<u>Handshake Problem - Time Complexity</u> <u>Using substitution method</u>

As we see the program , the return value is: return handshake (n-1)+(n-1) , there is two approaches:

First Approach:

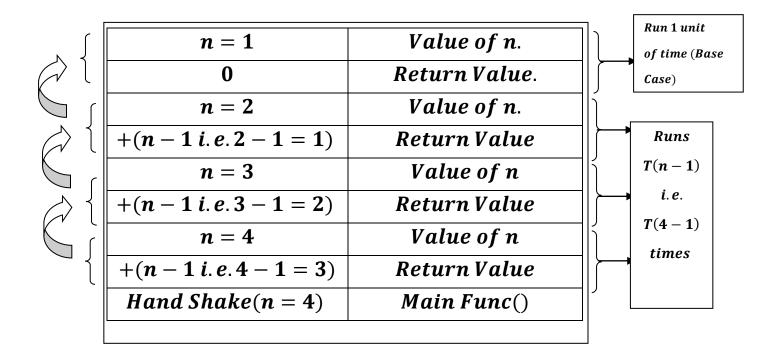
Take n = 4

Hence Base Cases:

n=0, has 0 handshake and it returns 0 takes 1 unit of time.

n=1 , has 0 handshake and it returns 0 takes 1 unit of time.

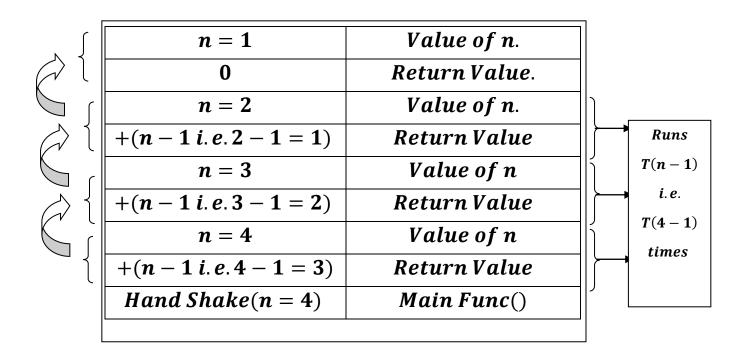
When n > 1, here 4, therefore total push inside stack:



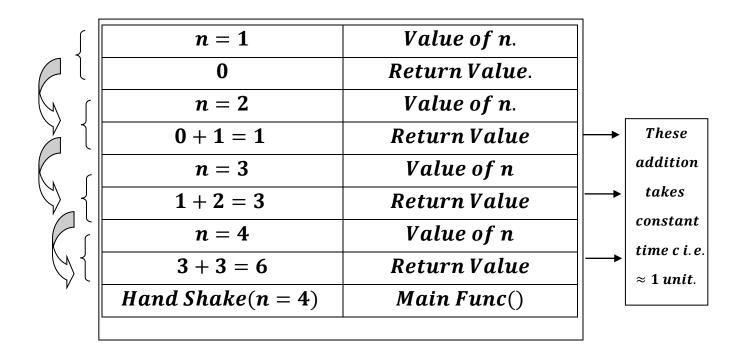
Hence we get our relation:

2nd Approach:

For base cases remain same, Lets take the stack representation of push:



Lets take the addition operation in pop:



$$T(n) = T(n-1) + 1 , for n > 1$$

$$To Compute H(n-1) H(n-1) by n-1$$

Which gives:

$$T(n) =$$

$$1 , for (n = 0)$$

$$for (n = 1)$$

$$T(n-1) + 1 , for (n > 1)$$

Moving with the 1st way:

And we know T(1) = 1, hence rewriting the linear recurrence equation as:

$$T(n) = T(n-1) + 1$$

Or 2nd way:

i. e. same
$$T(n) = T(n-1) + 1$$

Therefore, continuing with the equation: T(n-1) + 1,

$$T(n-1) = T(n-1-1) + 1 = T(n-2) + 1$$

Substituting this in T(n) we get:

$$T(n) = (T(n-2)+1)+1$$

Now,
$$T(n-2) = T(n-2-1) + 1 = T(n-3) + 1$$

Substituting this in T(n) we get:

$$T(n) = ((T(n-3)+1)+1)+1$$

Now if it runs upto `i` times we get:

Now,
$$T(n-i) = T(n-i-1) + 1$$

Substituting this in T(n) we get:

$$T(n) = T(n-i-1) + 1 + 1 + 1 + \cdots i \text{ times}$$

When i = n - 1 we get:

$$T(n) = T(n - (n - 1) - 1) + 1 + 1 + 1 + \dots + (n - 1)$$
times

$$= T(n-n+1-1)+1+1+1....(n-1)times$$

$$= T(0) + 1 + 1 + 1 \dots (n-1)times$$

And we know T(0) = 1 along with T(1) is 1 and as 1 is added to $1 + 1 + 1 + \cdots + (n-1)$ times it becomes:

$$= 1 + 1 + 1 + \cdots n times$$

i. e. $1 \times n = n$ and time complexity is:

$$= \mathbf{0}(\mathbf{n}).$$
