

## Solutions of Homework 5: Big $O, \Omega$

Important Note:  $\log n \leq \sqrt{n} \leq n$ .

**Q1.**

1.  $100n + 1 \leq 100n + n = 101n = O(n) \Rightarrow c = 101, n_0 = 1$ .
2.  $(10n + 1)^3 \leq (10n + n)^3 = (11n)^3 = 1331n^3 = O(n^3) \Rightarrow c = 1331, n_0 = 1$ .
3.  $3n^3 - 5n^2 - 100 \leq 3n^3 = O(n^3) \Rightarrow c = 3, n_0 = 1$ .
4.  $n^2 + n + \sqrt{n} + \log n \leq n^2 + n^2 + n^2 + n^2 = 4n^2 = O(n^2) \Rightarrow c = 4, n_0 = 1$ .

**Q.2**

- i.  $6n^2 - 2n \leq 6n^2 = O(n^2)$  ..... (1)  
 $6n^2 - 2n \geq 6n^2 - 2n^2 = 4n^2 = \Omega(n^2)$  .... (2)  
Form (1) and (2):  $6n^2 - 2n \log n = \Theta(n^2)$ .
- ii.  $\frac{6n^2}{\log^3 n + 1} \leq 6n^2 \leq n^3 = O(n^3)$ .
- iii.  $3n^3 + 44n^2 \geq n^2 = \Omega(n^2)$ .

**Q.3** Proof by contradiction: Assume that

$$\begin{aligned} & (\log n)^2 = O(\log n^2) \\ \rightarrow & (\log n)^2 \leq c \log n^2 = k \log n \text{ (note that } \log n^2 = 2 \log n) \\ \rightarrow & \log n \leq k. \end{aligned}$$

However, it is impossible to find such a constant  $k$  which is always greater than  $\log n$  for all possible values of  $n$  (taking into account that it is a monotonically increasing function). Therefore,  $(\log n)^2 \neq O(\log n^2)$ .

**Q.4** Note that:

- $2^{\log_2 n} = n$ .
- $2^{3 \log_2 n} = 2^{\log_2 n^3} = n^3$ .
- $\left(\frac{3}{2}\right)^n \leq 2^n$ .
- By sketching the graphs for  $\log^4 n$  and  $n$ , you can conclude that  $\log^4 n \leq n$ .
- A constant function (such as 100) does not grow with  $n$  as opposed to any other function and, therefore, it is upper bounded by any of these functions regardless the value of this constant.
- The exponential function  $a^n$  always dominates polynomial and logarithmic functions.

Hence, the required ascending order is:

$$100 \prec \log^4 n \prec 2^{\log_2 n} \prec 2^{3 \log_2 n} \prec n^3 \log^2 n \prec \left(\frac{3}{2}\right)^n \prec 2^n.$$