Cloud Computing Problem Formulation

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1 Problem Formulation

1.1 Center

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\mathbf{O}(t) = \text{all orders sequence in the queue in time slot t, e.g.}(o_i, o_{i+1}, ..., o_m)

\vec{a}(t) = \text{new arrival orders sequence in time slot t, e.g.}(o_j, o_{j+1}, ..., o_m)

\mathbf{O}(t+1) = \mathbf{O}(t) + \vec{a}(t) - \sum_{i=1}^{n} r_i(t)
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1.2 Locker

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\begin{aligned} \vec{\mathbf{u}_i} &= \text{available delivers sequence in locker i, e.g.}(u_i, u_{i+1}, ..., u_m) \\ \mathbf{o}_i(t) &= \text{all orders sequence in the queue in time slot t for locker i, e.g.}(o_i, o_{i+1}, ..., o_m) \\ r_i(t) &= \text{new arrival orders sequence for locker i in time slot t, e.g.}(o_j, o_{j+1}, ..., o_m) \\ h_i(t) &= \text{orders delivered for locker i in time slot t} \\ \mathbf{o}_i(t+1) &= \mathbf{o}_i(t) + r_i(t) - h_i(t) \\ h_i(t) &= \sum_{u \in \vec{\mathbf{u}_i}} b_u(t) \end{aligned}
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1.3 Deliver

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\begin{array}{l} \lambda_u(t) = \text{is deliver u active in time slot t, 0,1} \\ l_u = \text{locker position for deliver u} \\ \rho_u = \text{origin position for deliver u} \\ b_u(t) = \text{orders delivered by u in time slot t, e.g.}(o_i, o_{i+1}, ..., o_m) \\ d_u^{max} = \max \text{ distance for deliver u per time slot} \\ c_u^{max} = \max \text{ capacity of deliver u per time slot} \\ d_u(t) = \text{ distance for deliver u in time slot t} \\ c_u(t) = \text{ capacity for deliver u in time slot t} \\ \vec{v}_u(t) = \text{ paths for deliver u in time slot t, e.g. } (\rho_u, l_u, p_0, p_1, ..., p_n, \rho_u), \text{ where} \\ p_i \in \mathbf{V} \\ \mathbf{E}_u(t) = \text{ edges for deliver u in time slot t, and } \mathbf{E}_u(t) \subseteq \mathbf{E} \\ d_u(t) = \sum_{(u,v) \in \mathbf{E}_u(t)} w(u,v) \\ \sum_{o \in b_u(t)} g_o \leq c_u^{max} \\ \sum_{o \in b_u(t)} \mu_o \subseteq \vec{v}_u(t) \end{array}
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Order 1.4

 g_o = the capacity for order o

 σ_o = deadline of order o

 $\mu_o = \text{destination address of order o}$

1.5 Map

V = nodes in the city

 $\mathbf{E} = \text{edges}$ in the city

w(u, v) = the distance between position u and position v

1.6 **Payment**

 $\beta = \text{constant commission fee}$

 γ = payment for unit distance $\delta(t)$ = payment for time slot t

$$\Delta = \text{the whole payment}$$

$$\delta(t) = \sum_{u=1}^{n} \lambda_u(t) \cdot (\gamma d_u(t) + \beta)$$

$$\Delta = \sum_{t=1}^{T} \delta(t)$$

$$\Delta = \sum_{t=1}^{T} \delta(t)$$

1.7 Constraints

$$\sum_{n} g_o \le c_u^{max}$$

$$o \in b_u(t)$$

$$\sum_{o \in \vec{b_u(t)}} g_o \le c_u^{max}$$

$$\sum_{o \in \vec{b_u(t)}} \mu_o \subseteq \vec{v_u(t)}$$

$$o \in b_u(t)$$

$$\forall o \in \vec{h_i(t)}, t \leq \sigma_o$$