# System Programming

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# 1 MICROPROCESSOR

Masm: Microsoft Macro Assembler

To execute commands do

#### **Name Command Operand**

### 1.1 REGISTERS

#### 1.1.1 Data Registers

Four 32 bit data registers are EAX,EBX,ECX & EDX.

Four 16 bit data registers are there namely: AX, BX, CX, DX where

AX: 16 Bits is comprised of

АН	AL
8 bits	8 bits

Similarly, BX (16 bits) is comprised of

ВН	BL
8 bits	8 bits

Hence, in total

16 Bit	8 Bit Registers	Values	Remarks
AX	AH+AL	AH=AX/256	Accumulator;
		AL=AX%256	Ax is preffered
		AX=AH*256+AL	for Arithmetic
			and data
			transfer. I/O
BX	BH+BL	BH=BX/256	Base
		BL=BX%256	Serves as
			address for
			table lookup
CX	CH+CL	CH=CX/256	Count
		CL=CX%256	Looping
DX	DH+DL	DH=DX/256	Data
		DL=DX%256	Used in
			Mul/DIV. Used
			in I/O

#### 1.1.2 **Segment Registers**

**Four Segment Registers** 

Register name	Bits	Function
CS	16	Code Segments

		Stores address of Code
		segment
DS	16	Data Segments
SS	16	Stack Segments
ES	16	Extra Segment

### 1.1.3 Pointer And Index Register

All Registers are 16 Bits

Operand	Name	Remarks
SP	Stack Point	Used for accessing
		Stack Segment
BP	Base Pointer	Used for accessing
		data from stack
		segment
SI	Source Index	Used to point to
		memory location in
		the data segments
		addressed by DS
DI	Destination Index	Same as SI
IP	Instruction Pointer	Contains the offset
		address of the next
		instruction to be
		executed

# 1.1.4 Flag Register

16 Bit Register containing Various Flags like status and control flags

15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0
				OF	DF	IF	TF	SF	ZF		AF		PF		CF

Operand	Name	Remarks
CF	Carry Flag	CF=1, If there is any carry out
		from MSB on addition/sub
PF	Parity Flag	PF=1, IF the low byte of a
		result has an even no's of 1 bit,
		otherwise 0
AF	Auxiliary Flag	AF=1, If there is a carry out
		from bit on addition/
		subtractions otherwise AF=0
ZF	Zero Flag	ZF=1 for zero result otherwise
		ZF=0
SF	Sign Flag	SF=1 if MSB of result it 1
		Otherwise SF=0
OF	Overflow Flag	OF=1 if sign overflow occurs
		Otherwise OF=0
DF	Direction Flag	DF=0, SI & DI proceed in
		increasing memory address
		DF=1, decreasing

## 2 VARIABLES

#### Execution:

#### VAR\_NAME **COMMAND VALUES**

Command	Meaning	Bits
DB	Define Byte	8
DW	Define Word	16
DD	Define Double Word	32
DQ	Define QuadWord	64
DT	Define 10 bytes	80

Eg: ALPHA DB ? ;For Value Not initialized BY DB WRD DW 2;Initialized to 2 B\_arr DB 10H,20H,30H; Stored as B\_arr, B\_arr+1 & B\_arr+2 DW 10H,20H,30H; Stored as B\_arr, B\_arr+2 & B\_arr+4 B\_arr

> 'ABC'; A=string, B=String+1 ... string DB DB 'Hello',ODH,OAH,'\$'; H,E,L,L,O,ODH,OAH,\$

msq

### 2.1 CONSTANTS

Exectution:

VAR\_NAME **EQU VALUES** 

# 3 PROGRAM STRUCTURE

#### The Model Structure is as follows

```
.MODEL Small
.STACK 100H
.DATA
;Define Data
.CODE
somefn PROC
     ;Define fn
somefn ENDP
MAIN PROC
      ;Instructions
MAIN ENDP
END MAIN
```

### 3.1 .MODEL MEMORY\_MODEL

Memory_model	Remarks
SMALL	Code in 1 segment
	Data in 1 Segment
MEDIUM	Code in multiple segments
	Data in 1
COMPACT	Code in 1
	Data in multiple
LARGE	Code in multiple
	Data in multiple
	Array size max = 64kb
HUGE	Code in multiple
	Data in multiple
	Array size max >64kb

#### 3.2 .DATA

Variables are declared here, as previously discussed

### 3.3 .STACK

.STACK SIZE

EG .STACK 100H

#### 3.4 .CODE

For writing Code to solve any problem

## 4 PROCEDURES

Similar to the function in c++. Structure is as follows

```
Name PROC type
;Body of Procedure
RET
Name ENDP
;For calling
Call Name
```

#### Here,

- Type
  - Near: The statement that calls the procedure is in the same segment as the procedure. *Default* value.
  - o Far: Calling statement is in the different segment
- Return: Returns back to calling place.
- Call: Calls the function. It does 2 operations

- The return address of the calling proc is saved in top of the stack. This is the offset address of next instruction to be executed, after proc is completed.
- IP(instruction pointer) gets the offset address of the first instruction of the procedure.

### 5 ARRAYS

#### 5.1 1-D ARRAY

Eg are

- MSG DB 'HELLO\$'
- W DW 10,20,30,40,50
- GAMMA DW 100 DUP(?)
- DELTA DB 212 DUP(?)
- LINE DB 5,4,2 DUP(2,3DUP(0),1) ;{5,4,2,0,0,0,1,2,0,0,0,1}

#### 5.1.1 Addressing Modes

#### 5.1.1.1 Register Mode

Is used to call the reference. However, this is only restricted to few registers.

BX,SI,DI	DS
BP	SS

Eg:

- MOV AX, [SI] ; Moves the content addressed by SI to AX
- MOV BX,[BX] ;Content of address BX into BX

#### 5.1.1.2 Based and Indexed Addressing mode

```
[reg.+displacement]
[displacement+reg.]
[reg]+displacement
Displacement+[reg]
```

#### 5.1.2

Registered that can be used is same as previous.

Eg: Let W be a variable. All examples mean the same

- MOV AX, W[BX]
- MOV AX, [W+BX]
- MOV AX, W+[BX]
- MOV AX, [BX]+W

#### 5.1.2.1 PTR

type PTR address\_expression

Here Types are either BYTE or WORD.

It is used to override the declared type of expression-

Eg

DOLLAR DB 1AH
 CENT DB 52H

...

MOV AX, DOLLAR ;Wrong MOV AX, WORD PTR Dollar ;Right

### 5.1.2.2 Accessing elements in the stack

Eg

- MOV BP, SP
- MOV AX, [BP]
- MOV BX, [BP+2]
- MOV CX, [BP+4]

### 5.2 2-D ARRAY [A]<sub>MXN</sub>

Eg

• B DW 10,20,30,40

DW 50,60,70,80

DW 90,100,311,121 ;Row Major order A+[(i-1)N+(j-1)]S; S is the size (2 for W, 1 for B)

B DW 10,20

DW 50,60

DW 90,100 ;Column Major order A+[(i-1)+(j-1)M]S

#### 5.2.1 Addressing mode

#### 5.2.1.1 Base-Index mode

```
Variable[base_reg][index_reg]
```

What it does is gets info, after shifting the origin to Variable[Base\_reg] and adds offset of index\_reg. So A[4][3]==A[7]

### 6 STRING

Instruction	Command	Destination	Source
Move String	MOVSB	ES:DI	DS:SI
Compare String	CMPSB	ES:DI	DS:SI
Load String	LODSB	AL or AX	DS:SI
Scan String	SCASB	ES:DI	AL or AX

Store String STOSB ES:DI	AL or AX
--------------------------	----------

### 7 Macros

A procedure is called at execution time, control transfers to the procedure and returns after executing Its statements. A macro Is invoked at assembly time. The assembler copies the macro's statements into the program at the position of the invocation. When the program executes, there is no transfer of control.

```
Macro_name MACRO d1,d2 ... dn
;Statements

ENDM
;Bunch of shit
Macro_name d1,d2 ... dn
```

Where d1,d2 ... is the dummy list of arguments used by the macro

#### 7.1 LOCAL PSUEDO-OP

A macro with a loop or decision structure contains one or more labels. If such a macro Is Invoked more than once in a program, a duplicate label appears, resulting in an assembly error: This problem can be avoided by using local labels in the macro. To declare them, we use the LOCAL pseudo-op, whose syntax Is

```
LOCAL list_of__labels
```

#### 7.2 REPT MACRO

The REPT macro can be used to repeat a block of statements.

```
REPT expression
;Statements
ENDM
```

#### 7.3 Indefinite Repeat

#### 7.4 CONDITIONALS

```
CONDITIONAL
;Statements

ELSE
;Statements

ENDIF
```

#### 7.4.1 List of Conditional

Form	True if
IF exp	Constant expression is non-zero
IFE exp	Constant expression is zero
IFB <arg></arg>	Argument is missing
IFNB <arg></arg>	Argument is not missing
IFDEF <sym></sym>	Symbol sym is defined in the program (or as extern)
IFNDEF <sym></sym>	Symbol sym is not defined in the program (or as extern)
IFIDN <string1> <string2></string2></string1>	String1 and string2 are identical. Brackets are require.
IFDEF <string1>, <string2></string2></string1>	Are different
IF1	Assembler is making the first assembly pass
IF2	Assembler is making the second assembly
	pass

### 8 MEMORY MANAGEMENT

.COM: has one segment only. ORG 100h is used for stack no need for @data lines.

#### 8.1 Program Modules

#### 8.1.1 Near v/s Far

A procedure is NEAR if the statement that calls it is in the same segment as the procedure itself; a procedure is FAR if it is called from a different segment.

FAR procedure Is In a different segment from Its calling statement, the CALL instruction causes first CS and then IP to be saved on the stack, then CS:IP gets the segment offset of the procedure. To return, RET pops the stack twice to restore the original CS:IP.

#### 8.1.2 Extern

When assembling a module, the assembler must be Informed of names which are used In the module but are defined In other modules, otherwise these names will be flagged as undeclared.

EXTERN external\_name\_list

Here, external\_name\_list is a list of arguments of the form name:type where name is an external name, and type is one of the following: NEAR, FAR, WORD, BYTE, or DWORD.

#### 8.1.3 Public

A procedure or variable must be declared with the PUBUC pseudo-op if it is to be used in a different module

```
PUBLIC external_name_list
```

#### 8.2 Full Segment Definations

With the full segment definitions, the programmer can control how segments arc ordered, combined with each other, and aligned relative to each other In memory

#### 8.2.1 Segment Directive

```
name SEGMENT align combine class
;Statements
name ENDS
```

#### 8.2.1.1 Align Type

The align type of a segment declaration determines how the starting address of the segment is selected when the program is loaded in memory.

#### Table 14.1 Align Types

PARA	Segment begins at the next available paragraph (least significant hex digit of physical address is 0).
BYTE	Segment begins at the next available byte.
WORD	Segment begins at the next available word (least significant bit of physical address is 0).
PAGE	Segment begins at the next available page (two least significant hex digits of physical address are 0).

PARA is the default align type.

#### 8.2.1.2 Combine type

If a program contains segments of the same type, the combine type tells how they are to be combined when the program is loaded in memory.

#### Table 14.2 Combine Types

```
PUBLIC

Segments with the same name are concatenated (placed one after the other) to form a single, continuous memory block

COMIMON

Segments with the same name begin at the same place in memory: that is, are overlaid.

STACK

Has the same effect as PUBLIC, except that all offset addresses of instructions and data, in the segment are relative to the SS register. SP is initialized to the end of the segment.

AT paragraph

Indicates that the segment should begin at the specified paragraph.
```

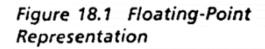
#### 8.2.1.3 Class Type

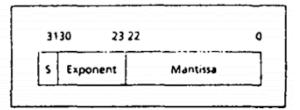
The *class* type of a segment declaration determines the order in which segments are loaded in memory. Class type declarations must be enclosed in single quotes. If two or more segments have

the same class, they are loaded in memory one after the other. If classes are not specified in segment declarations segments are loaded in the order, they appear in the source listing.

### 9 FLOATING POINT REPRESENTATION

In the floating-point representation, each number is represented in two parts: a mantissa, which contains the leading significant bits in a number, and an exponent, which is used to adjust the position of the binary point. For example, the number 2.5 in binary is 10.1b, and its floating-point representation has a mantissa of 1.1 and an exponent of 1. This is because 10.1b can he written as  $1.01 \times 21$ . Negative exponents are not represented as signed numbers. Instead, a number called the bais is added to the exponent to create a positive number. For example, if we use eight bits for the exponent, then the number  $(2^7 - 1)$  or 127 is chosen as the bias. To represent the number 0.0001b, we have a mantissa of 1.0 and an exponent of -4. After adding the bias of 127, we get 123 or 0111110 I b. It starts with a sign bit, followed by an 8-bit exponent, and a 23-bit mantissa. For short real figure 18.1 is represented





#### 9.1 8087 STACK

The 8087 has eight 80-bit data registers, and they function as a stack. Data may be pushed or popped from the stack. The top of the stack is addressed as ST or ST(0). The register directly beneath the top is addressed as ST(1). In general, the i<sup>th</sup> register in the stack is addressed as ST(i), where I must be a constant. Data types in 8087 are

Data		Precision	Most Significant Byte														
Formats	Range		7.	07	0	7	07	07	_ 0	7	07	07		7	0	7	0
Word integer	104	16 Bits	115		l <sub>o</sub>	Two	's Co	mplé	ment								
Short integer	109	32 Bits	J <sub>31</sub>					I <sub>O</sub>	īwo's	Con	npien	nent		•			
Long integer	1018	64 Bits	163										<u></u>	] c	omp	Two	
Packed BCD	1018	18 Digits	S	D1,	D16									_		D <sub>1</sub> D	2]
Short Real	10±38	24 Bits	SE,		Eo F,			F <sub>23</sub> F	o Imp	licit							
Long Real	104308	53 Bits	S E 10		Εo	F,				_			F <sub>62</sub>	F <sub>o</sub>	Imp	licit	
Ten:porary Real	10=4532	64 Bits	S E 14		E	o Fo										F <sub>6</sub>	3
Integer: I	WD 1	``															
Packed BCD: (-1) Kcal· (-1)^(2E-8im	(Fa'F1	.) .															
bias = 127 for Sh	ort Real																
. 1023 for L	ong Rea Temp Re							•								•	

Figure 18.2 8087 Data Types

### 10 Intel instruction Format

TYPE	BYTES REQUIRED	EXAMPLE
PREFIX	0-4 Bytes	Deals With Instruction behaviour Operand Size Override Address Size Override
O/P CODE	1 Byte	Not: 0011011w
MOD R/M	1 Byte	Modregreg Instructions like ADD ADDregreg
SIB	1 Byte	Scale Index Base  Scale * Index + Base  1 (5)  Scale   index   base
DISPLACEMENT	1 or 2 Bytes	Initial address
IMMEDIATE	1 or 2 Bytes	Numerical input MOV AX, <u>25</u>

### 10.1 Data Storage Buffer

Data Structure

1. Location counter (LC) - toinds to the next location when
code mill be blaced.

2. Op-code translation table - this table is pere-defined. Contains
the code for isach instruction

3. Symbol table

4. Others Stonege Buffer (SCB) - Contains ascer characters for the
string.

5. Forward Reference Table (FDT) - this contains pointers to the string
in SCB and offset ader when its
value mill be inserted in the object ad

### 10.2 Types of Record

ТҮРЕ	PREFIX	USE
HEADER	Н	Header type record, kind of stores the name of the file
REFERENCE	R	External Reference storage

DAT TYPE	D	Is used when the address has to be changed relative to initial address
MEMORY TYPE	M	A variable or the procedure which has been used in the code but memory is unknown

# 11 IMPORTANT COMMANDS

Command Name	Syntax	Eg	Remarks
ASCII Adjusted Addition	AAA	MOV AL,6 ADD AL,7 AAA	Used to check for carry in BCD values. Values is stored in AL. It does AH=AL/10 AL=AL%10
ASCII Adjusted Multiplication	AAD	AAD	AL=AL+10*AH AH=0
ASCII Adjusted Multiplication	AAM	AAM	Same as above
ASCII Adjusted Subtraction	AAS	MOV AL,6 SUB AL,7 AAS	Used to check for carry in BCD values. Values is stored in AL. It does AH=AL/10 AL=AL%10
Add	ADD dest, src/val	ADD AH,2 ADD AH,BX ADD Word1,AX	Dest=dest+src/val
Add with Carry	ADC dest, src/val	ADC AX,1 ADC AX,BX	Adds the carry bit along with normal ops
And	AND dest, src/val	AND AL,7FH {Clears sign bit of AL}	Can be used to clear specific Bits while reserving others. 7FH stands for DEL. Changes the destination.
Convert Byte to word	CBW	CBW	Converts byte to word. Used in DIV/IDIV
Clear Direction Flag	CLD	CLD	Clears the direction flag. If 0 then pointer to the data is incremented
Compare	CMP dest, src	CMP AX,BX CMP AX, '\$' CMP AX, 5	Sets flag CF=0 if not equal, CF=1 if equal
Compare String byte Compare String Word	CMPSB CMPSW		subtracts the byte with address ES:DI

			from the byte with
			address DS:SI, and
			sets the flags. The
			result is not stored.
Convert Word to	CMD	CWD	Converts word to
double word	CWD	CWD	double word. Used for
double word			IDIV
Decrement	DEC dest	DEC BX	Dest -=1
Decrement	DEC dest	DEC BX	Dest1
Division	DIV divisor	DIV BX	For <b>DIV</b> DX=0 (AH=0)
DIVISION	IDIV divisor	DIV BL	For <b>IDIV</b> , DX(AH)
	וטוע מועוגטו	DIV DE	should be signed
			extension. Use CWD.
			CATCHISION: OSC CVVD.
			Byte: Divisor is 8-bit
			register.
			AL = AX/divisor
			AH = AX%divisor
			Word Form: Divisor is
			16 bit. Dividend is 32
			bit DX:AX
			AX = DX:AX/div
			DX = DX:AX%Div
DUP	DUP(values)	DUP(?)	Fills random value
		DUP(42)	Fills 42
Load onto 8087 stack	FLD (load real)	FLD source	Where Source is a
	FILD (int load)	FILD source	memory location
	FBLD (packed BCD	FBLD source	
	load)		
Store 8087 into	FST (store real)	FST destination	Where destination is a
destination	FIST (store int)	FIST destination	memory location
	FSTP (Store real pop)	FSTP destination	
	FISTP (store int pop)	FISTP destination	
	FBSTP (store BCDpop)	FBSTP destination	
Arithmetic on 8087	FADD [[dest,] src]	Add	Each opcode can take zero,
stack	FSUB [dest, src]	Subtract	one, or two operands. An
	FMUL [dest, src]	Mul	instruction with no
		Divide	operands assumes ST(0) as the source and ST(1) as the
	FDIY [dest, src]		destination; the instruction
	FIADD source	Integer add	also pops the stack. In an
	FISUB source	Integer subtract	instruction with one
	FIMUL source	Integer mul	operand, the operand specifies a memory
	FIDIV source	Int div	location as the source; the
			destination. is assumed to
Increment	INC dest	INC BX	be ST(0).  Dest +=1
IIICI CIIICIIL	inc dest	INC AX	De3( +-1
Interrupt	INT 21h	Int 21h	Does when:
	211		AH=1: gets ASCII Code
			of key pressed
	<u> </u>		-:, p

Jump if greater	JG/JNLE	JG K3	Input: AL=ASCII of Key pressed AH=2: prints single char output Input: DL=ascii value of char to be printed AH=9: String output till it encounters '\$' Input: DX=offset address of the string Condition for jump:
(signed)  Jump if greater than	JGE/JNL	JGE K4	ZF=0 & SF=OF SF=OF
equal	/		C5 O5
Jump if less	JL/JNGE	JL 15	SF<>OF
Jump if less than & equal	JLE/JNG	JLE L4	ZF=1 & SF <> OF
Jump if greater (unsigned) i.e Jump if Above	JA/JNBE	JA K4	Condition for jump: CF=0 & ZF=0
Jump if above or equal	JAE/JNB	JAE	Condition for jump: CF=0
Jump if below	JB/JNAE	JB	Condition for jump: CF=1
Jump if below of equal/not above	JBE/JNA	JBE K3	Condition for jump: CF=1 or ZF=1
Jump if equal (Single flag Jump)	JE/JZ	JE	Condition for Jump: ZF=1
Jump if not equal	JNE/JNZ		ZF=0
Jump if carry Jump if not carry	JC JNC		CF=1 CF=0
Jump if overflow Jump if not overflow	JO JNO		OF=1 OF=0
Jump if signed Jump if not signed	JS JNS		SF=1 SF=0
Jump if even parity Jump if odd parity	JP/JPE JNP/JPO		PF=1 PF=0
Unconditional Jump	JMP	JMP k4	
Load Effective Address	LEA dest,src	LEA DX, MSG LEA SI, AX	Loads address of source to Destination
Load String	LODSB LODSW		moves the byte addressed by DS:SI into AL. SI is then incremented if DF = 0 or decremented if DF.= 1.
Move	MOV dest, src	MOV AH, WORD1 MOV AX,BX MOV AH, 'A'	Both destination and source cannot be memory variable.  MOV Word1, Word2

			Won't work.
Move string Byte	MOVSB REP MOVSB	MOVSB REP MOVSB	Moves a single byte from DS:SI to ES:DI. Repeats the command CX number of bytes
Move String Word	MOVSW REP MOVSW	MOVSW REP MOVSW	SI, DI increase by 2. CX number of words.
Multiplication	MUL source IMUL source	MUL BX MUL BL	Unsigned Signed In byte form: AX=source*AL Word Form: DX:AX=source*AX
Negation	NEG destination	NEG AX NEG AH	To negate destination
NOT	NOT Destination	NOT K	Changes bit from 0 to 1 and 0 from 1. {If K =1101b NOT K=0010b}
OR	OR	OR AL, 81H {Sets the msb and lsb of AL}	OR is used to set specific bit while reserving the others. Changes the source
Pop	POP destination	POP AX POP BX	Destination must be 16 bit. Does 2 operations 1) Pop - >destination 2) SP+=2
PopF	POPF	POPF	Content of top of the stack will be moved to flag register
Push	PUSH source	PUSH AX PUSH BX	Source must be 16 bit. It does 2 operations 1) SP-=2 2) SS:SP
PushF	PUSHF	PUSHF	Pushes the content of flag register on top of the stack
Rotate and Carry Left	RCL		Shifts CF into higher order word
Repeat until not equal to Repeat until equal to	REPNE SCASB REPE SCASB REPE CMPSB REPNE CMPSB	REPNE SCASB	will repeatedly subtract each string byte from AL, update DI, and decrement CX until there is a zero result (the target is found) or CX = 0 (the string ends).
Roll Left	ROL Dest, bits_shifted	Same as SHL	The ROL instruction shifts each bit to the

			1 6 11 1 1 1 1
			left, with the highest
			bit copied in the Carry
			flag and also into the
			lowest bit
Roll Right	ROR Dest, bits_shifted	Same as SHL	The ROL instruction
			shifts each bit to the
			Right, with the lowest
			bit copied in the Carry
			flag and also into the
			highest bit
Shift Arithmetic Left	SAL Dest, bits_shifted	SAL reg,CL	Identical to SHL.
Subtract with borrow	SBB Dest, Src/val	SBB AX, BX	Subtracts the carry bit
		SBB AX, 13	too along with normal
		,	subtraction
Scan Byte	SCASB	SCASB	SCASB subtracts the
Scan Word	SCASW		string byte pointed to
			by ES:DI from the
			contents
			of AL and uses the
			result to set the flags.
			ZF=0 ;If diff
			ZF=1; if same char
Shift Arithmetic Right	SHR Dest,Bits_shifted	SAR reg,CL	The lowest bit is
Shirt Arthinetic Night	STIN Dest, Dits_stillted	Same as SHI	copied to carry flag.
		Jame as Sm	And signed bit is
			copied to right
Shift Left	SHL Dest, bits_shifted	SHL reg, CL	The highest bit is
Silit Left	SHE Dest, bits_silited	SHL mem, CL	moved to carry flag.
		SHL reg, imm8	Fills lowest bit with 0.
		SHL mem, imm8	MSB->CF.
Shift Right	SHR Dest, bits_shifted	SHR reg, CL	The lowest bit is
Silit Right	SHK Dest, bits_stillted	SHR mem, CL	moved to carry flag.
		SHR reg, imm8	Fills highest bit with 0.
		_	MSB->CF.
Sot Direction Floa	STD	SHR mem, imm8 STD	Sets the direction flag.
Set Direction Flag	טוט	טוט	Pointer is
Stora String	CTOCD		decremented  Moves the content of
Store String	STOSB		Moves the content of
	STOSW		AL to ES:DI, and
			inc/dec the pointer
			according to
Subtract	CLID doct and had	CLID VIT 3	byte/word.
Subtract	SUB dest, src/val	SUB AH,2	Dest=dest-src/val
		SUB AH,BX	
Tast	TECT D	SUB Word1,AX	Composition to
Test	TEST Dest, src	TEST AH, 01H	Same as and but
			without changing the
	VOLIC 1	VOLIC ALL DI	source.
Exchange	XCHG dest, src	XCHG AH,BL	Both destination and
		XCHG AX,WORD1	source cannot be
		XCHG AX,BX	memory variable.

			XCHG Word1,Word2 Won't work.
XLAT	XLAT	MOV AL, 5 LEA BX, T1 XLAT ;This will do AL=[T1+5]	Low operand instruction which is used to convert byte value into another value that comes from byte table. It adds the content of AL to the address in BX and retrieves the value of the address to AL
XOR	XOR		XOR is used to Compliment specific bits while reserving the others

# 12 LINKER LOADER

## 12.1 Prog A

Code	Memory Address	Command	Values at Pass 1	Value at Pass 2	Value after Linking
PROG A: START 0					
EXTRN SI,BU,SU					
LDA #128	0000	29	01 28	01 28	01 28
LDA SI	0003	02	Na na	Na na	30 00
LDA #1	0006	29	00 01	00 01	00 01
LDX #0	0009	05	00 00	00 00	00 00
L1: STA BU,X	000C	OF	Na na	Na na	30 06
ADD #1	000F	19	00 01	00 01	00 01
TIZ SI	0012	2C	Na na	Na na	30 00
JEQ L2	0015	30	Na Na	00 1B	10 1B
J l1	0018	3C	00 OC	00 OC	10 0C
L2: JSUB SU	001B	48	Na Na	Na Na	20 00
RSUB	001E	4E	00 00	00 00	00 00
END					

### 12.1.1 Tables for Prog A

•				
Type	Name	Address		
Н	PROG A			
R	SI			
R	BU			
R	SU			

D		000C
D		0015
M	SI	0004
M	BU	000D
M	SI	0013
M	SU	001B

# 12.2 PROG B

Code	Memory Address	Command	Values at Pass 1	Value at Pass 2	Value after Linking
PROG B: START					
0					
PUBLIC SU					
EXTRN					
SI,BU,TO					
SU: LDA #0	0000	29	00 00	00 00	00 00
LDX #0	0003	05	00 00	00 00	00 00
L3: ADD BU,X	0006	1B	Na na	Na na	30 06
TIX SI	0009	2C	Na na	Na na	30 00
TEQ L4	000C	30	Na na	00 12	20 12
TE L3	000F	30	00 06	00 06	20 06
L4: STA TO	0012	0C	Na na	Na na	30 03
RSUB	0015	4F	00 00	00 00	00 00
END					

## 12.2.1 Tables for Prob B

Type	Name	Address
Н	PROG B	
R	SI	
R	BU	
R	ТО	
D		000D
D		0010
D	SU	0000
M	BU	0007
M	SI	000A
M	ТО	0013

## 12.3 PROG C

Code	Memory	Command	Values at Pass	Value at	Value after
	Address		1	Pass 2	Linking
PROG C: START					
0					
PUBLIC SI, BU,					
то					
SI DB 1	0000	29	Na na	Na na	Na na

TO DB 1	0003	05	Na na	Na na	Na na
BU DB 50	0006	1B	Na na	Na na	Na na
END					

# 12.3.1 Tables for Prog C

Туре	Name	Address
Н	PROG C	
D	SI	0000
D	ТО	0003
D	BU	0006

# 12.4 LINKER TABLE

Program Name	Load Point	Label	Absolute Adress
PROG A	1000		
PROG B	2000		
		SU	2000
PROG C	3000		
		SI	3000
		ТО	3003
		BU	3006