

Alternative Approach of Stack – Time Complexity

1. ISFULL function to check the stack is full or not as discussed before:

```
int isFull(Stack st)
{
    if (st.top == st.size - 1)
    {
        return 1;
    }
    return 0;
}
```

2. Double Stack function to double the stack if the stack is full :

```
void doubleStack(Stack *st)
{
    st->size = st->size * 2;
    st->s = (int *)realloc(st->s, st->size *
sizeof(int));
}
```

ReAlloc will reallocate the memory again with a new size of the array created.

3. In Push , when the element is pushed in Stack then its checked that if its full then double the stack , else just increment the top and push the element.

```
void push(Stack *st, int x)
{
    if (isFull(*st))
    {
        doubleStack(st);
        st->top++;
        st->s[st->top] = x;
        return;
    }

    else{
        st->top++;
        st->s[st->top] = x;
    }
}
```

Now, Coming back to execution:

At first we create object of Stack and declare a variable to decide what should be capacity of the stack in main function.

```
int main()
{
    Stack st;
    int cap;
    cout << "Enter the capacity of the stack" << endl;

    cin >> cap;
}
```

4. We know during push it will take $2n$ elements as as the size gets doubled:

```
int main()
{
    int elem;
    for (int i = 0; i < 2* cap; i++)
    {
        cout << "Enter the element to be pushed :" << endl;
        cin >> elem;
        push(&st, elem);
    }

}
```

5. Now pop function must pop the first set of elements that are (inserted before the array got doubled) :

```
int pop(Stack *st)
{
    if (isEmpty(*st))
    {
        cout << "Stack Underflow" << endl;
        return -1;
    }

    return st->s[st->top--];
}

int main()
{
    if (!isEmpty(st))
    {
        for (int i = 0; i < cap; i++)
        {
            cout << "Popped element is: " << pop(&st) << endl;
        }
    }
    else
    {
        cout << "Stack Underflow" << endl;
    }
}
```

Else , every other functions are all same earlier discussed .

Now coming to Time Complexity

1. Time Complexity of Push Operation:

Ans: As discussed earlier, Push Operation here takes Worst Case Complexity: $T(n) = O(n)$, where as, amortized time complexity = $O(1)$.

Q) Why we need amortized time complexity?

Ans) When an operation sometimes requires copying many elements, hence we need to resize then amortized analysis becomes necessary.

2. Time Complexity of Pop Operation:

Ans : As we are popping from each stack each at $O(1)$ time , now we are using for loop making: $n \times O(1)$ time = $O(n)$ time at worst case.

Q) Does here amortization time complexity require?

Ans: No , as no copying of elements and no resizing take place.

And Other Time Complexity of all types of Operations remains same that is $O(1)$.

Now coming to Space Complexity

As the stack takes `n` times push takes n units of space in memory , hence space complexity = $T(n) = O(n)$.

Hence, we sum up:

<i>Space complexity</i>	<i>$O(n)$</i>
<i>Creation of Stack</i>	<i>$O(1)$</i>
<i>Push</i>	<i>$O(1)$ Average</i>
<i>Pop()</i>	<i>pop func: $O(1)$, as for loop used to pop `n` times in main func making , $n \times O(1) = O(n)$.</i>
<i>Peek()</i>	<i>$O(1)$</i>
<i>IsEmpty()</i>	<i>$O(1)$</i>
<i>IsFull()</i>	<i>$O(1)$</i>
<i>DeleteStack()</i>	<i>$O(1)$</i>

