

IS5311-Discrete Mathematics
Chapter 1-Introduction to Discrete Mathematics

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Introduction

Key areas of discrete mathematics include:

- ① Sets and Relations
- ② Logic
- ③ Combinatorics
- ④ Graph Theory
- ⑤ Number Theory
- ⑥ Discrete Probability

Importance of Discrete Mathematics

- ① Foundation for Computer Science
- ② Cryptography and Information Security
- ③ Networks and Communication Systems
- ④ Operations Research and Optimization
- ⑤ Artificial Intelligence and Machine Learning
- ⑥ Combinatorial Optimization
- ⑦ Logic and Reasoning

Propositional Logic

Proposition

A proposition is a declarative sentence (that is, a sentence that declares a fact) that is either true or false, but not both.

Example

- ① Wi-Fi is available in the building. (T)
- ② Instagram is older than Facebook. (F)
- ③ $2 + 2 = 4$. (T)
- ④ Every smartphone has a headphone jack. (F)

Propositional Logic

Some examples of sentences that are not propositions.

- ① Did the program run? (Question)
- ② Charge your phone! (Command)
- ③ x is an even number (Open sentence)
- ④ $x + y = z$.

Key Concepts

- Propositional Variables
- Truth Values
- Propositional Calculus/Logic

Many mathematical statements are constructed by combining one or more propositions. New propositions, called **compound propositions**, are formed from existing propositions using logical operators.

Logical Operators

Negation(\neg)

Let p be a proposition. The negation of p , denoted by $\neg p$ (also denoted by \bar{p}), is the statement “It is not the case that p .” The proposition $\neg p$ is read “**not** p .”

Negation

Example

① Proposition: "The program is running."

Negation: "It is **not** the case that the program is running."

Simpler English: "The program is **not** running."

Negation

p	$\neg p$
T	F
F	T

Table: Truth table for the negation of a proposition p .

The negation operator constructs a new proposition from a single existing proposition. We will now introduce the logical operators that are used to form new propositions from two or more existing propositions. These logical operators are also called **connectives**.

Conjunction

Conjunction(\wedge)

Let p and q be propositions. The conjunction of p and q , denoted by $p \wedge q$, is the proposition “ p **and** q .”

Conjunction

Example

- A: The sensor is working.
- B: The sensor is connected.

Then $A \wedge B$: The sensor is working and the sensor is connected.

Conjunction

p	q	$p \wedge q$
T	T	T
T	F	F
F	T	F
F	F	F

Table: The Truth Table for the Conjunction of Two Propositions.

Disjunction

Disjunction (\vee)

Let p and q be propositions. The disjunction of p and q , denoted by $p \vee q$, is the proposition “ p or q .”

Disjunction

Example

- p: "The user pressed Enter."
- q: "The user clicked the button."

$p \vee q$: "The user pressed Enter **or** clicked the button."

Disjunction

p	q	$p \vee q$
T	T	T
T	F	T
F	T	T
F	F	F

Table: The Truth Table for the Disjunction of Two Propositions.

Exclusive

Exclusive(\oplus)

Let p and q be propositions. The exclusive or of p and q , denoted by $p \oplus q$, is the proposition that is true when exactly one of p and q is true and is false otherwise.

Exclusive

Example

- p: "The system uses a wired network."
- q: "The system uses a wireless network."

$p \oplus q$: "The system uses either a wired or wireless network, but **not both**.

Exclusive

p	q	$p \oplus q$
T	T	F
T	F	T
F	T	T
F	F	F

Table: The Truth Table for the Exclusive Or of Two Propositions

Conditional Statements

Conditional Statements(\rightarrow)

Let p and q be propositions. The conditional statement $p \rightarrow q$ is the proposition “**if** p , **then** q . ”

Conditional Statements

A conditional statement, often referred to as an "**if-then**" statement, expresses a relationship between two propositions where the truth of one proposition depends on the truth of the other.

Example

- ① "**If** it is raining outside, **then** I will bring an umbrella."
- ② "**If** I am elected, **then** I will lower taxes."
- ③ "**If** you get 100 on the final, **then** you will get an A."

Conditional Statements

p	q	$p \rightarrow q$
T	T	T
T	F	F
F	T	T
F	F	T

Table: The Truth Table for the Conditional Statement $p \rightarrow q$.

Conditional Statements

Exercise:

- ① Let p be the statement "The battery is low.", and q be the statement "The system sends a warning." Express $p \rightarrow q$ in English ?

Conditional Statements

We can form some new conditional statements starting with a conditional statement $p \rightarrow q$. In particular, there are three related conditional statements that occur so often that they have special names.

- The proposition $q \rightarrow p$ is called the **converse** of $p \rightarrow q$.

q	p	$q \rightarrow p$
T	T	
T	F	
F	T	
F	F	

Conditional Statements

- The **contrapositive** of $p \rightarrow q$ is the proposition $\neg q \rightarrow \neg p$.

$\neg q$	$\neg p$	$\neg q \rightarrow \neg p$
T	T	
T	F	
F	T	
F	F	

- The proposition $\neg p \rightarrow \neg q$ is called the **inverse** of $p \rightarrow q$.

$\neg p$	$\neg q$	$\neg p \rightarrow \neg q$
T	T	
T	F	
F	T	
F	F	

Conditional Statements

When two compound propositions always have the same truth value we call them **equivalent**.

- The conditional statement and its contrapositive are equivalent.
- The converse and the inverse of a conditional statement are also equivalent.

Conditional Statements

Example: What are the contrapositive, the converse, and the inverse of the conditional statement

“If the firewall is active, then the network is protected.”

Conditional Statements

Solution:

- The contrapositive of this conditional statement is
“If the network is not protected, then the firewall is not active.”
- The converse is
“If the network is protected, then the firewall is active.”
- The inverse is
“If the firewall is not active, then the network is not protected.”

Biconditional Statement

Biconditional Statement(\leftrightarrow)

Let p and q be propositions. The biconditional statement $p \leftrightarrow q$ is the proposition " **p if and only if q .**"

Note that the statement $p \leftrightarrow q$ is true when both the conditional statements $p \rightarrow q$ and $q \rightarrow p$ are true and is false otherwise.

Biconditional Statement

Example:

Let p be the statement "You attend the meeting," and let q be the statement "You receive the meeting minutes."

Biconditional Statement ($p \leftrightarrow q$):

"You can attend the meeting if and only if you receive the meeting minutes."

Truth Conditions:

- The statement is true if:
 - You attend the meeting and you receive the meeting minutes.
 - You do not attend the meeting and you do not receive the meeting minutes.
- The statement is false if:
 - You attend the meeting but do not receive the meeting minutes.
 - You do not attend the meeting but receive the meeting minutes.

Biconditional Statement

p	q	$p \leftrightarrow q$
T	T	T
T	F	F
F	T	F
F	F	T

Table: The Truth Table for the Biconditional $p \leftrightarrow q$.

Biconditional Statement

Example: Determine whether each of the following bi-conditional statement is true or false.

No.	Statement	True/False
1	A number is divisible by 10 if and only if it ends in 0.	
2	A triangle is equilateral if and only if it has three equal angles.	
3	An animal is a mammal if and only if it lays eggs.	
4	$2 + 2 = 5$ if and only if $3 + 3 = 6$.	
5	A person is a bachelor if and only if they are an unmarried man.	
6	A square is a rhombus if and only if all its angles are right angles.	

Biconditional Statement

Solution:

No.	Statement	True/False
1	A number is divisible by 10 if and only if it ends in 0.	True
2	A triangle is equilateral if and only if it has three equal angles.	True
3	An animal is a mammal if and only if it lays eggs.	False
4	$2 + 2 = 5$ if and only if $3 + 3 = 6$.	False
5	A person is a bachelor if and only if they are an unmarried man.	True
6	A square is a rhombus if and only if all its angles are right angles.	False

Truth Tables of Compound Propositions

Example: Construct the truth table for the compound proposition $(p \wedge q) \vee r$.

p	q	r	$(p \wedge q)$	$((p \wedge q) \vee r)$
T	T	T	T	T
T	T	F	T	T
T	F	T	F	T
T	F	F	F	F
F	T	T	F	T
F	T	F	F	F
F	F	T	F	T
F	F	F	F	F

Table: The Truth Table of $((p \wedge q) \vee r)$.

Truth Tables of Compound Propositions

Example: Construct the truth table for the compound proposition $(p \wedge \neg r) \vee (q \rightarrow r)$.

Solution:

p	q	r	$\neg r$	$p \wedge \neg r$	$q \rightarrow r$	$(p \wedge \neg r) \vee (q \rightarrow r)$
T	T	T	F	F	T	T
T	T	F	T	T	F	T
T	F	T	F	F	T	T
T	F	F	T	T	T	T
F	T	T	F	F	T	T
F	T	F	T	F	F	F
F	F	T	F	F	T	T
F	F	F	T	F	T	T

Table: The Truth Table of $(p \wedge \neg r) \vee (q \rightarrow r)$.

Exercise: Construct a truth table for each of these compound propositions.

- ① $p \vee \neg p$
- ② $p \wedge q$
- ③ $p \vee q \wedge p$
- ④ $(p \vee q) \rightarrow (q \wedge r)$
- ⑤ $\neg(p \wedge q) \leftrightarrow (\neg p \vee \neg q)$

Precedence of Logical Operators

Operator	Precedence
\neg	1
\wedge	2
\vee	3
\rightarrow	4
\leftrightarrow	5

Table: Precedence of Logical Operators.

Logic and Bit Operations

Computers represent information using bits. A bit is a symbol with two possible values: 0 (zero) and 1 (one). A bit can represent a truth value because there are two truth values: true and false. We use 1 to represent true (T) and 0 to represent false (F).

Truth Value	Bit
T	1
F	0

A variable is called a **Boolean variable** if its value is either true or false. Consequently, a Boolean variable can be represented using a bit.

Logic and Bit Operations

p	q	$\text{AND}(p \wedge q)$
1	1	1
1	0	0
0	1	0
0	0	0

p	q	$\text{OR}(p \vee q)$
1	1	1
1	0	1
0	1	1
0	0	0

p	q	$\text{XOR}(p \oplus q)$
1	1	0
1	0	1
0	1	1
0	0	0

Table: Tables for the Bit Operators OR, AND, and XOR.

Logic and Bit Operations

Bit String

A **bit string** is a sequence of zero or more bits. The length of this string is the number of bits in the string.

Example: 11011000 is a bit string of length eight.

Example: Given two bit strings 0110110110 and 1100011101, find their bitwise OR, bitwise AND, and bitwise XOR.

Logic and Bit Operations

Solution: Bitwise OR:

$$\begin{array}{r} 01110110 \\ \vee\quad 11001101 \\ \hline 11111111 \end{array}$$

Bitwise AND:

$$\begin{array}{r} 01110110 \\ \wedge\quad 11001101 \\ \hline 01000100 \end{array}$$

Bitwise XOR:

$$\begin{array}{r} 01110110 \\ \oplus\quad 11001101 \\ \hline 10111011 \end{array}$$

Logic and Bit Operations

Exercise: Evaluate each of these expressions.

- ① $(101011 \oplus 110011) \wedge (011010 \vee 100110)$.
- ② $(10101 \wedge 11011) \vee (10010 \oplus 01010)$.
- ③ $(111001 \wedge 100011) \oplus (110110 \vee 101010)$.
- ④ $(1111 \vee 0101) \wedge (1010 \oplus 0110)$.

Applications of Propositional Logic: Translating Statements

Translating English sentences into expressions with propositional variables and logical connectives serves several purposes. Once translated, we can analyze these expressions to determine their truth values.

Example:: Translate the following English sentence into a logical expression:

"If you upload your document and submit the form, then your assignment will be accepted."

Translating Statements

Solution: Let's represent:

- a : "You upload the document."
- b : "You submit the form."
- c : "Your assignment will be accepted"

So, we can translate the sentence into the logical expression

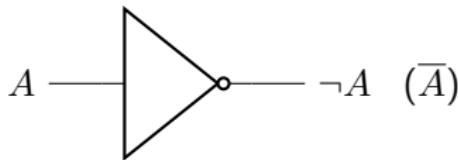
$$(a \wedge b) \rightarrow c.$$

Logic Circuits

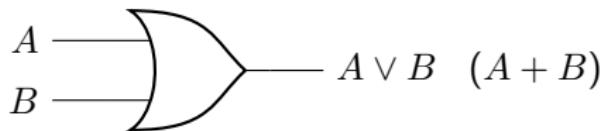
In a logic circuit, also known as a digital circuit, input signals p_1, p_2, \dots, p_n are received, each representing a bit (0 for "off" and 1 for "on"). The circuit processes these inputs and produces output signals s_1, s_2, \dots, s_n , also represented as bits.

Complicated digital circuits can be constructed from **three basic circuits**, called **gates**. These gates are fundamental building blocks for constructing complex digital circuits. Here are the three basic types of gates typically used:

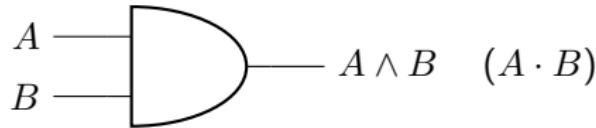
Logic Circuits



- ① This gate, also known as an inverter, produces an output that is the complement of its input.



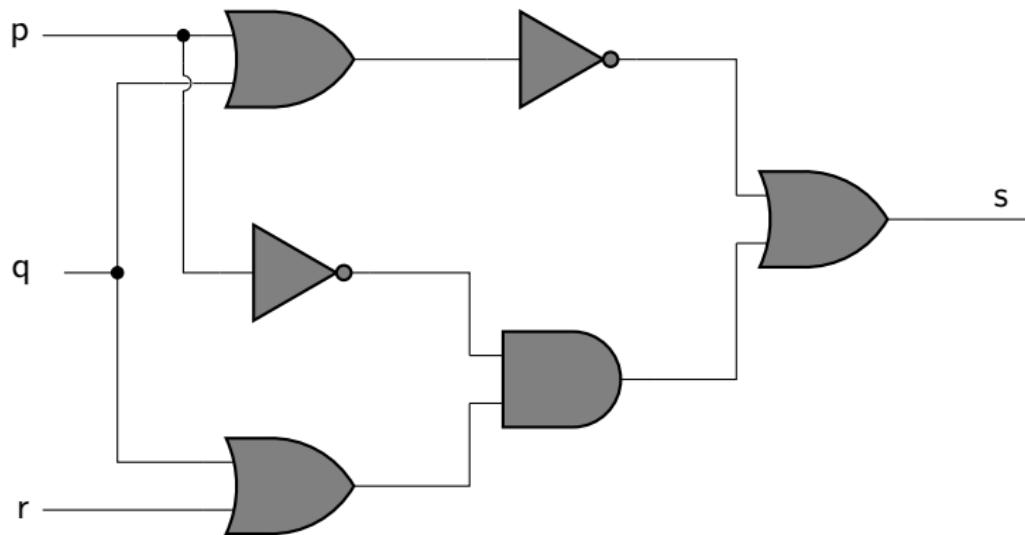
- ② This gate produces a high output (1) if any of its inputs are high (1). It produces a low output (0) only if all its inputs are low (0).



- ③ This gate produces a high output (1) only if all of its inputs are high (1). Otherwise, it produces a low output (0).

Logic Circuits

Example: Determine the output for the combinatorial circuit in the following figures?



Logic Circuits

Let us consider the logical flow of the circuit to determine the output s based on the inputs p , q , and r .

- First OR Gate (ORa):
 - Inputs: p and q
 - Output: $p \vee q$
- NOT Gate (Nob):
 - Input: Output of ORa, which is $p \vee q$
 - Output: $\neg(p \vee q)$
- Second OR Gate (ORb):
 - Inputs: q and r
 - Output: $q \vee r$
- NOT Gate (Noa):
 - Input: q
 - Output: $\neg q$

Logic Circuits

- AND Gate (ANDa):
 - Inputs: Output of Noa ($\neg q$) and Output of ORb ($q \vee r$)
 - Output: $\neg q \wedge (q \vee r)$
- Third OR Gate (ORc):
 - Inputs: Output of Nob ($\neg(p \vee q)$) and Output of ANDa ($\neg q \wedge (q \vee r)$)
 - Output: $\neg(p \vee q) \vee (\neg q \wedge (q \vee r))$

Thus, the output of the circuit s is:

$$s = \neg(p \vee q) \vee (\neg q \wedge (q \vee r))$$

Logic Circuits

Example: Consider the following logical expression: $(p \vee \neg r) \wedge (\neg p \vee (q \vee \neg r))$. Design a digital circuit that produces this output given input bits p , q , and r .

Propositional Equivalences

Tautology

A **tautology** is a compound proposition that is always true.

Example:

- ① The proposition $p \vee \neg p$ is a tautology because it is always true.
- ② You are online or not online.

Propositional Equivalences

Contradiction

In contrast, a **contradiction** is a compound proposition that is always false.

Example:

- ① The proposition $p \wedge \neg p$ is a contradiction because it can never be true; it asserts that p is both true and false simultaneously.
- ② You are online and not online at the same time.

Propositional Equivalences

Contingency

A **contingency** is a compound proposition that is neither a tautology nor a contradiction. Its truth value depends on the specific truth values assigned to its constituent propositional variables.

Example:

- ① The proposition $p \vee q$ is a contingency because its truth value depends on the truth values of p and q . If both p and q are true, the proposition is true; if both are false, the proposition is false.
- ② The sun is shining.

Logical Equivalences

Logically Equivalent

The compound propositions p and q are called **logically equivalent** if $p \leftrightarrow q$ is a tautology. The notation $p \equiv q$ denotes that p and q are logically equivalent.

Table: De Morgan's Laws

Expression	Equivalent Expression
$\neg(p \wedge q)$	$\neg p \vee \neg q$
$\neg(p \vee q)$	$\neg p \wedge \neg q$

Logical Equivalences

Example: To show that $\neg(p \vee q)$ and $\neg p \wedge \neg q$ are logically equivalent, we can use a truth table.

p	q	$p \vee q$	$\neg(p \vee q)$	$\neg p$	$\neg p \wedge \neg q$
T	T	T	F	F	F
T	F	T	F	F	F
F	T	T	F	T	F
F	F	F	T	T	T

Table: Truth Table for $\neg(p \vee q)$ and $\neg p \wedge \neg q$

Exercise: Show that $p \rightarrow q$ and $\neg p \vee q$ are logically equivalent.

Logical Equivalences

Equivalence	Name
$p \wedge T \equiv p$ $p \vee F \equiv p$	Identity laws
$p \vee T \equiv T$ $p \wedge F \equiv F$	Domination laws
$p \vee p \equiv p$ $p \wedge p \equiv p$	Idempotent laws
$\neg(\neg p) \equiv p$	Double negation law
$p \vee q \equiv q \vee p$ $p \wedge q \equiv q \wedge p$	Commutative laws
$(p \vee q) \vee r \equiv p \vee (q \vee r)$ $(p \wedge q) \wedge r \equiv p \wedge (q \wedge r)$	Associative laws

Table: Logical Equivalences

Logical Equivalences

Equivalence	Name
$p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$ $p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$	Distributive laws
$\neg(p \wedge q) \equiv \neg p \vee \neg q$ $\neg(p \vee q) \equiv \neg p \wedge \neg q$	De Morgan's laws
$p \vee (p \wedge q) \equiv p$ $p \wedge (p \vee q) \equiv p$	Absorption laws
$p \vee \neg p \equiv T$ $p \wedge \neg p \equiv F$	Negation laws

Table: Logical Equivalences

Logical Equivalences

- ① Using De Morgan's Laws show that $\neg(p \rightarrow q)$ is logically equivalent to $p \wedge \neg q$.
- ② Show that $\neg(p \vee (\neg p \wedge q))$ and $\neg p \wedge \neg q$ are logically equivalent by developing a series of logical equivalences.
- ③ Show that $(p \wedge q) \rightarrow (p \vee q)$ is a tautology.

Logical Equivalences

1

$$\begin{aligned}\neg(p \rightarrow q) &\equiv \neg(\neg p \vee q) && \text{by definition of implication} \\ &\equiv \neg(\neg p) \wedge \neg q && \text{by the second De Morgan's law} \\ &\equiv p \wedge \neg q && \text{by the double negation law}\end{aligned}$$

2

$$\begin{aligned}\neg(p \vee (\neg p \wedge q)) &\equiv \neg p \wedge \neg(\neg p \wedge q) && \text{by the second De Morgan's law} \\ &\equiv \neg p \wedge [\neg(\neg p) \vee \neg q] && \text{by the first De Morgan's law} \\ &\equiv \neg p \wedge (p \vee \neg q) && \text{by the double negation law} \\ &\equiv (\neg p \wedge p) \vee (\neg p \wedge \neg q) && \text{by the second distributive law} \\ &\equiv F \vee (\neg p \wedge \neg q) && \text{because } \neg p \wedge p \equiv F \\ &\equiv (\neg p \wedge \neg q) \vee F && \text{by the commutative law for disjunction} \\ &\equiv \neg p \wedge \neg q && \text{by the identity law for } F\end{aligned}$$

Predicates and Quantifiers

Consider the statement

" x is greater than 3"

It consists of two parts:

- ① The variable x , representing the subject,
- ② And the predicate "is greater than 3".

We denote this statement as $P(x)$, where

- ① P represents the predicate "is greater than 3,"
- ② x is the variable.

$P(x)$ is also known as the value of the propositional function P at x . Only when a specific value is assigned to x , $P(x)$ becomes a proposition with a definite truth value.

Predicates

Example: Let $Q(x)$ denote the statement " x is even." What are the truth values of $Q(4)$ and $Q(3)$?

Solution: For $Q(x)$ denoting the statement " x is even," the truth values are:

$Q(4)$: True, because 4 is an even number.

$Q(3)$: False, because 3 is not an even number.

Predicates

Example: Let $A(x)$ denote the statement

"Computer x is under attack by an intruder."

Suppose that of the computers on campus, only CS2 and MATH1 are currently under attack by intruders. What are the truth values of $A(\text{CS1})$, $A(\text{CS2})$, and $A(\text{MATH1})$?

Quantifiers

Universal Quantifier

The universal quantification of $P(x)$ is the statement

“ $P(x)$ for all values of x in the domain.”

The notation $\forall xP(x)$ denotes the universal quantification of $P(x)$. Here \forall is called the **universal quantifier**. We read $\forall xP(x)$ as “for all $xP(x)$ ” or “for every $xP(x)$.”

Quantifiers

Example: Let $C(x)$ denote “ x is a cat,” and let $A(x)$ denote “ x is an animal.” The statement “All cats are animals” can be written as:

$$\forall x(C(x) \rightarrow A(x))$$

Explanation:

- $\forall x$: This is the universal quantifier, which means “for all x .”
- $C(x) \rightarrow A(x)$: This is a conditional statement that reads as “if x is a cat, then x is an animal.”

Quantifiers

Example: Let $Q(x)$ be the statement " $x < 2$." The truth value of the quantification $\forall x Q(x)$, where the domain consists of all real numbers, is as follows:

Since $Q(x)$ is not true for every real number x

$Q(3)$ is false,

$x = 3$ serves as a counterexample for the statement $\forall x Q(x)$.

Hence, $\forall x Q(x)$ is false.

Quantifiers

Existential Quantifier

The existential quantification of $P(x)$ is the proposition

“There exists an element x in the domain such that $P(x)$.”

We use the notation $\exists x P(x)$ for the existential quantification of $P(x)$. Here, \exists is called the **existential quantifier**.

Quantifiers

Example: Let $R(x)$ be the statement "Device x has more than 8GB of RAM." What is the truth value of the quantification $\exists x R(x)$, where the domain consists of all the computing devices in a laboratory?

Quantifiers

Solution: Let $R(x)$ denote the statement "Device x has more than 8GB of RAM." The truth value of the quantification $\exists x R(x)$, where the domain consists of all computing devices in a lab, is as follows: If there exists at least one device in the lab that has more than 8GB of RAM (e.g., a 16GB workstation), then the statement $\exists x R(x)$ is **true**.

Observe that $\exists x R(x)$ is **false** only if *no* device in the lab has more than 8GB of RAM. That is, it is false if every device in the lab has RAM less than or equal to 8GB.

Exercise: Let $S(x)$ denote the statement "Device x never runs out of battery." What is the truth value of $\forall x S(x)$, where the domain consists of all mobile devices in the building?

Quantifiers

Statement	When True?	When False?
$\forall x P(x)$	$P(x)$ is true for every x .	There is an x for which $P(x)$ is false.
$\exists x P(x)$	There is an x for which $P(x)$ is true.	$P(x)$ is false for every x .

Table: Quantifiers

Negating Quantified Expressions

Consider the statement:

"Every IoT device in the building is connected to the secure network."

This is a universal quantification, represented as:

$$\forall x C(x), \text{ where } C(x)$$

is the statement:

" x is connected to the secure network,"

and the domain consists of all IoT devices in the building.

The negation of this statement is:

"It is not the case that every IoT device in the building is connected to the secure network,"

Negating Quantified Expressions

which is equivalent to:

"There is at least one IoT device in the building that is not connected to the secure network."

This is the existential quantification of the negation:

$$\exists x \neg C(x)$$

This example illustrates the equivalence:

$$\neg \forall x C(x) \equiv \exists x \neg C(x)$$

Negating Quantified Expressions

Quantifier	Equ. Statement	When is Negation True?	When is Negation False?
$\neg \forall x P(x)$	$\exists x \neg P(x)$	There exists an x for which $P(x)$ is false	Every x satisfies $P(x)$
$\neg \exists x P(x)$	$\forall x \neg P(x)$	Every x for which $P(x)$ is false	There exists an x for which $P(x)$ is true

Table: De Morgan's Laws for Quantifiers

Example: What are the negations of the statements “Every software update fixes at least one known bug.” ?

Negating Quantified Expressions

Solution:

Let $U(x)$ denote "software update x fixes at least one known bug."

Original statement: $\forall x U(x)$

Negation: $\neg(\forall x U(x)) \equiv \exists x \neg U(x)$

English interpretation of the negation:

There exists at least one software update that does not fix any known bugs.

Negating Quantified Expressions

Exercise: What is the negation of the statement:

$$\forall x \exists y (T(x, y) \rightarrow E(x, y))$$

where $T(x, y)$ denotes “device x transmits packet y ” and $E(x, y)$ denotes “packet y is encrypted by x ”.

Hint: Use logical equivalences to express the negation in simplified form.

Solution: Negating Quantified Expressions

Original: $\forall x \exists y (T(x, y) \rightarrow E(x, y))$

Negation:

$$\begin{aligned}\neg \forall x \exists y (T(x, y) \rightarrow E(x, y)) &\equiv \exists x \forall y \neg(T(x, y) \rightarrow E(x, y)) \\ &\equiv \exists x \forall y (T(x, y) \wedge \neg E(x, y))\end{aligned}$$

Interpretation: There is a device that transmits every packet without encrypting any of them.