| Character                   | Binary<br>code | Character | Binary<br>code |
|-----------------------------|----------------|-----------|----------------|
| A                           | 100 0001       | 0         | 011 0000       |
| В                           | 100 0010       | 1         | 011 0001       |
| C                           | 100 0011       | 2         | 011 0010       |
| D                           | 100 0100       | 3         | 011 0011       |
| E                           | 100 0101       | 4         | 011 0100       |
| F                           | 100 0110       | 5         | 011 0101       |
| G                           | 100 0111       | 6         | 011 0110       |
| Н                           | 100 1000       | 7         | 011 0111       |
| I                           | 100 1001       | 8         | 011 1000       |
| J                           | 100 1010       | 9         | 011 1001       |
| K                           | 100 1011       |           |                |
| L                           | 100 1100       |           |                |
| M                           | 100 1101       | space     | 010 0000       |
| N                           | 100 1110       | •         | 010 1110       |
| O                           | 100 1111       | (         | 010 1000       |
| P                           | 101 0000       | ÷         | 010 1011       |
| Q                           | 101 0001       | \$        | 010 0100       |
| $\widetilde{\widetilde{R}}$ | 101 0010       | *         | 010 1010       |
| S                           | 101 0011       | )         | 010 1001       |
| T                           | 101 0100       | _         | 010 1101       |

 TABLE 3-4 American Standard Code for information Interchange (ASCII)

bits stored in registers so that operations are performed on operands of the same type. In inspecting the bits of a computer register at random, one is likely to find that it represents some type of coded information rather than a binary number.

010 1111

010 1100

011 1101

101 0101

101 0110

101 0111

101 1000

101 1001

101 1010

Binary codes can be formulated for any set of discrete elements such as the musical notes and chess pieces and their positions on the chessboard. Binary codes are also used to formulate instructions that specify control information for the computer. This chapter is concerned with *data* representation. Instruction codes are discussed in Chap. 5.

## 3-2 Complements

U

V

W

X

Y

Z

Complements are used in digital computers for simplifying the subtraction operation and for logical manipulation. There are two types of complements for each base r system: the r's complement and the (r-1)'s complement.

When the value of the base r is substituted in the name, the two types are referred to as the 2's and l's complement for binary numbers and the 10's and 9's complement for decimal numbers.

## (r-1)'s Complement

9's complement

Given a number N in base r having n digits, the (r-1)'s complement of N is defined as  $(r^n-1)-N$ . For decimal numbers r=10 and r-1=9, so the 9's complement of N is  $(10^n-1)-N$ . Now,  $10^n$  represents a number that consists of a single 1 followed by n 0's.  $10^n-1$  is a number represented by n 9's. For example, with n=4 we have  $10^4=10000$  and  $10^4-1=9999$ . It follows that the 9's complement of a decimal number is obtained by subtracting each digit from 9. For example, the 9's complement of 546700 is 999999-546700=453299 and the 9's complement of 12389 is 99999-12389=87610.

1's complement

For binary numbers, r = 2 and r - 1 = 1, so the 1's complement of N is  $(2^n - 1) - N$ . Again,  $2^n$  is represented by a binary number that consists of a 1 followed by n 0's.  $2^n - 1$  is a binary number represented by n 1's. For example, with n = 4, we have  $2^4 = (10000)_2$  and  $2^4 - 1 = (1111)_2$ . Thus the 1's complement of a binary number is obtained by subtracting each digit from 1. However, the subtraction of a binary digit from 1 causes the bit to change from 0 to 1 or from 1 to 0. Therefore, the 1's complement of a binary number is formed by changing 1's into 0's and 0's into 1's. For example, the 1's complement of 1011001 is 0100110 and the 1's complement of 0001111 is 1110000.

The (r-1)'s complement of octal or hexadecimal numbers are obtained by subtracting each digit from 7 or F (decimal 15) respectively.

## (r's) Complement

The r's complement of an n-digit number N in base r is defined as  $r^n - N$  for  $N \neq 0$  and 0 for N = 0. Comparing with the (r - 1)'s complement, we note that the r's complement is obtained by adding 1 to the (r - 1)'s complement since  $r^n - N = [(r^n - 1) - N] + 1$ . Thus the 10's complement of the decimal 2389 is 7610 + 1 = 7611 and is obtained by adding 1 to the 9's complement value. The 2's complement of binary 101100 is 010011 + 1 = 010100 and is obtained by adding 1 to the 1's complement value.

10's complement

Since  $10^n$  is a number represented by a 1 followed by n 0's, then  $10^n - N$ , which is the 10's complement of N, can be formed also be leaving all least significant 0's unchanged, subtracting the first nonzero least significant digit from 10, and then subtracting all higher significant digits from 9. The 10's complement of 246700 is 753300 and is obtained by leaving the two zeros unchanged, subtracting 7 from 10, and subtracting the other three digits from 9. Similarly, the 2's complement can be formed by leaving all least significant 0's and the first 1 unchanged, and then replacing l's by 0's and 0's by l's in all other higher, significant bits. The 2's complement of 1101100 is 0010100 and is obtained by leaving the two low-order 0's and the first 1 unchanged, and then replacing l's by 0's and 0's by l's in the other four most significant bits.

2's complement

In the definitions above it was assumed that the numbers do not have a radix point. If the original number N contains a radix point, it should be removed temporarily to form the r's or (r-1)'s complement. The radix point is then restored to the complemented number in the same relative position. It is also worth mentioning that the complement of the complement restores the number to its original value. The r's complement of N is  $r^n - N$ . The complement of the complement is  $r^n - (r^n - N) = N$  giving back the original number.

## Subtraction of Unsigned Numbers

The direct method of subtraction taught in elementary schools uses the borrow concept. In this method we borrow a 1 from a higher significant position when the minuend digit is smaller than the corresponding subtrahend digit. This seems to be easiest when people perform subtraction with paper and pencil. When subtraction is implemented with digital hardware, this method is found to be less efficient than the method that uses complements.

The subtraction of two n-digit unsigned numbers  $M - N(N \neq 0)$  in base r can be done as follows:

- **1.** Add the minuend M to the r's complement of the subtrahend N. This performs  $M + (r^n N) = M N + r^n$ .
- **2.** If  $M \ge N$ , the sum will produce an end carry  $r^n$  which is discarded, and what is left is the result M N.
- **3.** If M < N, the sum does not produce an end carry and is equal to  $r^n (N M)$ , which is the r's complement of (N M). To obtain the answer in a familiar form, take the r's complement of the sum and place a negative sign in front.

Consider, for example, the subtraction 72532 - 13250 = 59282. The 10's complement of 13250 is 86750. Therefore:

```
M = 72532

10's complement of N = +86750

Sum = 159282

Discard end carry 10^5 = -\frac{100000}{59282}

Answer = \frac{100000}{59282}
```

Now consider an example with M < N. The subtraction 13250 - 72532 produces negative 59282. Using the procedure with complements, we have

$$M = 13250$$
10's complement of  $N = +27468$ 

$$Sum = 40718$$

subtraction

end carry