

and then the duel problem of LASSO is:

man property

with  $Q = \frac{1}{2} I_{\nu}$   $A = \begin{pmatrix} X^{T} \\ -X^{T} \end{pmatrix}$ 

P= y7

2) we want to helve the problem min vtav+ ptv

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There is no equality constraints; the associated contening protein is.

min  $t(v^T Q V + P^T V) - \sum_{i=1}^{2d} \log(b_i - A_i V)$ 

where A; is the inthe row of A.

we denote  $\phi(v) := -\frac{ed}{2}\log(b_1 - A_2 v)$ , and  $g_{\epsilon}(v) = t(v \circ v + e^{\tau}v) + \phi(v)$ 

 $\nabla \phi(v) = \frac{2d}{i=1} \frac{A_i^T}{b_i - A_i v} \Rightarrow \nabla g(v) = t((Q + Q^T) v + P) + \frac{2d}{i=1} \frac{A_i^T}{b_i - A_i v}$ 

 $\nabla^2 \phi(v) = \sum_{i=1}^{2d} \frac{A_i^T A_i}{(A_i v - b_i)^2} \Rightarrow \nabla^2 g(v) = t(0 + 0^T) + \sum_{i=1}^{2d} \frac{A_i^T A_i}{(A_i v - b_i)^2}$ 

We use Newton's method for uncontrained extimization to solve min give

E) we observe that when  $\mu$  is small ( $\mu$ =2), we do small centering steps and it takes a lot of Newton iterations to get to a solution with precision  $\varepsilon$ .

For higher values of  $\mu$ , we do songer steps on the central path and the total number of Newton iterations is much less than in the case with a small  $\mu$ . We also notice that the number of Newton iterations for  $\mu$  > 10 stays fairly constant when we irrepease  $\mu$ . Thus, it's sufficient to choose  $\mu$ =10

Here is the figure obtained for a precision  $\epsilon\text{=}10^{\text{-}6},$  for different values of  $\mu$ 

