## 20.5. TIME COMPLEXITY CALCULATION FOR LOOP (EG-4).

```
//outer loop executed n times for(i = 1; i \le n; i + 3) \{ k = k + 1; // \ constant \ time. \}
```

## **SOLUTION**

Here if we notice i+3 is increment factor, where 1 is lower bound and n is lower bound.

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We can write it as : c_1 \times 1 \le f(n) \le c \times n, Focusing on upper bound or Big 0 time complexity: f(n) \le c \times n, where n is g(n).
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Note: At every incement of i + 3, k = k + 1 prints.

Hence when:

$$f(1) \le c \times n \Longrightarrow when \ i = 1$$
  
 $k = k + 1 \ executes \ in \ 1 \ unit \ of \ time$   
 $i \ increments \ 1 + 3 = 4$ 

$$f(4) \le c \times n$$
 
$$k = k + 1 \text{ executes in 1 unit of time}$$
 
$$i \text{ increments } 4 + 3 = 7$$

$$f(7) \le c \times n$$
 
$$k = k + 1 \ executes \ in \ 1 \ unit \ of \ time$$
 
$$i \ increments \ 7 + 3 = 10$$

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Here the loop runs either : 
$$1 + (1+3) + (4+3) + \cdots + n$$
  
or :  $1 + (1+3) + (4+3) + \cdots + n - 1$ 

Here we will use [] which represents floor value:

$$.\,\emph{i.\,e.}\,\,\lfloor 1.23\,\rfloor = 1, \lfloor 1\,\rfloor = 1\,, \lfloor 2.26\,\rfloor = 2$$

## WHEN IT RUN UP TO n time.

for n=1, k=k+1 will print at i=1,  $\left(\left\lfloor \frac{n}{3}\right\rfloor\right)+1$ 

$$\Rightarrow \left( \left[ \begin{array}{c} \frac{1}{3} \right] + 1 \right) = 0 + 1 = 1$$

for n=2 , k=k+1 will print at i=1,  $\left(\left\lfloor \frac{n}{3}\right\rfloor\right)+1$ 

$$\Rightarrow \left( \left| \begin{array}{c} \frac{2}{3} \right| + 1 \right) = \mathbf{0} + \mathbf{1} = \mathbf{1}$$

for n=3, k=k+1 will print at  $i=1,3,\left(\left\lceil \frac{n}{3}\right\rceil\right)+1$ 

$$\Rightarrow \left( \left| \begin{array}{c|c} \frac{3}{3} & +1 \end{array} \right) = 1 + 1 = 2$$

for n=4 , k=k+1 will print at  $i=1,3,\left(\left\lceil \frac{n}{3}\right\rceil\right)+1$ 

$$\Rightarrow \left( \left[ \begin{array}{c} \frac{4}{3} \right] + 1 \right) = 1 + 1 = 2$$

for n=5, k=k+1 will print at  $i=1,3,\left(\left\lfloor \frac{n}{3}\right\rfloor\right)+1$ 

$$\Rightarrow \left( \left| \begin{array}{c} \frac{5}{3} \right| + 1 \right) = 1 + 1 = 2$$

for n = 6, k = k + 1 will print at  $i = 1, 3, 6, (\lfloor \frac{n}{3} \rfloor) + 1$ 

$$\Rightarrow \left( \left| \begin{array}{c} \frac{6}{3} \right| + 1 \right) = 2 + 1 = 3$$

for n=7 , k=k+1 will print at i=1,3,6,  $\left(\left\lfloor \frac{n}{3}\right\rfloor\right)+1$ 

$$\Rightarrow \left( \left| \begin{array}{c} \frac{7}{3} \right| + 1 \right) = 2 + 1 = 3$$

for n=8, k=k+1 will print at  $i=1,3,6,\left(\left\lfloor \frac{n}{3}\right\rfloor\right)+1$ 

$$\Rightarrow \left( \left| \begin{array}{c} \frac{8}{3} \right| + 1 \right) = 2 + 1 = 3$$

for n=9, k=k+1 will print at  $i=1,3,6,9,\left(\left\lfloor \frac{n}{3}\right\rfloor \right)+1$ 

$$\Rightarrow \left( \left| \frac{9}{3} \right| + 1 \right) = 3 + 1 = 4$$

for n = 10, k = k + 1 will print at  $i = 1, 3, 6, 9, \left( \left| \frac{n}{3} \right| \right) + 1$ 

$$\Rightarrow \left( \left| \begin{array}{c} \frac{10}{3} \right| + 1 \right) = 3 + 1 = 4$$

The innermost statement will be printed :  $O\left(\left(\left[\begin{array}{c} n \\ \overline{3} \end{array}\right]\right) + 1\right) = O(n)$ 

## WHEN IT RUN UP TO n-1 time.

If we notice for (n-1) time,

for 
$$n = 1$$
,  $k = k + 1$  will print at  $i = 1$ ,  $\left(\left|\frac{n-1}{3}\right|\right) + 1$ 

$$\Rightarrow \left( \left| \begin{array}{c} \mathbf{1} - \mathbf{1} \\ \mathbf{3} \end{array} \right| + 1 \right) = \mathbf{0} + \mathbf{1} = \mathbf{1}$$

for 
$$n=2$$
,  $k=k+1$  will print at  $i=1,\left(\left|\begin{array}{c}n-1\\3\end{array}\right|\right)+1$ 

$$\Rightarrow \left( \left\lfloor \frac{2-1}{3} \right\rfloor + 1 \right) = 0 + 1 = 1$$

for 
$$n=3$$
,  $k=k+1$  will print at  $i=1$ ,  $\left(\left\lfloor \frac{n-1}{3}\right\rfloor\right)+1$ 

$$\Rightarrow \left( \left\lfloor \frac{3-1}{3} \right\rfloor + 1 \right) = \mathbf{0} + \mathbf{1} = \mathbf{1}$$

for 
$$n=4$$
,  $k=k+1$  will print at  $i=1,4\left(\left|\begin{array}{c} n-1\\ \hline 3 \end{array}\right|\right)+1$ 

$$\Rightarrow \left( \left| \begin{array}{c} 4-1\\ \hline 3 \end{array} \right| + 1 \right) = 1 + 1 = 2$$

for 
$$n=5$$
,  $k=k+1$  will print at  $i=1,4\left(\left\lfloor \frac{n-1}{3}\right\rfloor\right)+1$ 

$$\Rightarrow \left( \left| \begin{array}{c} 5-1\\ \hline 3 \end{array} \right| + 1 \right) = 1 + 1 = 2$$

for 
$$n=6$$
,  $k=k+1$  will print at  $i=1,4\left(\left\lfloor \frac{n-1}{3}\right\rfloor\right)+1$ 

$$\Rightarrow \left( \left| \begin{array}{c} 6-1\\ \hline 3 \end{array} \right| + 1 \right) = 1 + 1 = 2$$

for 
$$n=7$$
 ,  $k=k+1$  will print at  $i=1,4,7\left(\left\lfloor \left\lfloor \frac{n-1}{3}\right\rfloor \right.\right)+1$ 

$$\Rightarrow \left( \left\lfloor \frac{7-1}{3} \right\rfloor + 1 \right) = 2 + 1 = 3$$

The innermost statement will be printed : 
$$O\left(\left(\left[\begin{array}{c}n-1\\3\end{array}\right]\right)+1\right)=O(n)$$