

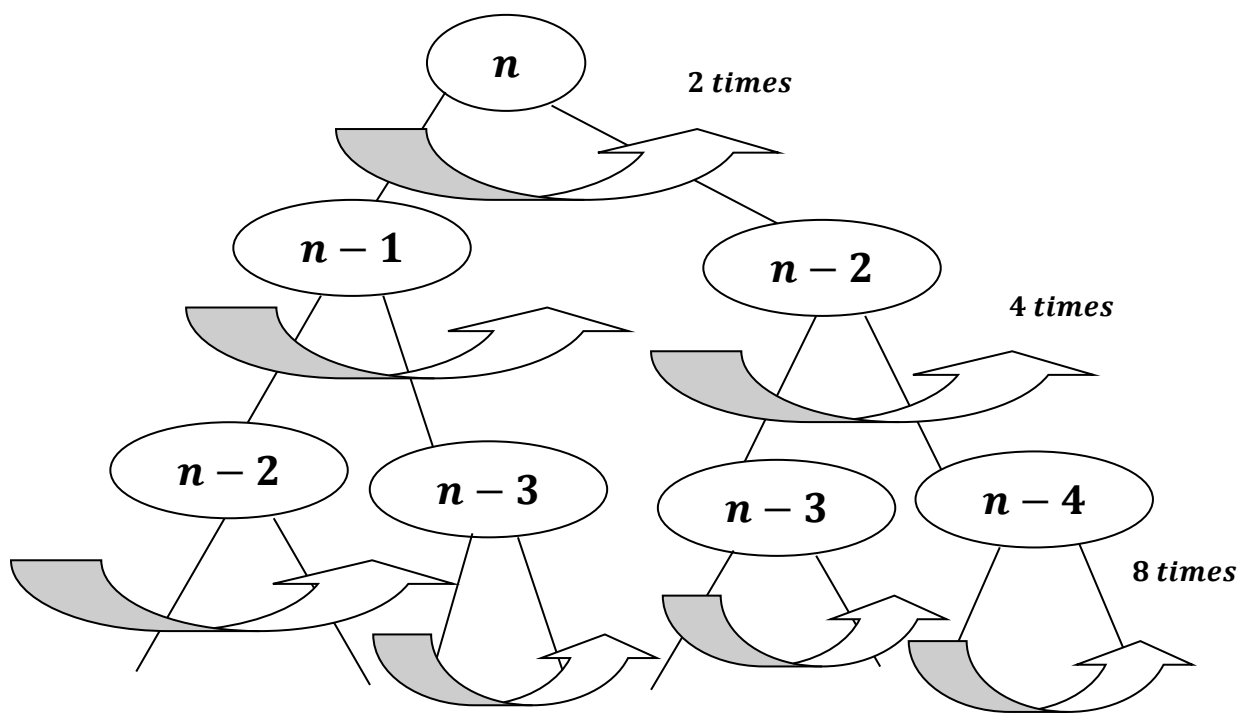
## ***Fibonacci Series Time Complexity with Recursion Tree***

*In fibonacci series :  $T(n) = T(n - 1) + T(n - 2)$ , where `n` is greater than 1 ,  
when `n` = 0 ,  $T(n) = 0$  and  $T(n) = 1$  , when `n` = 1, these are base cases for  
fibonacci series.*

$$T(n) = \begin{cases} 1 & \text{if}(n = 0) \\ 1 & \text{if}(n = 1) \\ T(n - 1) + T(n - 2) & \text{if}(n > 1) \end{cases}$$

***To start with: –***

***If the recursion goes from `n` to `1` times, considering every nodes, we get:***



<i>Level</i>	<i>No. of problems</i>	<i>Problem size</i>	<i>Work Done</i>
<b>0</b>	<b>1</b>	<b>1</b>	<b>1</b>
<b>1</b>	<b>2</b>	<b>1</b>	<b>2</b>
<b>2</b>	<b>4</b>	<b>1</b>	<b>4</b>
<b>.</b>	<b>.</b>	<b>.</b>	<b>.</b>
<b>.</b>	<b>.</b>	<b>.</b>	<b>.</b>
<b><math>n - 1</math></b>	<b><math>2^n</math></b>	<b>1</b>	<b><math>2^n</math></b>

*As we know:*

$$t_n = \begin{cases} 1 & \text{for } n = 1 \\ t_{n-1} + a & \text{for } n > 1 \end{cases}$$

*When  $a$  is 1 , problem size becomes 1 i. e. constant.*

*As we have same growth shown in the recurrence tree for  $T(n - 1)$  and  $T(n - 2)$  we have assumed :*

*$T(n - 1) \approx T(n - 2)$  for previous problem .*

*Now we can see the growth =  $1 + 2 + 4 + \dots 2^{n-1}$*

*We can rewrite it as*

*if we see it,  $1 + 2 + 2^2 + \dots$  is in geometric finite series:*

$$\Rightarrow \frac{x^{n+1} - 1}{x - 1} = \frac{2^{n+1} - 1}{2 - 1} = 2^{n+1} - 1 = O(2^n - 1) = O(2^n) \text{ is time complexity.}$$

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