

20.12. TIME COMPLEXITY CALCULATION

SINGLE FOR LOOP (SOME EXAMPLES BASED ON DECREMENT OPERATOR).

1. ($i--$)

```
for( $i = n; i \geq 1; i--$ ){  
     $k = k + 1;$   
}
```

Solution

$k = k + 1$ prints every 1 time taking 1 unit of time as i iterates

from $n, n - 1, \dots, 2, 1$ and $\sum_{i=1}^n 1 = (1 + 1 + 1 + \dots n \text{ times}) = n.$

Hence $O(n).$

2. $(i - 2)$

for($i = n; i \geq 1; i = i - 2$){

$k = k + 1;$

}

Solution

1. *if n is odd*(runs upto n times): $\left\lfloor \frac{n}{2} \right\rfloor + 1 =$ and $O \left(\left\lfloor \frac{n}{2} \right\rfloor + 1 \right) = O(n)$

2. *if n is even*(runs upto n-1 times): $\left\lfloor \frac{n-1}{2} \right\rfloor + 1$ and $O \left(\left\lfloor \frac{n-1}{2} \right\rfloor + 1 \right) = O(n)$

ANOTHER APPROACH

1. *if n is multiple of 2 i. e. an even number it will run $\frac{n}{2}$ times*

as every even number is divisible by 2 or we may say multiples of 2 .

hence $O \left(\frac{n}{2} \right) = O(n)$

2. *if n is not multiple of 2 i. e. an odd number it will run*

$\left\lfloor \frac{n}{2} \right\rfloor + 1$ times , *hence $O \left(\left\lfloor \frac{n}{2} \right\rfloor + 1 \right) = O(n)$*

2. $(i - 3)$

for($i = n; i \geq 1; i = i - 3$){

$k = k + 1;$

}

Solution

1. *if n is multiple of 3:* $\frac{n}{3} =$ and $O\left(\frac{n}{3}\right) = O(n)$

2. *if n is not multiple of 3:* $\left\lfloor \frac{n}{3} \right\rfloor + 1$ and $O\left(\left\lfloor \frac{n}{3} \right\rfloor + 1\right)$

$= O(n)$

3. *if n runs upto n - 1 times :* $\left\lfloor \frac{n-1}{3} \right\rfloor + 1$ and $O\left(\left\lfloor \frac{n-1}{3} \right\rfloor + 1\right) = O(n)$

3. *if (for loop runs from 0 to n - 1)*

for($i = n - 1; i > 0; i = i - 2$){

$k = k + 1;$

}

Solution

As it runs up to 0 to n - 1 times , then the approach will be:

1. *if n is multiple of 2 i. e. an even number it will run $\frac{n}{2}$ times*

as every even number is divisible by 2 or we may say multiples of 2 .

hence $O\left(\frac{n}{2}\right) = O(n)$

2. if n is not multiple of 2 i. e. an odd number it will run

$\left\lfloor \frac{n}{2} \right\rfloor + 1$ times , hence $O\left(\left\lfloor \frac{n}{2} \right\rfloor + 1\right) = O(n)$

4. if (for loop runs from 0 to $n - 1$)

$\text{for}(i = n - 1; i > 0; i = i - 3)\{$

$k = k + 1;$

$\}$

Solution

As it runs up to 0 to $n - 1$ times , then the approach will be:

1. if n is multiple of 3 i. e. an even number it will run $\frac{n}{3}$ times

as every even number is divisible by 3 or we may say multiples of 3 .

hence $O\left(\frac{n}{3}\right) = O(n)$

2. if n is not multiple of 3 i. e. an odd number it will run

$\left\lfloor \frac{n}{3} \right\rfloor + 1$ times , hence $O\left(\left\lfloor \frac{n}{3} \right\rfloor + 1\right) = O(n)$
