

## A1 Question 2

Lots of things had to be done to boyer moore to make it work the other way around:

### **(extended) Bad character rule:**

In Boyer moore, we process the bad character rule by iterating through the indices backwards, keeping a dictionary of the latest occurrence of each character and updating the bad character dictionary. The bad character dictionary is a dict that maps characters onto lists, and the lists contain at index  $i$  the leftmost occurrence of the character that isn't at index  $i$  or less.

At every index, we update the latest occurrence for the character at that index. It's quite similar to forwards preprocessing of the bad character rule but we iterate backwards instead.

### **Good prefix:**

To preprocess the good prefix rule, we first need `z_prefix`. Thankfully, `z_prefix` is the output of the `z` algorithm anyway, so we can just run `z` algorithm on the pattern to get what we need. Whereas In the original implementation detailed in the notes, the good suffix and matched prefix tables were both size 1 more than the length of the pattern, I didn't do that for the good prefix and matched suffix arrays because I didn't see the need.

Instead of iterating forwards, we iterate backwards through the pattern. This ensures that we get the leftmost occurrence of each prefix. Whenever we encounter a `Z` box that's bigger than 0, then we update the index in the `good_prefix` array corresponding to the size of the `z` box with the index that it occurred. This allows us to get the index of a re-occurrence of a prefix based on the size of the prefix, or check if a re-occurrence exists.

### **Matched Suffix:**

To preprocess Matched suffix, we need `z_suffix`. We do this by reversing the pattern, running `z` algorithm on it and then reversing that array.

We can then iterate through the indices, updating an index at matched suffix to a new number only if the `zsuff` value there (size of the `z` box ending there) is one more than the index, which means that that `z` box covers the entire prefix from 0 to  $i$ , and also matches the suffix, which means it's a new largest matched prefix.

We then iterate through again to fill in the gaps with their previous matched suffix, which makes sense since the biggest suffix from  $i+1$  to the end is going to be the biggest matched suffix from  $i$  to the end unless there's a new biggest suffix from  $i$  to the end.

### **Galil's Optimisation:**

Unfortunately I don't have the time to make Galil's optimisation work, but it works much the same way as in regular boyer moore. Have a range that you've already compared, and skip it for the next iteration