

friday

Notes



Topic Name : Time-Complexity 2

TODAY:

- => 1 - Comparing 2 Algorithms
- => 2 - Asymptotic Analysis : Big-O
Why neglect lower order terms?

AGENDA:

Why drop constant coefficient?

Lift

3 - Issues in Big-O notation

4 - Time & Space Complexities

5 - Mathematical and Graphical Representation
of Big-O. Problems.

6 - Why TLE? How to know before
submitting code.

7 - Worst Case, Best Case, Avg Case

Lift
Arrive

Big-O Notation

How to find Big-O Notation:

1. Calculate the no. of iterations based on the Input size.
- ✓ 2. Neglect the lower order terms.

$$\begin{matrix} N^2 \\ \times \end{matrix}, \begin{matrix} N \\ \times \end{matrix}, \begin{matrix} N^3 \\ \boxed{\quad} \end{matrix}$$

$$N^3 + N^2 + N = \underline{\underline{O(N^3)}}$$

Keep the highest order term.

- ✓ 3. Neglect the constant coefficients.

$$\begin{matrix} X \\ 4N^2 \\ \underline{\underline{= O(N^2)}} \end{matrix} \qquad \begin{matrix} X \\ 6N^2 \\ \underline{\underline{= O(N^2)}} \end{matrix}$$

Order of N^2

Big-O

Quid-1

$$F(N) = \boxed{N^3} + \begin{matrix} 10 \\ \uparrow \end{matrix} \begin{matrix} N^2 \\ \uparrow \end{matrix} + \begin{matrix} 8 \\ \uparrow \end{matrix}$$

$$P = 8 \cdot \underline{\underline{N^0}}$$

$$N^3 + \boxed{N^2 + N^0}$$

$$F(N) = \underline{c} \cdot N^0 = \underline{\underline{O(1)}}$$

$\times \quad c \text{ is constant.}$

Mathematical Relation : Common Complexities.

$$\boxed{N > 2} \quad \boxed{\log(N) < \sqrt{N} < N < N \log N < N\sqrt{N} < N^2 < 2^N}$$

$\underline{N=64}$ ||

$$\log_2(N) = \log_2 64 = 6$$

$$\sqrt{N} = \sqrt{64} = 8$$

$$N = 64$$

$$N \cdot \log(N) = 64 \cdot 6 = 384$$

$$N\sqrt{N} = 64 \cdot 8 = 512$$

$$N^2 = 64^2 = 4096$$

$$2^N = 2^{64}$$

Special Names

$O(1)$	=	<u>constant</u>	$\boxed{N^t}$	$t \geq 2$	$(N^2, N^3, N^4, N^5, \dots)$
$O(N)$	=	<u>Linear</u>			
$O(N^2)$	=	<u>Quadratic</u>			
$O(N \log N)$	=	<u>Linear Logarithmic</u>			
$O(\log N)$	=	<u>Logarithmic</u>			
$O(2^N)$	=	<u>Exponential</u>			

Ques-2

$$F(N) = N \log(N) + 7N^2 + 8$$

$$= \underline{\underline{O(N^2)}}$$

Time & Space = Resources.

Mohit Sharma

Swathi Yadav

Task: Given list containing N integers, sort them in ascending order.

Inp: [1, 7, -1, 2, 4, 6] Same Input.

[100, 2, --, --, --, 1]
100 elements.

out: [-1, 1, 2, 4, 6, 7]

↓

① TukTuk Sorting Algo (TSA)

② Ascending Algo (AA)

Execution
Time:

15 sec.

i) Windows 95

10 sec.

Macbook Pro M2

Dependency

9 sec.

Macbook Pro M1

ii)

7 sec.

10 sec.

6) Programming Lang.

C++
↓

Python.

Python

C++ is faster than Python.

iii)

12 sec.

10 sec.

c) Hardware

Top of a volcano.

↓

Canada.

Processor
Temperature

(External)

iv)

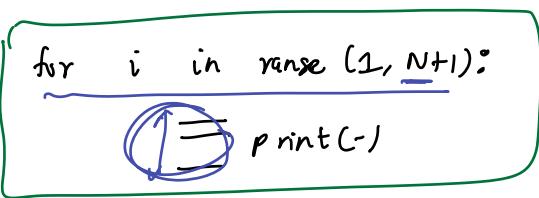
5 sec.

10 sec.

Observation:

\Rightarrow Execution Time depends on a lot of external factors.
Not a good measure of comparing time complexity (T_c)

(Q) What is a good comparison measure for T_c

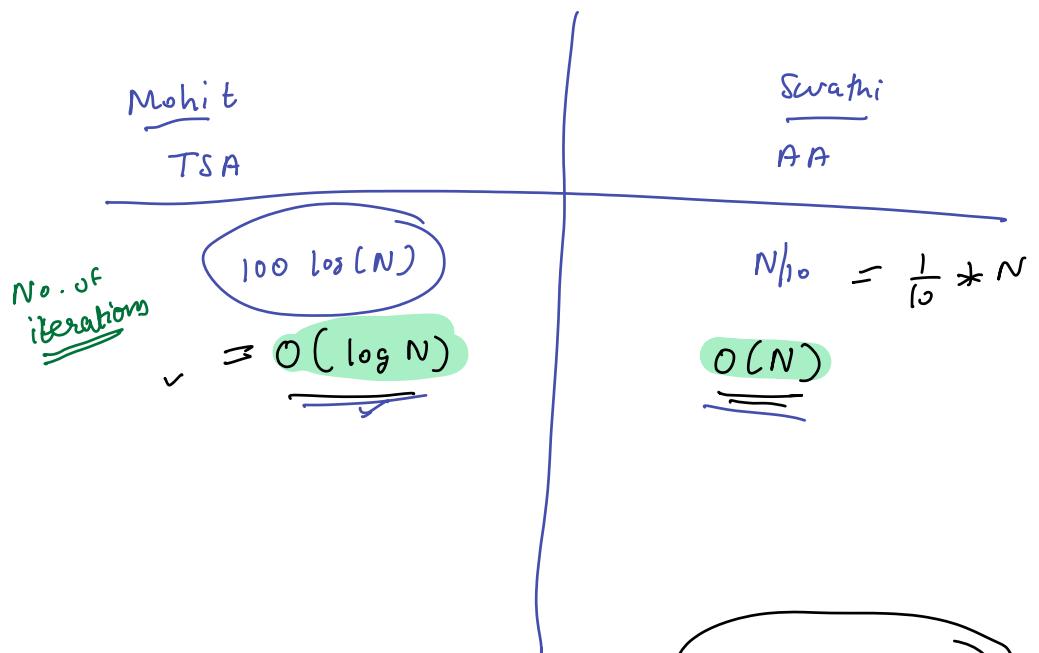
N iterations ← 
Iterations
is an independent factor.
as

$$\Rightarrow \underline{T(N)} \propto \underline{I(N)}$$

Time taken for input of size N is directly proportional
to the no. of iterations for input N .

$$I(N) = N$$

If $N \uparrow$ the no. of iterations increases
so does the time taken by the program.



\Rightarrow We will test their algorithm on diff values of N .

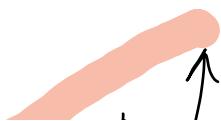
In general, \underline{N} values will be very Huge.

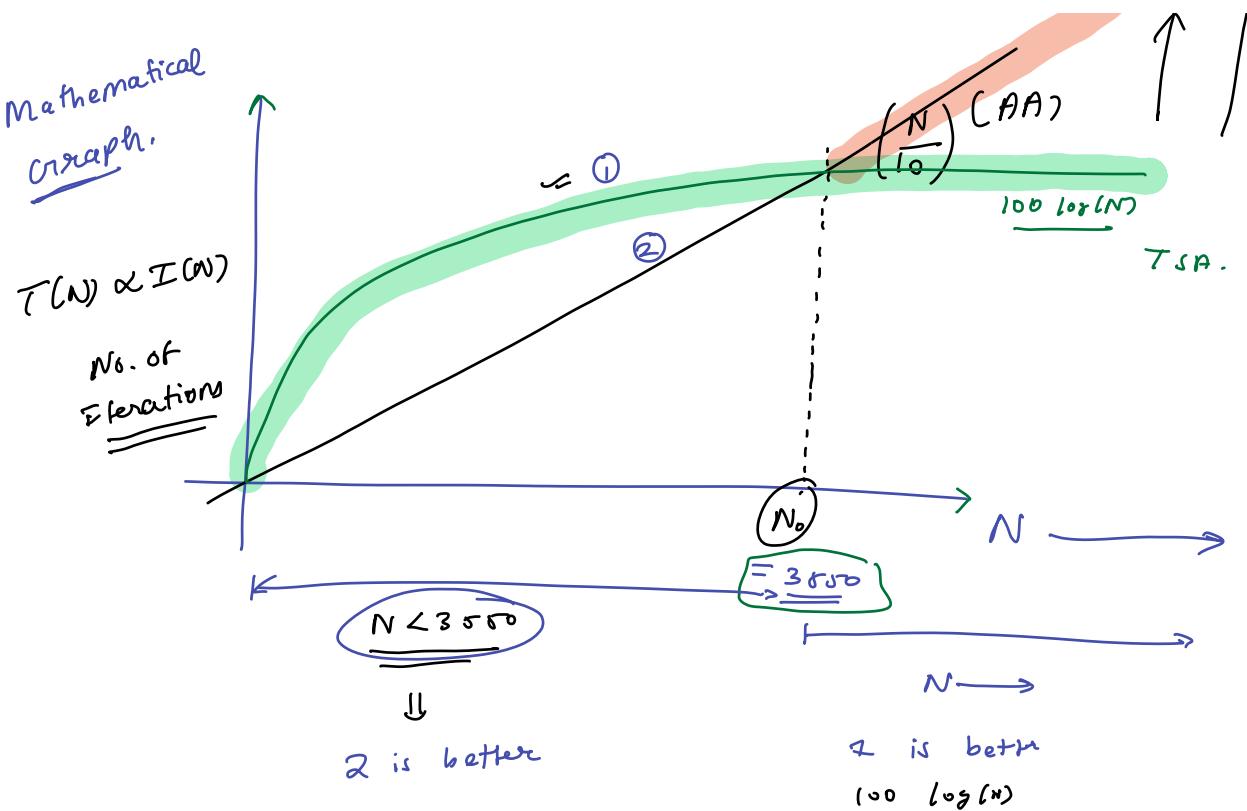
As $\underline{N \rightarrow \infty}$

\Rightarrow Asymptotic Analysis \Rightarrow Finding TC as N tends to infinity.

\Downarrow

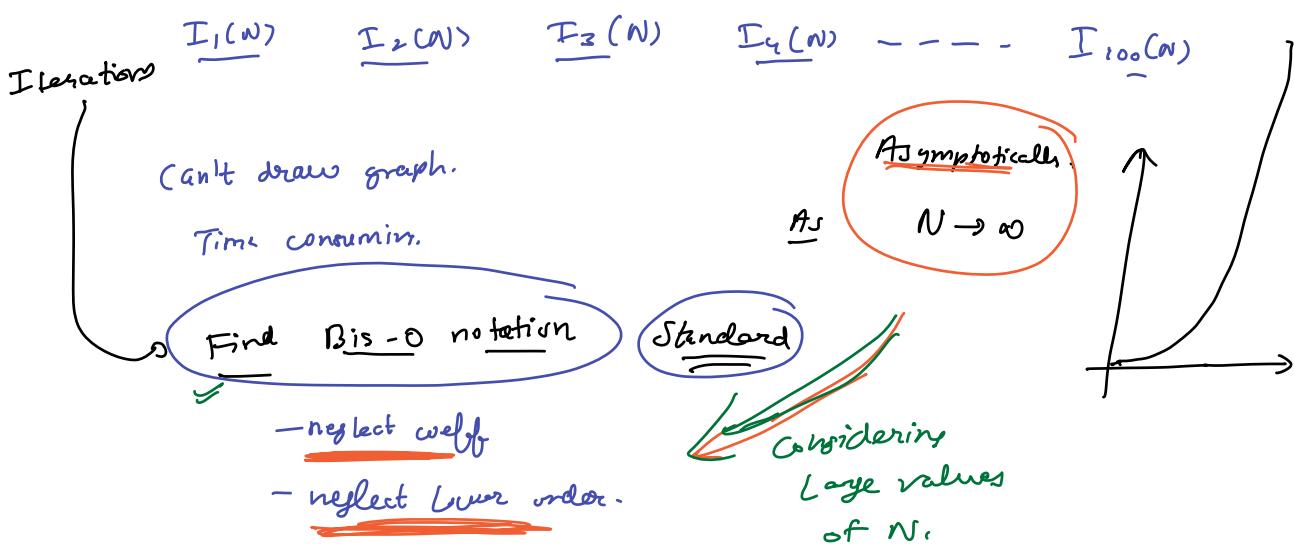
Bis-0





As per Asymptotic Analysis, we will prefer 1.
1 takes less time as $N \rightarrow \infty$.

(Q) Everyone in the class comes up with some sorting algo.



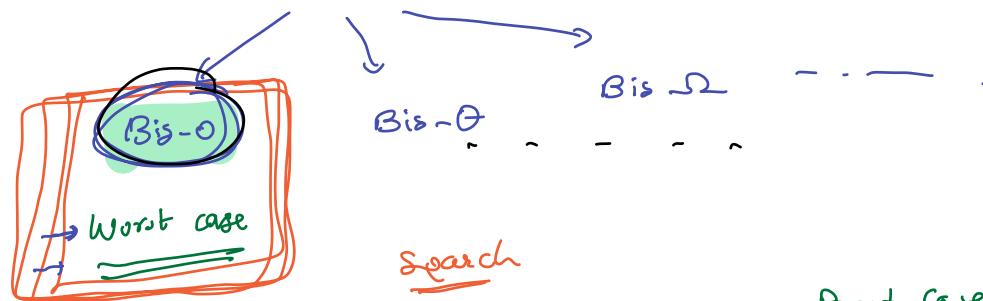
$$I_1(n) = \underline{n^2 + N \log n} + \underline{n \sqrt{n}} + \underline{n^3} + \underline{4n\sqrt{n}} = \underline{\underline{O(n^3)}} \quad \checkmark$$

$$I_2(n) = \underline{n^2 \sqrt{n}} + \underline{n^4} + \underline{n^2 \log n \sqrt{n}} + \underline{n^2} = \underline{\underline{O(n^4)}}$$

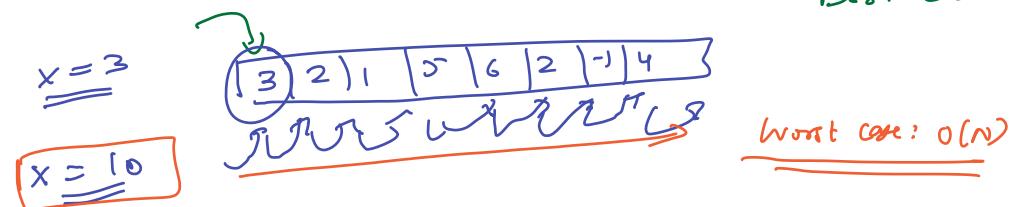
$$I_3(n) = \underline{n^2 + n} = \boxed{\underline{\underline{O(n^2)}}}$$

Performance of Algo as n becomes very large.

C Asymptotic Analysis



Best Case = $\Theta(1)$



Worst case: $O(n)$

Searching for x in list.

def search(l, x):

for ele in l:

if (ele == x):

return ele

Next class.

List size of N

[\leftarrow \rightarrow]

N elements.

return -1

No. of iterations = ? $\frac{1 \text{ or } 2 \text{ or } 3 \dots \text{ or } N}{N}$

$\downarrow \quad \rightarrow$

X

Best Case \uparrow

Worst Case \downarrow

Quiz

Mathematical Definition

$f(n)$ is $\underline{O}(g(n))$ if

there exist 2 constants c, n_0
such that

$$f(n) \leq c \cdot g(n) \quad \text{for all } n \geq n_0$$

$$\underline{f(n)} = \underline{N^2 + 10N + 8}$$

$$= \underline{\underline{O(N^2)}} \Rightarrow$$

$$= \underline{\underline{O(N^3)}} \quad \checkmark$$

$$= \underline{\underline{O(N^4)}}$$

$c \cdot g(n)$ is
Observation: an upper bound for $f(n)$

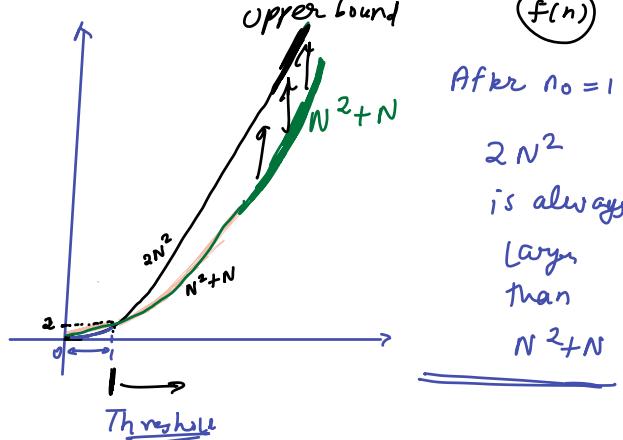
Es

$$\underline{f(n)} = \underline{N^2 + N} = \underline{\underline{O(N^2)}}$$

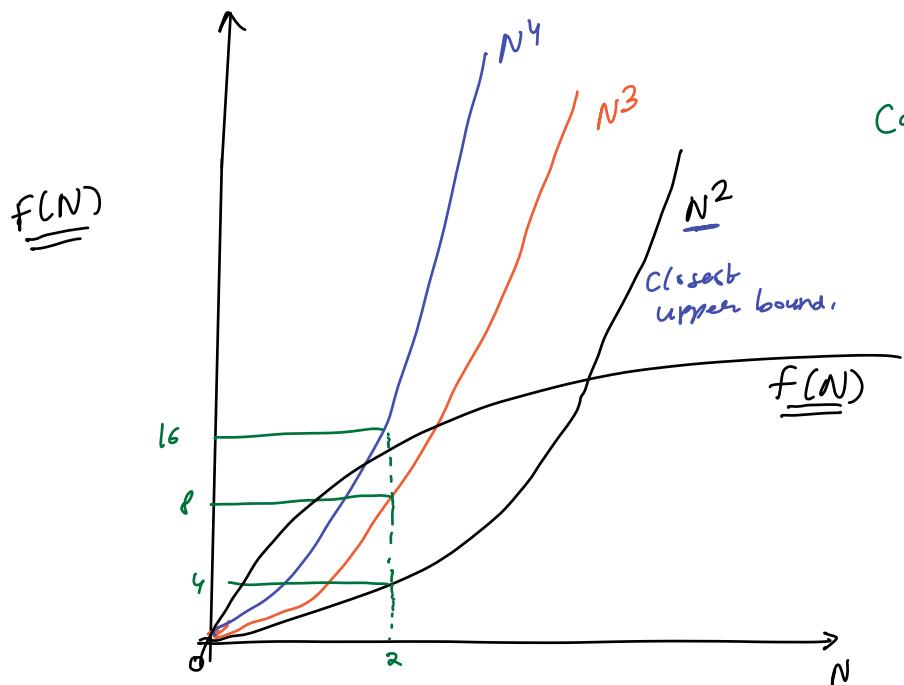
$$\underline{g(n)} = \underline{N^2}$$

$$c = 2$$

$$\underline{N^2 + N} \leq \underline{2N^2}$$



	$N^2 + N$	$2N^2$
$N=0$	0	0
$\left\{ \begin{array}{l} N=\frac{1}{2} \\ N=1 \end{array} \right.$	$\frac{1}{4} + \frac{1}{2} = \frac{3}{4} = 0.75$	$2 * \frac{1}{2} * \frac{1}{2} = \frac{1}{2} = 0.5$
$N=2$	2	2
$N=5$	$4+2=6$ $25+5=30$	$2*4=8$ $2*25=50$



Can I say $N^3 < N^4$
are both
Upper Bounds
for N^2 ?

$f(n)$ is $O(\underline{g(n)})$ if
there exist 2 constants c, n_0
such that

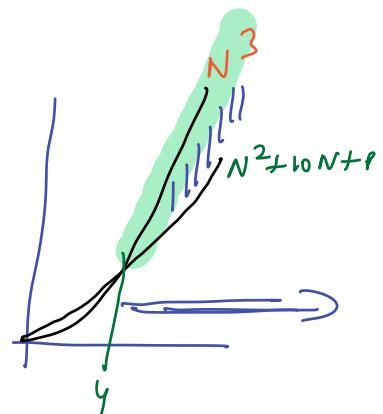
$$f(n) \leq c \cdot \underline{g(n)} \quad \text{for all } n \geq n_0$$

$$f(n) = \underline{n^2 + 10n + 8} \text{ is } O(n^3)$$

$$g(n) = N^3$$

$c = 1$
 $N_0 = 4$

$$\begin{aligned} \underline{n^2 + 10n + 8} &\leq N^3 \\ 16 + 40 + 8 &= 64 \end{aligned}$$



$$\Rightarrow f(n) = \underline{\underline{N^2 + 10n + 8}} \text{ is } O(N^2)$$

$$\underline{\underline{C = 100}}$$

$$\underline{\underline{C = 1000000}}$$

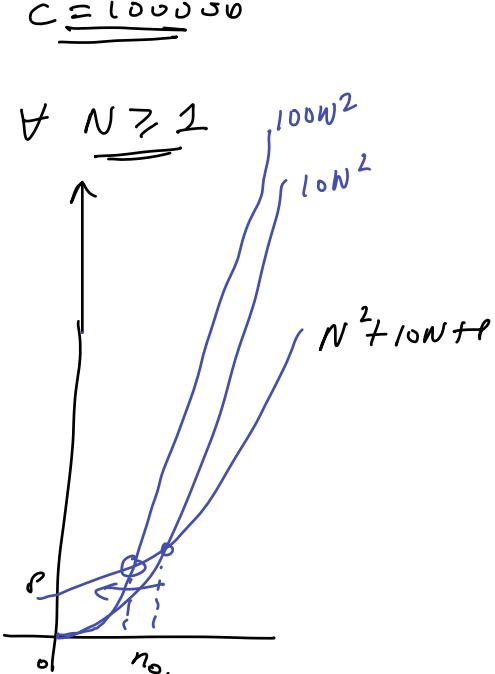
$$N^2 + 10n + 8 \leq 100 \underline{\underline{N^2}}$$

\downarrow

$$N=1 \quad \frac{1+10+8}{= 19} \quad 100$$

$$\underline{\underline{n_0 = 1}}$$

$$\boxed{f(n) \text{ is } O(N^2)}$$



✓ → Drop lower order terms

✓ → Drop constant coefficients

✓ $f(n) = \underline{\underline{N^2 + 4n + 8}} = \underline{\underline{O(N^2)}}$

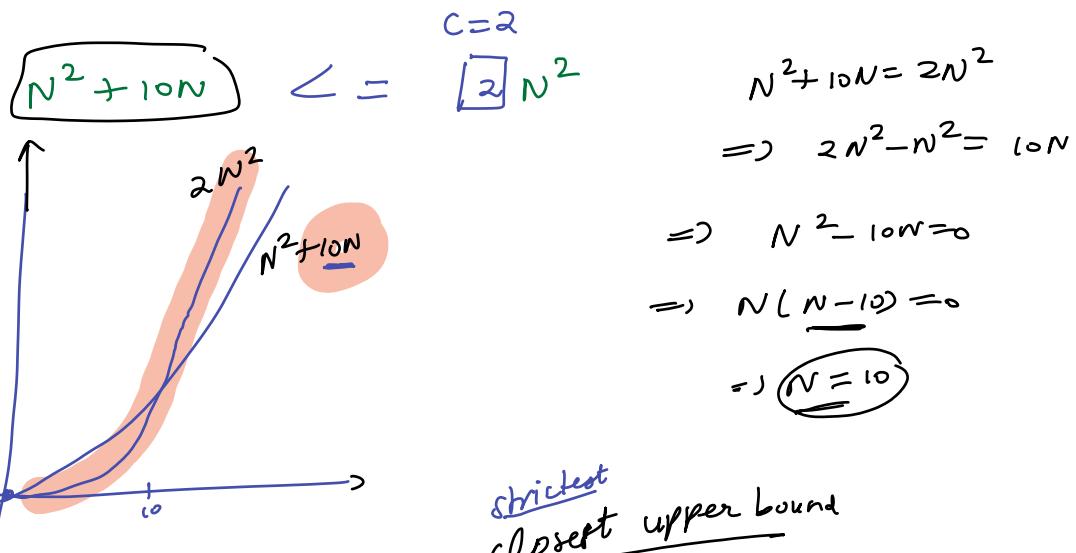
\Rightarrow Mathematical Definition

independent of constants & lower order terms.

$f(n)$ is $O(g(n))$ if

there exist 2 constants c, n_0
such that

$$f(n) \leq c \cdot g(n) \quad \text{for all } n \geq n_0$$



strictest closest upper bound

$$N^2 + 10N = O(N^2)$$

$$O(N^3)$$

$$O(N^4)$$

group
similar type.

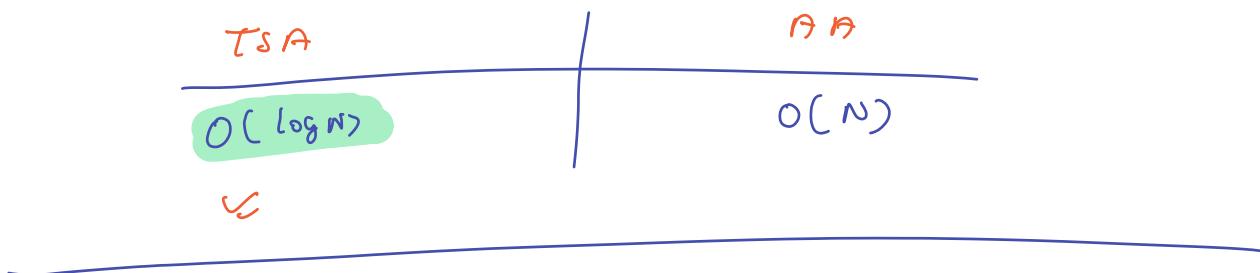
Comparing 2 algorithms

1) Find Big-O.

2) One which has highest order term is not efficient as compared with others.

Speed of Truck is less than

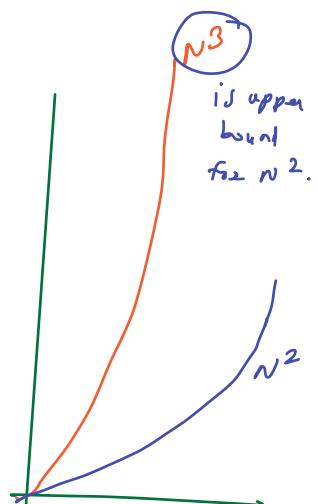
Speed of light.



Ques

✗ 1. $F(N) = \underline{n^3}$ is $O(\underline{n^2})$ ✗

✓ 2. $F(N) = n^3$ is $O(n^4)$



$$n^3 < \underline{C} \cdot \underline{n^2} \quad \forall \underline{n \geq n_0}$$

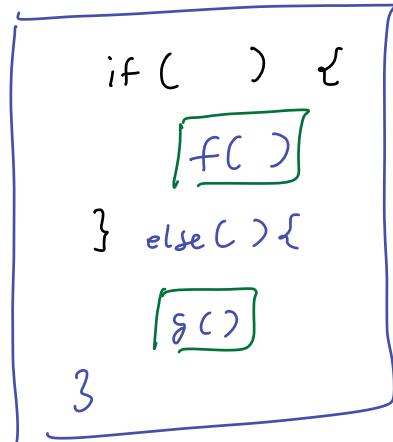
? ?

This will dominate after a certain threshold.

$$n^3 \leq n^4 \quad \forall n \geq 1$$

$$C=1, n_0=1.$$

Doubts



Worst Case

$$\begin{aligned}
 &\max(f, g) \\
 &\max(N, N^2) \\
 &= N^2
 \end{aligned}$$

Find Big-O for this.

$g(n)$

$$f(n) \leq c \cdot g(n)$$

$$\begin{aligned}
 &N^2 + 10N \\
 &= \underline{\underline{O(N^2)}}
 \end{aligned}$$

1. dropping lower order

2. dropping constant coefficients.

H.W - Q. 1



- 1) Try to contact TA.
- 2) Message on Slack.

Extra Reference
Session