## 20.8. TIME COMPLEXITY CALCULATION NESTED FOR LOOP (EG-7).

```
//outer loop executed n times for(i = 1; i \le n; i + +) \{ //inner loop executes n times for(j = 1; j \le n/2; j + +) \{ k = k + 1; // constant time. \}
```

## **SOLUTION:**

- 1. Inner most loop's statement  $\Rightarrow k = k + 1$  which runs at O(1) time.
- 2. No. of inputs in outer for loop takes 1 to n times. lets see the inner loop and runtime of inner loop's statement.

$$f(1) \le c \times n \Longrightarrow when i = 1$$

$$f(1) \le c \times n \Rightarrow when j = 1$$

c = c + 1 runs 1 unit of time.

$$f(2) \le c \times n \Longrightarrow when j = 2$$

c = c + 1 runs 1 unit of time.

$$f(3) \le c \times n \Longrightarrow when j = 3$$

c = c + 1 runs 1 unit of time.

... ... ....

$$f\left(\frac{n}{2}\right) \le c \times n \Longrightarrow when j = \frac{n}{2}$$

c = c + 1 runs 1 unit of time.

i.e.when i = 1, the inner most loop statement

$$run\left(1+1+1+1+\dots+\frac{n}{2}\right)=\frac{n}{2} times$$

$$f(2) \le c \times n \Longrightarrow when i = 2$$

$$f(1) \le c \times n \Longrightarrow when j = 1$$

c = c + 1 runs 1 unit of time.

$$f(2) \le c \times n \Longrightarrow when j = 2$$

c = c + 1 runs 1 unit of time.

$$f(3) \le c \times n \Longrightarrow when j = 3$$

c = c + 1 runs 1 unit of time.

... ... ....

$$f\left(\frac{n}{2}\right) \le c \times n \Longrightarrow when j = \frac{n}{2}$$

c = c + 1 runs 1 unit of time.

i. e. when i = 2, the inner most loop statement

$$run\left(1+1+1+1+\dots+\frac{n}{2}\right)=\frac{n}{2} times$$

••••

$$f(n) \le c \times n \Rightarrow when \ i = n$$
 $f(1) \le c \times n \Rightarrow when \ j = 1$ 
 $c = c + 1 \ runs \ 1 \ unit \ of \ time.$ 
 $f(2) \le c \times n \Rightarrow when \ j = 2$ 
 $c = c + 1 \ runs \ 1 \ unit \ of \ time.$ 
 $f(3) \le c \times n \Rightarrow when \ j = 3$ 
 $c = c + 1 \ runs \ 1 \ unit \ of \ time.$ 

... ... ....

$$f(n-1) \le c \times n \Longrightarrow when j = n-1$$
  
 $c = c + 1 runs 1 unit of time.$ 

i.e.when i = n, the inner most loop statement

$$run\left(1+1+1+1+\dots+\frac{n}{2}\right)=\frac{n}{2} times$$

Therefore printing the inner most statement (k = k + 1)

$$n\ times \frac{n}{2} = \frac{n}{2} + \frac{n}{2} + \dots + \frac{n}{2} = n\left(\frac{n}{2}\right) = \frac{n^2}{2}\ times\ ,$$
 hence  $O\left(\frac{n^2}{2}\right) = O(n^2)$