

DataStax Enterprise Architecture

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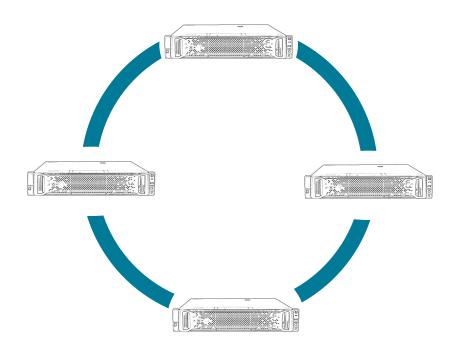
5. October 2017

Agenda

1	Topology and Data Structure
2	Request Handling
3	Lab1: Cassandra Access and Cassandra Stress



Design Goals and Objectives



- Continuously Available
- Master Less
- Fully Distributed
- Shared-Nothing Architecture
- Build In Replication
- Linear Scalability
- Scale out



Architecture

Apache CassandraTM Architecture

- Cluster layer
- Amazon DynamoDB paper
- masterless architecture

- Data-store layer
- Google Big Table paper
- Columns/columns family

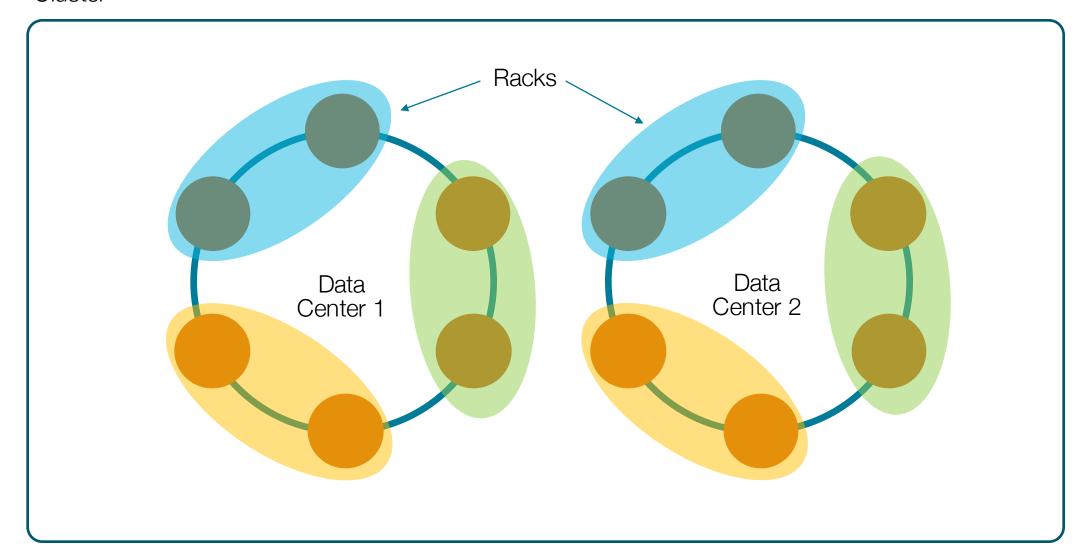




- All nodes are peers
 - Including seed nodes
 - No master
 - Discovery through gossip
- Built-in replication
 - Simplify your architecture!



Cluster



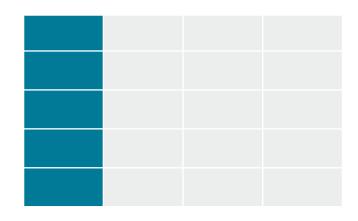


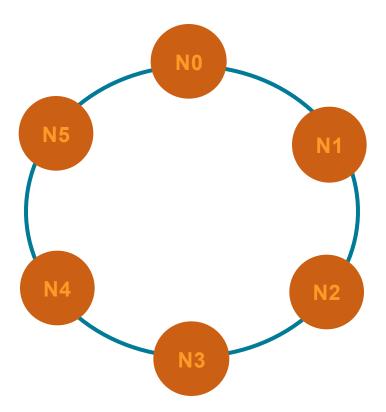
Tokens

Data is partitioned after its partition key

A unique token is allocated to a partition

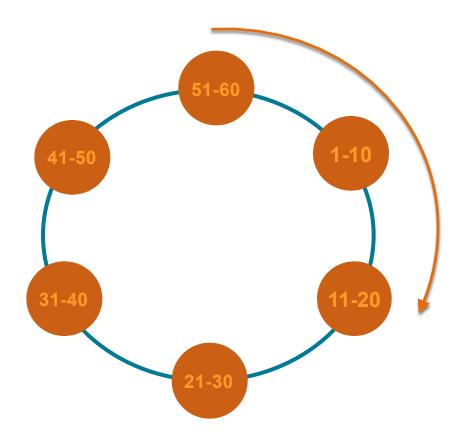
Token = random hash of #partition







Token Ranges



Token Range: 2⁻⁶³ - 2⁶³

Example with **Replication Factor 3**

N3 will own data for tokens 1 – 30

Token Range: 1-10, owned by N1,N2,N3 Token Range: 11-20, owned by N2,N3,N4 Token Range: 21-30, owned by N3,N4,N5



Primary key

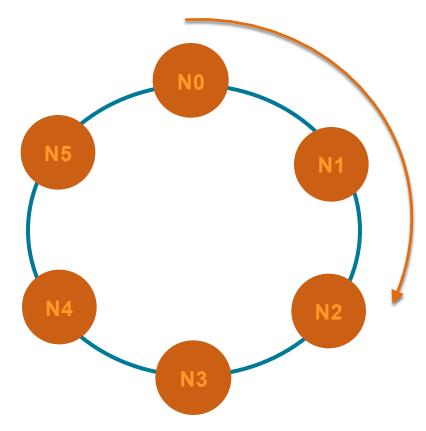
- Partition key
- Clustering columns
- Partitioner
 - Generates unique hash from partition key
- Replication strategy
 - Token hash determines starting point
 - Determines replica placement



Data Distribution

Token = hash of #partition → #node

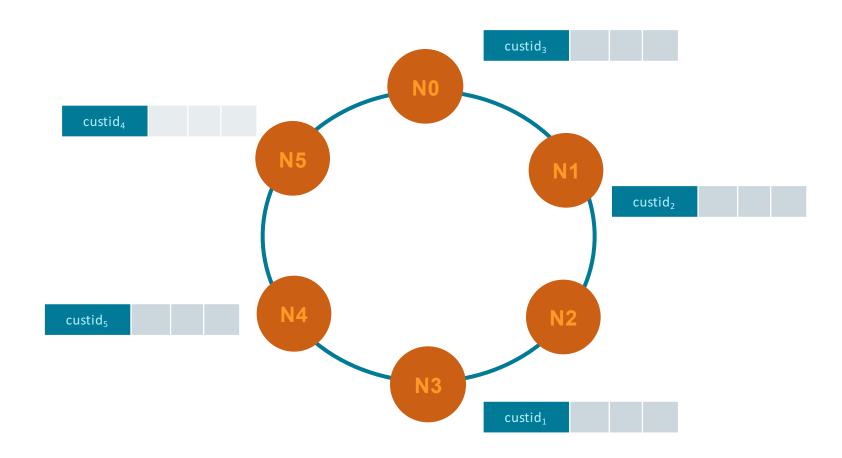
Token1	custid 1		
Token2	curstid 2		
Token3	curstid 3		
Token4	curstid 4		
Token4	curstid 5		



Data is evenly distributed and clock wise replicated



Data Distribution





Cassandra Query Language DDL / DML

```
CREATE KEYSPACE retailer WITH replication =
{'class': 'NetworkTopologyStrategy', 'DC1': '3'} ← Replication Factor
AND durable writes = true;
CREATE TABLE retailer.sales by customer (
   custid int,
   salesdt text,
   comment text,
                          Partition Key
   discount double,
   revenue double,
   PRIMARY KEY (custid, salesdt));
```

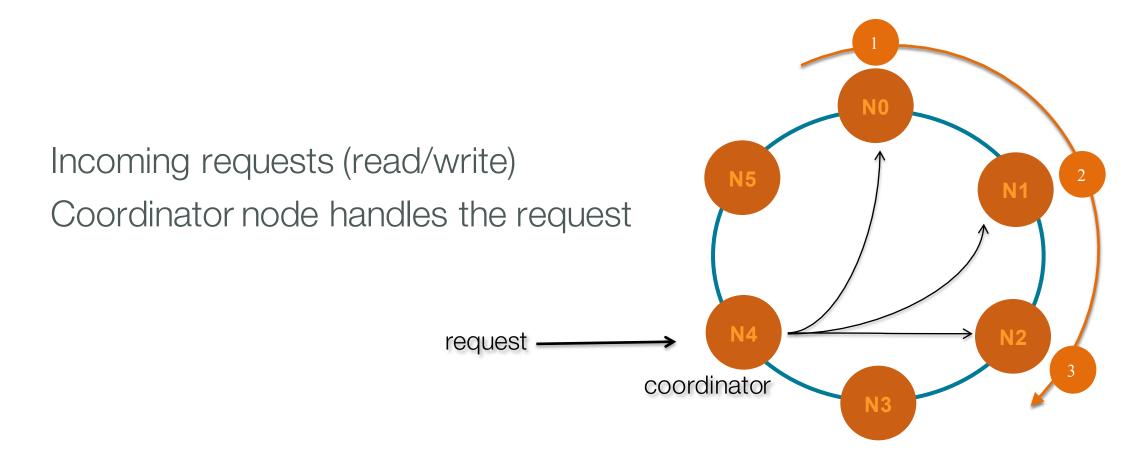
SELECT * FROM sales_by_customer where custid=1 OR custid=2 AND salesdt >=20160401;



Lab 1: Accessing the cluster

Tunable Consistency Read and write request handling

Coordinator node



Every node can be coordinator \rightarrow masterless



Consistency

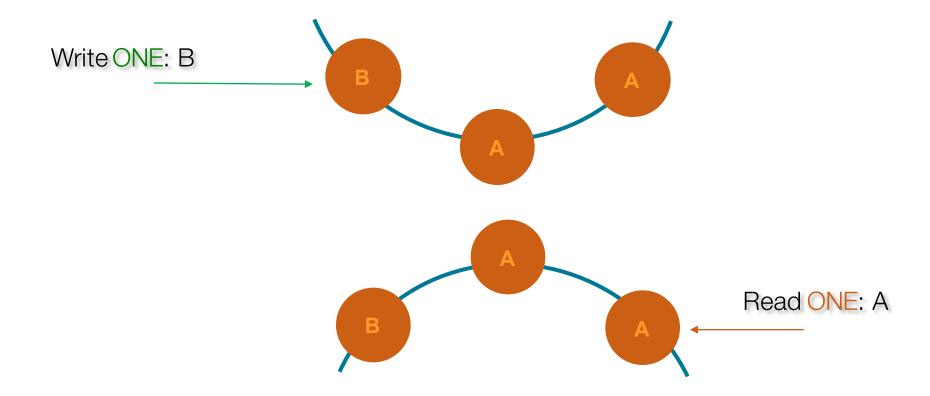
Tunable at runtime

- ONE
- QUORUM (strict majority w.r.t. RF)
- ALL

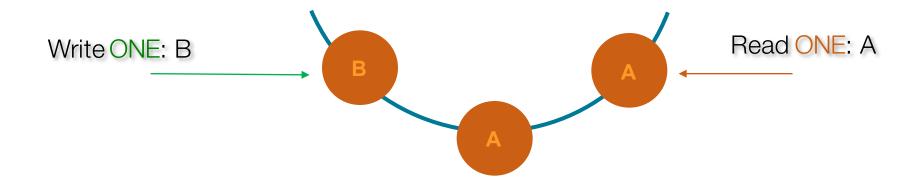
Apply both to read & write



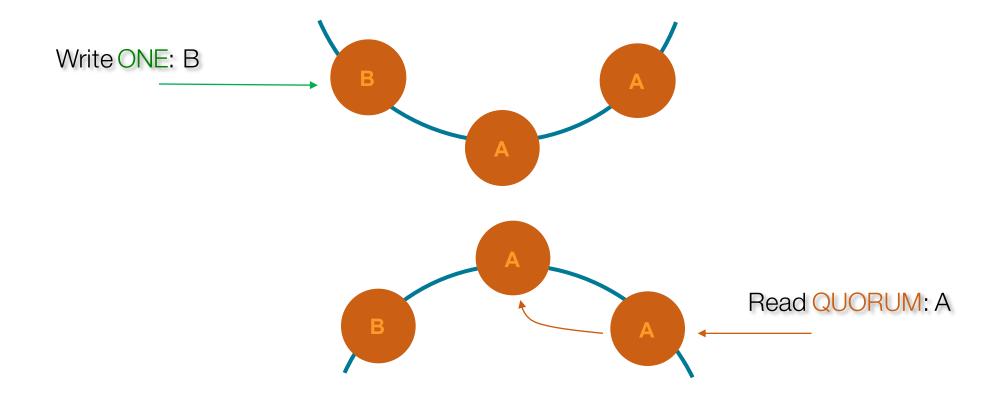
• RF = 3, Write ONE, Read ONE



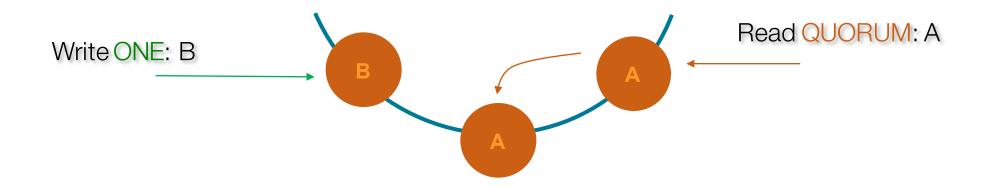
• RF = 3, Write ONE, Read ONE



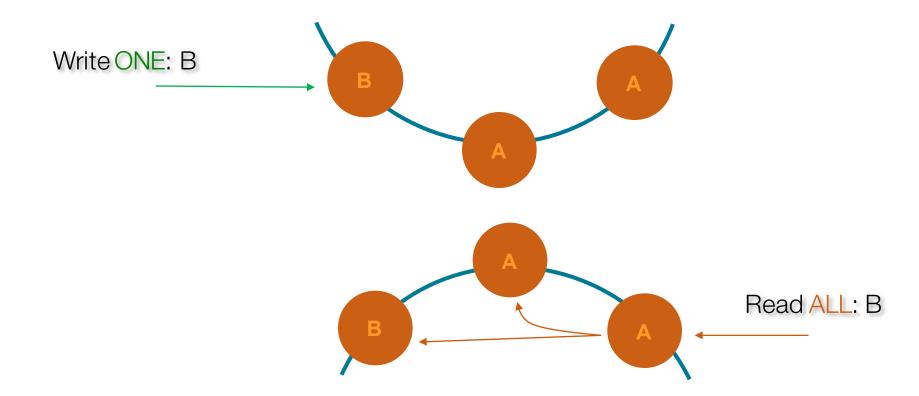
• RF = 3, Write ONE, Read QUORUM



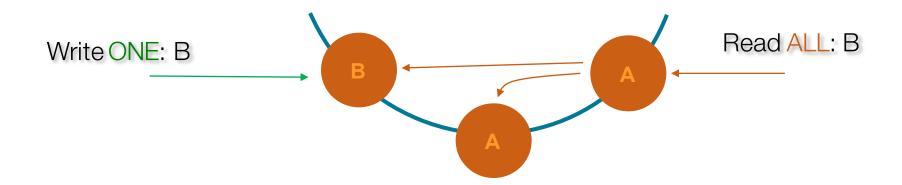
• RF = 3, Write ONE, Read QUORUM



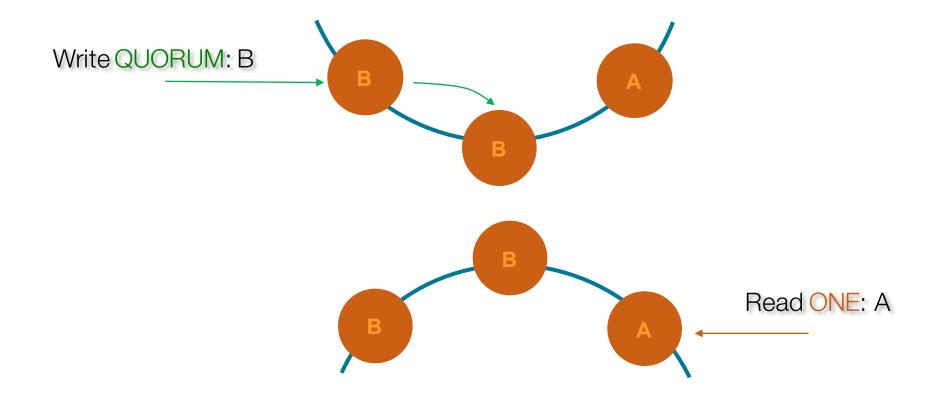
• RF = 3, Write ONE, Read ALL



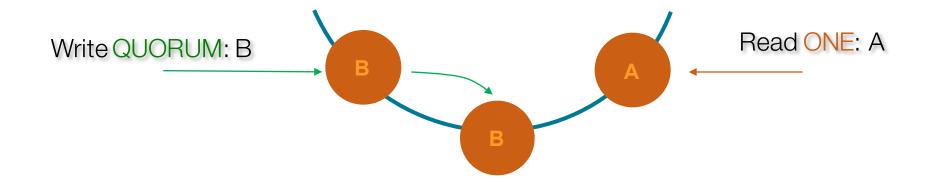
• RF = 3, Write ONE, Read ALL



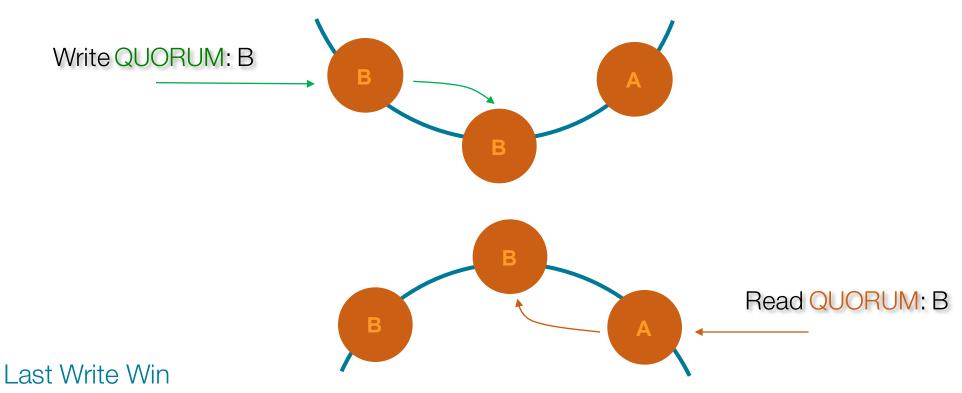
• RF = 3, Write QUORUM, Read ONE



• RF = 3, Write QUORUM, Read ONE

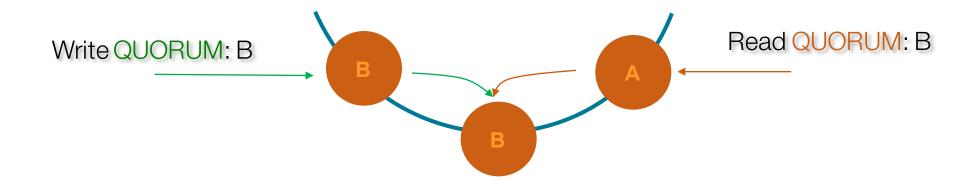


RF = 3, Write QUORUM, Read QUORUM



- R+W > RF = immediate consistency
- Background vs. foreground Read Repair. See more...

RF = 3, Write QUORUM, Read QUORUM



- Last Write Win
- R+W > RF = immediate consistency
- Background vs. foreground Read Repair. See more...

Consistency trade-off

Latency

Consistency



Consistency summary

ONE_{Read +} ONE_{Write}

available for read/write even (N-1) replicas down

QUORUM_{Read} + QUORUM_{Write}

available for read/write even 1+ replica down



- Background vs. Forground Read Repair
 - Compare digests
 - If any mismatch
 - re-request to same nodes (full data set)
 - compare full data sets, send update
 - block until out-of-date replicas respond
 - Return merged data set to the client
- Consistency Level
 - one, quorum, all
 - local vs. cluster wide



CRUD

```
INSERT INTO sales_by_customer (custid, salesdt, revenue) VALUES('1', '20160103', 799);

UPDATE sales_by_customer SET discount = 10 WHERE custid = 1;

DELETE sales_by_customer FROM sales_by_order WHERE custid = 1;

SELECT revenue FROM sales_by_customer WHERE custid = 1;
```



Time To Live

```
INSERT INTO users(...) VALUES(...) USING TTL = 3600;
UPDATE users USING TTL = 3600 SET age = xxx WHERE ...;
```



Time To Live

```
CREATE TABLE readings_by_sensor(
    sensor_id text,
    collect_time timestamp,
    electricity long,
    gas long,
    health double,
    PRIMARY KEY (sensor_id, collect_time)
)
WITH CLUSTERING ORDER BY (collect_time ASC)
AND default_time_to_live = 86400; // TTL 1 day
```

INSERT INTO readings_by_sensor(...) VALUES(...) USING TTL = 86400;



Collections

```
CREATE TABLE users (
    login text,
    name text,
    age int,
    friends set<text>,
    hobbies list<text>,
    languages map<int, text>,
    ...
    PRIMARY KEY(login));
```



User Defined Types (UDT)

```
CREATE TYPE address (
         street_number int,
         street_name text,
         postcode int,
         country text);

CREATE TABLE users (
        login text,
        ...
        location address,
        ...
        PRIMARY KEY(login));
```



UDT DML



Light Weight Transactions

```
INSERT INTO users(...) VALUES(...) IF NOT EXISTS;

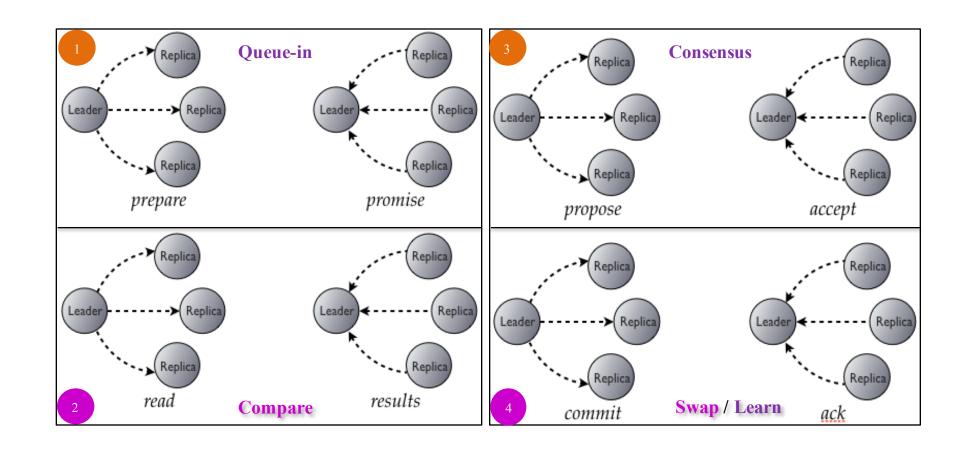
DELETE users WHERE ... IF EXISTS; UPDATE users

SET age = xxx WHERE ... IF age = 30;
```

Linearizable writes on a single partition



Lightweight Transaction (LWT) PAXOS implementation





Driver Code

```
Cluster cluster = Cluster.builder()
      .addContactPoint("127.0.0.1")
      .withLoadBalancingPolicy(new TokenAwarePolicy(DCAwareRoundRobinPolicy.builder()
          .withLocalDc("myLocalDC")
          .build())
      .build();
PreparedStatement prepared = session.prepare
      ( "insert into sales by customer(custid, salesdt) values (?, ?)");
BoundStatement bound = prepared.bind("1", "20170102");
session.execute(bound); // Throws <u>UnavailableException</u> If consistency doesn't met, downgrade is
                             possible with corresponding RetryPolicy. Read More...
```



Lab 2: Hands-on DSE CQL

Vielen Dank!

Eventuell Bootstrap, ReBalance, Num Tokens VNodes

- All nodes are peers
 - Including seed nodes
 - No master
 - Discovery through gossip
- Built-in replication
 - Simplify your architecture!

