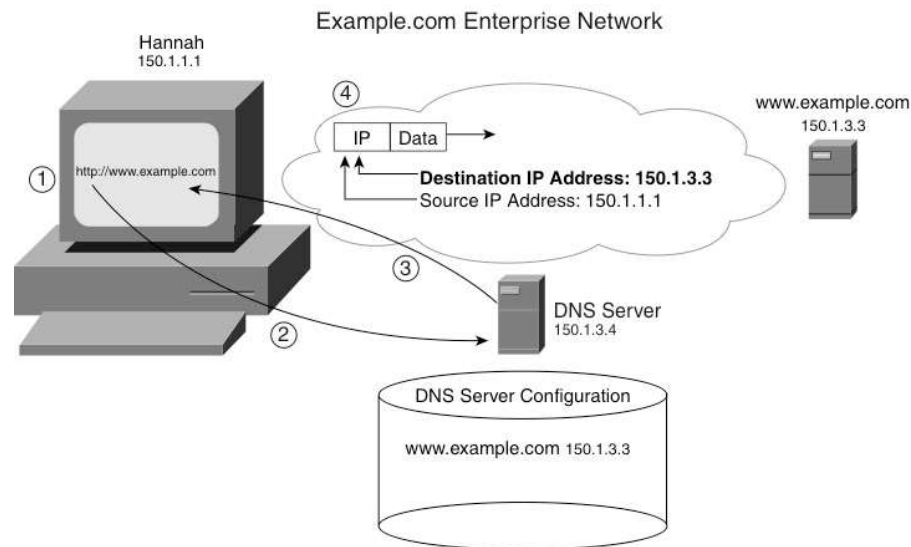
Name resolution

- ◆ IP addresses are unwieldy for humans
 - Many computers are therefore also identifiable by hostname
 - But IP headers must contain IP addresses
- ◆ Name resolution is process of converting from hostname to IP address
 - One approach is to look it up in a local host file
 - ◆ But host files can't possibly be comprehensive and are maintenance intensive
 - ◆ So seldom used in practice



Domain Name Service (DNS)

- ◆ Most hosts use DNS for name resolution
 - Defines protocols for name resolution and the structure of hostnames
 - Most enterprise networks and ISPs include at least one DNS server
 - ◆ Every host on the network is configured with the IP address of one or more DNS servers
 - May be manually configured or automatically configured via DHCP
 - ◆ DNS server is essentially a shared host file for the network
 - It knows the hostname and IP address of every host on that network
 - Hosts send a DNS resolution request to the DNS server whenever they require name resolution
 - An administrator must keep the list on the DNS server up to date for all the hosts on the network



Root DNS servers

- ◆ Each enterprise DNS server is responsible only for resolving host names for its network
 - Doesn't necessarily know how to resolve hostnames outside its network
- ◆ Each enterprise DNS server is configured with the IP addresses of one or more root DNS servers
 - When asked to resolve a hostname it doesn't know, a DNS server will contact a root DNS server for assistance
 - Root DNS servers don't resolve individual hostnames, they resolve domain names
 - ◆ Domain names identify enterprise networks, each with their own enterprise DNS server(s)
 - ◆ Each root DNS server maintains a list associating a domain name with the IP address(es) of the DNS server(s) for that domain
 - ◆ When a root server is asked to resolve a hostname, it returns the IP address of an enterprise DNS server that should know how to resolve it
 - ◆ The original DNS server must then send another name resolution request to that domain's enterprise DNS server
 - ◆ Most servers will store the answer in a local cache for future reference

Root DNS servers illustrated

