

Additional type rules to "Extensional Normalisation and Type-Directed Partial Evaluation for Typed Lambda Calculus with Sums"

$$\frac{}{\cdot \text{ctx}} \text{ctx-EMP} \quad \text{Empty context}$$

Making "lower" variables

$$\frac{\Gamma \vdash A \text{ type}}{\Gamma :: (x:A) \text{ ctx}} \text{ctx-EXT1}$$

$$\Gamma :: (x:A) \text{ ctx}$$

concatenation variable not in Γ

Making type variable

$$\frac{\Gamma \text{ ctx}}{\Gamma :: (A \text{ type}) \text{ ctx}} \text{ctx-EXT2}$$

A context Γ is hence a list
of membership judgements $b:B$
and type judgements $A \text{ type}$

corresponds to membership judgement

$$A:U_0 \text{ in HoTT}$$

0th level universe

$$\frac{\Gamma \vdash A : \text{type} \quad \Gamma \vdash B : \text{type}}{\Gamma \vdash A \times B : \text{type}} \quad \times \text{ intro}$$

$$\frac{}{\Gamma \vdash 1 : \text{type}} \quad \text{Unit intro}$$

$$\frac{}{\Gamma \vdash 0 : \text{type}} \quad \begin{array}{l} \text{empty} \\ \text{type} \end{array} \text{ intro}$$

$$\frac{\Gamma \vdash A : \text{type} \quad \Gamma \vdash B : \text{type}}{\Gamma \vdash A + B : \text{type}} \quad \begin{array}{l} \text{sum type} \\ \text{intro} \end{array}$$

$$\frac{\Gamma \vdash A : \text{type} \quad \Gamma \vdash B : \text{type}}{\Gamma : A \rightarrow B \text{ type}} \quad \begin{array}{l} \text{function} \\ \text{type} \\ \text{intro} \end{array}$$

$$\frac{\Gamma :: (x : A) :: \Gamma' \text{ ctx}}{\Gamma :: (x : A) :: \Gamma' \text{ vble}}$$

$$\Gamma \vdots (\lambda x. A) :: \Gamma \vdash x : A$$

$$\frac{\Gamma :: (B \text{ type}) :: \Gamma' \text{ ctx}}{\Gamma :: (B \text{ type}) :: \Gamma' \vdash B \text{ type}} \text{Vble}_2$$

We may also want ^{optional} to include the following substitution & weakening rules from the HoTT book

$$\frac{\Gamma \vdash A \text{ type} \quad \Gamma :: \Delta \vdash b : B}{\Gamma :: (x : A) :: \Delta \vdash b : B} \text{Wk}_1$$

$$\frac{\Gamma \vdash A \text{ type} \quad \Gamma :: \Delta \vdash K \text{ type}}{\Gamma :: (x : A) :: \Delta \vdash K \text{ type}} \text{Wk}_2$$

$$\frac{\Gamma \vdash a : A \quad \Gamma :: (x : A) :: \Delta \vdash b : B}{\Gamma :: \Delta[a/x] \vdash b[a/x] : B} \text{Subst.}$$

← replace x with a
in term b

$$\Gamma \vdash J \text{ type} \quad \Gamma :: (L \text{ term}) :: \Delta \vdash b : B$$

$$\frac{\Gamma :: \Delta[J/L] \vdash b[J/L]}{\text{Subst}_2}$$

$$\frac{\Gamma \vdash a:A \quad \Gamma :: (x:A) :: \Delta \vdash V \text{ type}}{\Gamma :: \Delta[a/x] \vdash V \text{ type}} \text{Subst}_3$$

$$\frac{\Gamma \vdash J \text{ type} \quad \Gamma :: (L \text{ type}) :: \Delta \vdash V \text{ type}}{\Gamma :: \Delta[J/L] \vdash b[J/L]} \text{Subst}_4$$

Last modified: 17:03