

Delete from BST; DELETE(X)

- ①  $x$  has no children; trivial deletion —  $O(1)$
- ②  $x$  has one child; change one pointer —  $O(1)$
- ③  $x$  has 2 children; find successor, remove successor (recursive delete call), put successor value in  $x$  location —  $O(n)$

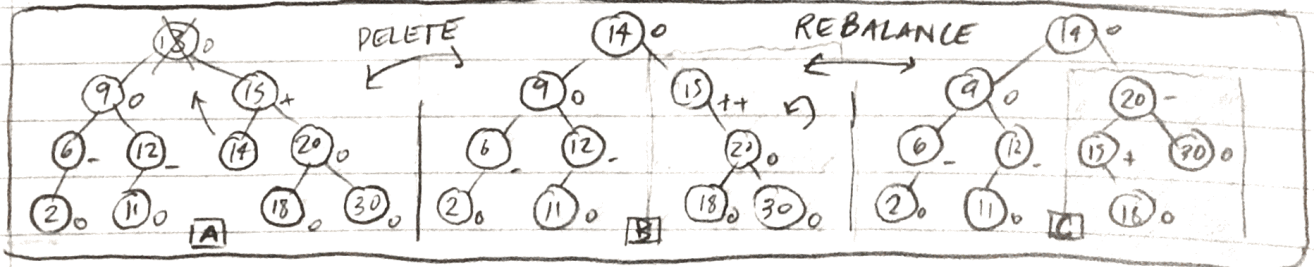
height of  
BST not  
guaranteed  $\log n$

- Note: there will be questions that test your knowledge of the difference between BST and balanced trees

Delete X from AVL tree:

- find node  $k$  where  $x$  is stored, delete node  $x$  as in regular BST (outlined above), then rebalance the tree
- in all three cases, we end up deleting a leaf;
  - ①  $x$  has no children,  $x$  is leaf, delete leaf  $x$
  - ②  $x$  has one child  $x'$ ,  $x'$  must be leaf to preserve balance when  $x'$  replaces  $x$ , leaf  $x'$  is eliminated
  - ③  $x$  has two children, find successor (which by definition has at most one child), deleting successor is case ① or case ②
- every time deletion is made, go from deleted node to root and rebalance whenever unbalanced node encountered

EX:



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