

***Handshake Time complexity
using Recursion Tree Method***

$$T(n) = \begin{cases} 1 & , for (n = 0) \\ 1 & , for (n = 1) \\ T(n-1) + T(1) & , for (n > 1) \end{cases}$$

Or for

$$T(n) = \begin{cases} 1 & , for (n = 0) \\ 1 & , for (n = 1) \\ T(n-1) + 1 & , for (n > 1) \end{cases}$$

We know :

$$t_n = \begin{cases} 1 & for\ n = 1 \\ t_{n-1} + a & for\ n > 1 \end{cases}$$

When a is 1 , problem size becomes 1 i. e. constant.

Therefore, the recurrence tree for this recurrence equation would be shown below:

<i>Level</i>	<i>No. of problems</i>	<i>Problem Size</i>	<i>Work done = Problem Size \times No. of Problems</i>	
0	1	1	$1 \times 1 = 1$	1
1	1	1	$1 \times 1 = 1$	1
2	1	1	$1 \times 1 = 1$	1
.	.	.	.	1
.	.	.	.	
$n - 1$	1	1	$1 \times 1 = 1$	1

One can observe that the work done at every level is 1 and the size of the problem is reduced by a factor of 1 at every level .

Therefore, at the level $n - 1$, the problem size and work done would be 1. The total number of levels is n (as the level of the root is 0). Therefore, the final cost can be estimated as follows:

$$\textit{Total Cost} = \sum_{i=1}^n 1 = 1 + 1 + 1 + \cdots n \textit{ times} = 1 \times n = n.$$

Therefore, the asymptotic complexity would be $O(n)$.
