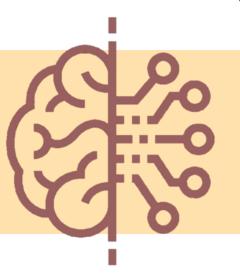


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Exhaustive Search Algorithm

(Exhaustive – Uninformed - Blind Search)



Artificial Intelligence

School of Computing Universiti Teknologi Malaysia





Outline

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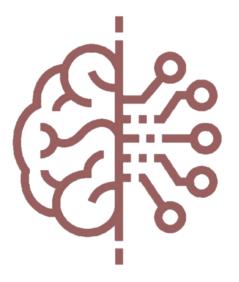
- 1. Selection of a Search Strategy
 - Breadth-first
 - Depth-first
 - Comparison of both search strategy
- 2. Using State Space to Represent Reasoning
 - And/or graph
 - Hypergraph





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Selection of a Search Strategy



- Breadth-first
- Depth-first
- Comparison of both search strategy





Selection of a Search Strategy

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Most of the effort is often spent on the selection of an appropriate search strategy for a given problem

- ✓ Uninformed search (blind / exhaustive search)
 - number of steps, path cost unknown
 - knows when it reaches a goal
- ✓ Informed search (heuristic search)
 - has background information about the problem
 - map, costs of actions





Search Strategy: The order

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Uninformed Search

- Breadth-first
- Depth-first
- Iterative deepening
- Uniform-cost search
- Depth-limited search
- Bi-directional search
- Constraint satisfaction

Informed Search

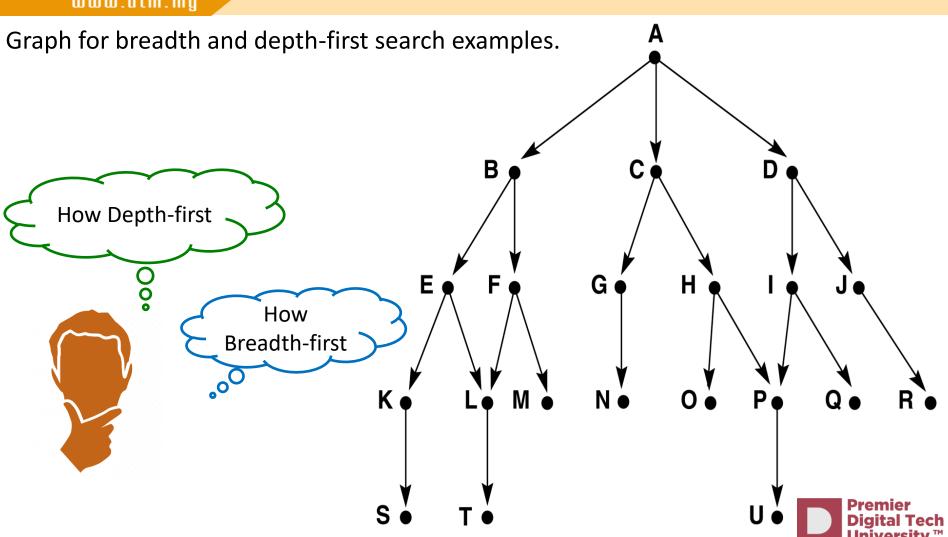
- Best-first search
- Search with heuristics
- Memory-bounded search
- Iterative improvement search





Search Strategy

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- All the nodes reachable from the current node are explored first
 - Achieved by the TREE-SEARCH method by appending newly generated nodes at the end of the search queue

function BREADTH-FIRST-SEARCH(problem) returns solution

return TREE-SEARCH(problem, FIFO-QUEUE())

Time Complexity	b ^{d+1}
Space Complexity	b ^{d+1}
Completeness	yes (for finite b)
Optimality	yes (for non-negative path costs)

b branching factor

d depth of the tree

Solution is found





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Function Breadth-first search algorithm

```
function breadth_first_search;
begin
                                                                             % initialize
  open := [Start];
  closed := [];
  while open ≠ [] do
                                                                       % states remain
    begin
       remove leftmost state from open, call it X;
         if X is a goal then return SUCCESS
                                                                          % goal found
           else begin
             generate children of X;
              put X on closed;
             discard children of X if already on open or closed;
                                                                          % loop check
              put remaining children on right end of open
                                                                               % queue
           end
    end
                                                                        % no states left
  return FAIL
end.
```

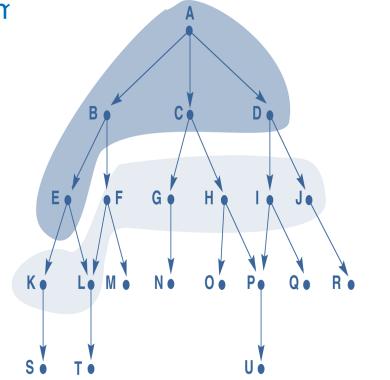




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Function Breadth-first search algorithm

- 1. open = [A]; closed = []
- 2. open = [B,C,D]; closed = [A]
- 3. open = [C,D,E,F]; closed = [B,A]
- 4. open = [D,E,F,G,H]; closed = [C,B,A]
- 5. open = [E,F,G,H,I,J]; closed = [D,C,B,A]
- 6. open = [F,G,H,I,J,K,L]; closed = [E,D,C,B,A]
- 7. open = [G,H,I,J,K,L,M] (as L is already on open); closed = [F,E,D,C,B,A]
- 8. open = [H,I,J,K,L,M,N]; closed = [G,F,E,D,C,B,A]
- 9. and so on until either U is found or **open** = []





Closed

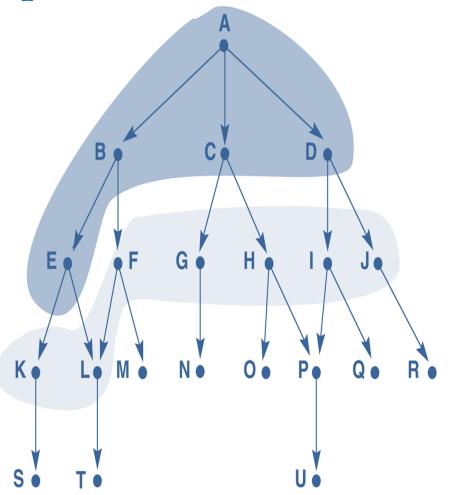
Open



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Function Breadth-first search algorithm

- States on open and closed are highlighted.
- Explore level by level
- Use list open in backtrack states not yet evaluated
- Use list close in backtrack states already evaluated
- Which state is removed from open determines the order of search- BFS adds at the right list and removes from the left (queue-FIFO)



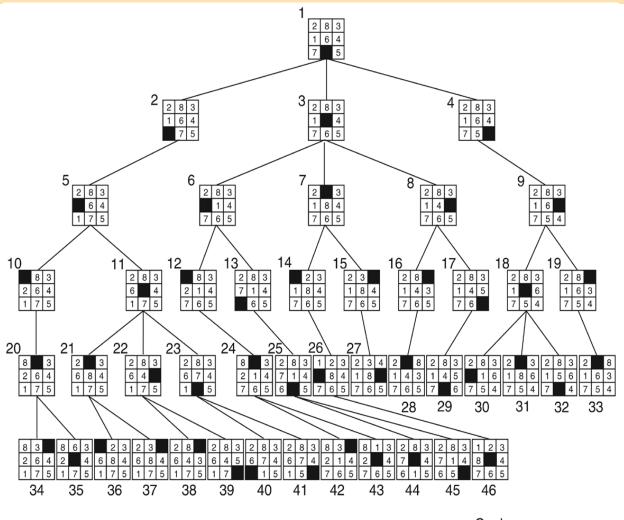






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Breadth-first search of the 8-puzzle, showing order in which states were removed from open



Goal



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Function Depth-first search algorithm

```
begin
                                                                             % initialize
  open := [Start];
  closed := [];
  while open ≠ [] do
                                                                       % states remain
    begin
      remove leftmost state from open, call it X;
      if X is a goal then return SUCCESS
                                                                          % goal found
         else begin
           generate children of X;
           put X on closed;
           discard children of X if already on open or closed;
                                                                          % loop check
           put remaining children on left end of open
                                                                                % stack
         end
    end:
  return FAIL
                                                                        % no states left
end.
```



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Function Depth-first search algorithm

- 1. open = [A]; closed = []
- 2. open = [B,C,D]; closed = [A]
- 3. **open = [E,F,C,D]**; **closed = [B,A]**
- 4. open = [K,L,F,C,D]; closed = [E,B,A]
- 5. **open = [S,L,F,C,D]**; **closed = [K,E,B,A]**
- 6. open = [L,F,C,D]; closed = [S,K,E,B,A]
- 8. open = [F,C,D]; closed = [T,L,S,K,E,B,A]
- 9. open = [M,C,D], as L is already on closed; closed = [F,T,L,S,K,E,B,A]
- 10. open = [C,D]; closed = [M,F,T,L,S,K,E,B,A]
- 11. open = [G,H,D]; closed = [C,M,F,T,L,S,K,E,B,A]



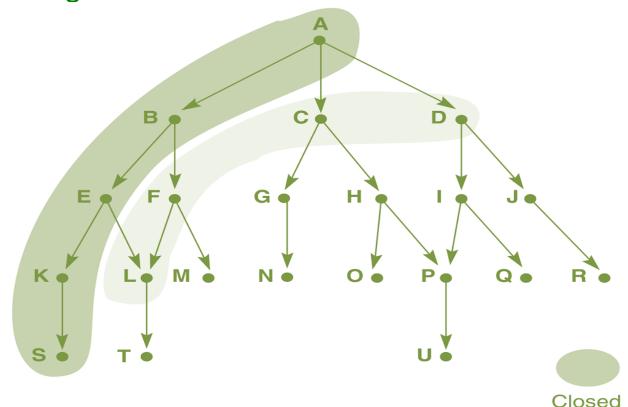


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Function Depth-first search algorithm

States on open and closed are highlighted.

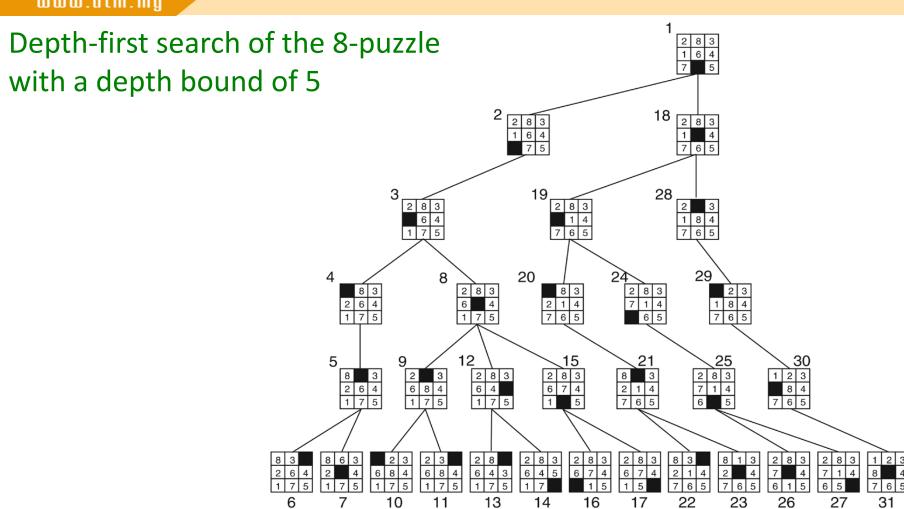
- Child and descendants are evaluated before siblings
- Goes deeper into search space whenever possible
- Which state is removed from open determines the order of search-DFS adds and removes from the left end (stack-LIFO)





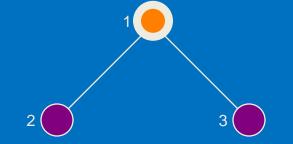


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Initial
Visited
Fringe
Current
Visible

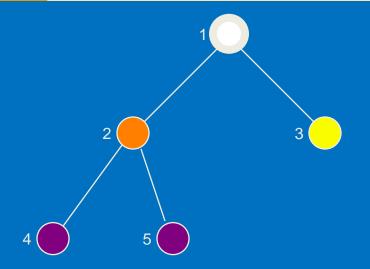
Goal

Fringe = nodes waiting in queue to be explored





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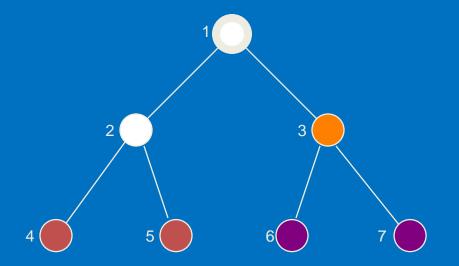
Initial
Visited
Fringe
Current
Visible
Goal

Fringe = nodes
waiting in queue to be explored





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Initial
Visited
Fringe
Current
Visible
Goal

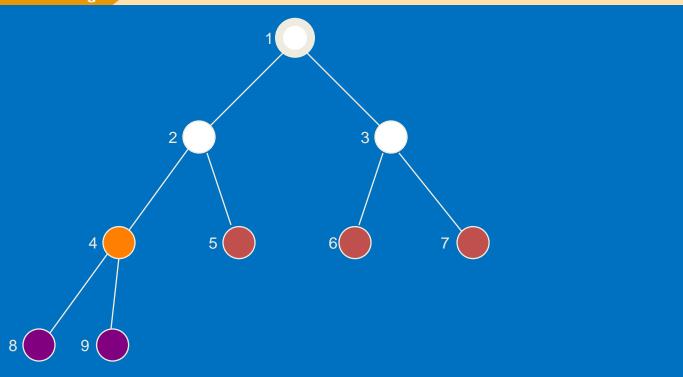
Fringe = nodes
waiting in queue to be explored



Fringe: [4,5] + [6,7]



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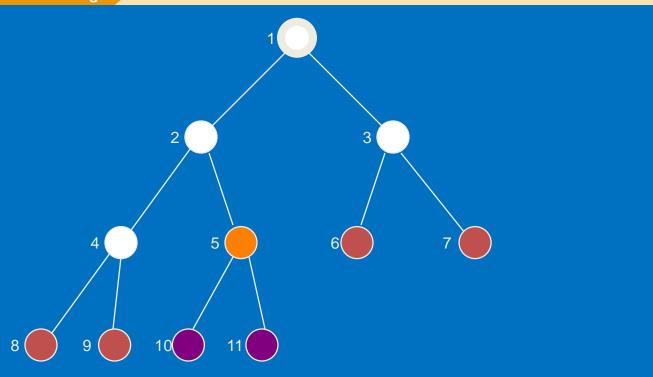
Initial
Visited
Fringe
Current
Visible
Goal

Fringe: [5,6,7] + [8,9]





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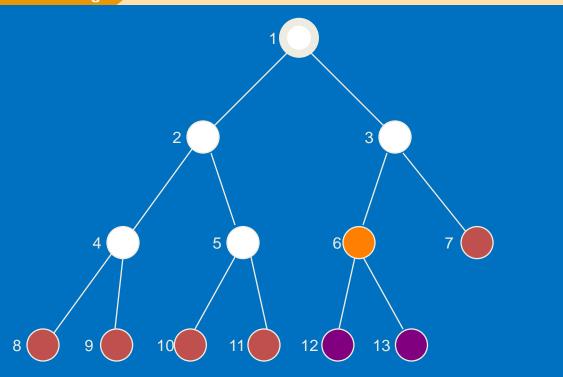
Initial
Visited
Fringe
Current
Visible
Goal

Fringe: [6,7,8,9] + [10,11]





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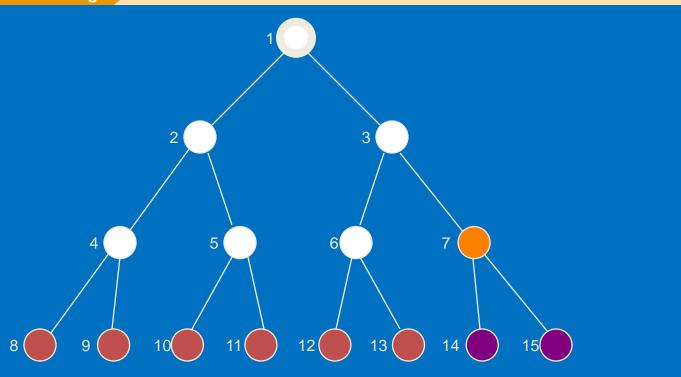
Initial
Visited
Fringe
Current
Visible
Goal

Fringe: [7,8,9,10,11] + [12,13]





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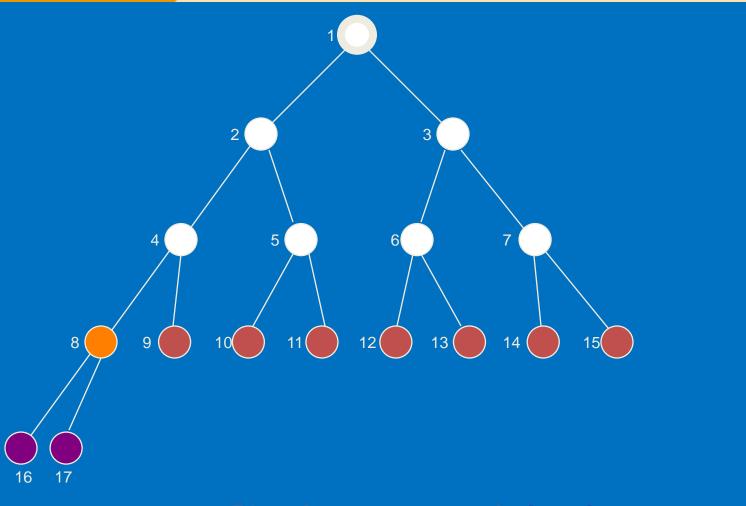
Initial
Visited
Fringe
Current
Visible
Goal

Fringe: [8,9.10,11,12,13] + [14,15]





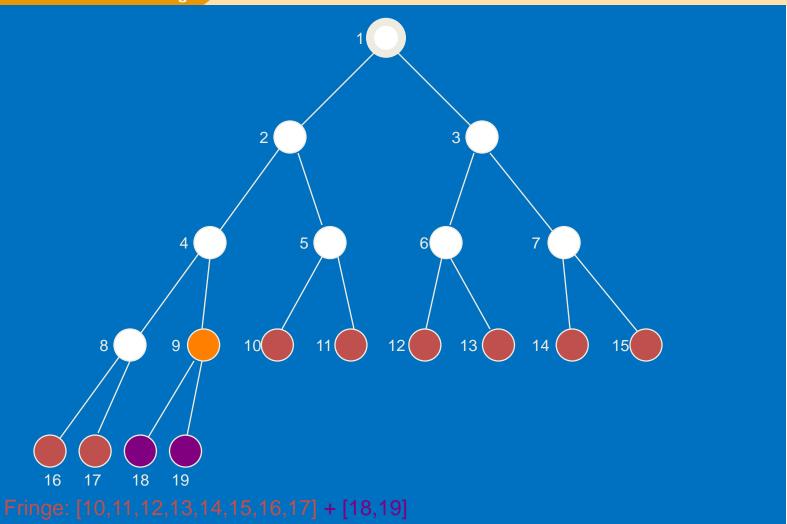
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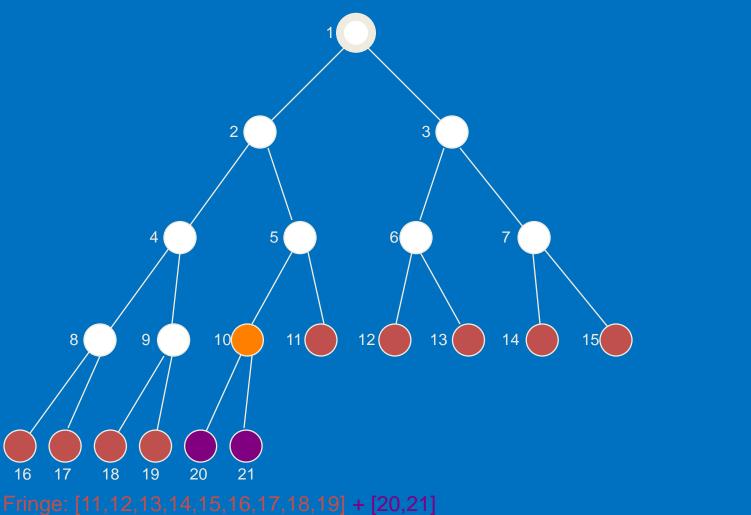
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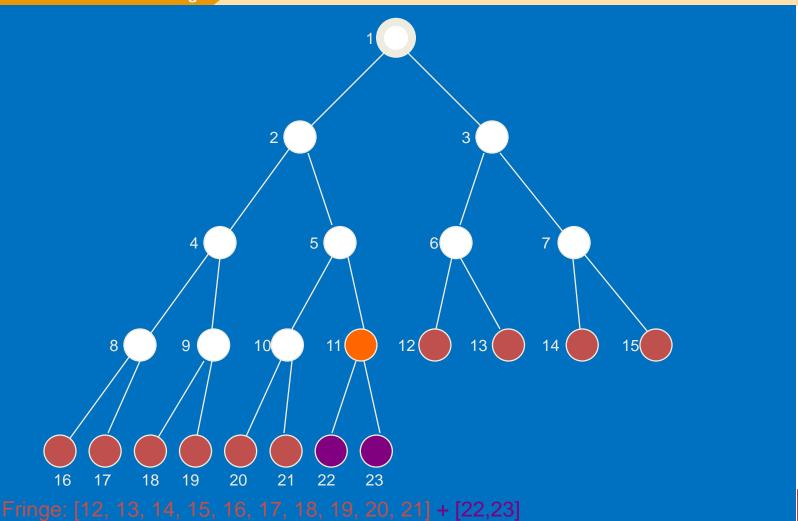
www.utm.my







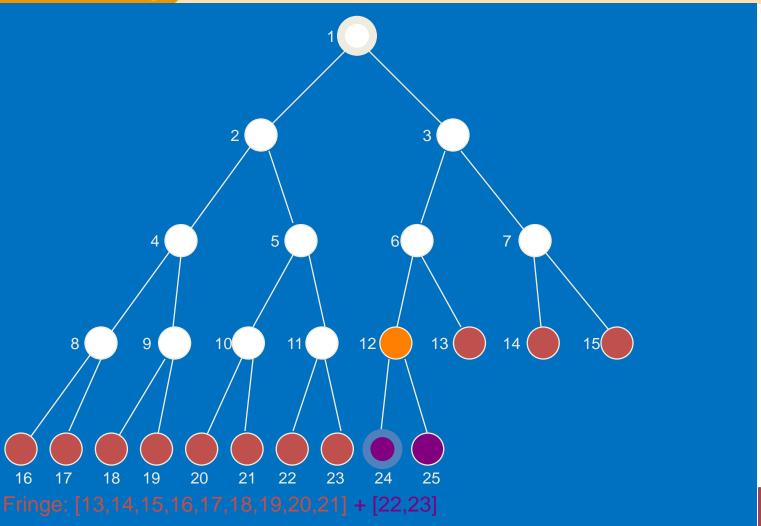
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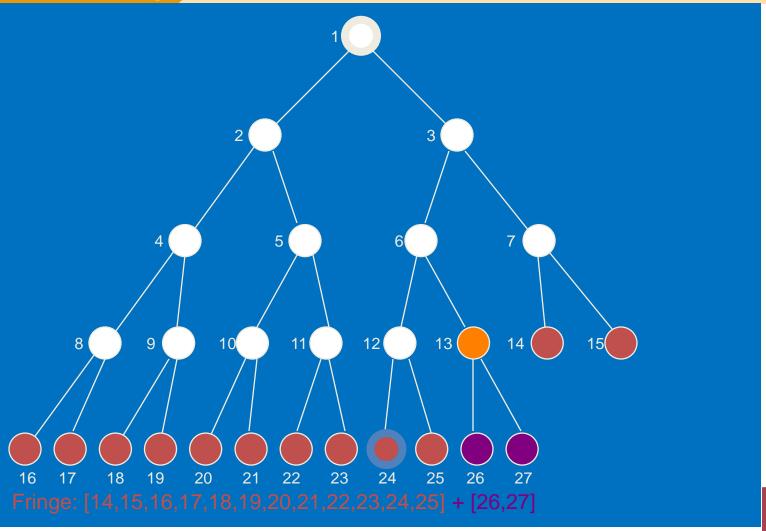
Initial
Visited
Fringe
Current
Visible
Goal

Note: The goal node is "visible" here, but we can not perform the goal test yet.





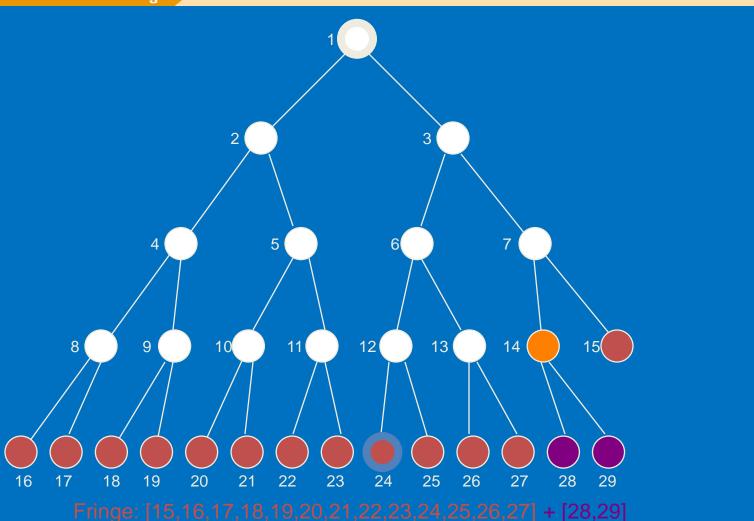
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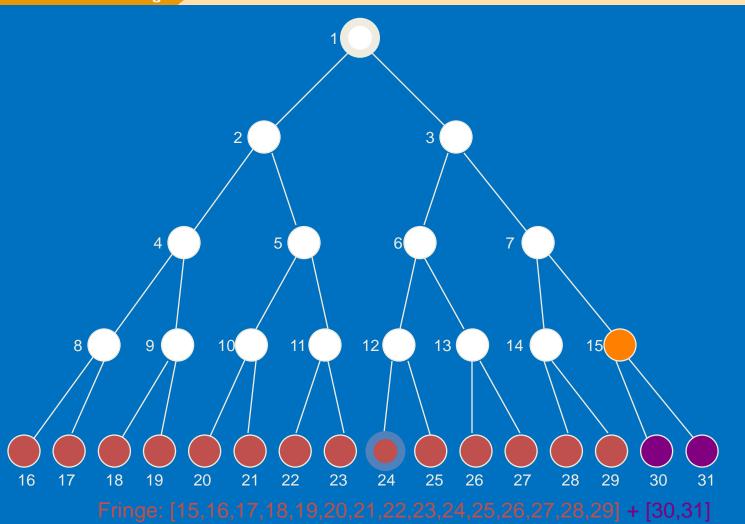
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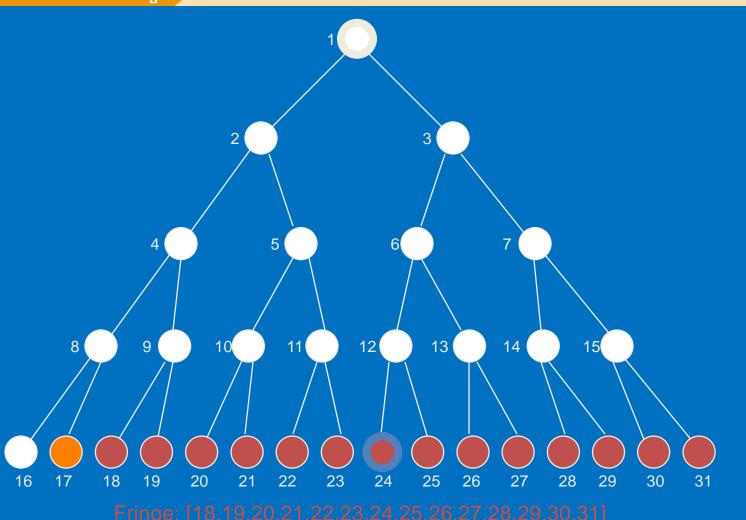
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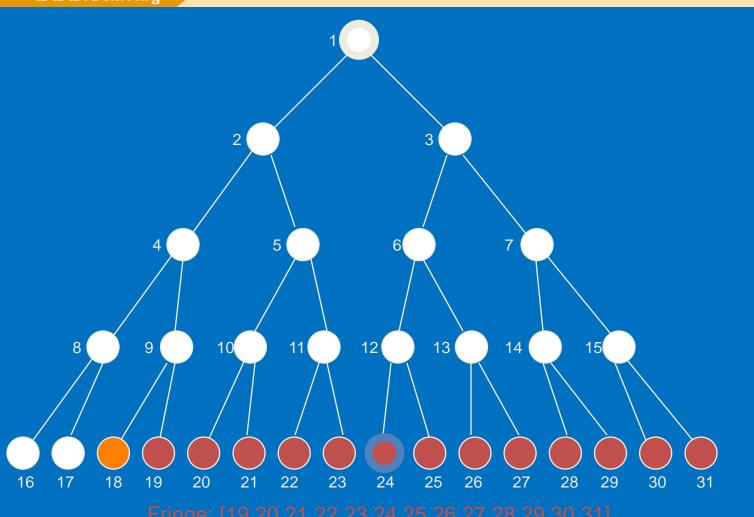
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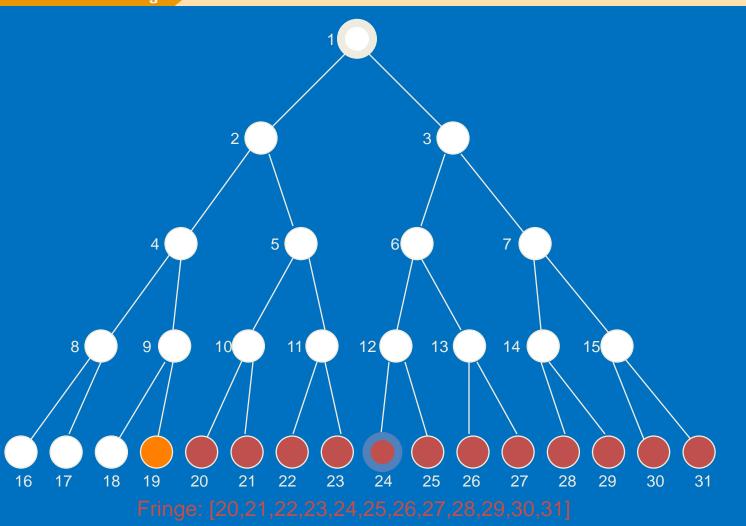
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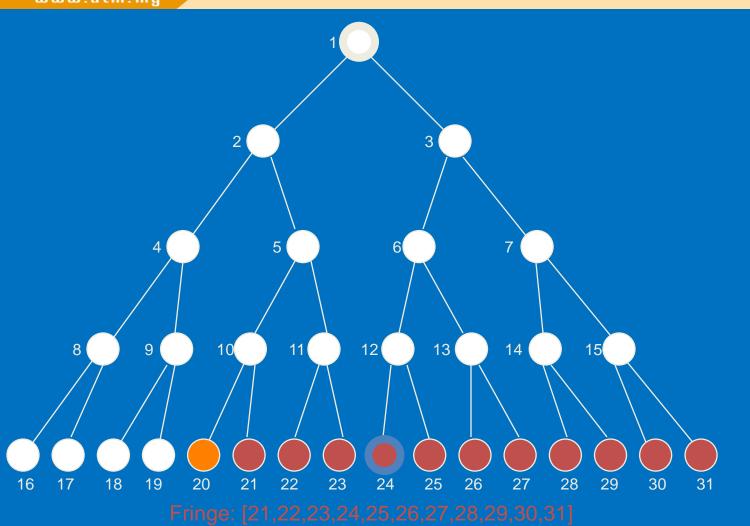
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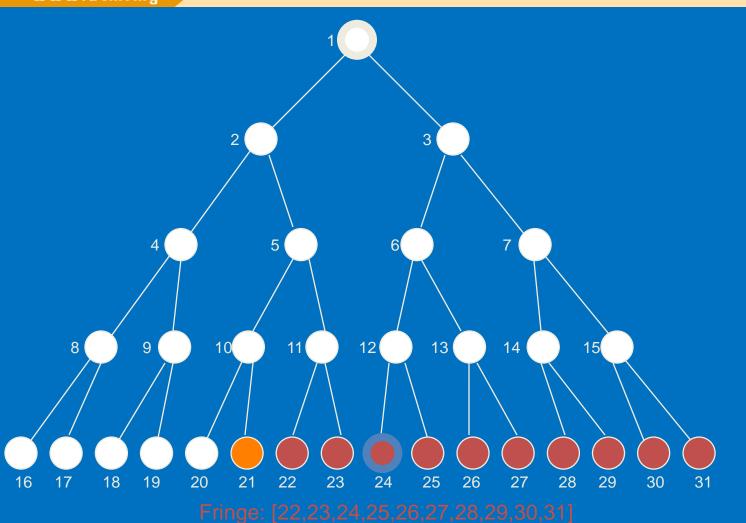
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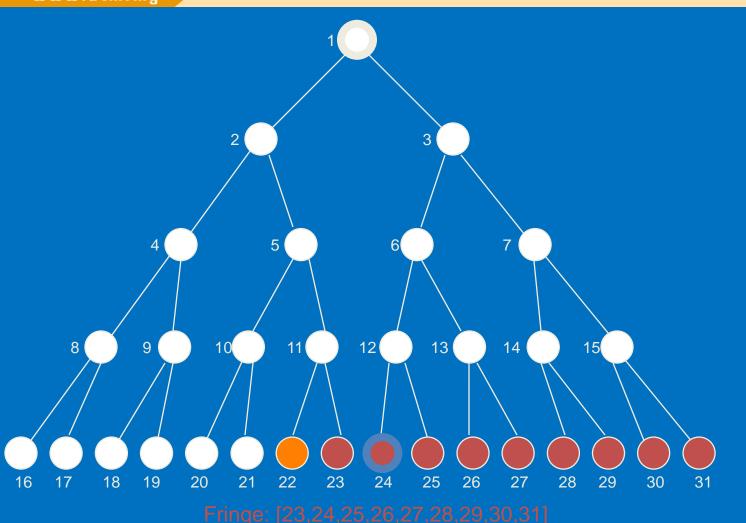
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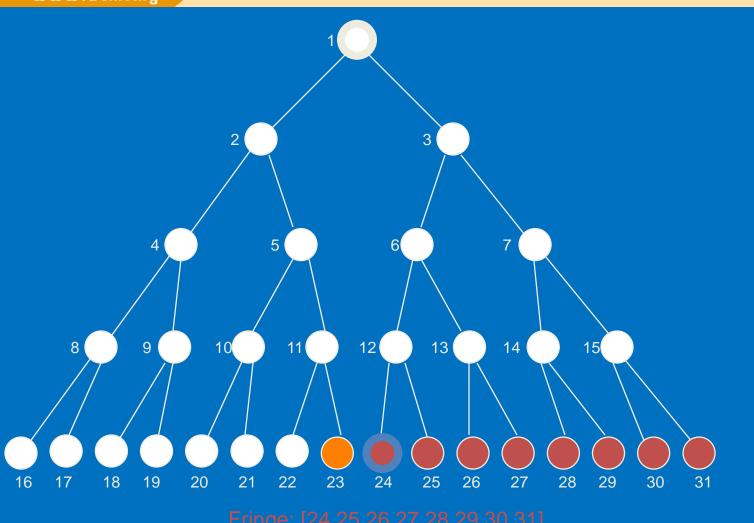






Breadth-First Snapshot 23

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Initial
Visited
Fringe
Current
Visible
Goal

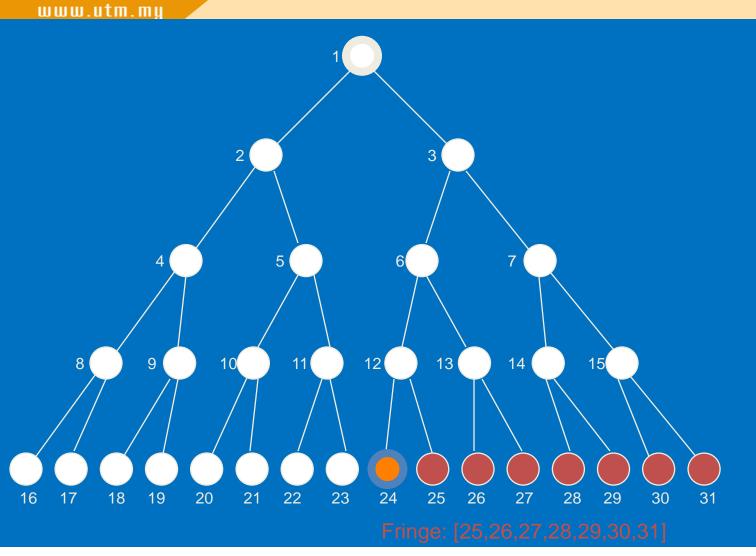
Fringe = nodes
waiting in queue to be explored





Breadth-First Snapshot 24





Initial Visited Fringe Current Visible Goal

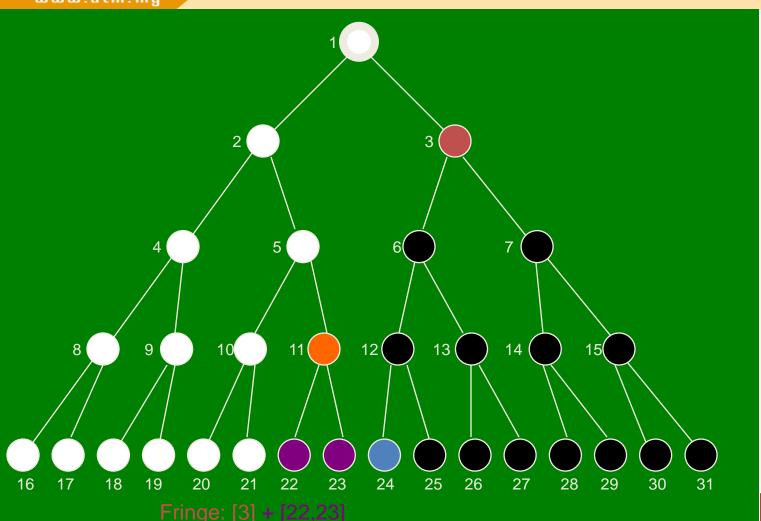
Note: The goal test is positive for this node, and a solution is found in 24 steps.





Depth-First Snapshot

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Initial
Visited
Fringe
Current
Visible
Goal

Fringe = nodes
waiting in queue to be explored





Activity 1: Depth-First vs. Breadth-First

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- Depth-first goes off into one branch until it reaches a leaf node
 - Not good if the goal is on another branch
 - Neither complete nor optimal
 - Uses much less space than breadth-first

Question 1: Give reasons why Depth First search uses less space that breadth first?

- Breadth-first is more careful by checking all alternatives
 - Complete and optimal
 - Under most circumstances
 - But ..Very memory-intensive !!

Question 2: Give a reason why Breadth First search uses very memory-intensive?





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The following state space graph below (Figure 1) represents the gas piping network. The node in the graph denotes the gas gate while the arc denotes the pipe that connects the entrance to the gate. Based on this scenario, answer the following questions:

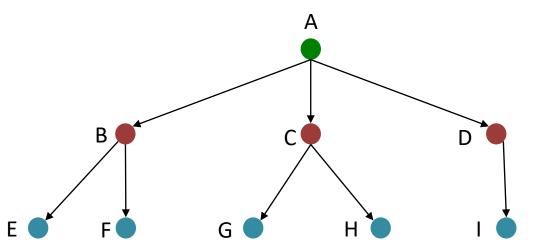


Figure 1. Gas piping network

1. Using graph theory, identify the suitable node/nodes for each description in Table below.

Description	Node (Gas gate)
Start node	А
The child node of A	
The parent node of G. H	
The child node of B	
The parent node of I	





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2. Based on Figure 1, let node A is a start node, and the goal node is E. Perform a **breadth-first search**, then list down the order of nodes to be visited (OPEN and CLOSED list) from the starting node to the goal node in Table below.

Iteration	OPEN	CLOSED
0		
1		
2		

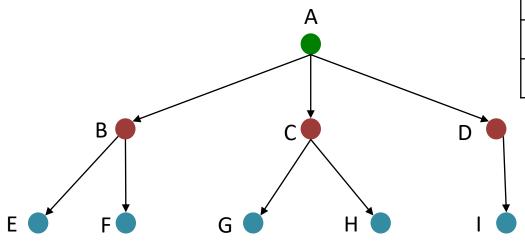


Figure 1. Gas piping network

3. Produce the sequence of nodes for the gas can be flowed from gate A to gate E.





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2. Based on Figure 1, let node A is a start node, and the goal node is E. Perform a breadth-first search, then list down the order of nodes to be visited (OPEN and CLOSED list) from the starting node to the goal node in Table below.

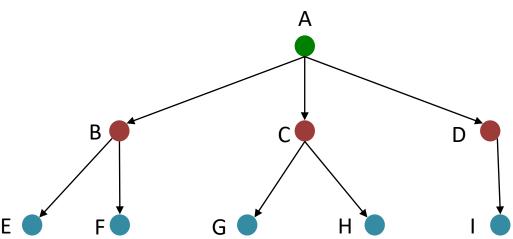


Figure	1.	Gas	piping	networ	k

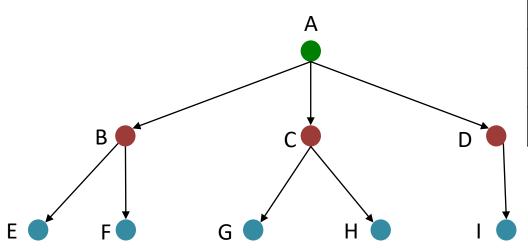
Iteration	OPEN	CLOSED
0	А	0
1	B,C,D	А
2	C,D,E,F	В,А
3	D,E,F,G,H	С,В,А
4	E,F,G,H,I	D,C,B,A
5	E is found	





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1. Based on Figure 1, let node A is a start node, and the goal node is E. Perform a **depth-first search**, then list down the order of nodes to be visited (OPEN and CLOSED list) from the starting node to the goal node in Table below.



Iteration	OPEN	CLOSED
0		
1		
2		
3		
4		
5		

Figure 1. Gas piping network

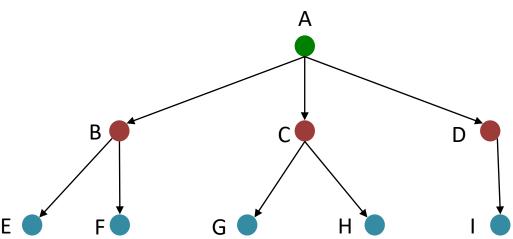
2. Produce the sequence of nodes for the gas can be flowed from gate A to gate E.





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3. Based on Figure 1, let node A is a start node, and the goal node is E. Perform a depth-first search, then list down the order of nodes to be visited (OPEN and CLOSED list) from the starting node to the goal node in Table below.



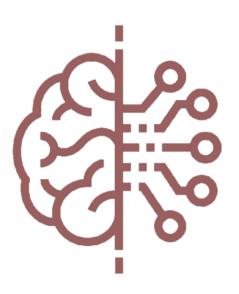
G 🔵		Н		
Figure	1. Gas	piping	netv	vork

Iteration	OPEN	CLOSED
0	А	[]
1	B,C,D	А
2	E,F,C,D	В,А
3	E is found	
4		
5		





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Using State Space to Represent Reasoning

- And/or graph
- Hypergraph





Using State Space to Represent Reasoning: Predicate calculus

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- Predicate calculus map nodes/some states of a graph onto state space
- Inference rules describe the arc between states
- Eg. Is expression a logical consequence of given assertion?
 Use search to solve
- Guarantee correctness of conclusion derived- a formal proof of integrity of solution





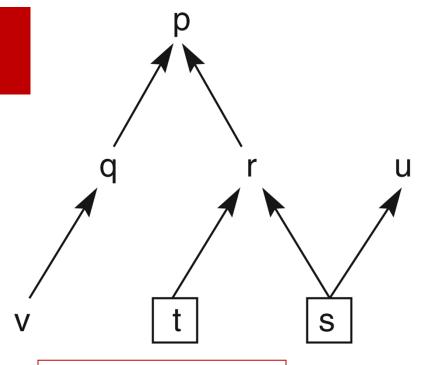
Draw a graph: Propositional calculus

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Example: Define a graph from propositional calculus to view the logical relationships.

$$q \rightarrow p$$
 $r \rightarrow p$
 $v \rightarrow q$
 $s \rightarrow r$
 $t \rightarrow r$
 $s \rightarrow u$
 s

Translate into graph?



- State space graph of a set of implications in the propositional calculus.
- Reasoning:- How to infer p using moden ponens?

s , s -->r yields r

r, r-->p yields p

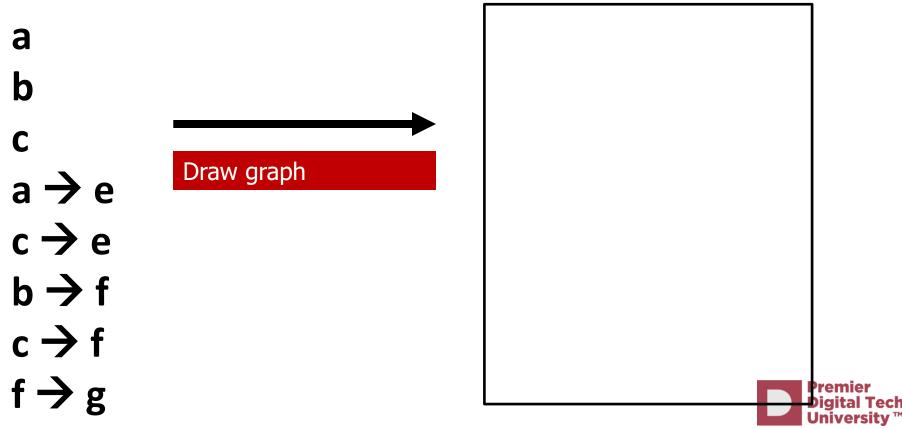




Activity 2: From propositional calculus to a graph

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Give a set of propositional calculus expressions, you are required to draw a graph.

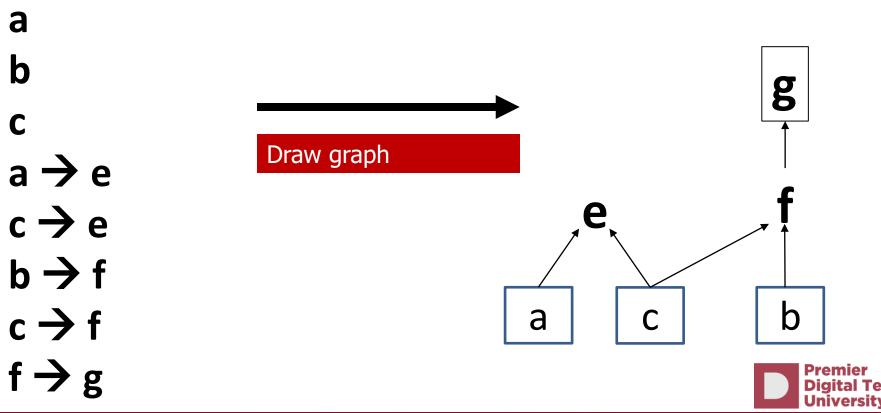




Activity 2: From propositional calculus to a graph

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Give a set of propositional calculus expressions, you are required to draw a graph.

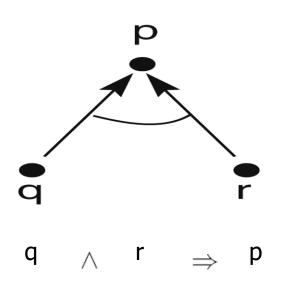


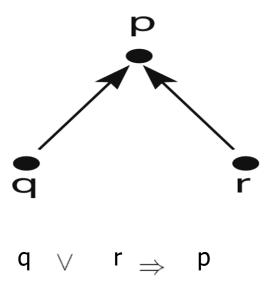


And/or graph

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- To represent logical operators AND and OR in predicate calculus
- extension of basic graph model known as and/or graph
- To differ the relationships graphically, curve link indicates operator AND.









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AND/OR graph is a specialization of a type of graph known as *hypergraph*, which connects nodes by sets of arcs rather than a single arcs

DEFINITION

HYPERGRAPH

A hypergraph consists of:

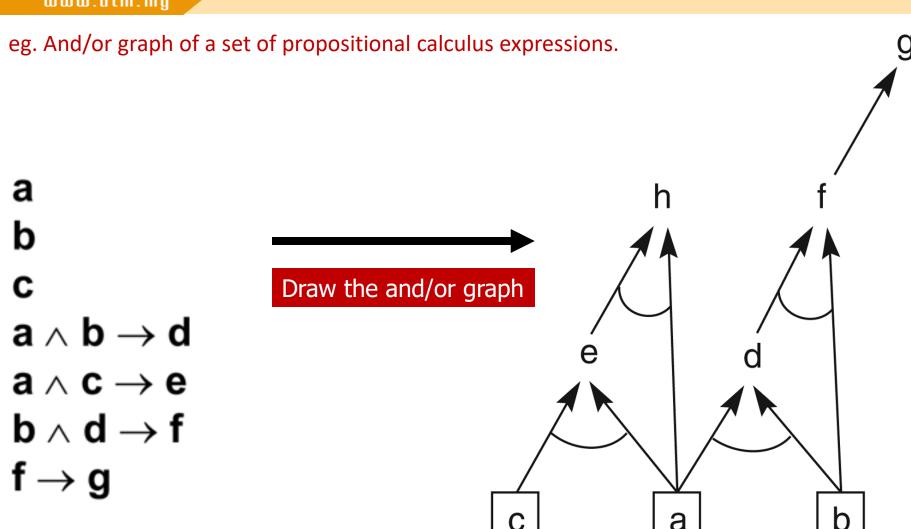
N, a set of nodes.

 \mathbf{H} , a set of hyperarcs defined by ordered pairs in which the first element of the pair is a single node from \mathbf{N} and the second element is a subset of \mathbf{N} .

An ordinary graph is a special case of hypergraph in which all the sets of descendant nodes have a cardinality of 1.

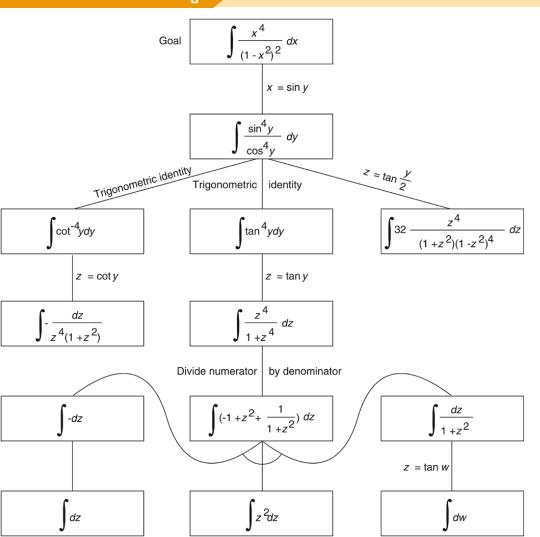


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More example:

And/or graph of part of the state space for integrating a function, from Nilsson (1971).

Goal-directed search





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- Fred is a collie.
 collie(fred).
- Sam is Fred's master. master(fred,sam).
- 3. The day is Saturday. day(saturday).
- It is cold on Saturday.
 ¬ (warm(saturday)).
- 5. Fred is trained. trained(fred).
- Spaniels are good dogs and so are trained collies.
 ∀ X[spaniel(X) ∨ (collie(X) ∧ trained(X)) → gooddog(X)]
- 7. If a dog is a good dog and has a master then he will be with his master. \forall (X,Y,Z) [gooddog(X) \land master(X,Y) \land location(Y,Z) \rightarrow location(X,Z)]
- 8. If it is Saturday and warm, then Sam is at the park.

 (day(saturday) ∧ warm(saturday)) → location(sam,park).
- 9. If it is Saturday and not warm, then Sam is at the museum. $(day(saturday) \land \neg (warm(saturday))) \rightarrow location(sam,museum).$

More example:

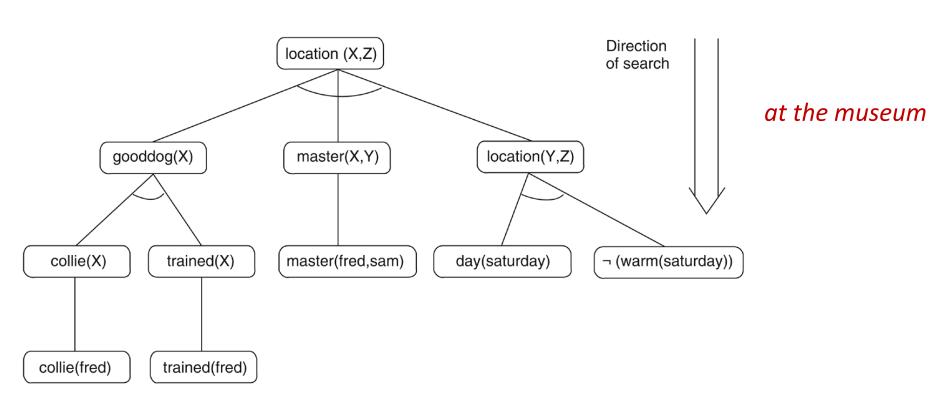
- Where is Fred??
- The facts and rules of this example are given as English sentences followed by their predicate calculus equivalents:





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Where is Fred???..The solution subgraph showing that fred is ???



Substitutions = {fred/X, sam/Y, museum/Z}

Goal-driven search using predicate calculus

