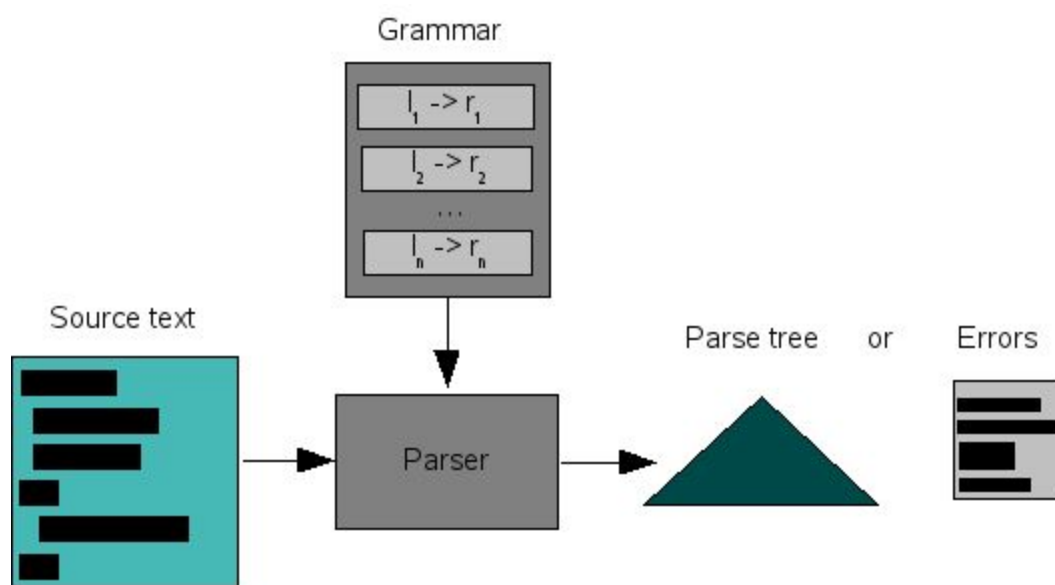


Parser

Phase #2



Submitted to : Ms. Meghana Pujar

Submitted by : Abhijith Anilkumar (13CO102)

Chirag Jamadagni (13CO117)

George CM (13CO119)

Table of Contents

❑ Abstract	3
❑ Introduction	4
❑ Syntactical Analyzer	4
❑ Yacc Script	5
❑ Design	8
❑ Code	8
❑ Test Cases	21
❑ Implementation	24
❑ Extensive Test Cases	25
❑ Results/Future work	30
❑ References	30

List of Tables

● Test Cases	21
--------------------	----

List of Figures

● Flowchart of a parser	5
● Syntactical Analyzer with Yacc	7
● Parser Code - 1	11
● Parser Code - 2	12
● Parser Code - 3	12
● Parser Code - 4	13
● Parser Code - 5	13
● Parser Code - 6	14
● Parser Code - 7	14
● Parser Code - 8	15
● Parser Code - 9	15
● Parser Code - 10	16
● Parser Code - 11	16
● Parser Code - 12	17
● Parser Code - 13	17
● Symbol Table - 1	18
● Symbol Table - 2	19
● Symbol Table - 3	19
● Run file	20
● Example - 1 Input	25
● Example - 1 Output	26
● Example - 1 Symbol table	26
● Example - 1 Constant table	27
● Example - 2 Input	27
● Example - 2 Output	28

Abstract

Compiler design for a language involves two phases - the analysis phase, and the synthesis phase. The aim of the second phase of the course project of building a compiler is to build a parser, using a tool called Yacc. The lexical analyzer generated in the first phase reads the source program and generates tokens that are given as input to the parser which then creates a syntax tree in accordance with the grammar, consequently leading to the generation of intermediate code that is fed into the synthesis phase, to obtain the correct, equivalent machine level code. Various test cases are also demonstrated, along with code snippets and output snapshots, in order to demonstrate the functioning of the parser.

Introduction

Syntactical Analyzer

Also known as the parser. Syntactic analysis is the second phase of a compiler. Parsing or syntactic analysis is the process of analysing a string of symbols, in computer languages, conforming to the rules of a formal grammar. It takes the tokens generated by the lexical analyser and a context free grammar and gives a parse tree corresponding to the grammar, with the tokens inputted, as the output. It basically gives a structural representation of the input, checking for correct syntax in the process.

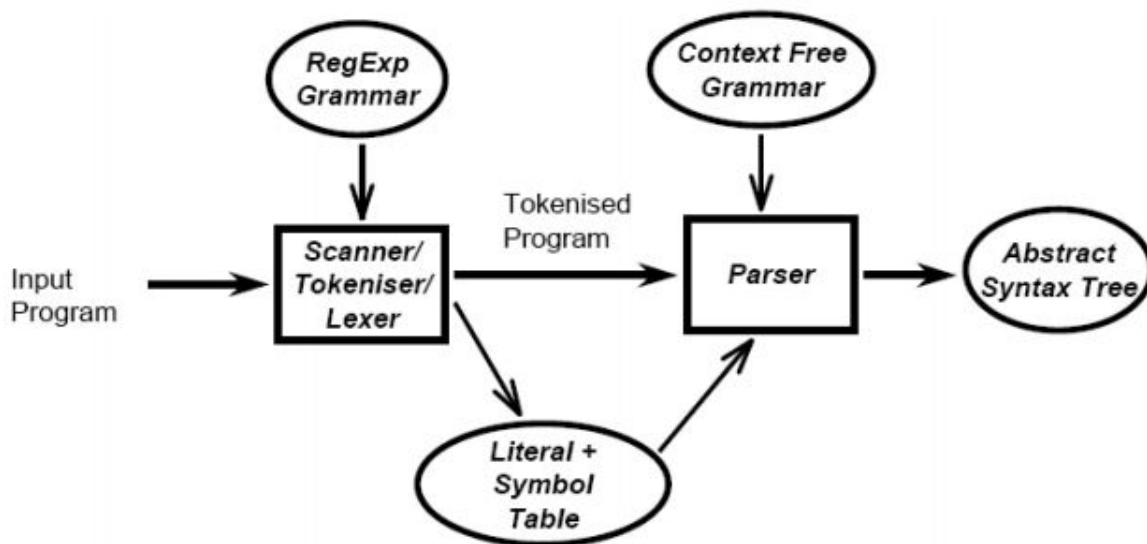


Fig 1: Flowchart of a parser

Yacc Script

Yacc(Yet Another Compiler-Compiler) provides a general tool for describing the input to a computer program. The Yacc user specifies the structures of his input, together with code to be invoked as each such structure is recognized. Yacc turns such a specification into a subroutine that handles the input process; frequently, it is convenient and appropriate to have most of the flow of control in the user's application handled by this subroutine. The input subroutine produced by Yacc calls a user-supplied routine to return the next basic input item. Thus, the user can specify his input in terms of individual input characters or in terms of higher level constructs such as names and numbers. The user-supplied routine may also handle idiomatic features such as comment and continuation conventions, which typically defy easy grammatical specification.

Yacc provides a general tool for imposing structure on the input to a computer program. The Yacc user prepares a specification of the input process; this includes rules describing the input structure, code to be invoked when these rules are recognized, and a low-level routine to do the basic input. Yacc then generates a function to control the input process. This function, called a parser, calls the user-supplied low-level input routine (the lexical analyzer) to pick up the basic items (called tokens) from the input stream. These tokens are organized according to the input structure rules, called grammar rules; when one of these rules has been recognized, then user code supplied for this rule, an action, is invoked; actions have the ability to return values and make use of the values of other actions. Yacc is written in a portable dialect of C[1] and the actions, and output subroutine, are in C as well. Moreover, many of the syntactic conventions of Yacc follow C.

Yacc requires token names to be declared as such. In addition, it is often desirable to include the lexical analyzer as part of the specification file; it may be useful to include other programs as well. Thus, every specification file consists of three sections:

declarations
%%
rules
%%
programs

The sections are separated by double percent “%%” marks. The percent “%” symbol is generally used in Yacc specifications as an escape character.

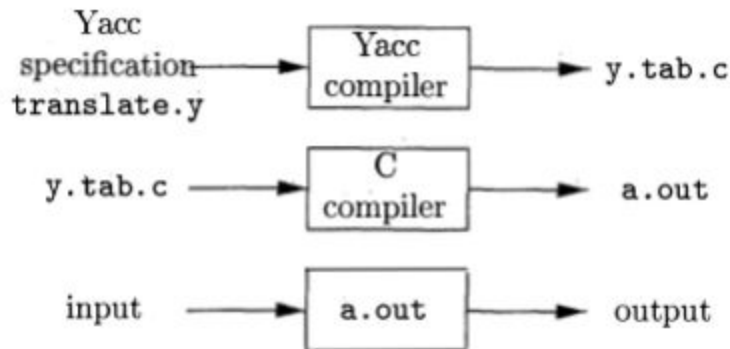


Fig 2: Creating a syntactical analyzer with Yacc

Design

Code

Lex can be used in two ways. The first one is where it is used for simple transformations, or for analysis and statistics gathering on a lexical level. It can also be used with a parser generator to perform the lexical and syntactical analysis phase, because it is particularly easy to interface Lex and Yacc. Lex programs recognize only regular expressions; Yacc writes parsers that accept a large class of context free grammars, but it requires a low level analyzer to recognize input tokens. When the tasks are divided into lexical and syntactical phases, Yacc is used to assign structure.

The rules section contains multiple grammar rules that specify the syntactical constructs of the source language. We have associated different actions to different rules, like returning values or obtaining values returned from previous actions.

In order to get the tokens from the lexical analyzer, the integer valued function `yylex()` must be specified. This function returns an integer which is the token number representing the kind of token that is read. If this token is associated with a value, it should be assigned to the external variable `yylval`.

The following are the key points in order to link the Lex file with the Yacc file:

- The token names in the Lex file should match with those of the Yacc file.
- In order that the Yacc file receives each of the tokens from the Lex file, we must also see to it that every token generated by the lexical analyzer is returned to the Yacc file, for it to check its rules and parse.
- The main function of the Lex file must also be modified, and a `yyparse()` function must be called from it, in order to return all the generated tokens to the Yacc file.

The following are the cases that arise while running the program:

- If the source program is free of syntactical errors, then the tokens generated by the lexical analyzer are generated and displayed.
 - In this case, along with the generated tokens, a table called the symbol table, is also generated, updated, and displayed which contains the token type and the token name.
 - Along with this, a table called the constant table, containing the various constant variables is also displayed.

- On the other hand, if the source program does contain syntax errors, then the tokens are properly generated up to the position of the error, after which an error message, generated by `yyerror()` appears, stating the line number at which the error has occurred.
 - The symbol table and constant table are not displayed in this case.

- *Errors:* For displaying the line numbers for the errors, we used a global variable `lineno` initialized to 1 and which keeps incrementing whenever a newline is scanned.

- *Symbol table and Constant table:* We have created a header file, "symbol.h" where there are functions to insert the lexemes into the symbol tables. When we see the <EOF> we print both the tables.

- *Shift-reduce conflicts:* There is only one shift-reduce conflict present in our grammar and no reduce-reduce conflicts at all. This conflict is due to the classic "dangling else" problem. The problem can be fixed by assigning precedence to the rules that causes conflict. The rule causing conflict is:

```
selection_statement : IF '(' expression ')' statement
                    | IF '(' expression ')' statement ELSE
statement
```

In order to fix this, we start by making ELSE and LOWER_THAN_ELSE (a pseudo-token) non associative:

```
%nonassoc LOWER_THAN_ELSE
%nonassoc ELSE
```

This gives ELSE more precedence over LOWER_THAN_ELSE simply because LOWER_THAN_ELSE is declared first.

Then in the conflicting rule you have to assign a precedence to either the shift or reduce action:

```
selection_stmt : IF '(' expression ')' statement %prec
LOWER_THAN_ELSE ; | IF '(' expression ')' statement ELSE statement ;
```

Here, higher precedence is given to shifting.

Parser.y

This is the main parser code. This is a yacc file which contains the declarations, rules and programs and defines the actions which should be taken for various cases.

```
1 %{
2 #include <stdio.h>
3 #include <string.h>
4 #include <stdlib.h>
5 #include "symbol.h"
6 #include "y.tab.h"
7 int errorFlag=0;
8 extern FILE *yyin;
9 extern FILE *yyout;
10 extern int line,cnt;
11 extern symrec *sym_table;
12 extern symrec *const_table;
13 extern char *tempid;
14 %}
15 %token IDENTIFIER CONSTANT STRING_LITERAL SIZEOF
16 %token PTR_OP INC_OP DEC_OP LEFT_OP RIGHT_OP LE_OP GE_OP EQ_OP NE_OP
17 %token AND_OP OR_OP MUL_ASSIGN DIV_ASSIGN MOD_ASSIGN ADD_ASSIGN
18 %token SUB_ASSIGN LEFT_ASSIGN RIGHT_ASSIGN AND_ASSIGN
19 %token XOR_ASSIGN OR_ASSIGN TYPE_NAME
20
21 %token TYPEDEF EXTERN STATIC AUTO REGISTER
22 %token CHAR SHORT INT LONG SIGNED UNSIGNED FLOAT DOUBLE CONST VOLATILE VOID
23 %token STRUCT UNION ENUM ELLIPSIS
24
25 %token CASE DEFAULT IF ELSE SWITCH WHILE DO FOR GOTO CONTINUE BREAK RETURN
26
27 %nonassoc LOWER_THAN_ELSE
28 %nonassoc ELSE
29
30 %start translation_unit
31 %%
32
33 primary_expression
34 : IDENTIFIER { putsym(tempid, 'v', line); }
35 | CONSTANT { putsym(tempid, 'c', line); }
36 | STRING_LITERAL { putsym(tempid, 'c', line); }
37 | '(' expression ')'
```

1,120 Top

Figure 3

```

38 ;
39
40 postfix_expression
41 : primary_expression
42 | postfix_expression '[' expression ']'
43 | postfix_expression '(' argument_expression_list ')'
44 | postfix_expression '.' IDENTIFIER { putsym(tempid, 'v', line); }
45 | postfix_expression PTR_OP IDENTIFIER { putsym(tempid, 'v', line); }
46 | postfix_expression INC_OP { putsym(tempid, 'o', line); }
47 | postfix_expression DEC_OP { putsym(tempid, 'o', line); }
48 ;
49
50
51 argument_expression_list
52 : assignment_expression
53 | argument_expression_list ',' assignment_expression { putsym(" ", 'p', line); }
54 ;
55
56 unary_expression
57 : postfix_expression
58 | INC_OP unary_expression { putsym("++", 'o', line); }
59 | DEC_OP unary_expression { putsym("--", 'o', line); }
60 | unary_operator cast_expression
61 | SIZEOF unary_expression { putsym("sizeof", 'o', line); }
62 | SIZEOF '(' type_name ')' { putsym("sizeof", 'o', line); }
63 ;
64
65 unary_operator
66 : '&' { putsym("&", 'o', line); }
67 | '*' { putsym("*", 'o', line); }
68 | '+' { putsym("+", 'o', line); }
69 | '-' { putsym("-", 'o', line); }
70 | '~' { putsym("~", 'o', line); }
71 | '!' { putsym("!", 'o', line); }
72 ;
73
74 cast_expression

```

39,104

8%

Figure 4

```

75 : unary_expression
76 | '(' type_name ')' cast_expression
77 ;
78
79 multiplicative_expression
80 : cast_expression
81 | multiplicative_expression '*' cast_expression { putsym("*", 'o', line); }
82 | multiplicative_expression '/' cast_expression { putsym("/", 'o', line); }
83 | multiplicative_expression '%' cast_expression { putsym("%", 'o', line); }
84 ;
85
86 additive_expression
87 : multiplicative_expression
88 | additive_expression '+' multiplicative_expression { putsym("+", 'o', line); }
89 | additive_expression '-' multiplicative_expression { putsym("-", 'o', line); }
90 ;
91
92 shift_expression
93 : additive_expression
94 | shift_expression LEFT_OP additive_expression { putsym("<<", 'o', line); }
95 | shift_expression RIGHT_OP additive_expression { putsym(">>", 'o', line); }
96 ;
97
98 relational_expression
99 : shift_expression
100 | relational_expression '<' shift_expression
101 | relational_expression '>' shift_expression
102 | relational_expression LE_OP shift_expression { putsym("<=", 'o', line); }
103 | relational_expression GE_OP shift_expression { putsym(">=", 'o', line); }
104 ;
105
106 equality_expression
107 : relational_expression
108 | equality_expression EQ_OP relational_expression { putsym("==", 'o', line); }
109 | equality_expression NE_OP relational_expression { putsym("!=", 'o', line); }
110 ;
111

```

111,116

17%

Figure 5

```

112 and_expression
113 : equality_expression
114 | and_expression ' & ' equality_expression { putsym("&", 'o', line); }
115 ;
116 exclusive_or_expression
117 : and_expression
118 | exclusive_or_expression '^' and_expression { putsym("^", 'o', line); }
119 ;
120 inclusive_or_expression
121 : exclusive_or_expression
122 | inclusive_or_expression '|' exclusive_or_expression { putsym("|", 'o', line); }
123 ;
124 logical_and_expression
125 : inclusive_or_expression
126 | logical_and_expression AND_OP inclusive_or_expression { putsym("&&", 'o', line); }
127 ;
128 logical_or_expression
129 : logical_and_expression
130 | logical_or_expression OR_OP logical_and_expression { putsym("||", 'o', line); }
131 ;
132 conditional_expression
133 : logical_or_expression
134 | logical_or_expression '?' expression ':' conditional_expression { putsym("?:", 'o', line); }
135 ;
136 assignment_expression
137 : conditional_expression
138 | unary_expression assignment_operator assignment_expression
139 ;
140 assignment_operator
141 : '=' { putsym("=", 'o', line); }

```

148,113-116 25%

Figure 6

```

149 | MUL_ASSIGN { putsym("*= ", 'o', line); }
150 | DIV_ASSIGN { putsym("/= ", 'o', line); }
151 | MOD_ASSIGN { putsym("%= ", 'o', line); }
152 | ADD_ASSIGN { putsym("+= ", 'o', line); }
153 | SUB_ASSIGN { putsym("-= ", 'o', line); }
154 | LEFT_ASSIGN { putsym("<=<=", 'o', line); }
155 | RIGHT_ASSIGN { putsym(">=>=", 'o', line); }
156 | AND_ASSIGN { putsym("&= ", 'o', line); }
157 | XOR_ASSIGN { putsym("^= ", 'o', line); }
158 | OR_ASSIGN { putsym("|= ", 'o', line); }
159 ;
160 expression
161 : assignment_expression
162 | expression ',' assignment_expression { putsym(",", 'p', line); }
163 ;
164 constant_expression
165 : conditional_expression
166 ;
167 declaration
168 : declaration_specifiers ';' { putsym(";", 'p', line); }
169 | declaration_specifiers init_declarator_list ';' { putsym(";", 'p', line); }
170 ;
171 declaration_specifiers
172 : storage_class_specifier
173 | storage_class_specifier declaration_specifiers
174 | type_specifier
175 | type_specifier declaration_specifiers
176 | type_qualifier
177 | type_qualifier declaration_specifiers
178 ;
179 init_declarator_list
180 : init_declarator

```

185,113-116 34%

Figure 7

```

186 | init_declarator_list ',' init_declarator { putsym(" ", 'p', line); }
187 ;
188
189 init_declarator
190 : declarator
191 | declarator '=' initializer { putsym("=", 'o', line); }
192 ;
193
194 storage_class_specifier
195 : TYPEDEF { putsym("typedef", 'k', line); }
196 | EXTERN { putsym("extern", 'k', line); }
197 | STATIC { putsym("static", 'k', line); }
198 | AUTO { putsym("auto", 'k', line); }
199 | REGISTER { putsym("register", 'k', line); }
200 ;
201
202 type_specifier
203 : VOID { putsym("void", 'k', line); }
204 | CHAR { putsym("char", 'k', line); }
205 | SHORT { putsym("short", 'k', line); }
206 | INT { putsym("int", 'k', line); }
207 | LONG { putsym("long", 'k', line); }
208 | FLOAT { putsym("float", 'k', line); }
209 | DOUBLE { putsym("double", 'k', line); }
210 | SIGNED { putsym("signed", 'k', line); }
211 | UNSIGNED { putsym("unsigned", 'k', line); }
212 | struct_or_union_specifier
213 | enum_specifier
214 | TYPE_NAME
215 ;
216
217 struct_or_union_specifier
218 : struct_or_union IDENTIFIER '{' struct_declaration_list '}'
219 | struct_or_union '{' struct_declaration_list '}'
220 | struct_or_union IDENTIFIER
221 ;
222

```

222,124 42%

Figure 8

```

223 struct_or_union
224 : STRUCT { putsym("struct", 'k', line); }
225 | UNION { putsym("union", 'k', line); }
226 ;
227
228 struct_declaration_list
229 : struct_declaration
230 | struct_declaration_list struct_declaration
231 ;
232
233 struct_declaration
234 : specifier_qualifier_list struct_declarator_list ';' { putsym(";", 'p', line); }
235 ;
236
237 specifier_qualifier_list
238 : type_specifier specifier_qualifier_list
239 | type_specifier
240 | type_qualifier specifier_qualifier_list
241 | type_qualifier
242 ;
243
244 struct_declarator_list
245 : struct_declarator
246 | struct_declarator_list ',' struct_declarator { putsym(" ", 'p', line); }
247 ;
248
249 struct_declarator
250 : declarator
251 | ':' constant_expression { putsym(":", 'p', line); }
252 | declarator ':' constant_expression { putsym(":", 'p', line); }
253 ;
254
255 enum_specifier
256 : ENUM '{' enumerator_list '}' { putsym("enum", 'k', line); }
257 | ENUM IDENTIFIER '{' enumerator_list '}' { putsym("enum", 'k', line); putsym(tempid, 'v', line); }
258 | ENUM IDENTIFIER { putsym("enum", 'k', line); putsym(tempid, 'v', line); }
259 ;

```

259,113-116 51%

Figure 9


```

260
261 enumerator_list
262 : enumerator
263 | enumerator_list ',' enumerator { putsym(" ", 'p', line); }
264 ;
265
266 enumerator
267 : IDENTIFIER { putsym(tempid, 'v', line); }
268 | IDENTIFIER '=' constant_expression { putsym("=", 'o', line); putsym(tempid, 'v', line); }
269 ;
270
271 type_qualifier
272 : CONST { putsym("const", 'k', line); }
273 | VOLATILE { putsym("volatile", 'k', line); }
274 ;
275
276 declarator
277 : pointer direct_declarator
278 | direct_declarator
279 ;
280
281 direct_declarator
282 : IDENTIFIER { putsym(tempid, 'v', line); }
283 | '(' declarator ')'
284 | direct_declarator '[' constant_expression ']'
285 | direct_declarator '[' ']'
286 | direct_declarator '(' parameter_type_list ')'
287 | direct_declarator '(' identifier_list ')'
288 | direct_declarator '(' ')'
289 ;
290
291 pointer
292 : '*' { putsym(" ", 'o', line); }
293 | '*' type_qualifier_list { putsym(" ", 'o', line); }
294 | '*' pointer { putsym(" ", 'o', line); }
295 | '*' type_qualifier_list pointer { putsym(" ", 'o', line); }
296 ;

```

296,125-128 59%

Figure 10

```

297
298 type_qualifier_list
299 : type_qualifier
300 | type_qualifier_list type_qualifier
301 ;
302
303
304 parameter_type_list
305 : parameter_list
306 | parameter_list ',' ELLIPSIS { putsym(" ", 'p', line); putsym("::", 'o', line); }
307 ;
308
309 parameter_list
310 : parameter_declaration
311 | parameter_list ',' parameter_declaration { putsym(" ", 'p', line); }
312 ;
313
314 parameter_declaration
315 : declaration_specifiers declarator
316 | declaration_specifiers abstract_declarator
317 | declaration_specifiers
318 ;
319
320 identifier_list
321 : IDENTIFIER { putsym(tempid, 'v', line); }
322 | identifier_list ',' IDENTIFIER { putsym(tempid, 'v', line); putsym(" ", 'p', line); }
323 ;
324
325 type_name
326 : specifier_qualifier_list
327 | specifier_qualifier_list abstract_declarator
328 ;
329
330 abstract_declarator
331 : pointer
332 | direct_abstract_declarator
333 | pointer direct_abstract_declarator

```

333,121-124 68%

Figure 11

```

334 ;
335
336 direct_abstract_declarator
337 : '(' abstract_declarator ')'
338 | '[' ']'
339 | '[' constant_expression ']'
340 | direct_abstract_declarator '[' ']'
341 | direct_abstract_declarator '[' constant_expression ']'
342 | '(' ')'
343 | '(' parameter_type_list ')'
344 | direct_abstract_declarator '(' ')'
345 | direct_abstract_declarator '(' parameter_type_list ')'
346 ;
347
348 initializer
349 : assignment_expression
350 | '[' initializer_list ']'
351 | '[' initializer_list ',' ']'
352 ;
353
354 initializer_list
355 : initializer
356 | initializer_list ',' initializer { putsym(":", 'p', line); }
357 ;
358
359 statement
360 : labeled_statement
361 | compound_statement
362 | expression_statement
363 | selection_statement
364 | iteration_statement
365 | jump_statement
366 ;
367
368 labeled_statement
369 : IDENTIFIER ':' statement { putsym(tempid, 'v', line); }
370 | CASE constant_expression ':' statement { putsym(":", 'p', line); putsym("case", 'k', line); }

```

370,117-120 76%

Figure 12

```

371 | DEFAULT ':' statement { putsym(":", 'p', line); putsym("default", 'k', line); }
372 ;
373
374 compound_statement
375 : '{' '}'
376 | '{' statement_list '}'
377 | '{' declaration_list '}'
378 | '{' declaration_list statement_list '}'
379 ;
380
381 declaration_list
382 : declaration
383 | declaration_list declaration
384 ;
385
386 statement_list
387 : statement
388 | statement_list statement
389 ;
390
391 expression_statement
392 : ';' { putsym(";", 'p', line); }
393 | expression ';' { putsym(";", 'p', line); }
394 ;
395
396 selection_statement
397 : IF '(' expression ')' statement %prec LOWER_THAN_ELSE { putsym("if", 'k', line); }
398 | IF '(' expression ')' statement ELSE statement { putsym("if", 'k', line); putsym("else", 'k', line); }
399 | SWITCH '(' expression ')' statement
400 ;
401
402 iteration_statement
403 : WHILE '(' expression ')' statement { putsym("while", 'k', line); }
404 | DO statement WHILE '(' expression ')' ';'
405 | FOR '(' expression_statement expression_statement ')' statement { putsym("for", 'k', line); }
406 | FOR '(' expression_statement expression_statement expression ')' statement { putsym("for", 'k', line); }
407 ;

```

407,117-120 85%

Figure 13

Symbol.h

Contains functions to input a token into the symbol or constant table and output in on the standard output. Basically, this file contains the hash implementation of the symbol and constant table.

```
1 #include <stdlib.h>
2 #define t_void 1
3 #define t_char 2
4 #define t_int 3
5 #define t_float 4
6 struct symrec
7 {
8     char *name,type[20],line[100];
9     struct symrec *next;
10 };
11 typedef struct symrec symrec;
12 symrec *sym_table = (symrec *)0;
13 symrec *const_table = (symrec *)0;
14
15 void *putsym(char *sym_name,char sym_type, int linec)
16 {
17     symrec *ptr,*newptr;
18     char line[39],linec[19];
19     snprintf(linec, 19, "%d", linec);
20     strcpy(line,"");
21     strcat(line,linec);
22     char type[20];
23     switch(sym_type)
24     {
25         case 'c':
26             strcpy(type,"Constant");
27             break;
28         case 'v':
29             strcpy(type,"Identifier");
30             break;
31         case 'p':
32             strcpy(type,"Punctuator");
33             break;
34         case 'o':
35             strcpy(type,"Operator");
36             break;
37         case 'k':
```

6,128 [Top](#)

Figure 16

```

38         strcpy(type,"Keyword");
39         break;
40     case 's':
41         strcpy(type,"String Literal");
42         break;
43     case 'd':
44         strcpy(type,"Preprocessor Statement");
45         break;
46     }
47     if(sym_type == 'c')
48     {
49         for(newptr=const_table;newptr!=(symrec *)0;newptr=(symrec *)newptr->next)
50             if(strcmp(newptr->name,sym_name)==0)
51             {
52                 strcat(newptr->line,line);
53                 return;
54             }
55         ptr=(symrec *)malloc(sizeof(symrec));
56         ptr->name=(char *)malloc(strlen(sym_name)+1);
57         strcpy(ptr->name,sym_name);
58         strcpy(ptr->type,type);
59         strcpy(ptr->line,line);
60         ptr->next=(struct symrec *)const_table;
61         const_table=ptr;
62     }
63     else
64     {
65         for(newptr=sym_table;newptr!=(symrec *)0;newptr=(symrec *)newptr->next)
66             if(strcmp(newptr->name,sym_name)==0)
67             {
68                 strcat(newptr->line,line);
69                 return;
70             }
71         ptr=(symrec *)malloc(sizeof(symrec));
72         ptr->name=(char *)malloc(strlen(sym_name)+1);
73         strcpy(ptr->name,sym_name);
74         strcpy(ptr->type,type);

```

74,110-116 74%

Figure 17

```

51     {
52         strcat(newptr->line,line);
53         return;
54     }
55     ptr=(symrec *)malloc(sizeof(symrec));
56     ptr->name=(char *)malloc(strlen(sym_name)+1);
57     strcpy(ptr->name,sym_name);
58     strcpy(ptr->type,type);
59     strcpy(ptr->line,line);
60     ptr->next=(struct symrec *)const_table;
61     const_table=ptr;
62 }
63 else
64 {
65     for(newptr=sym_table;newptr!=(symrec *)0;newptr=(symrec *)newptr->next)
66         if(strcmp(newptr->name,sym_name)==0)
67         {
68             strcat(newptr->line,line);
69             return;
70         }
71         ptr=(symrec *)malloc(sizeof(symrec));
72         ptr->name=(char *)malloc(strlen(sym_name)+1);
73         strcpy(ptr->name,sym_name);
74         strcpy(ptr->type,type);
75         strcpy(ptr->line,line);
76         ptr->next=(struct symrec *)sym_table;
77         sym_table=ptr;
78     }
79 }
80 void *getsym(char *sym_name)
81 {
82     symrec *ptr;
83     for(ptr=sym_table;ptr!=(symrec *)0;ptr=(symrec *)ptr->next)
84         if(strcmp(ptr->name,sym_name)==0)
85             return ptr;
86     return 0;
87 }

```

87,112 Bot

Figure 18

Run

This is a shell script which automates the compilation and execution of the code.

7 lines (6 sloc) 117 Bytes	
1	#!/bin/sh
2	yacc -d parser.y
3	lex scanner.l
4	gcc lex.yy.c y.tab.c -w -g
5	./a.out \$1
6	rm -rf lex.yy.c y.tab.c y.tab.h a.out

Figure 19

Test Cases

Serial No.	Test Case	Expected Output	Status
1	<pre>#include<stdio.h> void main() { int a,,b; }</pre>	line no.4:syntax error	Working
2	<pre>#include<stdio.h> void main() { int a,c,b c=a+b; }</pre>	line no.4:syntax error	Working
3	<pre>#include<stdio.h> void main() { int a,b,c; if(a>b); a = b-c; else a = b+c; }</pre>	line no.8:syntax error	Working
4	<pre>#include<stdio.h> void main() { int a,b,c; for(i=0;i++) a = a+b; }</pre>	line no.6:syntax error	Working

5	<pre>#include<stdio.h> void main() { int a,b,c; for(i=0;i<b;i++) { for(j=0;j<c;j--) { a=a+b; } } }</pre>	line no.12:syntax error	Working
6	<pre>#include<stdio.h> void main() { int a,b,c; for(i=0;i<b;i++) { a=b*; } }</pre>	line no.8:syntax error	Working
7	<pre>#include<stdio.h> void main() { int a,b,c[10; //syntax error }</pre>	ine no.4:syntax error	Working

8	<pre>#include<stdio.h> void main() { int a[10],i,j; while(i<10) { scanf("%d",&a[i]); } for(i=0;i<10;i++) { for(j=0;j<9-i;j++) { if(a[j]>a[j+1]) { swap(&a[j],&a[j+1]); } } } void swap(int *a,int *b) { *a=*a+*b; *b=*a-*b; *a=*a-*b; }</pre>	Prints the symbol and constants table	Working
---	--	---	---------

Implementation

The implementation is as follows:

- Firstly, a file is created with a “.y” extension, which represents a yacc file. The file created here is “*parser.y*”.
- The file, like all yacc files, contains three main sections: Declarations, Rules, and Programs.
- All the definitions used throughout the program are defined under this section, including the header file that is created for the implementation of the symbol table, i.e., “*symbol.h*”.
- The Rules section contains the various grammatical rules that are to be followed while parsing.
- The Programs section is where the *yyparse()* and the *yyerror()* functions are called, in order to generate the error message at the line where the syntax error has occurred.
- The header file “*symbol.h*” contains functions to insert a token into the symbol table and also to output it on stdout. “*symbol.h*” essentially consists of the hash implementation of the symbol and constant tables.
- Once the file is made, it is run using the shell script by giving the following command: *./run*

Extensive Test Cases

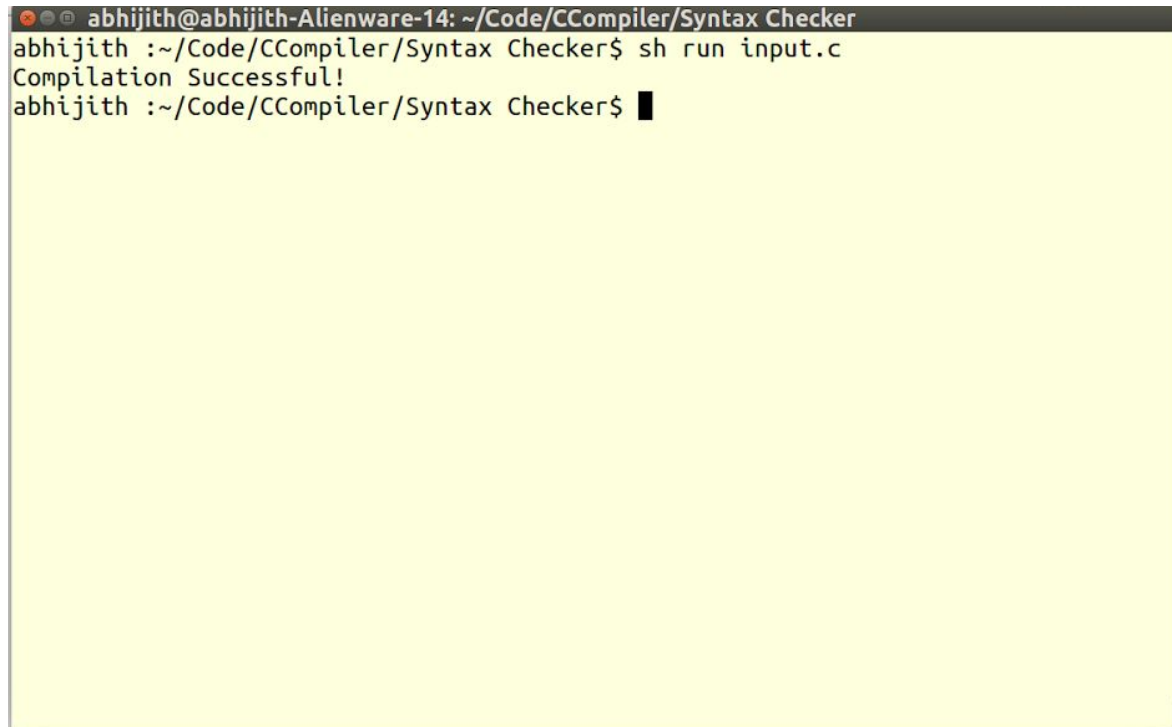
Case 1 : Positive test case with no errors

Input

```
1  #include<stdio.h>
2
3  void main()
4  {
5      int a[10],i,j;
6
7      while(i<10)
8      {
9          scanf("%d",&a[i]);
10     }
11
12
13     for(i=0;i<10;i++)
14     {
15         for(j=0;j<9-i;j++)
16         {
17             if(a[j]>a[j+1])
18             {
19                 swap(&a[j],&a[j+1]);
20             }
21         }
22     }
23 }
24
25 void swap(int *a,int *b)
26 {
27     *a=*a+*b;
28     *b=*a-*b;
29     *a=*a-*b;
30 }
```

Figure 20

Output

A terminal window with a dark title bar and a light yellow background. The title bar contains the text 'abhijith@abhijith-Alienware-14: ~/Code/CCompiler/Syntax Checker'. The terminal shows the following commands and output:

```
abhijith :~/Code/CCompiler/Syntax Checker$ sh run input.c
Compilation Successful!
abhijith :~/Code/CCompiler/Syntax Checker$
```

Figure 21

Symbol Table		
Name	Token Class	Line Number
}	Punctuator	22
return	Keyword	21
if	Keyword	21
/	Operator	17
*	Operator	17 17
=	Operator	17
&	Operator	8 10 16
scanf	Identifier	8 10 16
printf	Identifier	7 9 15 18 20
;	Punctuator	6 7 8 9 10 15 16 17 18 20 21
]	Punctuator	6
[Punctuator	6
num	Identifier	6
ee	Identifier	6
si	Identifier	6 17 18
time	Identifier	6 16 17
,	Punctuator	6 6 6 6 6 8 10 16 18
rate	Identifier	6 10 17
amount	Identifier	6 8 17
{	Punctuator	5
}	Punctuator	5 7 8 9 10 15 16 17 18 20 21 21
(Punctuator	5 7 8 9 10 15 16 17 18 20 21 21
main	Identifier	5
int	Keyword	5 6

Figure 22

Constant Table	
Value	Line Number
0	21
"Hello"	20
1	19
"\nSimple Interest : %d"	18
"\nEnter Period of Time : "	15
"\nEnter Rate of Interest : "	9
"%d"	8 10 16
"\nEnter Principal Amount : "	7
100	6 17

Figure 23

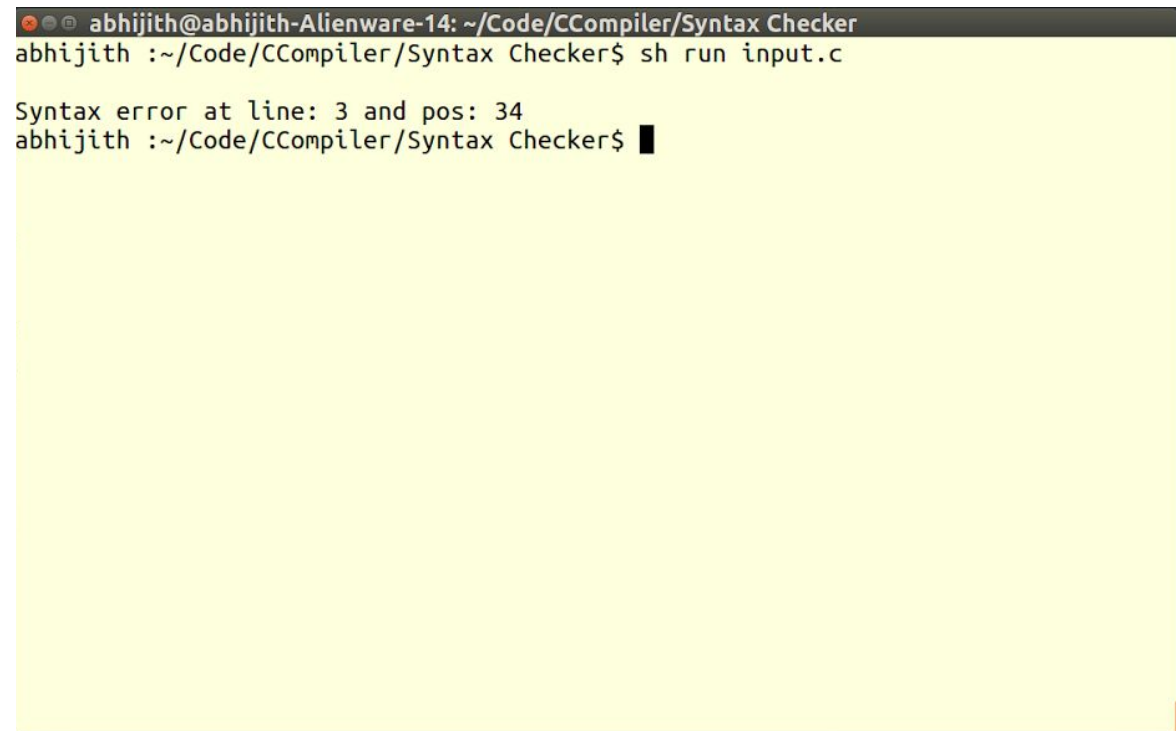
Case 2 : Negative test case with syntax errors

Input

```
1  #include<stdio.h>
2  /* This is a multi
3  line comment*/
4  int main() {
5      int amount, rate, time, si, 56ee;
6      //This is a single line comment
7      printf("\nEnter Principal Amount : ");
8      scanf("%d", &amount);
9
10     printf("\nEnter Rate of Interest : ");
11     scanf("%d", &rate);
12
13     printf("\nEnter Period of Time : ");
14     scanf("%d", &time);
15     scanf("",);
16     si = (amount * rate * time / 100;
17     printf("\nSimple Interest : %d", si);
18     printf("",);
19     return(0);
20 }
21 }
22 "I am d
```

Figure 24

Output

A terminal window with a dark title bar showing the user 'abhijith' on a machine named 'abhijith-Alienware-14' in the directory '~/Code/CCompiler/Syntax Checker'. The user has executed the command 'sh run input.c'. The output shows a syntax error at line 3, position 34. The prompt returns to the shell.

```
abhijith@abhijith-Alienware-14: ~/Code/CCompiler/Syntax Checker
abhijith :~/Code/CCompiler/Syntax Checker$ sh run input.c

Syntax error at line: 3 and pos: 34
abhijith :~/Code/CCompiler/Syntax Checker$
```

Figure 25

Results/ Future Work

The outcome of the second phase of this project is a syntactical analyzer that takes in the tokens generated by the lexical analyzer built in the first phase as input, and displays syntax error on the corresponding line number, and also symbol table and constant table.

Our next aim would be to implement a add a semantic checker to the existing compiler, and build the symbol table to be used during checking and code generation, in order to work our way towards making a full-fledged compiler.

References

1. <http://www.tldp.org/HOWTO/Lex-YACC-HOWTO-6.html>
2. <http://easywaysnow.blogspot.in/2012/10/running-lex-and-yacc-program-inubuntu.html?m=1>
3. <https://www.lysator.liu.se/c/ANSI-C-grammar-y.html>
4. <http://dinosaur.compilertools.net/yacc/>