19. ANALYSIS OF CONSTANT VARIABLE IN A PROGRAM.

TILL NOW WE HAVE LEARNT THAT:

Big - Oh or Worst Case Complexity:

 $f(n) \le cg(n)$, where c and n_0 are constants and $n \ge n_0$, then f(n) = O(g(n))

Big - Omega or Best Case Complexity:

 $f(n) \geq cg(n)$, where c and n_0 are constants and $n \geq n_0$, then $f(n) = \Omega(g(n))$

<u>Big - Theta or Average Case Complexity:</u>

 $c_1g(n) \le f(n) \le c_2g(n)$, where c, n_1 and n_2 are constants and $n \ge \{n_1, n_2\}$,

then
$$f(n) = \Theta(g(n))$$

APPROACH TO FIND TIME COMPLEXITY

Approach to find time complexity of a program is to find the worst complexity i.e., to find 'n' times run of a particular code of a program.

FINDING TIME COMPLEXITY OF CONSTANT TERMS.

REWRITING IT IN ASYMPTOTIC NOTATION

Therefore we write it as:

1. int a=1; It will take one unit of time to run. Where n is size of input (no. of input), here no. of input = 1.It will be also written as:

$$f(n) \le c \times g(n)$$

$$\Rightarrow f(1) \le c \times g(1)$$

- \Rightarrow Hence O(1) i.e. O(c) = O(1), where c is any constant.
- i.e. It will take one unit of time to run.
- 2. float b= 2.9f; It will take one unit of time to run.

 Where n is size of input (no. of input), here no. of input = 1.

 It will be also written as:

$$f(n) \le c \times g(n)$$

$$\Rightarrow f(1) \le c \times g(1)$$

- \Rightarrow Hence O(1) i.e. O(c) = O(1), where c is any constant.
- i.e. It will take one unit of time to run.

3. char c= 'a;'. Where n is size of input (no. of input), here no. of input = 1.

It will be also written as:

$$f(n) \le c \times g(n)$$

 $\Rightarrow f(1) \le c \times g(1)$

 \Rightarrow Hence O(1) i.e. O(c) = O(1), where c is any constant.

 $i.\,e.\,It\,will\,take\,one\,unit\,of\,time\,to\,run\,.$

THAT IS FOR ANY CONSTANT VARIABLE IT WILL RUN O(1) times.