

# CSC236 - Week 1, Lecture 2

Cristyn Howard

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- abstract data type (ADT) - an object and its operations
- data structure - specific implementation of some ADT
- One example of an ADT is a Priority Queue (PQ).
  - object: maintains a set  $S$  of elements with keys that can be compared
  - operations:
    - \*  $\text{INSERT}(x, S)$ : insert item  $x$  into set  $S$  such that the order is maintained
    - \*  $\text{MAX}(S)$ : return a value with the maximum priority
    - \*  $\text{EXTRACTMAX}(S)$ : find an element with max priority and remove it from the set
- PQ application example: might be used by an operating system to organize processes that need to be run
- Some data structures for PQs:

	Data Structure	Worst-case run time of:	
		INSERT	EXTRACTMAX
–	unsorted linked list	$O(1)$	$O(n)$
	sorted linked list	$O(n)$	$O(1)$
	heaps	$O(\log n)$	$O(\log n)$

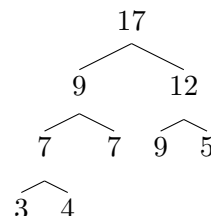
- Note that heaps are superior! So what are they, and how do they work?

Recall: Complete Binary Tree (CBT): start at the leftmost node, every node has at most two children, fill every level before moving to the next

Recall: Height of binary tree: length of longest path from the root to any leaf of the tree.

- For CBT of size  $n$ , height =  $\lfloor \log_2 n \rfloor$

- heaps: store set  $S$  of size  $n$  in an  $n$ -node complete binary tree
  - elements in a heap are stored such that the priority of any node is greater than the priority of its children
  - heaps for set  $S$  are not unique, i.e. there are many valid heap configurations of set  $S$
  - Example:  $S = \{3, 4, 5, 7, 7, 9, 9, 12, 17\}$ , A heap containing  $s$ :



- In a computer, heaps are stored in a (1-indexed) array such that the elements represent the nodes of the heap ordered from top-to-bottom, left-to-right.

– An array storing our heap of S:

Index:	1	2	3	4	5	6	7	8	9
Value:	17	9	12	7	7	9	5	3	4

- How do we relate parents to children (and vice versa) when the heap is stored in an array?

Parents to Children	Children to Parents
right child index = parent index $\times$ 2	parent index = $\lfloor \text{child index} / 2 \rfloor$
left child index = (parent index $\times$ 2) + 1	

- How are the PQ operations implemented on a heap stored as an array?

- INSERT(): (A) add new element to last space in array; (B) get parent index, get parent value; (C) if parent value is smaller than new element value, swap element locations, go to B.

\* worst case run time of insert: compare & swap take constant time, the upper bound on the number of swaps is the height of the tree, thus  $\text{Insert} \in \Theta(\log n)$ .

·  $O(\log n)$ : every insert takes at most  $c \times \log n$  steps, for some constant  $c$ .

·  $\Omega(\log n)$  :  $\exists$  some element for which insert takes at least  $c \times \log n$  steps, (e.g.  $x > \text{root}$ ).

- MAX():  $\Theta(1)$ , get first element of the array

- EXTRACTMAX(): (A) swap first & last elements of the array; (B) return last element & reduce heap size by one; (C) compare first element value to values of both children, swap element with larger of two children; (D) continue until element is larger than both children, or until the indices of both children are greater than heap size (i.e. element has no children)

\*  $\Theta(\log n)$  for reasons similar to insert

- SORTING: extract max until all elements have been removed from the array,  $\Theta(n \times \log n)$