

# differential privacy

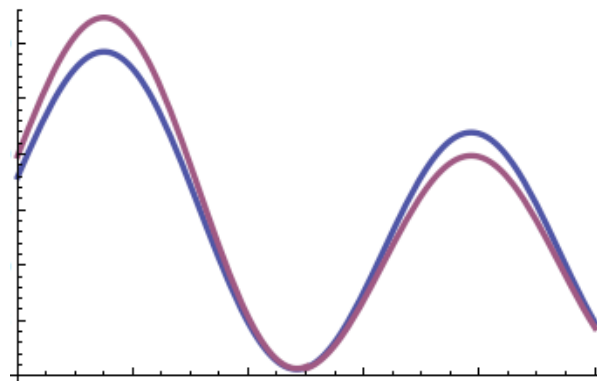
[DinurNissim03, DworkNissimMcSherrySmith06, Dwork06]

$\epsilon$ -Differential Privacy for algorithm  $M$ :

for any two neighboring data sets  $x_1, x_2$ , differing by the addition or removal of a single row

any  $S \subseteq \text{range}(M)$ ,

$$\Pr[M(x_1) \in S] \leq e^\epsilon \Pr[M(x_2) \in S]$$



# sensitivity of a function $f$

$$\Delta f = \max_{x_1, x_2} |f(x_1) - f(x_2)|_1$$

for neighboring data sets  $x_1, x_2$

measures how much one person can affect output

sensitivity is  $1/|x|$  for queries returning the average value of count queries mapping  $X$  to  $\{0,1\}$

linear queries :  $X \rightarrow [0,1]$  over the dataset (think statistical queries)

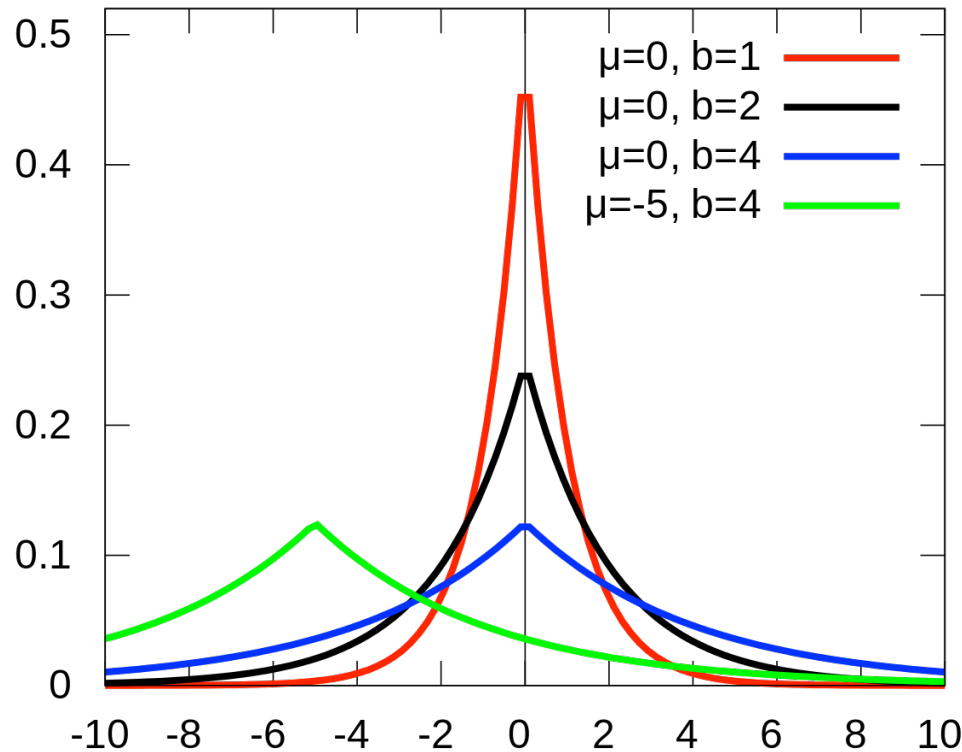
# scale noise with sensitivity

[DworkNissimMcSherrySmith06]

$$\Delta f = \max_{x_1, x_2} |f(x_1) - f(x_2)|_1$$

for neighboring data sets  $x_1, x_2$

# Laplace distribution $\text{Lap}(b)$



$$\text{pdf}(z) = \exp(-|z|/b) / 2b$$
$$\text{variance} = 2b^2$$

For  $Y \sim \text{Lap}(b)$ ,  $\Pr[|Y| \geq bt] = \exp(-t)$

# Laplace mechanism

Def. Given  $f : \mathbb{N}^{|X|} \rightarrow \mathbb{R}^k$  the Laplace Mechanism is defined as

$$M_L(x, f(\cdot), \epsilon) = f(x) + (Y_1, \dots, Y_k)$$

where the  $Y_i$  are iid random draws from  $\text{Lap}(b)$  with  $b = \Delta f / \epsilon$ .

(If we want discrete output space, subsequently round accordingly.)

# Laplace mechanism: Privacy

Thm. The Laplace Mechanism preserves  $(\epsilon, 0)$ -differential privacy.

## Laplace Mechanism: Privacy

Thm The Laplace Mechanism preserves  $(\epsilon, 0)$ -dp

pf Let  $x, x'$  be neighboring databases, so  $\|x - x'\|_1 \leq 1$

Let  $f: \mathcal{N}^{|x|} \rightarrow \mathbb{R}^k$

let  $p_x$  be prob. density function of  $M_L(x, f, \epsilon)$   
 $p_{x'}$  " " " " "  $M_L(x', f, \epsilon)$

Let  $z \in \mathbb{R}^k$

$$\begin{aligned} \frac{p_x(z)}{p_{x'}(z)} &= \prod_{i=1}^k \left[ \frac{\exp\left(-\frac{\epsilon |f(x)_i - z_i|}{\Delta f}\right)}{\exp\left(-\frac{\epsilon |f(x')_i - z_i|}{\Delta f}\right)} \right] = \prod_{i=1}^k \exp\left[\frac{\epsilon |f(x')_i - z_i| - |f(x)_i - z_i|}{\Delta f}\right] \\ &\leq \prod_{i=1}^k \exp\left(\frac{\epsilon |f(x)_i - f(x')_i|}{\Delta f}\right) = \exp\left(\frac{\epsilon \|f(x) - f(x')\|_1}{\Delta f}\right) \\ &\leq \exp(\epsilon) \end{aligned}$$

# Laplace mechanism: Accuracy

Thm. The Laplace Mechanism preserves...



## Laplace Mechanism - Accuracy

Thm Let  $f: \mathcal{N}^{|\mathcal{X}|} \rightarrow \mathbb{R}^K$ ,  $y = M_L(x, f, \varepsilon)$ . Then  $\forall \delta \in [0, 1]$ :

$$\Pr \left[ \|f(x) - y\|_\infty \geq \ln\left(\frac{K}{\delta}\right) \left(\frac{\Delta f}{\varepsilon}\right) \right] \leq \delta$$

PF

$$\begin{aligned} \Pr \left[ \|f(x) - y\|_\infty \geq \ln\left(\frac{K}{\delta}\right) \left(\frac{\Delta f}{\varepsilon}\right) \right] &= \Pr \left[ \max_{i \in [K]} |Y_i| \geq \ln\left(\frac{K}{\delta}\right) \left(\frac{\Delta f}{\varepsilon}\right) \right] \\ &\leq K \cdot \Pr \left[ |Y_i| \geq \ln\left(\frac{K}{\delta}\right) \left(\frac{\Delta f}{\varepsilon}\right) \right] \\ &= K \left( \frac{\delta}{K} \right) \\ &= \delta \end{aligned}$$

$\leftarrow \Pr[|Y| \geq tb] = \exp(-t)$   
 $b = \Delta f / \varepsilon$

## Notes

1. Could replace Laplacian by gaussian noise  
add noise scaled to  $N(0, \sigma^2)$ ,  $\sigma \sim \Delta f \ln(1/s) / \epsilon$   
gives  $(\epsilon, s)$ -dp
2. The simpler randomized response algorithm  
is local  
However overall its accuracy is worse.

# applying the Laplace mechanism

single counting query: how many people in the database satisfy predicate  $P$ ?

sensitivity = 1

can add noise  $\text{Lap}(1/\epsilon)$

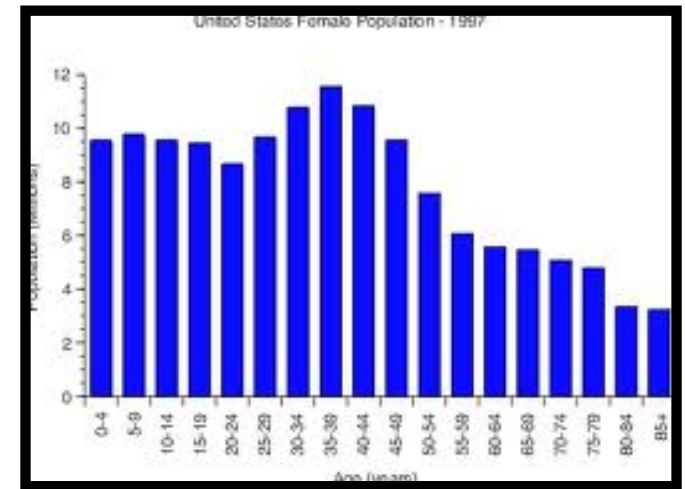
# applying the Laplace mechanism

vector-valued queries of dimension  $d$

Can apply composition and add noise  $\text{Lap}(d \Delta f / \epsilon)$  in each component of output, where  $\Delta f$  is sensitivity of each component.

(Could also split the noise unevenly. Could also treat the queries separately and apply composition.)

# applying the Laplace mechanism



histogram queries

could again use noise  $\text{Lap}(d/\epsilon)$

but actually only need  $\text{Lap}(1/\epsilon)$ , since  
sensitivity generalizes as  $\max \ell_1$  distance

What is the sensitivity of the query, “how many people in the database are both over age 50 AND smokers?”

# example



Suppose we wanted to determine the most commonly-“liked” Facebook page, subject to DP

could give a DP count of the number of likes for each page, but sensitivity would grow with the max number of “likes” a person could give (bad)

but we only want to know the max, not every count—could that be easier?

# reportNoisyMax

For  $m$  count queries add noise  $\text{Lap}(1/\epsilon)$  to each, and report the index of the largest noised query.

Claim: reportNoisyMax is  $(\epsilon, 0)$ -differentially private.



# reportNoisyMax

For  $m$  count queries add noise  $\text{Lap}(1/\epsilon)$  to each, and report the index of the largest noised query.

What about accuracy?