

Advanced Programming 2 - Scalability

DR. ELIAHU KHALASTCHI

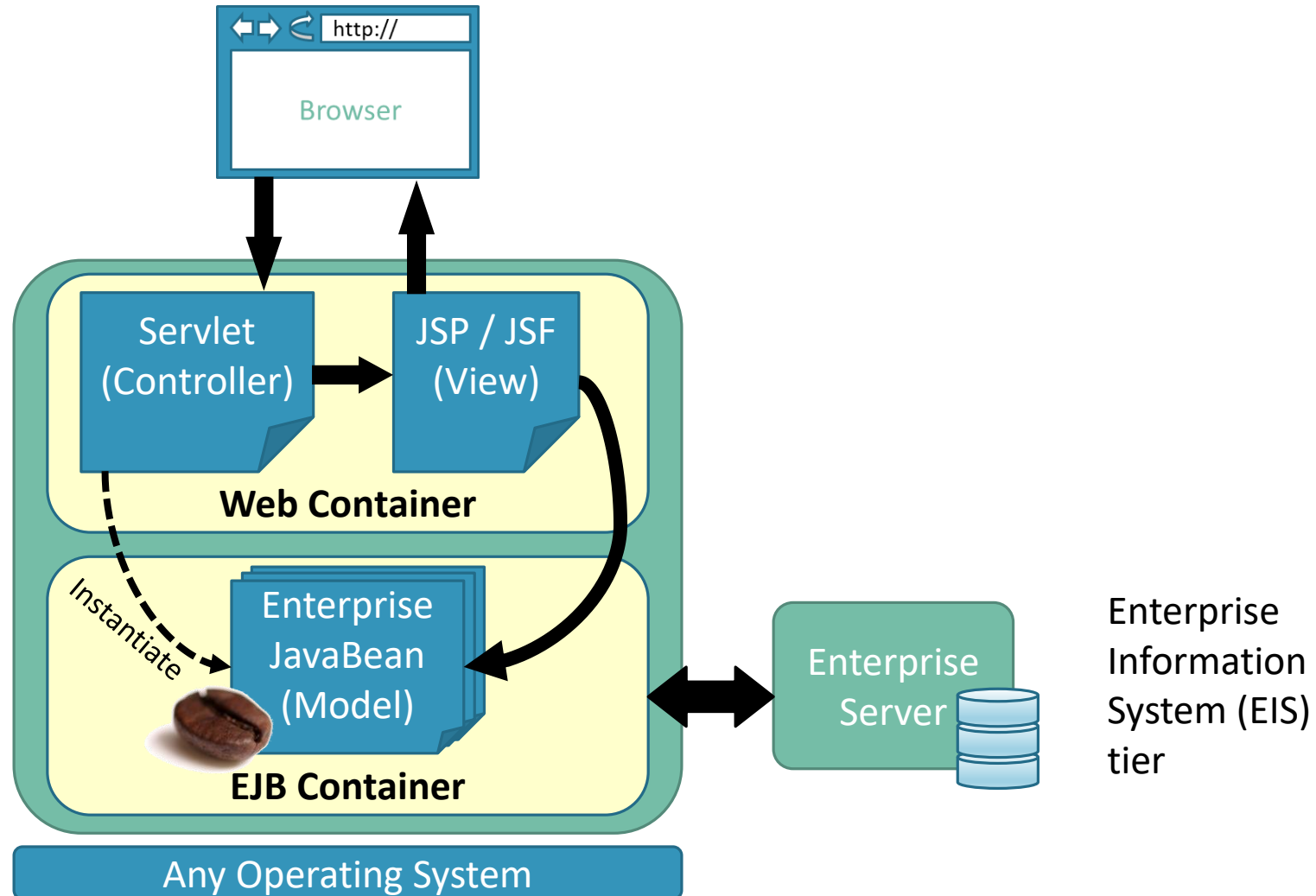
2016

A solid teal horizontal bar spanning the width of the slide at the bottom.

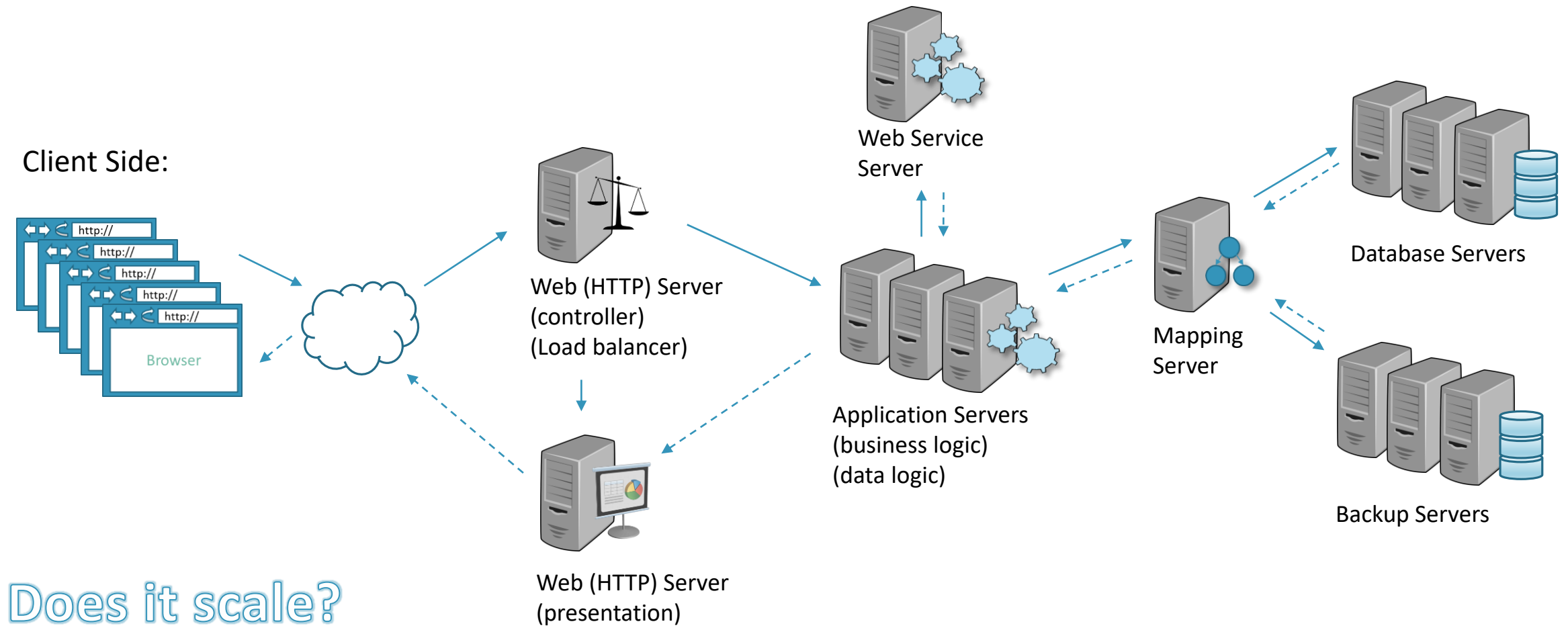
In the last lessons...

Client Side:
(client tier)

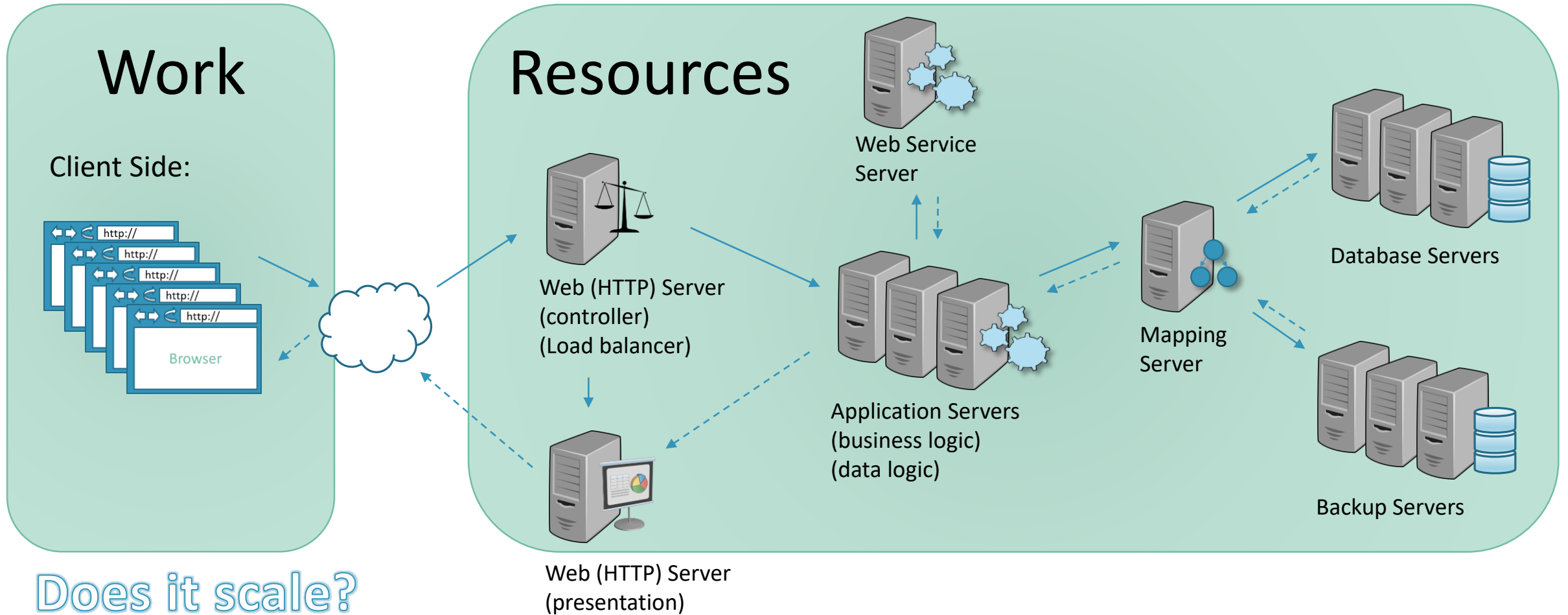
Java2 EE
Application Server
(middle tier)



In reality, it is more complex...



In reality, it is more complex...

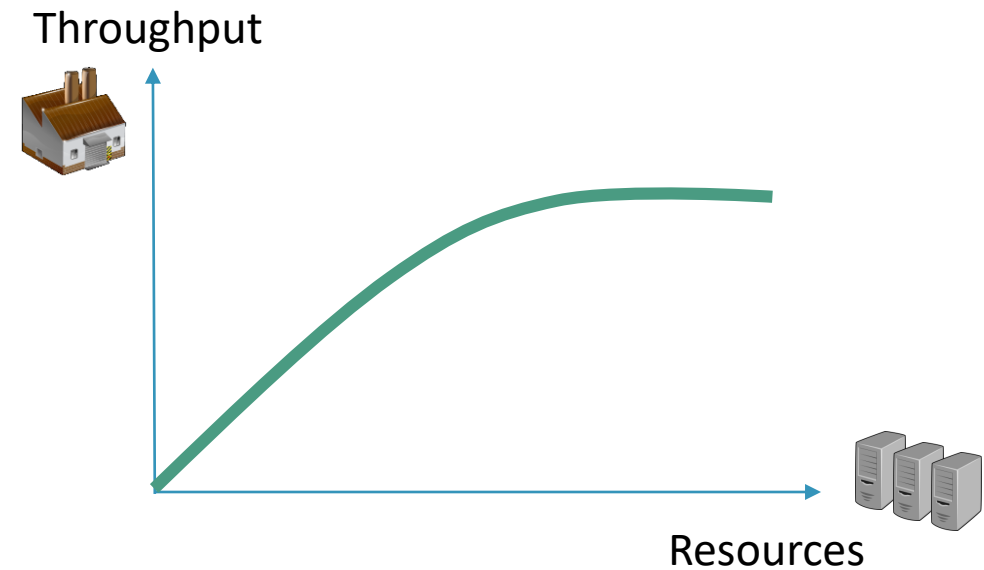
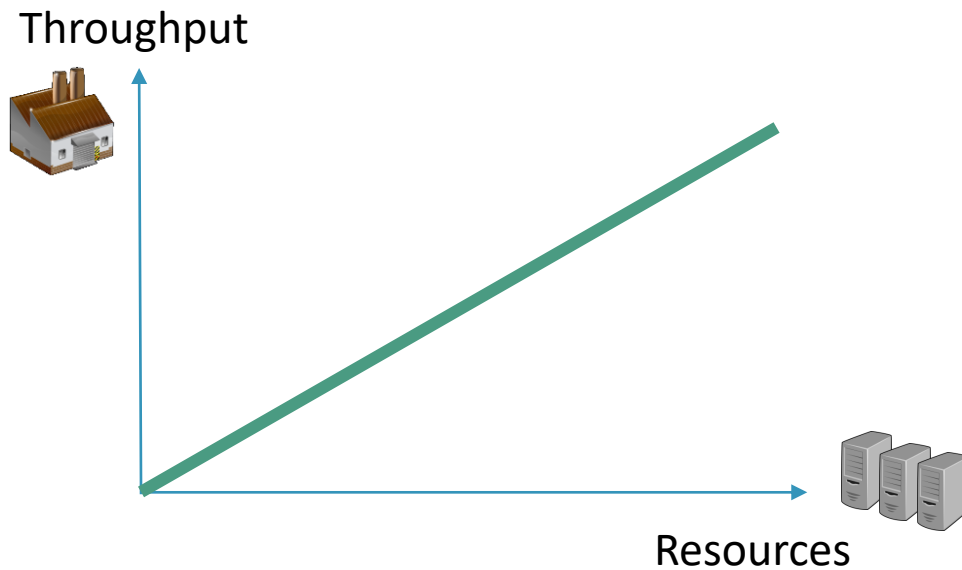


Definitions

- Scalability

- The capability of a system to handle a growing amount of work,
- or its potential to be enlarged in order to accommodate that growth

- Linear Scalability vs. None-Linear Scalability



Definitions (in the Web App context)

- **Scalability** – the required effort to increase the capacity to handle greater amounts of load
 - How much additional traffic can the system handle?
 - How easy it is to add more storage?
 - How many more transactions can be handled?
- **Performance** – optimized utilization of resources for fast response and low latency
- **Availability** – the uptime of the application
 - Redundancy for key components
 - Rapid recovery for partial system failure
 - Graceful degradation when a problem occurs
- Other parameters – reliability, manageability, cost, etc.

Let's start small... one server



A single machine cannot be tuned differently

Application server needs

- More CPU
- Optimized TCP/IP connections
- None blocking I/O, etc.

Database Server needs more RAM

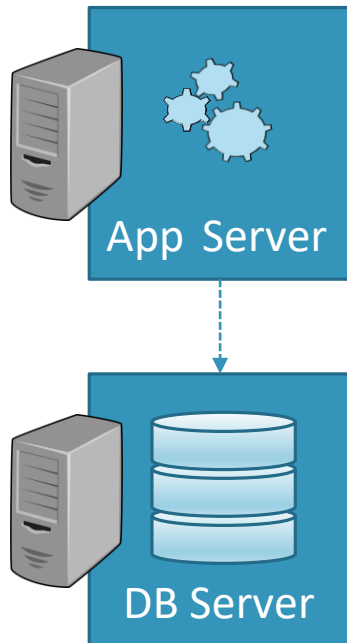
- And different optimizations

After some time... (more users)

Scale-up:

- Put more CPU
- Put more RAM

A Better Idea: split the nodes



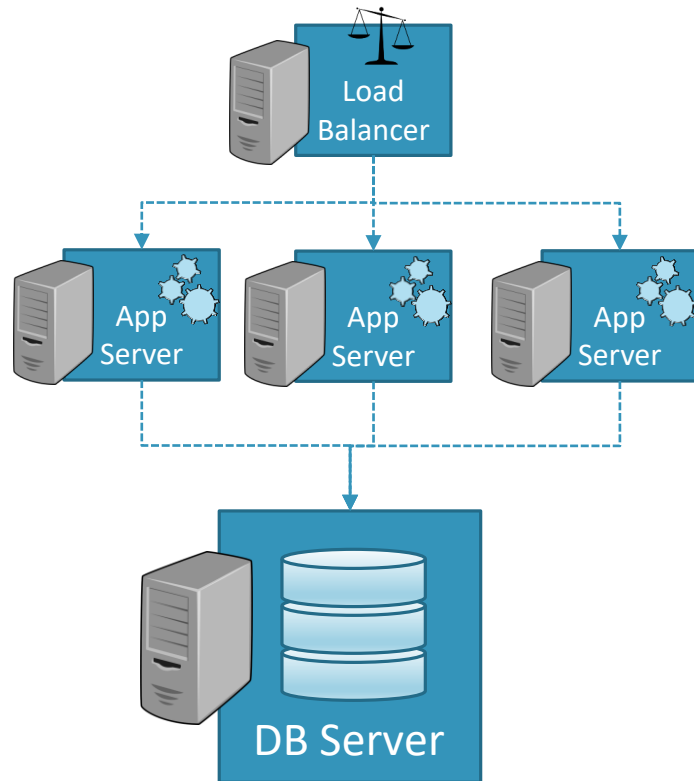
We can optimize each node differently

Downsides:

- The availability is not increased
- If one node fails the entire system fails
- Scalability is not increased
- More users →
 - more CPU for the App Server
 - more RAM for the DB Server
- Scale-up...
 - They are still bottlenecks...

We need to scale-out, use more nodes

A Better Idea: Scale-out



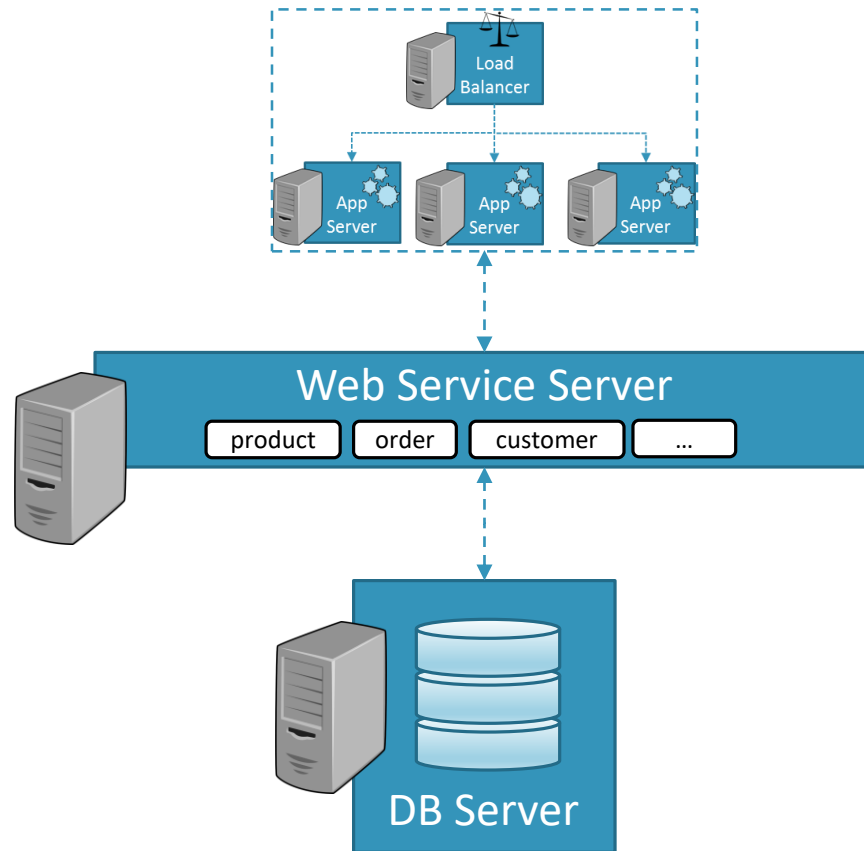
The load balancer takes requests
and forward them to the app servers
while balancing the load between them

This is scalable,
we can always add more app servers

The load balancer = a single point of failure

The DB Server = a single point of failure

Using Services



Services can be independent

Dedicated teams can work on them in parallel

Using more services doesn't increase availability

- Unless we use redundancy...

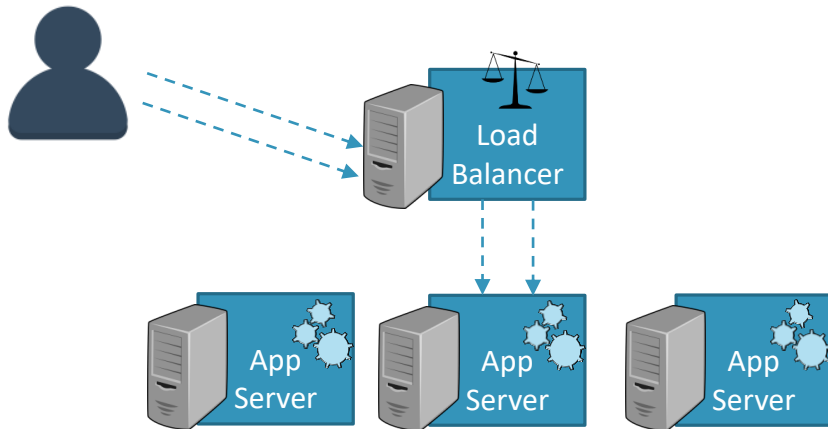
Performance issue

- The data needs to be cross-service-join

Handling a sticky session...

A sticky session

- A number of related stateful requests
- needs to be handled by the same machine
- Does not scale



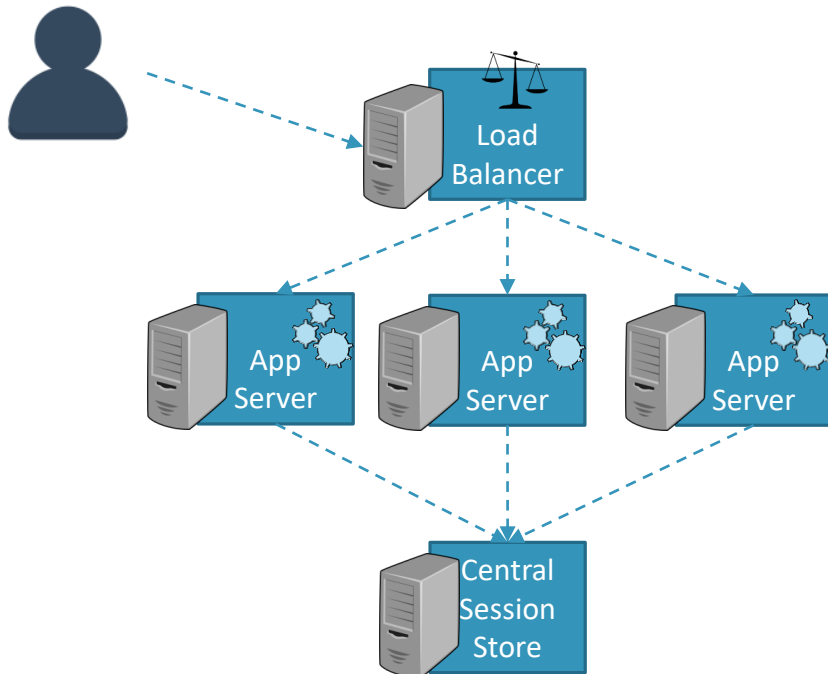
Handling a sticky session...

A sticky session

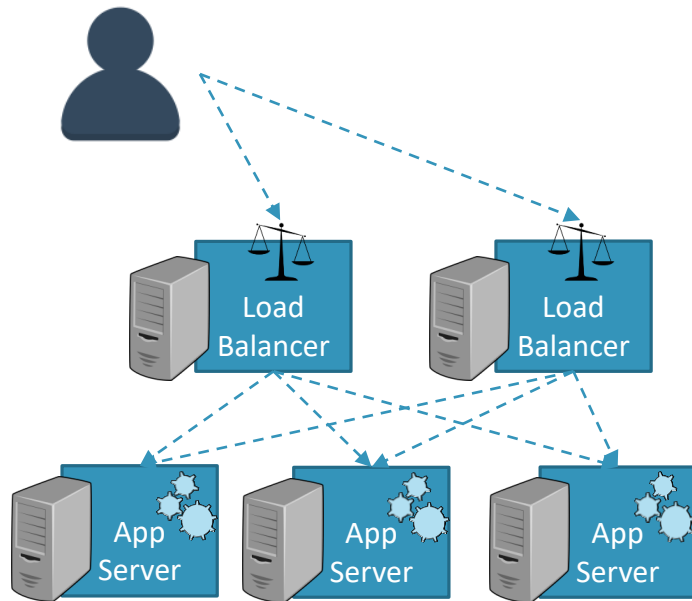
- A number of related stateful requests
- needs to be handled by the same machine
- Does not scale

We can use another server to store the session state

- A database, mem-cache
- Performance issue – another call is made



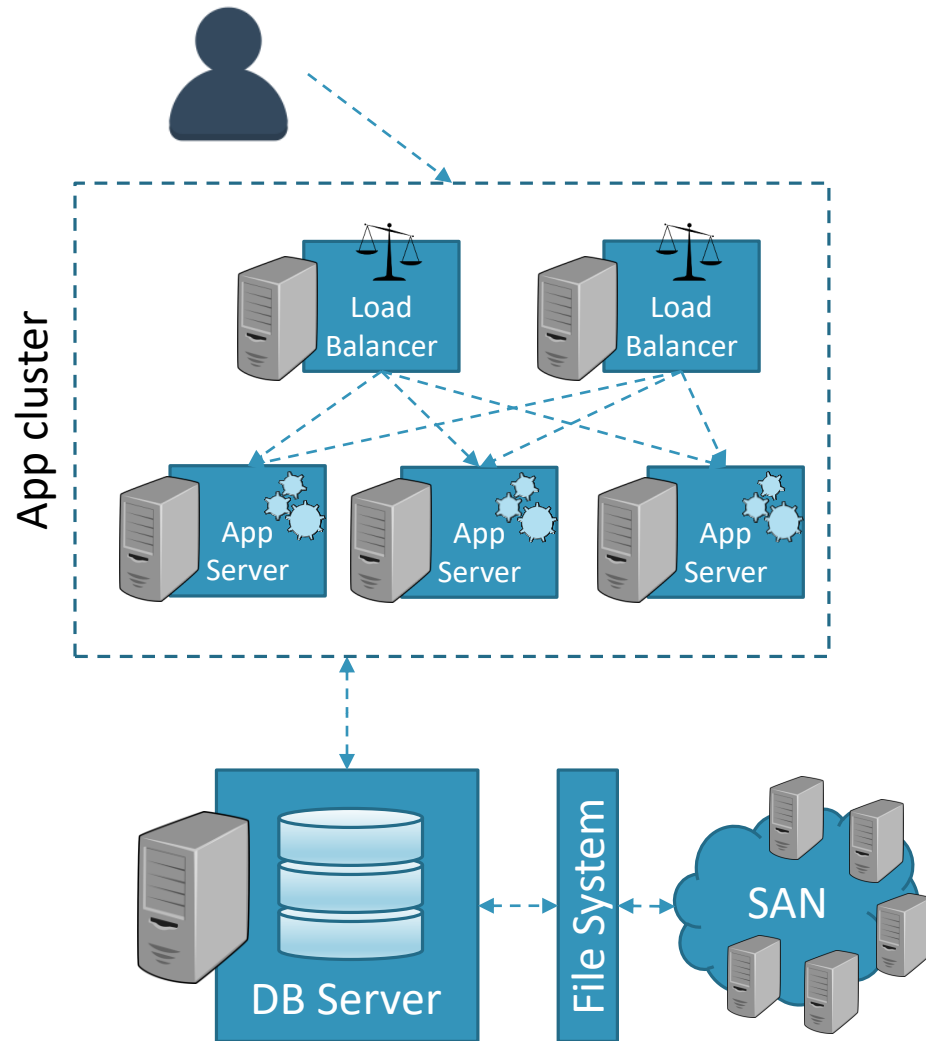
Load Balancer Redundancy



Can be used as a backup

Or be used together

What about the database?



The database can use a SAN

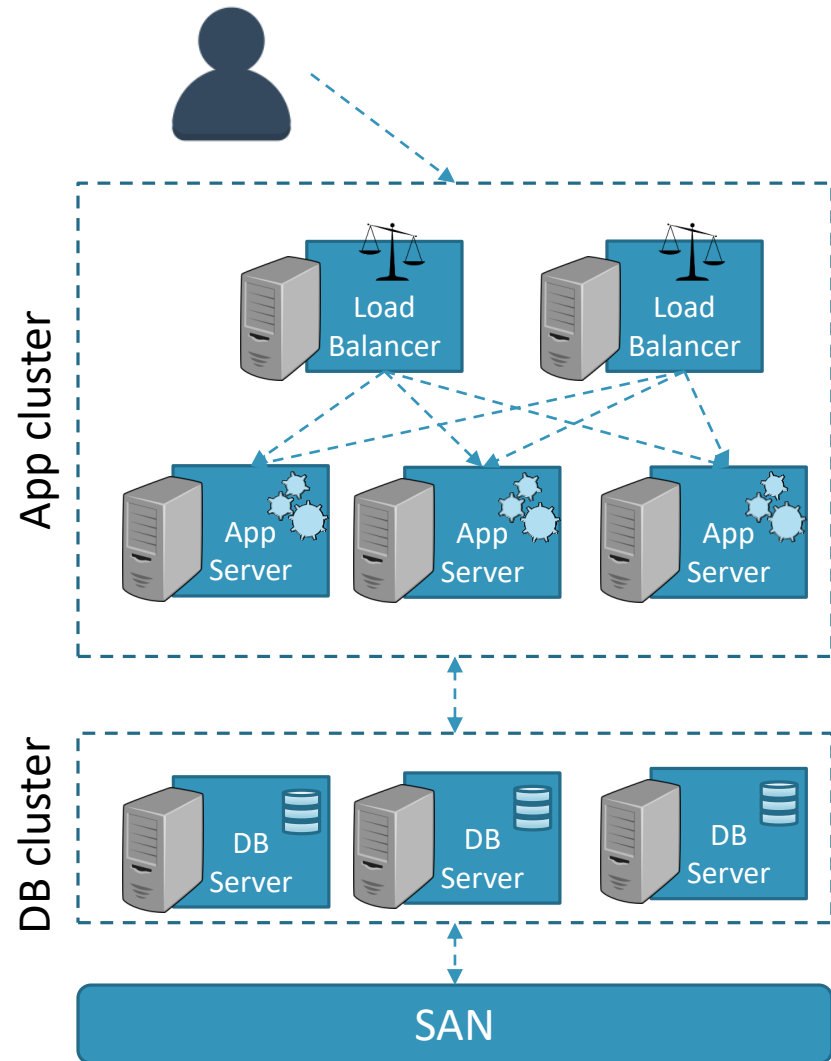
SAN – Storage Area Network

A SAN can be **scaled up**

- More RAM, discs in the network

But still the DB is a single point of failure

Scaling out the database



No single point of failure to the DB

Can retrieve data from either DB server

If we use replicas... (8, 16)

- Wait for the backup – consistent, performance issue
- Don't wait – nonconsistency, better performance

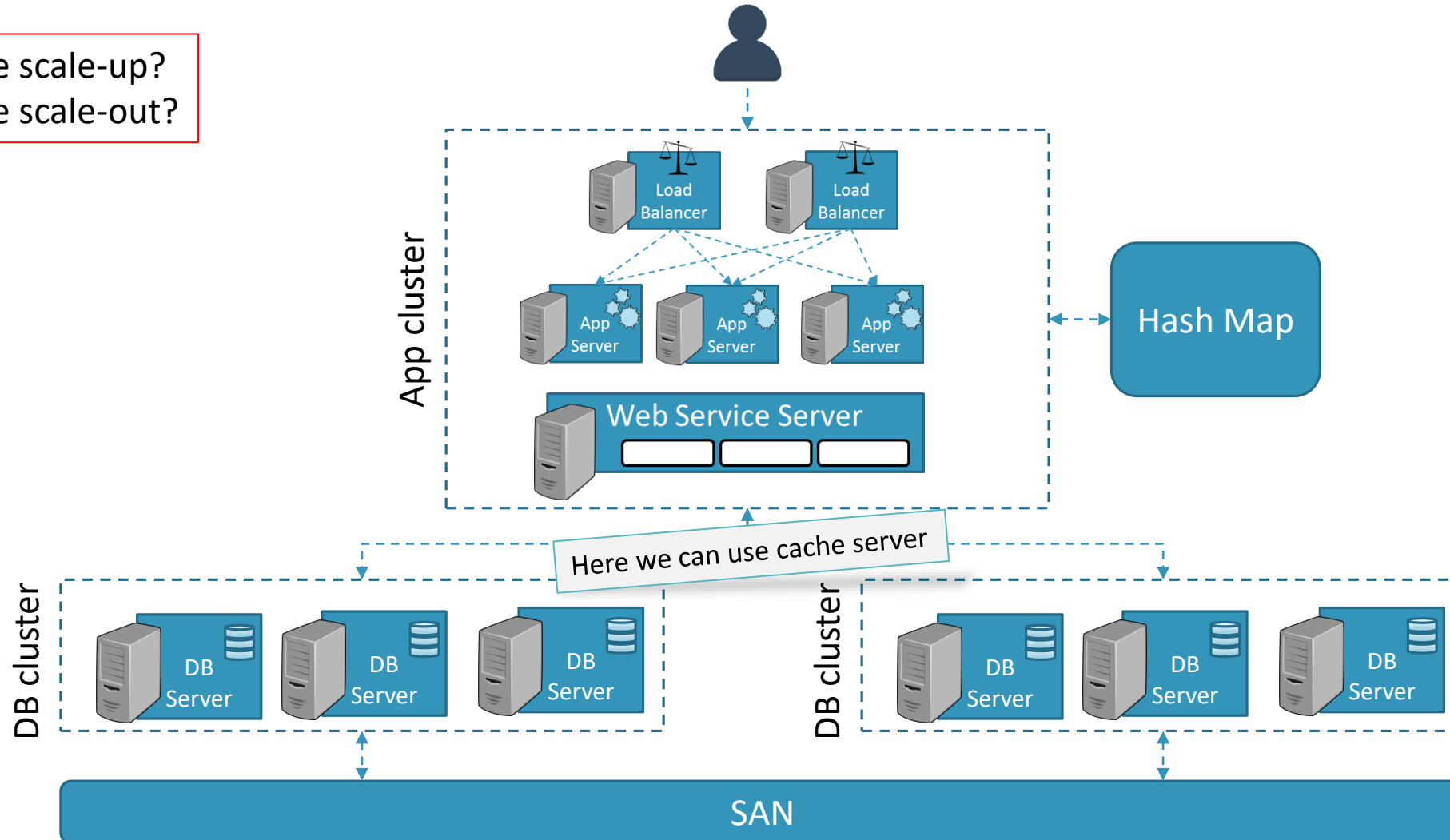
Scaling out a DB is difficult (horizontal partitioning)

- We lose the advantage of ACID transactions
- Atomicity, Consistency, Isolation, Durability

We can set different responsibilities for different DBs (vertical partitioning)

Overview

What can we scale-up?
What can we scale-out?

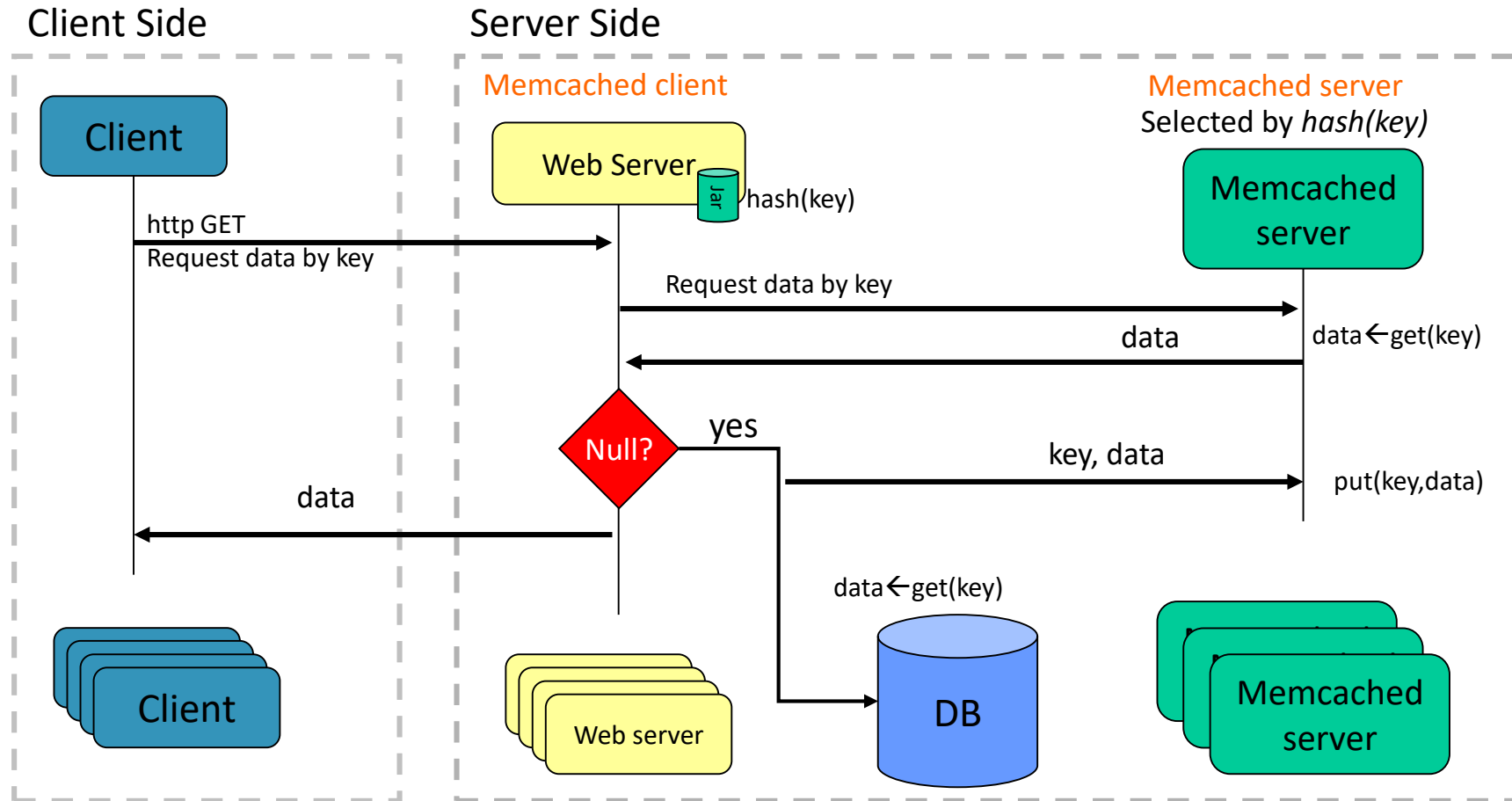


Memcached



A CASE STUDY

Memcached Overview



Pseudo code example

```
function get_foo(int userid) {  
    data = db_select("SELECT * FROM users WHERE userid = ?", userid);  
    return data;  
}
```

```
function get_foo(int userid) {  
    /* first try the cache */  
    data = memcached_fetch("userrow:" + userid);  
    if (!data) {  
        /* not found : request database */  
        data = db_select("SELECT * FROM users WHERE userid = ?", userid);  
        /* then store in cache until next get */  
        memcached_add("userrow:" + userid, data);  
    }  
    return data;  
}
```

```
function update_foo(int userid, string dbUpdateString) {  
    /* first update database */  
    result = db_execute(dbUpdateString);  
    if (result) {  
        /* database update successful : fetch data to be stored in cache */  
        data = db_select("SELECT * FROM users WHERE userid = ?", userid);  
        /* the previous line could also look like data = createDataFromDBString(dbUpdateString); */  
        /* then store in cache until next get */  
        memcached_set("userrow:" + userid, data);  
    }  
}
```

Memcached scalability

- The scalability is **linear**
- As the number of keys increases
- We can add a linear number of memcached server

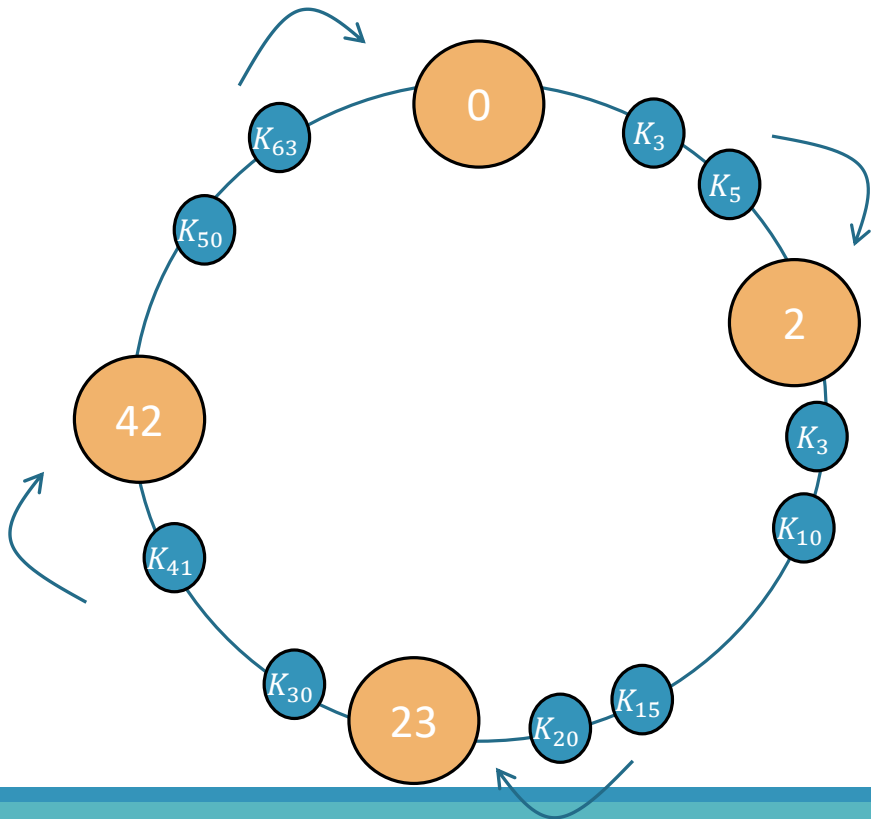
- Data is retrieved from the RAM
 - Faster than from the DB
 - Less load on the DB

If a Memcached Server crashes...

- One less server keys can be referred to...
- The hash map size is reduced from n to $n - 1$
- Most mapping algorithms require the **remapping of all keys**
 - E.g., put object o in machine number $hash(o) \bmod n$
- This is a catastrophe...
 - All the data from the DB needs to be reloaded to the memcached servers
- The same problem exists when adding a server... (scale-out)
- Unless we use **consistent hashing**

Consistent Hashing

- If a node is removed (or added) – only K/n keys need to be remapped
 - Where K is the number of keys, and n is the number of slots
- Main idea: **keys** and **nodes** are mapped to the same (ring-like) ID space



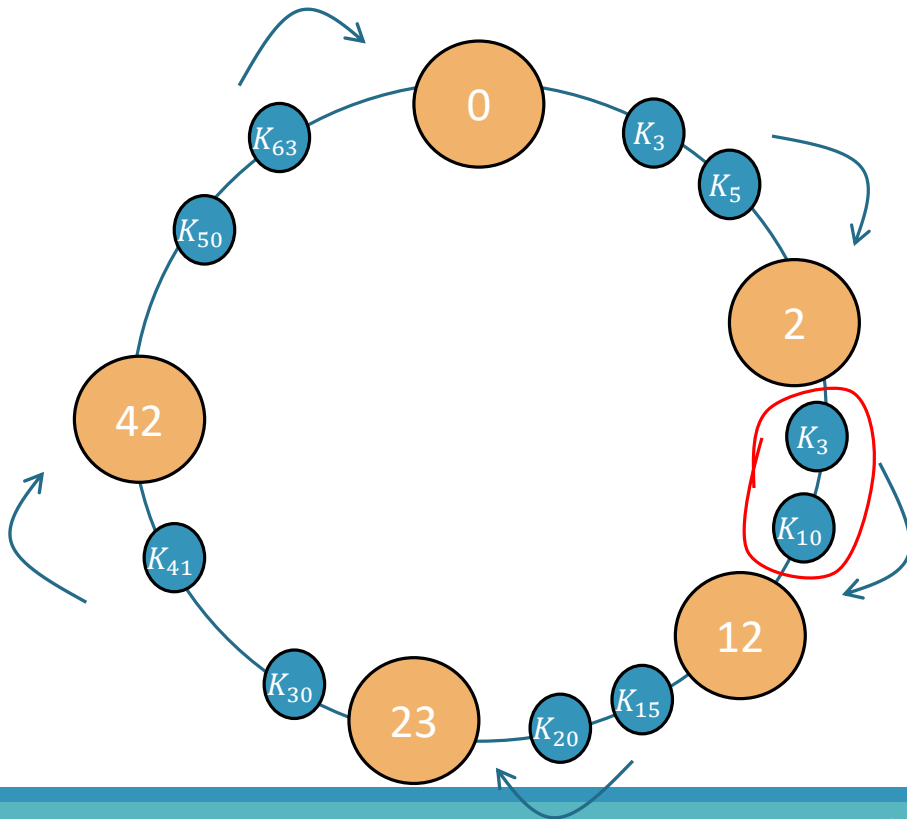
using $hash(IP)$ a node is mapped to a place on the ring – its ID

$hash(Key)$ also points to a place on the ring – the key ID

We need to map a key ID to a node ID
- We go clockwise and assign the key to the next node

Consistent Hashing

- If a node is removed (or added) – only K/n keys need to be remapped
 - Where K is the number of keys, and n is the number of slots
- Main idea: **keys** and **nodes** are mapped to the same (ring-like) ID space



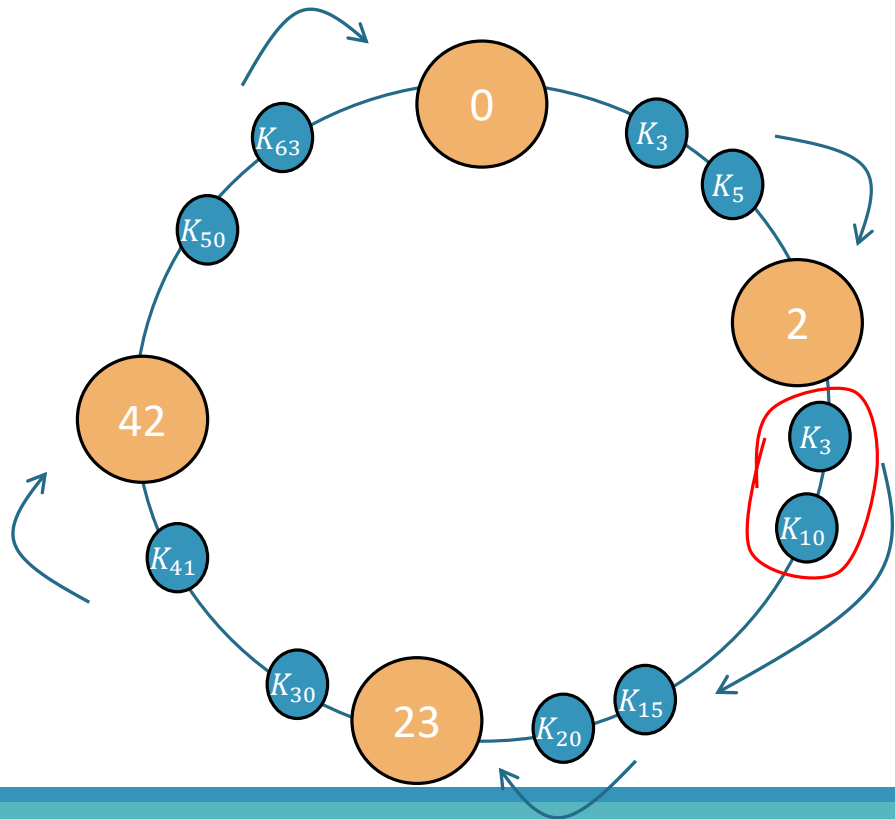
using $hash(IP)$ a node is mapped to a place on the ring – its ID

$hash(Key)$ also points to a place on the ring – the key ID

We need to map a key ID to a node ID
- We go clockwise and assign the key to the next node

Consistent Hashing

- If a node is removed (or added) – only K/n keys need to be remapped
 - Where K is the number of keys, and n is the number of slots
- Main idea: **keys** and **nodes** are mapped to the same (ring-like) ID space



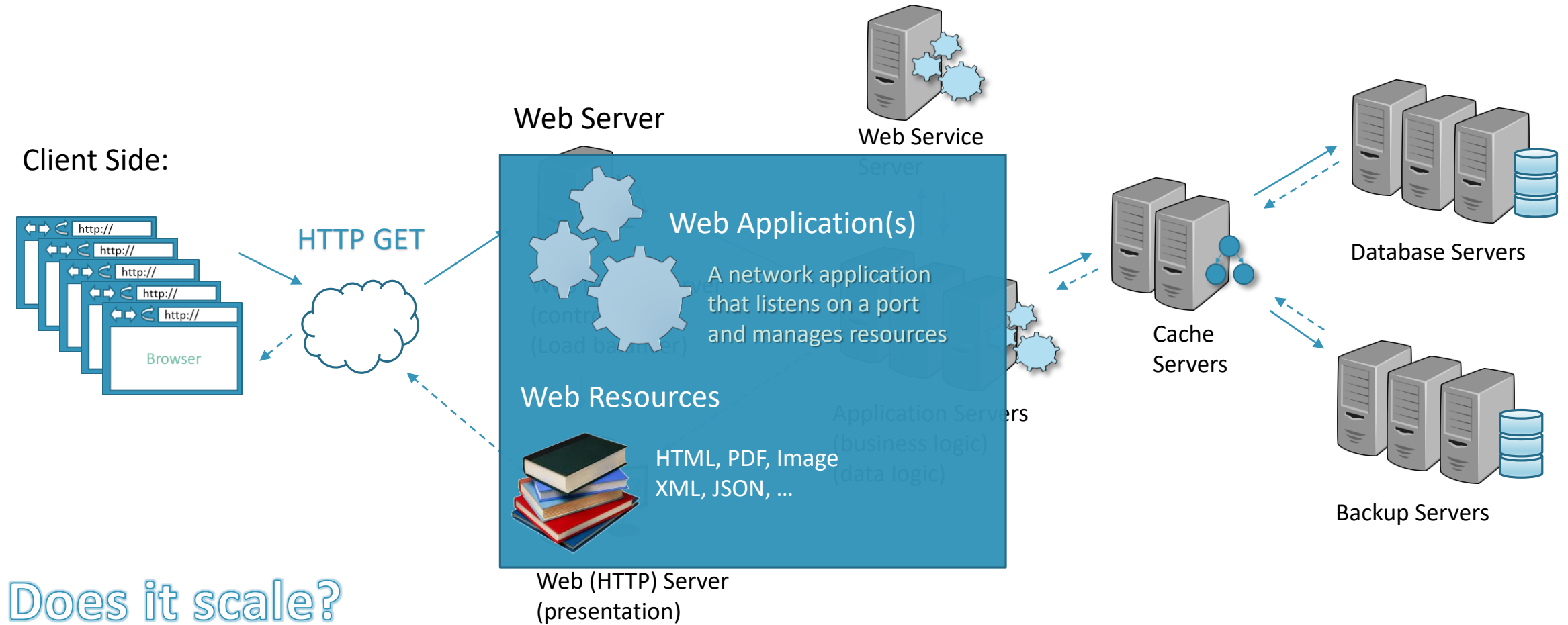
using $hash(IP)$ a node is mapped to a place on the ring – its ID

$hash(Key)$ also points to a place on the ring – the key ID

We need to map a key ID to a node ID
- We go clockwise and assign the key to the next node

Summary

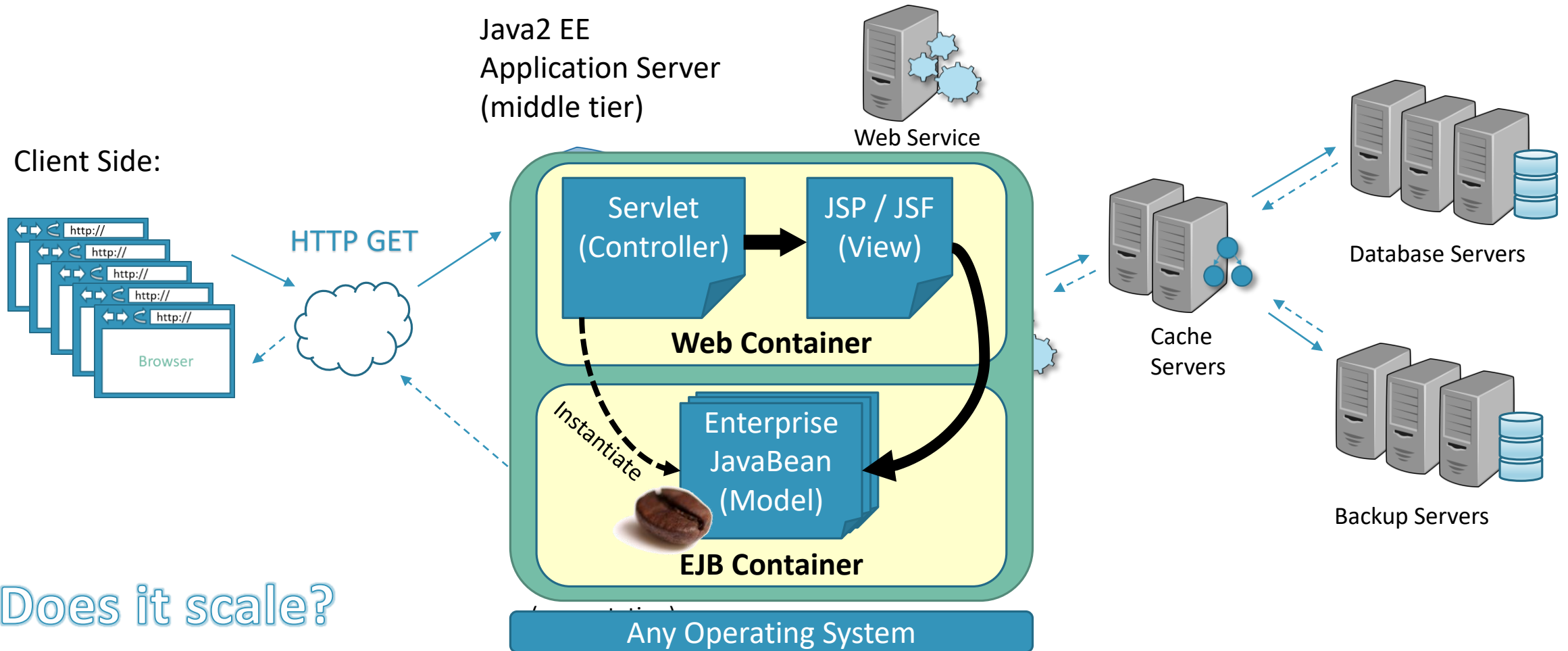
Our Enterprise



Does it scale?

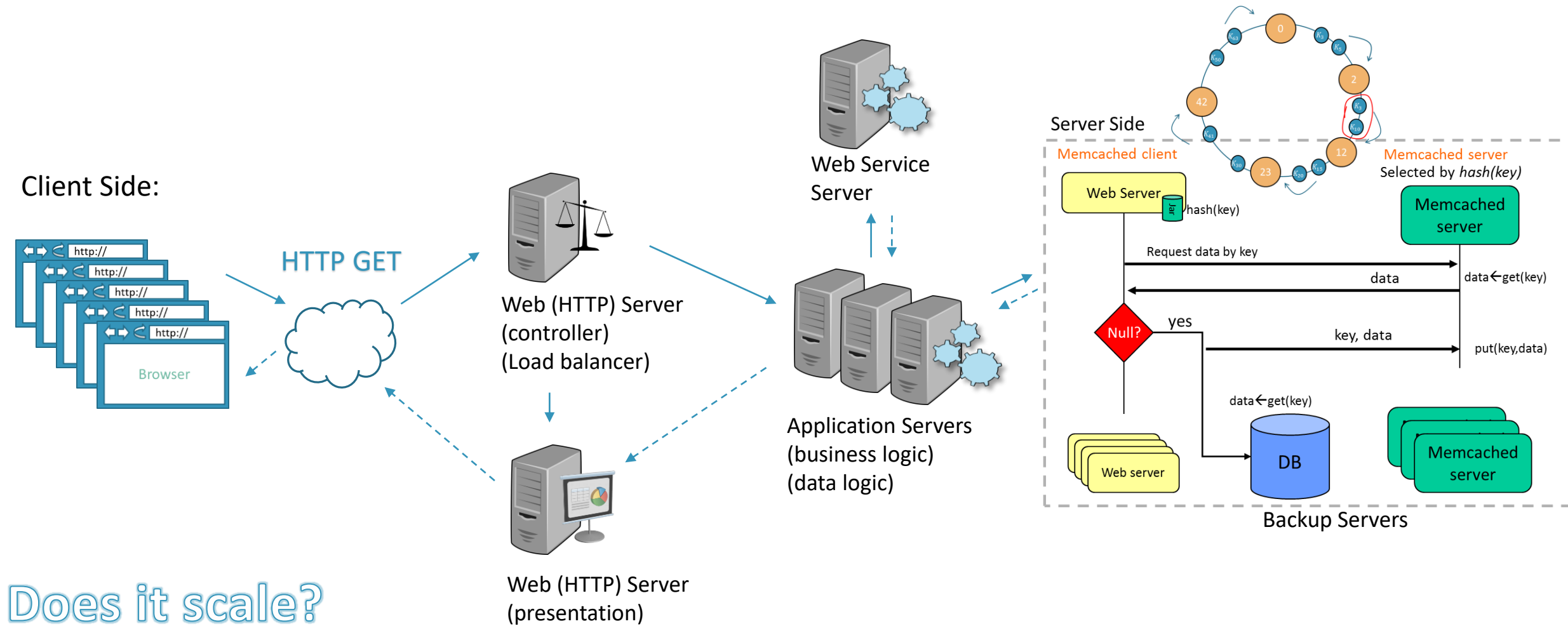
Our Enterprise

Java2 EE
Application Server
(middle tier)



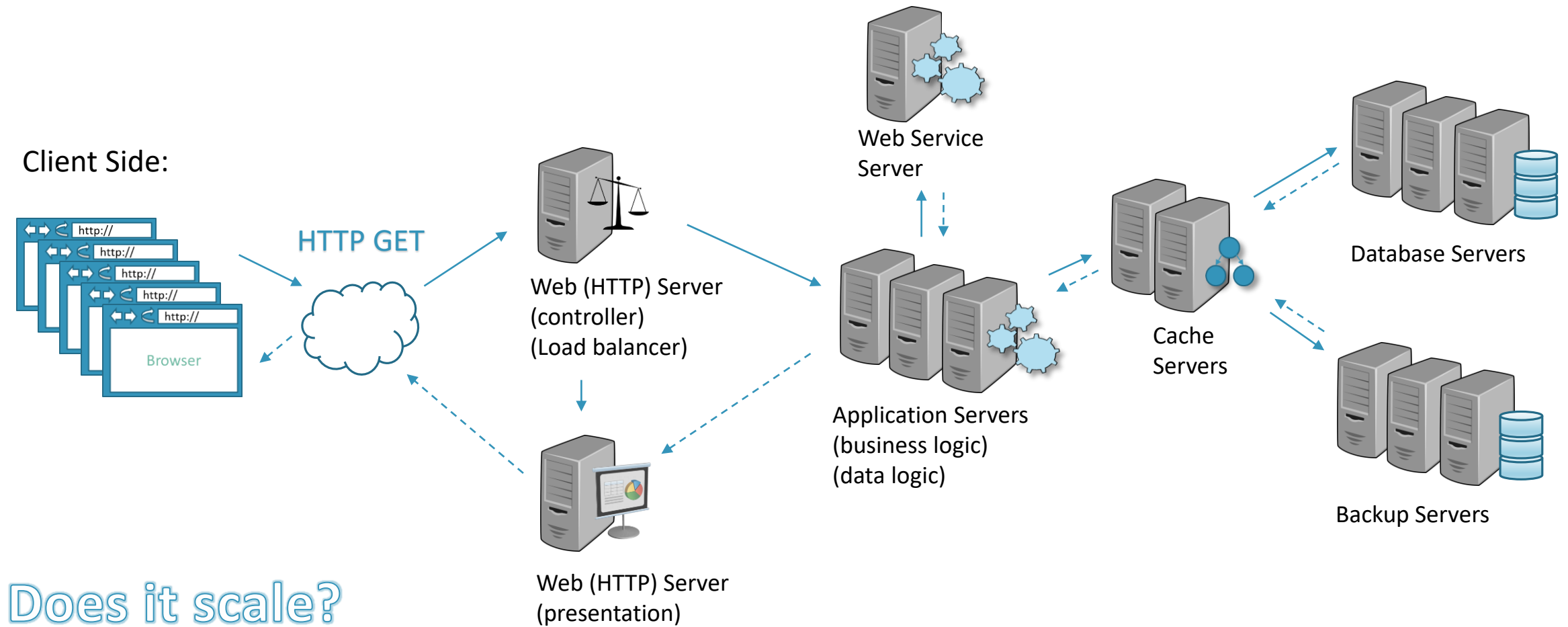
Does it scale?

Our Enterprise

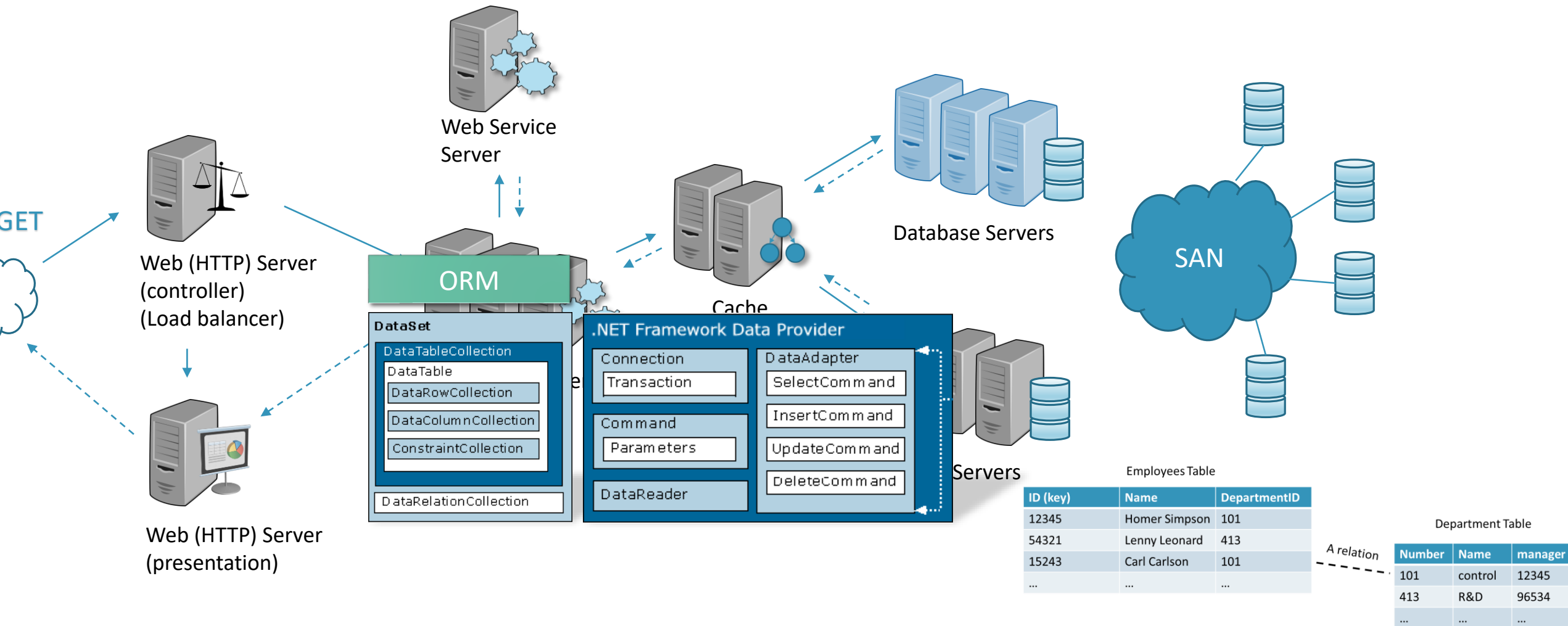


Does it scale?

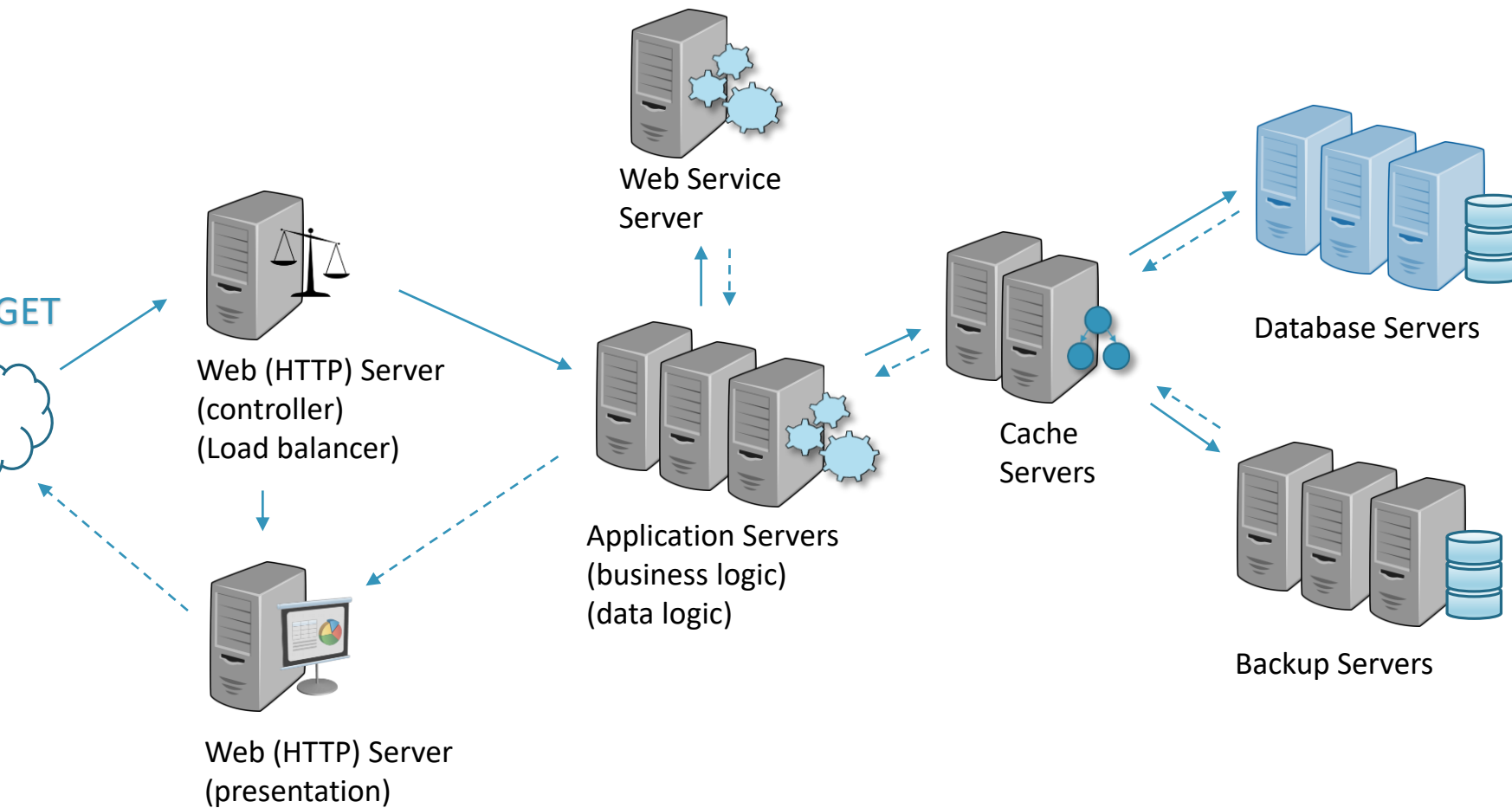
Our Enterprise



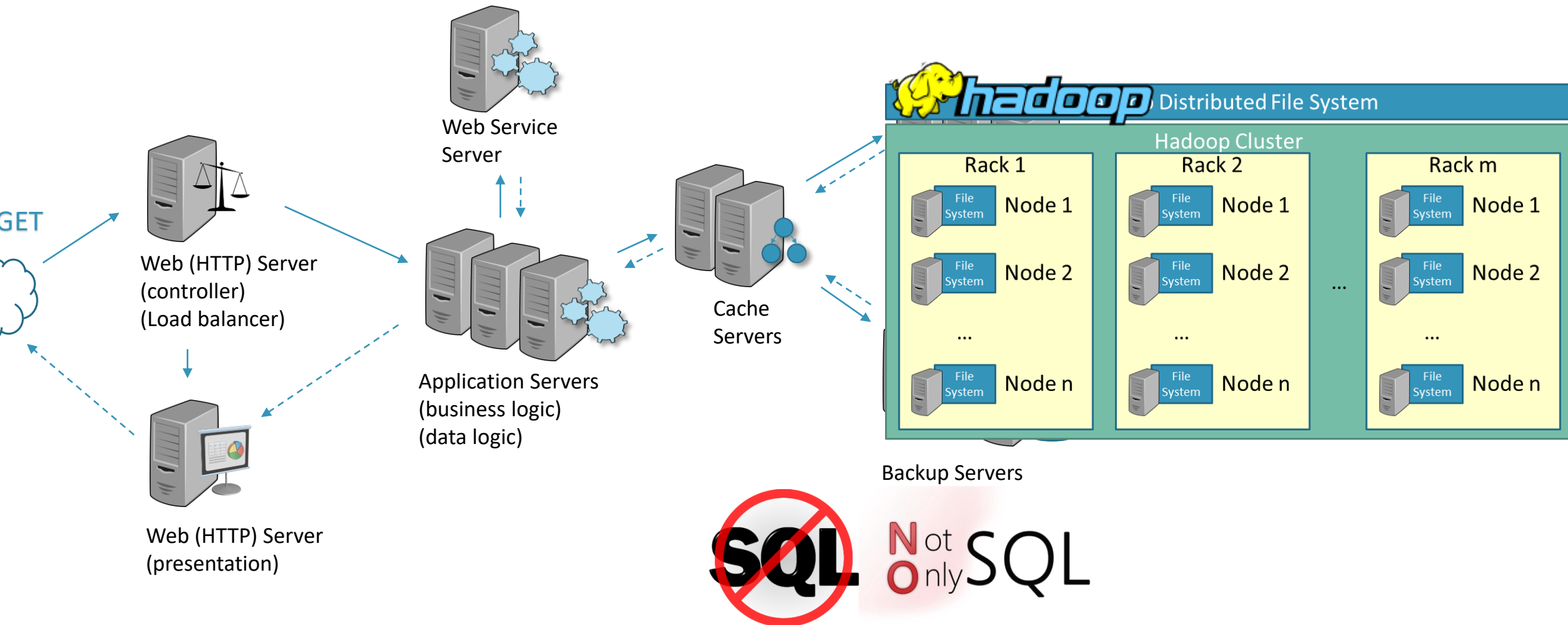
Our Enterprise



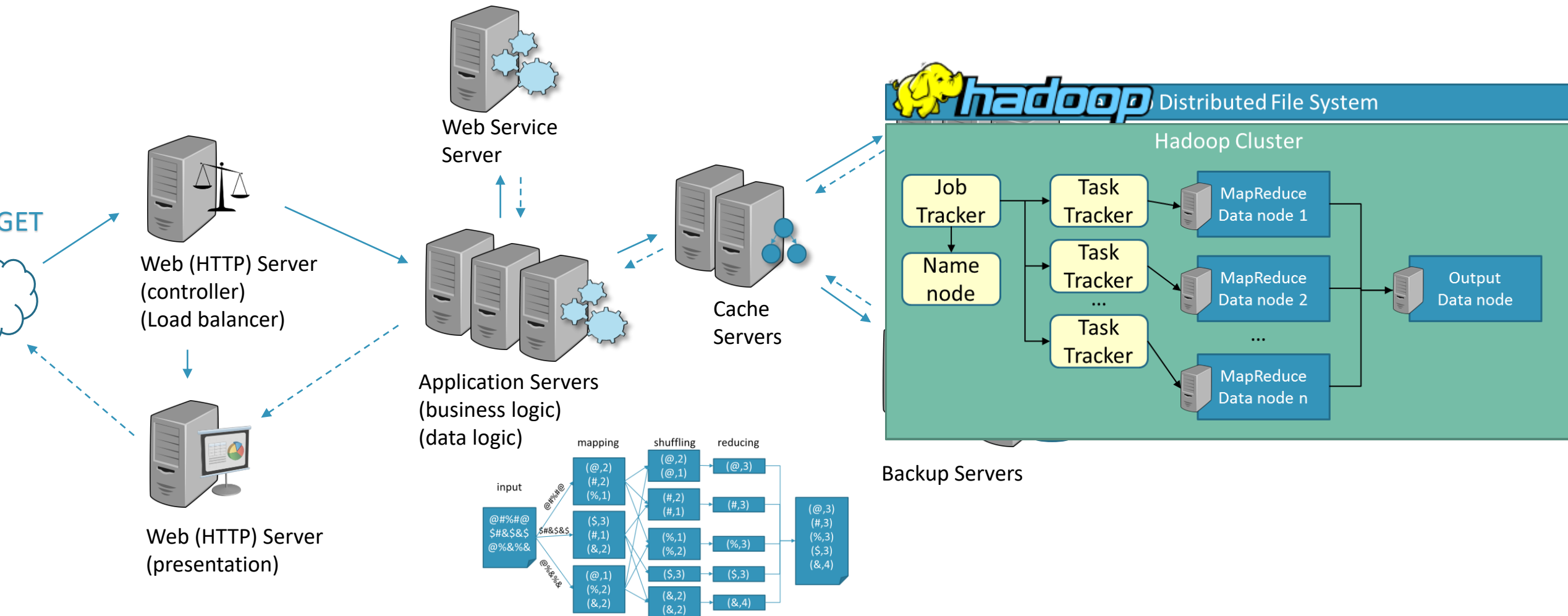
Our Enterprise



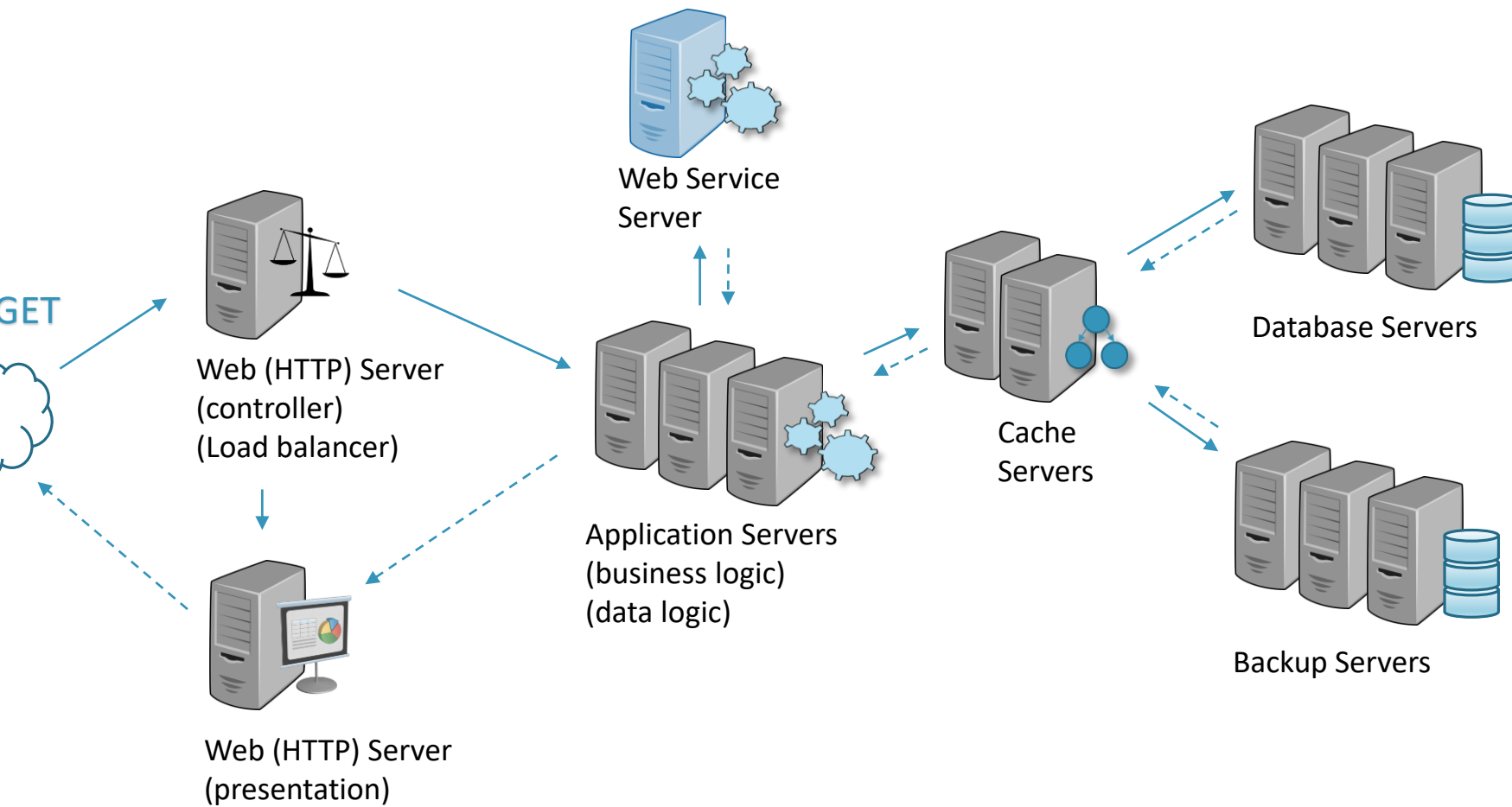
Our Enterprise



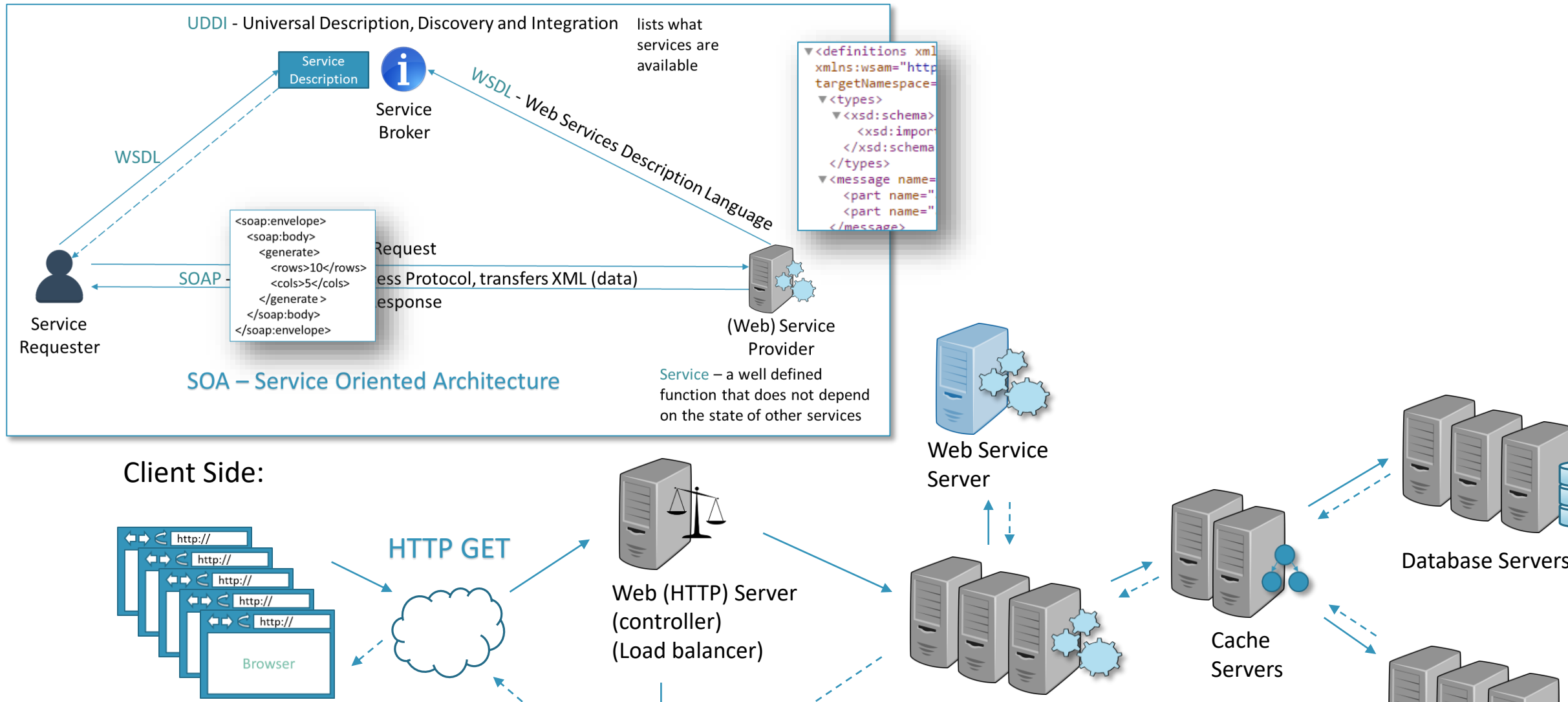
Our Enterprise



Our Enterprise



Our Enerprise



Our Enterprise

SOAP

```
<?xml version="1.0"?>
<soap:Envelope
xmlns:soap="http://www.w3.org/2001/12/soap-envelope"
soap:encodingStyle="http://www.w3.org/2001/12/soap-encoding">
  <soap:body pb="http://www.acme.com/phonebook">
    <pb:GetUserDetails>
      <pb:UserID>12345</pb:UserID>
    </pb:GetUserDetails>
  </soap:Body>
</soap:Envelope>
```

REST

```
http://www.acme.com/phonebook/UserDetails/12345
```

Client Side:

