Amortized Time Complexity - Theory(Introduction)

It applies not to a single run of algorithm but rather to a sequence of operations performed on the same data structure. As we see in the stack previously. As the stack gets 2 it gets doubled.

We can term such type of Stack as Augmented Stack where Augmented means greater in size.

Such doubling size, we see in dynamic hash table data structures, where each time of overflow the size increases.

According Cormen, Amortized Analysis can be defined as follows: "An amortized analysis is any strategy for analysing a sequence of operations to show that the average cost per operation is small, even though a single operation within the sequence might be expensive."

Real World Example

Let us assume that X gets a salary of 10,000. He saves 2000 per month, which goes on for 10 months.

In Worst Case Scenario: if the expenditure amounts to $\fill \fill \fil$

In short, like banks keep the account details of a customer, one is required to maintain the account of the cost of the sequence of operations.

OBSERVATION

- 1. $Amortized\ Cost > Actual\ Cost$.
- 2. It gives us the tight bound and tight bound is associated with Average Case Time Complexity.

Classification of Calculation on Amortized Analysis


