Mandatory Exercise - Backwards-chaining

Anders Wind - awis@itu.dk

Task 1

Backwards-chaining is goal driven and can have a better time complexity than forward-chaining. This is due to the fact that Backwards chaining only "searches" the relevant part of the knowledge base, but forward-chaining has to search all different implications of the given facts. Therefore situations where we have a lot of facts that do not have any impact on the query, backwards chaining will be advantageous.

Task 2

I assume that the KB is in CNF.

I note that only one of the symbols in a clause's premise needs to be false for us to say that the clause is false, but only one of the clauses for a propositional symbol needs to be infered for the propositional symbol to be entailed.

```
function PL-BC-Entails? (KB, q) returns true or false
                 inputs: KB, the knowledge base, a set of propositional definite clauses
                                                       q, the query, a propositional symbol
                 Value-Map ← Check-Clause(KB, q, Map.Empty)
                  if q is in Value-Map
                                     return Value-Map[q]
                 return false
function Check-Clause (KB, q, Value-Map) returns map
                 inputs: KB, the knowledge base, a set of propositional definite clauses
                                                        \mathbf{q}\text{,} the query, a propositional symbol
                                                       Value-Map, a map of propositional symbol keys mapping to boolean values, empty to begin with
                  if q is in Value-Map
                                     return Value-Map
                 Value-Map.add(q, false)
                  for each clause c in KB which implies q
                                    if c is fact
                                                       Value-Map[q] ← true
                                                       return Value-Map
                                    else
                                                        result ← true
                                                       Value-Map ← Update-Map(KB, c.PREMISE, Value-Map)
                                                       for each (key, value) in Value-Map where key is in c.PREMISE
                                                                           if !value
                                                                                              result ← false
                                                                                             break
                                                       if result
                                                                           Value-Map[q] ← true
                                                                          return Value-Map
                  return Value-Map
function Update-Map(KB, symbols, Value-Map) returns map
                  inputs: KB, the knowledge base, a set of propositional definite clauses
                                                         symbols, a set of propositional symbols
                                                       Value-Map, a map of propositional symbol keys mapping to boolean value
                  for each Propositional-Symbol ps in symbols % \left( 1\right) =\left( 1\right) \left( 1\right) +\left( 1\right) \left( 1\right) \left( 1\right) +\left( 1\right) \left( 1\right) 
                                     Value-Map ← Check-Clause(KB, ps, Value-Map)
                  return Value-Map
```

Task 3

Yes but it returns false due to a cyclic case. My algorithm handles cyclic cases by checking if a Symbol has already been already exists in the Value-Map.

This is achieved by adding (q, false) to the Value-Map before recoursively checking the Premise of a clause. Then if q is reached through again in one of the premise symbols(or through any number of recoursions) it will just return the value-map(where the value of the symbol is false) and not try to update again.

Task 4

I assume that it is meant, that time complexity the algorithm runs in based on the amount of propositional symbols. The algorithm runs in linear time since every symbol will only be visited once due to the generation of the value-map and if a symbol already has been entailed it will not be entailed again.