

Delete Element at a given position in array – Mechanism

10. Shifting of Data in Array[Right Shift].

Static – Delete Element At Position In Array

<i>1 based indexing(logical)</i>	<i>0 based indexing(logical)</i>
$a[1] = 1$	$a[0] = 1$
$a[2] = 2$	$a[1] = 2$
$a[3] = 3$	$a[2] = 3$
$a[4] = 4$	$a[3] = 4$
$a[5] = 5$	$a[4] = 5$
$a[6] = 6$	$a[5] = 6$

If this is 1 based indexing :

Worst Case Scenario

1 – based indexing – deleting $a[1]$

$a[6] \rightarrow a[5] = 6$ [Left to right shift] \rightarrow 1 time

$a[5] \rightarrow a[4] = 5$ [Left to right shift] \rightarrow 2nd time

$a[4] \rightarrow a[3] = 4$ [Left to right shift] \rightarrow 3rd time

$a[3] \rightarrow a[2] = 3$ [Left to right shift] \rightarrow 4rth time

$a[2] \rightarrow a[1] = 2$ [Left to right shift] \rightarrow 5th time

Internal shifting – deleting $a[0]$

$a[i = 0] \rightarrow a[i = 0 + 1] = 2$ [Left to right shift] – 1 time

$a[i = 1] \rightarrow a[i = 1 + 1 = 2] = 3$ [Left to right shift] – 2nd time

$a[i = 2] \rightarrow a[i = 2 + 1 = 3] = 4$ [Left to right shift] – 3rd time

$a[i = 3] \rightarrow a[i = 3 + 1 = 4] = 5$ [Left to right shift] – 4rth time

$a[i = 4] \rightarrow a[i = 4 + 1 = 5] = 6$ [Left to right shift] – 5th time

Lower Bound : $i = pos - 1$

Upper Bound : $i < size - 1$

Operation : $a[i] = a[i + 1]$

```
for (int i = pos - 1; i < size - 1; i++)
{
    a[i] = a[i + 1];
}

size = size - 1;

cout << "The array after deletion is: " << endl;

for (int i = 0; i < size; i++)
{
    cout << "a[" << i << "]= " << a[i] << endl;
}
```

6	0x5014	$a[5]$ or $a + 5$
5	0x5010	$a[4]$ or $a + 4$
4	0x500C	$a[3]$ or $a + 3$
3	0x5008	$a[2]$ or $a + 2$
2	0x5004	$a[1]$ or $a + 1$
1	0x5000	$a[0]$ or $a + 0$

If we view it in this way:

6	5	4	3	2	1
0x5014	0x5010	0x500C	0x5008	0x5004	0x5000
$a[5] \text{ or } a + 5$	$a[4] \text{ or } a + 4$	$a[3] \text{ or } a + 3$	$a[2] \text{ or } a + 2$	$a[1] \text{ or } a + 1$	$a[0] \text{ or } a + 0$
$a[6]$	$a[5]$	$a[4]$	$a[3]$	$a[2]$	$a[1]$

Now we are using 1 – based indexing ,lets say we are deleting $a[1]$:

$i = pos - 1 = 1 - 1 = 0$ and $i = 0 < size - 1 = 6 - 1 = 5$:

$$a[i = 0] = a[0 + 1 = 1] = 2$$

$$a + 0 = 0x5000 + 0 \times 4 \text{ bytes} = 0x5000 + 0_{10} = 0x5000 + 0_{16} = 0x5000$$

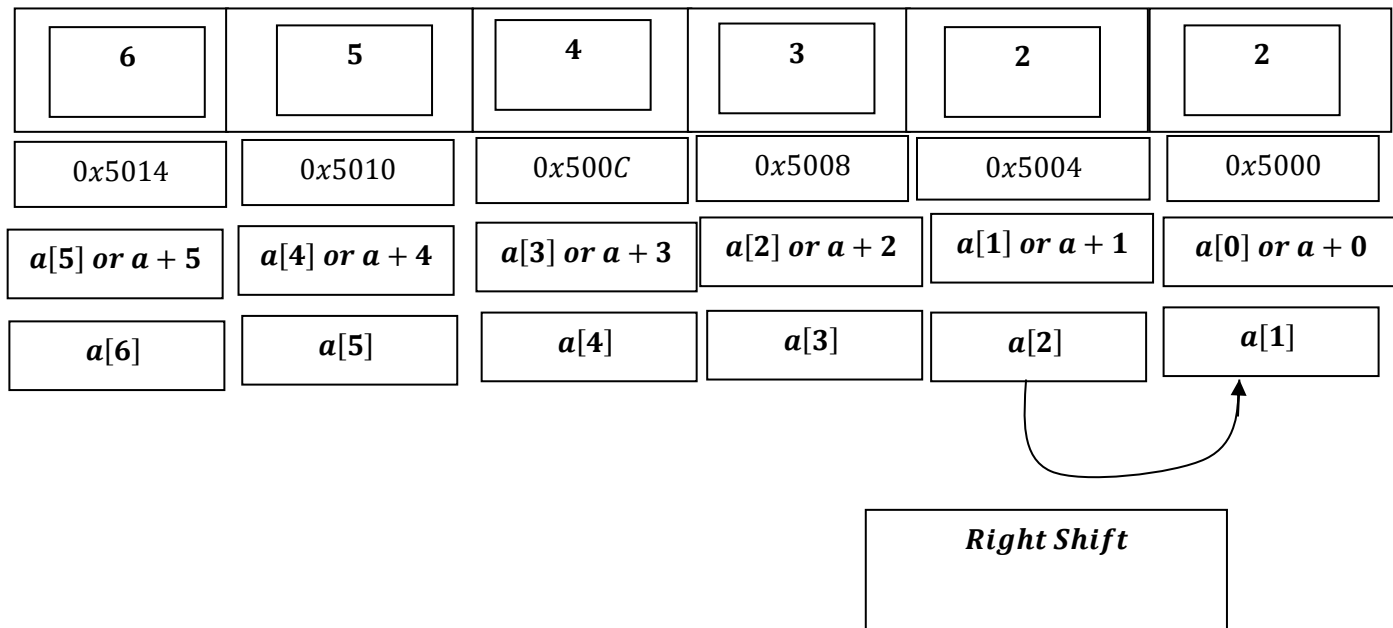
=

$$a + 1 = 0x5000 + 1 \times 4 \text{ bytes} = 0x5000 + 4_{10} = 0x5000 + 4_{16} = 0x5004$$

= 2

6	0x5014	$a[5] \text{ or } a + 5$
5	0x5010	$a[4] \text{ or } a + 4$
4	0x500C	$a[3] \text{ or } a + 3$
3	0x5008	$a[2] \text{ or } a + 2$
2	0x5004	$a[1] \text{ or } a + 1$
2	0x5000	$a[0] \text{ or } a + 0$

If we view it in this way:



Now, $i++ \Rightarrow i = i + 1 \Rightarrow i = 0 + 1 = 1$. [Post increment]

$i = 1$ and $i = 1 < size - 1 = 6 - 1 = 5$:

$$a[i = 1] = a[1 + 1 = 2] = 3$$

$$a + 1 = 0x5000 + 1 \times 4 \text{ bytes} = 0x5000 + 4_{10} = 0x5000 + 4_{16} = 0x5004$$

=

$$a + 2 = 0x5000 + 2 \times 4 \text{ bytes} = 0x5000 + 8_{10} = 0x5000 + 8_{16} = 0x5008$$

= 3

6	0x5014	$a[5] \text{ or } a + 5$
5	0x5010	$a[4] \text{ or } a + 4$
4	0x500C	$a[3] \text{ or } a + 3$
3	0x5008	$a[2] \text{ or } a + 2$
3	0x5004	$a[1] \text{ or } a + 1$
2	0x5000	$a[0] \text{ or } a + 0$

If we view it in this way:

6	5	4	3	3	2
0x5014	0x5010	0x500C	0x5008	0x5004	0x5000
$a[5] \text{ or } a + 5$	$a[4] \text{ or } a + 4$	$a[3] \text{ or } a + 3$	$a[2] \text{ or } a + 2$	$a[1] \text{ or } a + 1$	$a[0] \text{ or } a + 0$
$a[6]$	$a[5]$	$a[4]$	$a[3]$	$a[2]$	$a[1]$

Right Shift

Now, $i++ \Rightarrow i = i + 1 \Rightarrow i = 1 + 1 = 2$. [Post increment]

$i = 2$ and $i = 2 < size - 1 = 6 - 1 = 5$:

$$a[i = 2] = a[2 + 1 = 3] = 4$$

$$a + 2 = 0x5000 + 2 \times 4 \text{ bytes} = 0x5000 + 8_{10} = 0x5000 + 8_{16} = 0x5008$$

=


$$a + 3 = 0x5000 + 3 \times 4 \text{ bytes} = 0x5000 + 12_{10} = 0x5000 + C_{16} = 0x500C$$

= 4.

6	0x5014	$a[5] \text{ or } a + 5$
5	0x5010	$a[4] \text{ or } a + 4$
4	0x500C	$a[3] \text{ or } a + 3$
4	0x5008	$a[2] \text{ or } a + 2$
3	0x5004	$a[1] \text{ or } a + 1$
2	0x5000	$a[0] \text{ or } a + 0$

If we view it in this way:

6	5	4	4	3	2
0x5014	0x5010	0x500C	0x5008	0x5004	0x5000
$a[5] \text{ or } a + 5$	$a[4] \text{ or } a + 4$	$a[3] \text{ or } a + 3$	$a[2] \text{ or } a + 2$	$a[1] \text{ or } a + 1$	$a[0] \text{ or } a + 0$
$a[6]$	$a[5]$	$a[4]$	$a[3]$	$a[2]$	$a[1]$



Right Shift

Now, $i++ \Rightarrow i = i + 1 \Rightarrow i = 2 + 1 = 3$. [*Post increment*]

$i = 3$ and $i = 3 < \text{size} - 1 = 6 - 1 = 5$:

$$a[i = 3] = a[3 + 1 = 4] = 5$$

$$a + 3 = 0x5000 + 3 \times 4 \text{ bytes} = 0x5000 + 12_{10} = 0x5000 + C_{16} = 0x500C$$

=

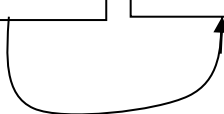
$$a + 4 = 0x5000 + 4 \times 4 \text{ bytes} = 0x5000 + 16_{10} = 0x5000 + 10_{16} = 0x5010$$

= 5.

6	0x5014	$a[5] \text{ or } a + 5$
5	0x5010	$a[4] \text{ or } a + 4$
5	0x500C	$a[3] \text{ or } a + 3$
4	0x5008	$a[2] \text{ or } a + 2$
3	0x5004	$a[1] \text{ or } a + 1$
2	0x5000	$a[0] \text{ or } a + 0$

If we view it in this way:

6	5	5	4	3	2
0x5014	0x5010	0x500C	0x5008	0x5004	0x5000
$a[5] \text{ or } a + 5$	$a[4] \text{ or } a + 4$	$a[3] \text{ or } a + 3$	$a[2] \text{ or } a + 2$	$a[1] \text{ or } a + 1$	$a[0] \text{ or } a + 0$
$a[6]$	$a[5]$	$a[4]$	$a[3]$	$a[2]$	$a[1]$



Right Shift

Now, $i++ \Rightarrow i = i + 1 \Rightarrow i = 3 + 1 = 4$. [*Post increment*]

$i = 4$ and $i = 4 < \text{size} - 1 = 6 - 1 = 5$:

$$a[i = 4] = a[4 + 1 = 5] = 6$$

$$a + 4 = 0x5000 + 4 \times 4 \text{ bytes} = 0x5000 + 16_{10} = 0x5000 + 10_{16} = 0x5010$$

=

$$a + 5 = 0x5000 + 5 \times 4 \text{ bytes} = 0x5000 + 20_{10} = 0x5000 + 14_{16} = 0x5014$$

= 6.

6	0x5014	$a[5] \text{ or } a + 5$
6	0x5010	$a[4] \text{ or } a + 4$
5	0x500C	$a[3] \text{ or } a + 3$
4	0x5008	$a[2] \text{ or } a + 2$
3	0x5004	$a[1] \text{ or } a + 1$
2	0x5000	$a[0] \text{ or } a + 0$

If we view it in this way:

6	6	5	4	3	2
0x5014	0x5010	0x500C	0x5008	0x5004	0x5000
$a[5] \text{ or } a + 5$	$a[4] \text{ or } a + 4$	$a[3] \text{ or } a + 3$	$a[2] \text{ or } a + 2$	$a[1] \text{ or } a + 1$	$a[0] \text{ or } a + 0$
$a[6]$	$a[5]$	$a[4]$	$a[3]$	$a[2]$	$a[1]$

Right Shift

Now, $i++ \Rightarrow i = i + 1 \Rightarrow i = 4 + 1 = 5$. [Post increment]

$i = 5$ and $i = 5 < size - 1 = 6 - 1 = 5$, the condition become false, hence loop exits.

Then:

```
size = size - 1;
```

size = 6 - 1 = 5, it becomes 5.

```
for (int i = 0; i < size; i++)  
{  
    cout << "a[" << i << "]=" << a[i] << endl;  
}
```

It prints:

$a[0] = 2$

$a[1] = 3$

$a[2] = 4$

$a[3] = 5$

$a[4] = 6$

For 0 – based indexing

$a[i = 0] \rightarrow a[i = 0 + 1] = 2$ [Left to right shift] – 1 time

$a[i = 1] \rightarrow a[i = 1 + 1 = 2] = 3$ [Left to right shift] – 2nd time

$a[i = 2] \rightarrow a[i = 2 + 1 = 3] = 4$ [Left to right shift] – 3rd time

$a[i = 3] \rightarrow a[i = 3 + 1 = 4] = 5$ [Left to right shift] – 4rth time

$a[i = 4] \rightarrow a[i = 4 + 1 = 5] = 6$ [Left to right shift] – 5th time

Lower Bound : $i = pos$

Upper Bound : $i < size - 1$

Operation : $a[i] = a[i + 1]$

```
for (int i = pos ; i < size - 1; i++)
{
    a[i] = a[i + 1];
}

size = size - 1;

cout << "The array after deletion is: " << endl;

for (int i = 0; i < size; i++)
{
    cout << "a[" << i << "]= " << a[i] << endl;
}
```

This also change input validation ,for 0 – based indexing:

```
if (pos < 0 || pos > size-1)
{
    cout << "Invalid position";
}
```

Also ,input validation ,for 1 – based indexing:

```
if (pos < 1 || pos > size)
{
    cout << "Invalid position";
}
```
