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# Databases Project – Spring 2021

Team No: G34

Members: Siran Li, Ke Wang, Ningwei Ma

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# Deliverable 3

# **Changes**

Based on the feedback, we have made some changes including ER model, DDL, data processing and query implementation of Deliverable 2. We have quoted some of the feedback and have responded to it.

# ER MODEL

Modifications respective to Feedback on M2 (the modifications are based on the feedback on ER:

-- 'Yes, I think that the "condition" entity could be merged with "case". It doesn't give you any specific advantage to keep these two entities as separate.'

Removed the 'condition' entity and moved all its attributes to the 'case' entity.

-- 'For example, you can create a new entity (and the corresponding table) for "road conditions", with road condition as an attribute and the case\_id acts as a foreign key referred from the collision table. Note that in this case, both the case\_id and the road condition serve as the PK.'

Changed the attributes of 'road\_en' entity into 'case\_id' and 'road\_con', the pairs of them serve as primary keys.

-- 'Shouldn't you have a foreign key reference of the primary keys of party into this table? Also, I don't think factors can exist without a party. Lastly, I don't think you need a separate table for the "have" relationship b/w party and other\_fac\_en, it is absorbed as a single table for 'other\_fac\_en'.'

Changed the attributes of 'other\_fac\_en' entity into 'party\_id' and 'other\_fac', both of them serve as primary keys. Replaced the attribute 've num' of 'vehicle' entity by 'party id'.

--'It is indeed correct to have a single "safety\_equipment" entity in the ER, however, during the translation I think you would need two separate tables: safetey\_equipment for party and victim. Similar to the previous point and with a very similar explanation, I don't think you need separate tables for the relationships "have ps" and "have vs".'

For this feedback, we decided to separate the 'other\_fac\_en' entity into two entities in the ER model as well.

Separated the 'safety\_equip\_en' entity into two entities 'safety\_p' (the safety equipment for party) and 'safety\_v' (the safety equipment for victim').

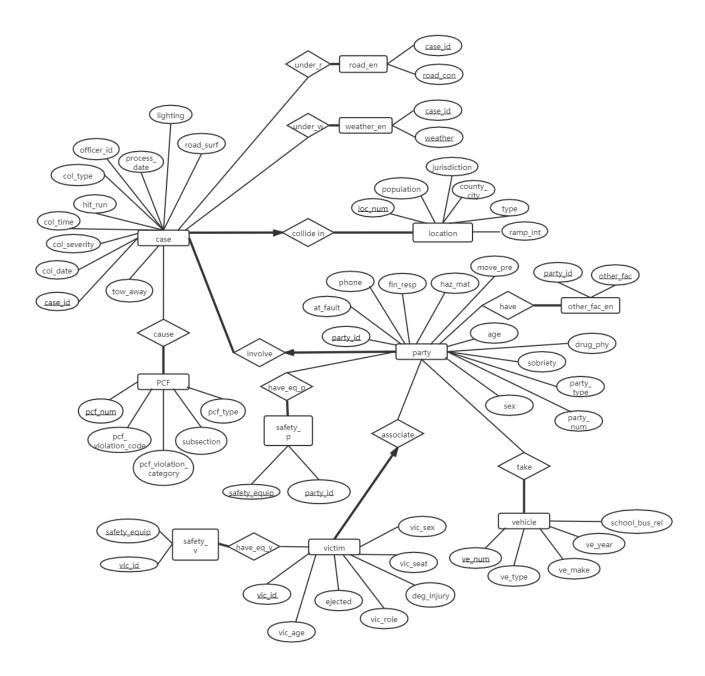
For 'safety\_p', its attributes are changed to 'safety\_equip' and 'party\_id', both of them serve as primary keys.

For 'safety\_v', its attributes are changed to 'safety\_equip' and 'vic\_id' both of them serve as primary keys.

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# DDL

In response to deliverable 2's feedback, we have modified the DDL.

--' I don't think you need a separate table for the "have" relationship b/w party and other\_fac\_en, have\_ps, have\_vs, under\_r, under\_w, take'

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We deleted the relational tables such as under\_w, under\_r, have, have\_ps, have\_vs, and merged them to the independent entity tables including weather\_en, other\_fac\_en, safety\_equipment, road\_en, respectively.

Now in the relational table, such conditions (weather, road condition, other\_fac, safety\_p, safety\_v, road\_con) will be like an attribute and the case\_id/party\_id/victim\_id will act as a foreign key referred from the corresponding table. We use the pairs of (xx\_id, xx\_condition) as the primary key to ensure that they are unique.

--'other\_fac\_en should be added the 'on delete cascade' constraint. '

In these tables (other\_fac\_en, road\_en, weather\_en, safety\_p, safety\_v), we add "on delete cascade" constraint on the foreign key, because these record cannot exist without party/case/victim. Once the corresponding party/case/victim is deleted from the database, the record related to them would also be deleted.

--'Also, shouldn't ve\_num be added as a foreign key in the party table? Lastly, similar to the above two points, I don't think you need a separate table for the "take" relationship.'

Each case happens under exactly one location and PCF and each party drives one kind of vehicle(by the data and our way of data processing). We deleted the relational table 'take' between party\_involve-take-vehicle, but just added a column "ve\_num" in the party\_involve table referring to the "ve\_num" of vehicle table. We also deleted the relation 'cause' table between PCF-casue-case and added a foreign key "pcf\_num" to the case table referring to the "pcf\_num" of PCF table.

--'Recall that party can/should not exist without a collision, and victim can/should not exist without a party.'

We added "on delete cascade" constraint on the "party\_id" foreign key in the associate\_victim table, and the "case\_id" foreign key in the party\_involve table, because party is the weak entity of the case, and the victim if the weak entity of the party (by our assumption and data processing).

--'Since party is a weak entity of case, shouldn't you have a foreign key reference of the primary keys of party into this table?'

We modeled the party a weak entity of the case entity. We put the foreign key reference of the case\_id in the party entity, based on this assumption: party has exactly one participation in this relationship, but one case may be related to many parties. So, we think it might be more reasonable to put the foreign key to case in the party table. And thus we did not put foreign key referring to party in the case table, to avoid a reference loop.

--'I just want to make sure that the duplicates are not because of the datatype issue (you are modeling case\_id as integer), and are genuine duplicates'

We changed the type of "case\_id" to VARCHAR2, because there may be several zeros in front of the number, which is also a part of the identity. Now there is no duplicated case\_id. (We did not change the type of party\_id and vict\_id, because they do not have duplicates)

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-- 'Yes, I think that the "condition" entity could be merged with "case". It doesn't give you any specific advantage to keep these two entities as separate.'

We have merged the condition table into the case table, with table condition's two attributes becoming the attributes of case now.

We have 11 tables now, and the DDL is attached below. Modified parts are marked in yellow.

\_\_\_\_\_\_

```
CREATE TABLE PCF(
pcf_num INTEGER,
pcf_violation_code INTEGER,
pcf_violation_category VARCHAR2(50),
subsection VARCHAR2(3),
pcf_type VARCHAR2(50),
PRIMARY KEY (pcf_num)
);
CREATE TABLE Vehicle(
ve_num INTEGER,
ve_type VARCHAR2(50),
ve_make VARCHAR2(30),
ve_year INTEGER,
school bus rel VARCHAR2(5),
PRIMARY KEY (ve_num)
);
CREATE TABLE Location(
loc_num INTEGER,
```

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```
population INTEGER,
county_city INTEGER,
loc_type VARCHAR2(20),
ramp_int VARCHAR2(10),
PRIMARY KEY (loc num)
);
CREATE TABLE Case(
case_id VARCHAR2(30),
loc_num INTEGER NOT NULL,
pcf_num INTEGER NOT NULL,
col_severity VARCHAR2(30),
col_time DATE,
col_date DATE,
hit_run VARCHAR2(30),
jurisdiction INTEGER,
officer_id VARCHAR2(10),
process_date DATE,
tow_away INTEGER,
col_type VARCHAR2(30),
lighting VARCHAR2(50),
road_surf VARCHAR2(10),
PRIMARY KEY (case id),
FOREIGN KEY (loc_num) REFERENCES Location(loc_num),
```

FOREIGN KEY(pcf\_num) REFERENCES PCF(pcf\_num)

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```
);
CREATE TABLE Party_involve(
party_id INTEGER,
case_id VARCHAR2(30) NOT NULL,
at_fault INTEGER,
phone VARCHAR2(3),
fin_resp VARCHAR2(3),
haz_mat VARCHAR2(3),
move_pre VARCHAR2(30),
age INTEGER,
drug_phy VARCHAR2(3),
sobriety VARCHAR2(3),
party_type VARCHAR2(15),
party_num INTEGER,
sex VARCHAR2(6),
ve_num INTEGER,
PRIMARY KEY (party_id),
FOREIGN KEY (case_id) REFERENCES Case(case_id)
ON DELETE CASCADE,
FOREIGN KEY (ve_num) REFERENCES Vehicle(ve_num)
);
CREATE TABLE Associate_victim(
vic_id INTEGER,
```

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```
party_id INTEGER NOT NULL,
vic_age INTEGER,
ejected INTEGER,
vic_role INTEGER,
deg_injury VARCHAR2(50),
vic_seat INTEGER,
vic_sex VARCHAR2(6),
PRIMARY KEY (vic_id),
FOREIGN KEY (party_id) REFERENCES Party_involve(party_id)
ON DELETE CASCADE
);
CREATE TABLE Other_fac_en(
party_id INTEGER,
other_fac VARCHAR2(3),
PRIMARY KEY (party_id, other_fac),
FOREIGN KEY (party_id) REFERENCES Party_involve(party_id)
ON DELETE CASCADE
);
CREATE TABLE road_en(
case_id VARCHAR2(30),
road con VARCHAR2(20),
PRIMARY KEY (case_id, road_con),
FOREIGN KEY (case_id) REFERENCES Case(case_id)
```

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# ON DELETE CASCADE ); CREATE TABLE Weather\_en( case id VARCHAR2(30), weather VARCHAR2(20), PRIMARY KEY (case\_id, weather), FOREIGN KEY (case\_id) REFERENCES Case(case\_id) ON DELETE CASCADE ); CREATE TABLE safety\_p( party\_id INTEGER, safety\_equip VARCHAR2(3), PRIMARY KEY (party\_id, safety\_equip), FOREIGN KEY (party\_id) REFERENCES Party\_involve(party\_id) ON DELETE CASCADE ); CREATE TABLE safety\_v( vic\_id INTEGER, safety\_equip VARCHAR2(3), PRIMARY KEY (vic\_id, safety\_equip), FOREIGN KEY (vic\_id) REFERENCES Associate\_victim(vic\_id)

ON DELETE CASCADE);

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# Query Implementation of Deliverable 2

#### Query 1:

#### **Description of logic:**

List the year-number of collisions per year. We use "group by" to group case by year (extracted from col\_date) and count the number of cases of each year.

#### SQL statement

SELECT EXTRACT (YEAR FROM col\_date) AS YEAR, count(\*) AS N\_collisions FROM case
GROUP BY EXTRACT (YEAR FROM col\_date)
ORDER BY YEAR ASC

#### Query result (if the result is big, just a snippet)

YEAR	N_COLLISIONS
2001	522562
2002	544741
2003	538954
2004	538295
2005	532725
2006	498850
2007	501908
2017	7
2018	21

#### Query 2:

#### **Description of logic:**

In the "take" table, group entries by "ve\_make" and count the number of parties of each "ve\_make", then find the max count and the corresponding "ve\_make". Before that we need to use "ve\_number" to know the "ve\_make", so we first join table "vehicle" and "party\_involve". To illustrate the whole row of the most popular, we fetch first 1 row only.

#### **SQL** statement

SELECT VE MAKE, COUNT(\*) AS N COLLISION

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FROM (VEHICLE INNER JOIN PARTY\_INVOLVE ON VEHICLE.VE\_NUM = PARTY\_INVOLVE.VE\_NUM)
GROUP BY VE\_MAKE
ORDER BY N\_COLLISION DESC
FETCH FIRST 1 ROWS ONLY

#### Query result (if the result is big, just a snippet)

VE_MAKE	N_VEHICLE
FORD	1129701

#### Query 3:

#### **Description of logic:**

In the lighting attribute of condition, find the description that contains "dark", and count the fraction of cases that occur in such condition, and keep 2 significant digits..

#### **SQL** statement

SELECT ROUND(NOM/(SELECT COUNT(\*) FROM CASE),2)
FROM
(SELECT COUNT(\*) AS NOM
FROM CASE
WHERE CASE.LIGHTING LIKE '%dark%')

#### Query result (if the result is big, just a snippet)

query resure (if the resure is big) just a simple
FRACTION
0.28

## Query 4:

### **Description of logic:**

Find the number of collisions that have occurred under snowy weather. We count the number of entries that have weather\_con = 'snowing' in the table "weather\_en"

#### **SQL** statement

SELECT COUNT(\*)
FROM WEATHER\_EN
WHERE WEATHER LIKE '%snowing%'

#### Query result (if the result is big, just a snippet)

Query result (i) the result is alg, just a simple
N_COLLISIONS
8530

#### Query 5:

Description of logic:

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Group by collisions by which day they are during a week, and count the total number of collisions of that day, then find the row of highest number of cases. We use TO\_CHAR (COL\_DATE, 'D') to extract the day of the week.

#### **SQL** statement

SELECT TO\_CHAR(COL\_DATE, 'D') AS WEEK\_DAY, COUNT(\*) AS N\_COLLISONS FROM CASE
GROUP BY TO\_CHAR(COL\_DATE, 'D')
ORDER BY N\_COLLISONS DESC
FETCH FIRST 1 ROWS ONLY

## Query result (if the result is big, just a snippet)

WEEK_DAY	N_COLLISIONS
6	614853

#### Query 6:

#### **Description of logic:**

List all weather types and their corresponding number of collisions in ascending order of the collisions. We group cases by "weather" and list "weather" and the count number.

## **SQL** statement

SELECT WEATHER, COUNT(\*) AS N\_COLLISION FROM WEATHER\_EN GROUP BY WEATHER ORDER BY N\_COLLISION

## Query result (if the result is big, just a snippet)

WEATHER	N_COLLISION
other	6960
snowing	8530
wind	13952
fog	21259
raining	223752
cloudy	548250
clear	2941042

#### Query 7:

#### **Description of logic:**

Count the number of parties that are at-fault, with financial responsibility and loose material. We first extract the "road\_num" of "road\_loose", and find which parties are associated with such road condition. We filter the "party\_id" table who is at fault and with financial responsibility. Finally we count the number of the selected parties.

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#### **SQL** statement

SELECT COUNT(\*) AS N\_PARTIES
FROM PARTY\_INVOLVE P, ROAD\_EN R
WHERE P.CASE\_ID = R.CASE\_ID AND P.AT\_FAULT = 1 AND P.FIN\_RESP = 'Y' AND R.ROAD\_CON LIKE '%loose material%'

# Query result (if the result is big, just a snippet)

-	N_PA	ARTIES	, ,
	4	803	

#### Query 8:

#### Description of logic:

Find the median victim age: we directly use the "MEDIAN" function of SQL from the associate\_victim table.

Find the most common victim seating position. We group the victims with seating position, and count the number of victims of each vic\_seat, order them in the descending order of this number and find the max.

# **SQL** statement

#### 8.a

SELECT MEDIAN(vic\_age) AS MEDIAN\_VIC\_AGE FROM ASSOCIATE\_VICTIM v2;

#### 8.b

SELECT VIC\_SEAT AS MOST\_COMMON\_SEAT\_POSITION FROM

(SELECT COUNT(vic\_seat) AS count, vic\_seat
FROM associate\_victim v2
GROUP BY vic\_seat
ORDER BY count DESC)

FETCH FIRST 1 ROWS ONLY;

# Query result (if the result is big, just a snippet)

#### 8.a

<b></b>
MEDIAN_VIC_AGE
25

#### 8.b

MOST_COMMON_SEAT_POSITION	
3	

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# Query 9:

#### **Description of logic:**

Fraction of all participants (victims + parties) that have been victims using a belt. All participants refer to both parties and victims, so our denominator is the sum of number of all victims and parties. We first extract the 'vic\_id's who use belt using table have\_vs and safety\_equip\_en. Then we count the unique 'vic\_id's and use this number as the numerator. Finally we get the fraction and keep 2 significant digits.

#### **SQL** statement

```
SELECT ROUND(A.FRACTION,3) AS fraction

FROM(SELECT DISTINCT

(SELECT COUNT(vic_id) AS count

FROM

(SELECT h1.vic_id as vic_id

FROM SAFETY_V H1

WHERE H1.SAFETY_EQUIP like '%C%') v_belt)/

((SELECT COUNT(party_id) FROM party_involve)

+(SELECT COUNT(vic_id) FROM associate_victim)) as fraction

FROM party involve) a
```

#### Query result (if the result is big, just a snippet)

FRACTION	
0.011	

#### Query 10:

#### Description of logic:

Compute the fraction of collisions happening for each hour of the day, and display the ratio as percentage for all the hours of the day. We first use cast(col\_time as timestamp) to extract the hour in which the case occurred. Then we group the cases by the specific hour and count the number of the cases, then order them by the number. We also directly calculate the count of total cases. Then we divide the count number of each hour by the total number to get each fraction.

#### **SQL** statement

**SELECT DISTINCT** 

EXTRACT(hour from cast(col\_time as timestamp)) as hour, CONCAT( ROUND((COUNT(\*)/(SELECT COUNT(\*) FROM CASE)\*100.0),2),'%') as FRACTION

FROM CASE

GROUP BY EXTRACT(hour from cast(col\_time as timestamp))

ORDER BY hour ASC

Query result (if the result is big, just a snippet)

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HOUR	
	FRACTION
0	1.91%
1	1.83%
2	1.81%
3	1.15%
4	0.98%
5	1.45%
6	2.62%
7	5.17%
8	5.23%
9	4.09%
10	4.23%
11	4.89%
12	5.78%
13	5.78%
14	6.55%
15	7.75%

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r	
16	7.33%
17	7.91%
18	6.30%
19	4.43%
20	3.49%
21	3.28%
22	2.86%
23	2.38%
(null)	0.81%

# Data processing

We made modifications based on the following feedbacks:

-- 'How are you generated identifiers: e.g. pcf\_num, loc\_num, con\_num, etc.?'

For identifiers of pcf\_num, ve\_num and loc\_num (now con\_num has been deleted), we used Pandas to generate indexes on the dataframes (reset\_index(drop=True)). Thus the indexes are all natural numbers and unique. We generated these indexes to serve as keys for respective entities.

-- 'How are you handling attributes with missing values, if any? Especially, the attributes that are important and shouldn't/can't be simply ignored.'

Case\_id, party\_id and victim\_id do not and should not have null values.

For road\_en, we first extracted all distinct road\_en values and dropped the null value, then joined it with the 2 columns of road condition in the *collisions2018.csv* to get the not null (case\_id, road\_con) pairs, and dropped the duplicates. In this way if a collision's road condition is unavailable, it will not appear in the road\_en table, and all the pairs are unique. The same method also works for weather\_en, other\_fac\_en, safety\_p and safety\_v.

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For vehicle, we first extracted all the distinct tuples of their attributes, and we did not drop the null values, so there will not be null value in the party\_involve table for ve\_num, but there is a row of "ve\_num | null | null | null | null | null | is the vehicle table. It is the same for location and PCF table.

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# Deliverable 3

# **Assumptions**

Based on our assumptions on M1 and M2, we made the following further assumptions.

-- 'I just want to make sure that the duplicates are not because of the datatype issue (you are modeling case\_id as integer), and are genuine duplicates. Please reanalyze this part.'

As stated in previous parts, we changed the type of case\_id into string now, which fixed this problem. The duplicates are indeed caused by the datatype issue, and now the problem has been fixed.

-- (Feedback on M1) "Party\_number refers to the specific party of a particular case, so party\_number + case\_id is unique for each party, playing the same role as party\_id." \* Have you verified this from the data?'

Furthermore, after fixing this problem, it supports our earlier assumption "case\_id+party\_num" uniquely identifies 'party id'.

# **Query Implementation**

# Query a:

# Description of logic:

We first select parties whose "party\_type" is 'driver' and the "age" is not null. Then we just project columns of "age" and "at\_fault" and classify them into different age groups. If the party is at fault, the value of "at\_fault" is 1, otherwise it is 0 (if it's not null), so we can just calculate the sum of "at\_fault" of each age group, and then divide the sum by the count of each age group to get the ratio

#### SQL statement

**SELECT** 

(CASE WHEN AGE<=18 THEN 'underage'

WHEN AGE BETWEEN 19 and 21 then 'young I'

when AGE BETWEEN 22 AND 24 THEN 'young II'

when AGE BETWEEN 24 AND 60 THEN 'adult'

when AGE BETWEEN 61 AND 64 THEN 'elder I'

else 'elder II' END) as AGE\_RANGE,

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CONCAT(100\*ROUND(SUM(AT\_FAULT)/COUNT (AT\_FAULT), 3), '%') as ratio

FROM (SELECT AGE, AT\_FAULT FROM PARTY\_INVOLVE WHERE PARTY\_TYPE like '%driver%' and AGE IS NOT NULL and AT\_FAULT IS NOT NULL)

group by (CASE WHEN AGE<=18 THEN 'underage'

WHEN AGE BETWEEN 19 and 21 then 'young I'

when AGE BETWEEN 22 AND 24 THEN 'young II'

when AGE BETWEEN 24 AND 60 THEN 'adult'

when AGE BETWEEN 61 AND 64 THEN 'elder I'

ELSE 'elder II' END)

order by RATIO DESC

# Query result (if the result is big, just a snippet)

AGE_RANGE	RATIO
underage	64.70%
young I	58%
young II	51.90%
elder II	50.50%
adult	41.00%
elder I	40.10%

We found the underage, young I and young II group has the 3 highest at\_fault ratio. If we were an insurance company, I will raise premiums for young people under 24 years old and the elder over 65 years old.

#### Query b:

# Description of logic:

We join the party\_involve table with vehicle on ve\_num to get the vehicle details for each party, then we join it with road\_en (whose road\_con contains "holes") on case\_id to get the road\_condition of the case that the party is involved. If in a case there are vehicles with the same type, we count 1 for how

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many collisions this vehicle\_type participates. Finally we group the cases based on the ve\_type and calculate the total number of cases of each group, order them by the count and fetch the top 5 type.

#### SQL statement

SELECT VE\_TYPE, COUNT(\*) AS COUNT

**FROM** 

(SELECT DISTINCT VEHICLE.VE\_TYPE, PARTY\_INVOLVE.CASE\_ID

FROM PARTY\_INVOLVE, VEHICLE, ROAD\_EN

WHERE ROAD\_EN.ROAD\_CON LIKE '%holes%'

AND VE\_TYPE IS NOT NULL

AND ROAD\_EN.CASE\_ID = PARTY\_INVOLVE.CASE\_ID

AND VEHICLE.VE\_NUM = PARTY\_INVOLVE.VE\_NUM)

GROUP BY VE\_TYPE

ORDER BY COUNT DESC

FETCH FIRST 5 ROWS ONLY

# Query result (if the result is big, just a snippet)

VE_TYPE	COUNT
passenger	6940
car	0340
pickup or	
panel	2017
truck	
motorcycle	432
or scooter	452
bicycle	417
truck or	
truck	351
tractor	331
with trailer	

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## Query c:

# Description of logic:

We first select the victims whose deg\_injury is like "severe injury" or "killed", and then join it with the table party\_involve on the party\_id and then table vehicle (select those whose ve\_make is not null) on ve\_num to get [vic\_id|ve\_make]. We group the table by ve\_make and count the number of victims of each ve\_make. Finally we order the table by the count and fetch the top-10 vehicle makes and their number of victims.

#### SQL statement

SELECT VE\_MAKE, COUNT(ASSOCIATE\_VICTIM.VIC\_ID) AS COUNT

FROM VEHICLE, PARTY\_INVOLVE, ASSOCIATE\_VICTIM

WHERE VEHICLE.VE\_NUM = PARTY\_INVOLVE.VE\_NUM

AND ASSOCIATE\_VICTIM.PARTY\_ID = PARTY\_INVOLVE.PARTY\_ID

AND (ASSOCIATE\_VICTIM.DEG\_INJURY LIKE '%severe injury%' OR ASSOCIATE\_VICTIM.DEG\_INJURY LIKE '%killed%')

AND VEHICLE.VE\_MAKE IS NOT NULL

GROUP BY VE\_MAKE

ORDER BY COUNT DESC

FETCH FIRST 10 ROWS ONLY

# Query result (if the result is big, just a snippet)

VE_MAKE	COUNT
FORD	13924
HONDA	12060
TOYOTA	10639
CHEVROLET	10418
NISSAN	3860
DODGE	3641

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NOT	3603	
STATED	3003	
HARLEY-	3410	
DAVIDSON	3410	
SUZUKI	2482	
YAMAHA	2105	

## Query d:

## Description of logic:

We assume that the fraction is calculated by: ( # victims with no injury and of one seat place) / (# victims of this seat place). We don't consider the number of collisions because if there are 2 victims seating in same/different position in one collision, it is ambiguous to count for 1 or 2 collisions, and the collision can be of several seating positions.

We partition the table by vic\_seat, and add a column for count of victims of one kind of vic\_seat. Then we select the tuples whose deg\_injury is like "no injury", and calculate the sum of 1/count of each vic\_seat group. The denominator is the total count and the nominator is the count of the no injured, so we get the safety\_index of each vic\_seat. To get the max and min safety\_index and their corresponding seat, we union two tables, one is for the min and one is for the max.

#### SQL statement

(SELECT vic\_seat, ROUND(SUM(1.0/denom),2) AS SAFETY\_INDEX

**FROM** 

(SELECT VIC\_SEAT, DEG\_INJURY, COUNT(DEG\_INJURY) OVER (PARTITION BY VIC\_SEAT) AS DENOM

FROM ASSOCIATE VICTIM)

WHERE DEG\_INJURY LIKE '%no injury%'

GROUP BY VIC SEAT

ORDER BY SAFETY\_INDEX DESC

FETCH FIRST 1 ROWS ONLY)

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**UNION ALL** 

(SELECT vic\_seat, ROUND(SUM(1.0/denom),2) AS SAFETY\_INDEX

**FROM** 

(SELECT VIC\_SEAT, DEG\_INJURY, COUNT(DEG\_INJURY) OVER (PARTITION BY VIC\_SEAT) AS DENOM

FROM ASSOCIATE VICTIM)

WHERE DEG\_INJURY LIKE '%no injury%'

**GROUP BY VIC SEAT** 

ORDER BY SAFETY\_INDEX

FETCH FIRST 1 ROWS ONLY)

# Query result (if the result is big, just a snippet)

VIC_SEAT	SAFETY_INDEX
5	0.839
1	0.009

#### Query e:

#### Description of logic:

We assume that the query is to find ve\_types that participate in at least 10 collisions of EACH city and in at least half of the cities. We join the vehicle and party to know what type of the vehicle every party is driving, and then join it with table case and location to know in which city the collisions happen. There may be two same ve\_type in the same collision, and we count 1 for the #collision of this ve\_type, so for each [ve\_type, county\_city], we calculate the total number of collisions as

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CASE\_NUM. After selecting [ve\_type, county\_city] tuples where the CASE\_NUM >= 10, we group the table by ve\_type and select the number of those whose count(county\_city) >= total number of cities / 2

#### SQL statement

SELECT COUNT(VE\_TYPE) AS NUM\_VE\_TYPE

**FROM** 

(SELECT VE\_TYPE, COUNT(COUNTY\_CITY) AS CITY\_NUM

**FROM** 

(SELECT DISTINCT VE\_TYPE, COUNTY\_CITY, COUNT(\*) AS CASE\_NUM

FROM VEHICLE, PARTY\_INVOLVE, CASE, LOCATION

WHERE VEHICLE.VE\_NUM=PARTY\_INVOLVE.VE\_NUM

AND PARTY\_INVOLVE.CASE\_ID=CASE.CASE\_ID

AND LOCATION.LOC\_NUM=CASE.LOC\_NUM

GROUP BY VE\_TYPE, COUNTY\_CITY)CITY

WHERE CASE\_NUM >=10

GROUP BY VE\_TYPE) COUNT\_CITY,

(SELECT COUNT(DISTINCT COUNTY\_CITY) AS TOTAL\_CITY FROM LOCATION) TOTAL

WHERE CITY NUM>= TOTAL CITY/2

# Query result (if the result is big, just a snippet)

NUM_	VE_	TYPE
13		

#### Query f:

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# Description of logic:

To find the most populated cities, we use the order by function defined by ourselves to order the population level of cities descendingly, and fetch the first 3 distinct county\_city and the corresponding population. In our model, each case has a loc\_num, but several location entries can have the same county\_city, so we need to join the 3-row table with the location and the case table to get [county\_city|population|case\_id], and then join it with party\_involve and associate\_victim table to get [county\_city|population|case\_id|vic\_d|vic\_age]. We group the table over case\_id and add a column of average victim age of each case. Then we partition the table over county\_city, and add a row\_number() column in the order of victim age, and select the youngest 10.

#### SQL statement

```
SELECT COUNTY CITY, POPULATION, CASE ID, AVG AGE
  (SELECT COUNTY_CITY, POPULATION, CASE_ID, AVG_AGE, ROW_NUMBER()
OVER(PARTITION BY COUNTY CITY ORDER BY AVG AGE ASC) AS ROW NUMBER
 FROM
    (SELECT COL_CITY.COUNTY_CITY, COL_CITY.POPULATION, CASE.CASE_ID,
AVG(VIC AGE) AS AVG AGE
   FROM
     (SELECT DISTINCT COUNTY_CITY, POPULATION
       FROM LOCATION
       order by (case population
         when 7 then 0
         when 6 then 1
         when 5 then 2
         when 4 then 3
         when 3 then 4
         when 2 then 5
         when 1 then 6
         when 9 then 7
         when 0 then 8
         else 9 end)
       FETCH FIRST 3 ROWS ONLY) COL CITY,
       LOCATION, CASE, PARTY_INVOLVE, ASSOCIATE_VICTIM
   WHERE COL_CITY.county_city=LOCATION.county_city
     AND CASE.LOC_NUM=LOCATION.LOC_NUM
     AND PARTY INVOLVE.CASE ID=CASE.CASE ID
     AND ASSOCIATE VICTIM.PARTY ID=PARTY INVOLVE.PARTY ID
     AND VIC AGE IS NOT NULL
     GROUP BY COL_CITY.COUNTY_CITY, COL_CITY.POPULATION,
CASE.CASE_ID)AGE)RANK_AGE
WHERE ROW NUMBER<=10
ORDER BY COUNTY_CITY, ROW_NUMBER
```

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# Query result (if the result is big, just a snippet)

COUNTY_CITY	POPULATION	CASE ID	AVG_AGE
109	7	1050440	0
109	7	2858865	0
109	7	2489862	0
109	7	1572294	0
109	7	739484	0
109	7	491424	0
109	7	3553733	0
109	7	3486455	0
109	7	3012670	0
109	7	2856537	0
3801	7	984042	0
3801	7	2945655	0
3801	7	2613850	0
3801	7	1686308	0
3801	7	1783897	0
3801	7	353302	0
3801	7	2903949	0
3801	7	1895167	0
3801	7	1490755	0
3801	7	994429	0
4313	7	1702989	0
4313	7	2530732	0
4313	7	2427381	0
4313	7	2034278	0
4313	7	1585055	0
4313	7	3391590	0
4313	7	2034270	0
4313	7	1506715	0
4313	7	1322740	0
4313	7	1239869	0

# Query g:

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# Description of logic:

We first select case with col\_type like "pedestrian" to get the "PE\_CASE" table, then we join PE\_CASE with party\_involve and associate\_victim table to get the corresponding victims of the case. We select the tuples that satisfies the min(vic\_age) > 100 and find the eldest victim age of each collision.

#### SQL statement

SELECT DISTINCT CASE\_ID, MAX(VIC\_AGE) OVER(PARTITION BY CASE\_ID) AS ELDEST

**FROM** 

(SELECT PE\_CASE.CASE\_ID, VIC\_AGE, MIN(VIC\_AGE) OVER (PARTITION BY PE\_CASE.CASE\_ID)AS MIN\_AGE

**FROM** 

(SELECT CASE\_ID

FROM CASE

WHERE COL TYPE LIKE '%pedestrian%') PE CASE, PARTY INVOLVE, ASSOCIATE VICTIM

WHERE PARTY\_INVOLVE.CASE\_ID=PE\_CASE.CASE\_ID

AND ASSOCIATE VICTIM.PARTY ID=PARTY INVOLVE.PARTY ID)CASE AGE

WHERE MIN\_AGE>100

ORDER BY CASE\_ID

# Query result (if the result is big, just a snippet)

CASE_ID	ELDEST
36446	110
69198	101

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439197	102
445265	101
566220	102
644226	103
784061	102
817210	102
820619	101
828116	102
851026	106
868472	103
1209166	101
1213340	121
1347636	101
1373664	101
1548445	102

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1847678	104
2472739	103
2531557	103

# Query h:

## Description of logic:

We first group the party\_involve table by the ve\_num to get [ve\_num | N\_COLLISION] table VE\_COL, then we join the VE\_COL with the vehicle table on the ve\_num to get how many collisions a kind of vehicle has participated in. We select those who has participated in more than 9 collisions. Because we assign each vehicle with same [ve\_type, ve\_make, ve\_year] a unique ve\_num, so we finally use [ve\_type, ve\_make, ve\_year] to denote the vehicle id.

#### SQL statement

```
WITH

VE_COL AS

(

SELECT VE_NUM, COUNT(case_id) AS N_COLLISION

FROM PARTY_INVOLVE

GROUP BY VE_NUM
)
```

SELECT VEHICLE.VE\_TYPE || ', ' || VEHICLE.VE\_MAKE || ', ' || VEHICLE.VE\_YEAR AS VE\_ID, VE\_COL.N\_COLLISION AS N\_COLLISION

FROM (VEHICLE INNER JOIN VE\_COL ON VE\_COL.VE\_NUM = VEHICLE.VE\_NUM)

WHERE N\_COLLISION > 9 AND VEHICLE.VE\_MAKE IS NOT NULL AND VEHICLE.VE\_TYPE IS NOT NULL AND VEHICLE.VE\_TYPE != 'pedestrain'

ORDER BY N\_COLLISION DESC

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# FETCH FIRST 20 ROWS ONLY

# Query result (if the result is big, just a snippet)

	1
VE_ID	N_COLLISION
passenger car, TOYOTA, 2000	52469
passenger car, FORD, 2000	51915
passenger car, HONDA, 2000	50255
passenger car, FORD, 1998	49148
passenger car, TOYOTA, 2001	47205
passenger car, HONDA, 2001	45243
passenger car, FORD, 2001	45203
passenger car, TOYOTA, 1999	42907
passenger car, HONDA, 1998	42063
passenger car, FORD, 1999	41918
passenger car, FORD, 1995	40225
passenger car, HONDA, 1997	39182
passenger car, FORD, 1997	38852
passenger car, HONDA, 1999	38528
passenger car, TOYOTA, 2002	38403
passenger car, TOYOTA, 1998	37982
passenger car, TOYOTA, 1997	37134
passenger car, TOYOTA, 2003	35919
passenger car, HONDA, 2002	35759
passenger car, FORD, 2002	35433

#### Observation:

The vehicles participated in most collisions are all passenger cars. Besides, these vehicles are all made by large car manufactures such as Totota, Ford and Honda.

# Query i:

# Description of logic:

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Because we assign each kind of location a loc\_num but several loc\_num have the same county\_city, so we first join location and case, then group the table by "county\_city" to count how many collisions happen in the same county\_city (if it is not null). Then we order them by number of collisions and fetch the top-10 county\_city.

#### SQL statement

SELECT COUNTY\_CITY, COUNT(CASE\_ID) as N\_COLLISION FROM

LOCATION, CASE

WHERE LOCATION.LOC\_NUM = CASE.LOC\_NUM and LOCATION.COUNTY\_CITY IS NOT NULL

**GROUP BY COUNTY\_CITY** 

ORDER BY N\_COLLISION DESC

FETCH FIRST 10 ROWS ONLY

# Query result (if the result is big, just a snippet)

COUNTY_CITY	N_COLLISION
1942	399582
1900	118446
3400	80191
3711	76867
109	72995
3300	61453
3404	58068
4313	57852
1941	53565
3801	48450

#### Query j:

#### Description of logic:

We focus on the "group by" function. First we group by lighting of the case, there are "DAY", "NIGHT" and others, in the dataset, lighting conditions of "dawn and dusk" are ambiguous, so we just distinguish them by the collaboration date and collaboration is not "DAY" or

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"NIGHT", we group them by the date and time, based on different seasons. After grouping, we count the number of collisions of each lighting condition, and order them by the count.

# SQL statement **WITH** COL\_PERIOD AS ( **SELECT** CASE WHEN CASE.LIGHTING LIKE '%day%' THEN 'DAY' WHEN CASE.LIGHTING LIKE '%dark%' THEN 'NIGHT' **ELSE** (CASE WHEN EXTRACT(MONTH FROM COL\_TIME)>8 OR EXTRACT(MONTH FROM COL TIME)<4 THEN (CASE WHEN EXTRACT(HOUR FROM CAST(COL TIME AS TIMESTAMP))>5 AND EXTRACT(HOUR FROM CAST(COL TIME AS TIMESTAMP))<8 THEN 'DAWN' WHEN EXTRACT(HOUR FROM CAST(COL TIME AS TIMESTAMP))>17 AND EXTRACT(HOUR FROM CAST(COL\_TIME AS TIMESTAMP))<20 THEN 'DUSK' WHEN EXTRACT(HOUR FROM CAST(COL TIME AS TIMESTAMP))>7 AND EXTRACT(HOUR FROM CAST(COL TIME AS TIMESTAMP))<18 THEN 'DAY' ELSE 'NIGHT' END) WHEN EXTRACT(MONTH FROM COL\_TIME)>3 OR EXTRACT(MONTH FROM COL TIME)<9 THEN (CASE

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WHEN EXTRACT(HOUR FROM CAST(COL\_TIME AS TIMESTAMP))>3 AND EXTRACT(HOUR FROM CAST(COL\_TIME AS TIMESTAMP))<6 THEN 'DAWN'

WHEN EXTRACT(HOUR FROM CAST(COL\_TIME AS TIMESTAMP))>19 AND EXTRACT(HOUR FROM CAST(COL\_TIME AS TIMESTAMP))<22 THEN 'DUSK'

WHEN EXTRACT(HOUR FROM CAST(COL\_TIME AS TIMESTAMP))>5 AND EXTRACT(HOUR FROM CAST(COL\_TIME AS TIMESTAMP))<20 THEN 'DAY'

ELSE 'NIGHT'

END)

**ELSE NULL** 

END)

**END AS PERIOD** 

FROM CASE

)

SELECT PERIOD, COUNT(\*) AS N\_COLLISION

FROM COL\_PERIOD

WHERE PERIOD IS NOT NULL

**GROUP BY PERIOD** 

ORDER BY N\_COLLISION DESC

# Query result (if the result is big, just a snippet)

PERIOD	N_COLLISION
DAY	2575153
NIGHT	1041614
DAWN	30317
DUSK	28092

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# Query Performance Analysis - Indexing

## Query 2:

<Initial Running time/IO: 3371ms/37073</p>

Optimized Running time/IO:1511ms/24804

**Explain the improvement:** 

Added the following index:

create index ve\_case\_ix on party\_involve(case\_id, ve\_num)

#### **Necessities:**

From the initial plan, we can find there are three tables joining: ROAD\_EN, VEHICLE, PARTY\_INVOLVE and each of them involves a full file scan. We only need [case\_id, ve\_num] of the party\_involve table, the whole content of the road\_en table, [ve\_num, ve\_type] of the vehicle table. Since road\_con and case\_id are both primary keys of the road\_en table, there is no need to add indexes in the road\_en tables. So we build indexes on [case\_id, ve\_num] of the party\_involve table to reduce full scan and optimize.

For this query, we added a joint index 've\_case\_ix' on <case\_id, ve\_num> on entity 'party\_involve'. In this case, when doing hash join, the system can do a fast full scan on 'party\_involve' instead of a full scan.

In our case, the index 've\_case\_ix' is relevant when we want to group the collisions according to the type of vehicles. So a fast full scan using a B-tree index is useful.

#### Influence on the cost:

After applying the two indexes, we could find that the total cost is reduced from 37073 to 24804. The cost of access of party\_involve is reduced from 19124 to 6855, and thus the hash join of party\_involve and road\_en is reduced, and then the hash join of vehicle and the party\_involve is also reduced due to the less cost of scanning the table party\_involve. Now the cost of scanning the party\_involve is still a big part of the total cost, but the cost of right semi join and the unique operation becomes the dominant part.

#### **Initial plan**

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OPERATION	OBJECT_NAME	OPTIONS	CARDINALITY COST	Γ
SELECT STATEMENT			5	37073
SORT		ORDER BY	5	37073
□ ■ VIEW	SYS. null		5	37072
⇔ <b>σ∀</b> Filter Predicates				
from\$_subquery\$_005.rowlimit_\$\$_rownumber<=5				
⇒- ● WINDOW		SORT PUSHED RANK	15	37072
⇒- <b>σÿ</b> Filter Predicates				
ROW_NUMBER() OVER ( ORDER BY COUNT(*) DESC )<=5				
⊟- ● HASH		GROUP BY	15	37072
⇒			301205	35412
⊟ ● HASH		UNIQUE	301205	35412
⊨ M HASH JOIN			301205	31478
□ On Access Predicates				
VEHICLE. VE_NUM=PARTY_INVOLVE. VE_NUM				
TABLE ACCESS	VEHICLE	FULL	21932	68
⊕ <b>O∀</b> Filter Predicates				
└── VE_TYPE IS NOT NULL				
⊞ ⋈ HASH JOIN		RIGHT SEMI	362485	31409
➡ <b>O</b> ™ Access Predicates				
ROAD_EN. CASE_ID=PARTY_INVOLVE. CASE_ID				
□ ■ TABLE ACCESS	ROAD EN	FULL	182607	3040
⊕ <b>O</b> ♥ Filter Predicates				
ROAD_EN.ROAD_CON_LIKE '%holes%'				
TABLE ACCESS	PARTY INVOLVE	FULL	7286606	19124

# Improved plan>

OPERATION	OBJECT_NAME	OPTIONS	CARDINALITY	COST	
B ● SELECT STATEMENT				5	24804
⊕ • SORT		ORDER BY		5	24804
UIEW VIEW	SYS. null			5	24803
⊕- <b>σ</b> ÿ Filter Predicates					
from\$_subquery\$_005.rowlimit_\$\$_rownumber<=5					
⊟- ● WINDOW		SORT PUSHED RANK		15	24803
Ģ- <b>σ∀</b> Filter Predicates					
ROW_NUMBER() OVER ( ORDER BY COUNT(*) DESC )<=5					
⊟- ● HASH		GROUP BY		15	24803
⇒ <b>■</b> VIEW				301205	23143
HASH		UNIQUE		301205	23143
⇒ ⋈ HASH JOIN				301205	19209
□ Om Access Predicates					
VEHICLE. VE_NUM=PARTY_INVOLVE. VE_NUM					
□ TABLE ACCESS	<u>VEHICLE</u>	FULL		21932	68
⊕ <b>ov</b> Filter Predicates					
- VE_TYPE IS NOT NULL					
⇒ MHASH JOIN		RIGHT SEMI		362485	19140
□ Om Access Predicates					
ROAD_EN. CASE_ID=PARTY_INVOLVE. CASE_ID					
TABLE ACCESS	ROAD EN	FULL		182607	3040
⊕ <b>O</b> ♥ Filter Predicates					
ROAD_EN.ROAD_CON LIKE '%holes%'					
□ o∉ INDEX	VE CASE IX	FAST FULL SCAN		7286606	6855

# Query 5:

< Initial Running time/IO: 21.494s/48986

Optimized Running time/IO: 14.703s /24096

# **Explain the improvement:**

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# Added the following index:

create index ve\_case\_ix on party\_involve(case\_id, ve\_num)
create index case\_loc\_ix on case(case\_id, loc\_num)
create index county\_city\_ix on location(loc\_num, county\_city)

#### **Necessities:**

From the initial plan, we can find there are four tables joining: PARTY\_INVOLVE, VEHICLE, CASE, LOCATION and each of them involves a full file scan. We only need [case\_id, ve\_num] of the party\_involve table, [ve\_num, ve\_type] of the vehicle table, [case\_id, loc\_num] of the case table and [loc\_num, county\_city] of the location table. So we build indexes on [case\_id, ve\_num] of the party\_involve table, indexes on [case\_id, loc\_num] of the case table, and indexes on [loc,num, county\_city] of the location table to reduce full scan and optimize.

For this query, we added a joint index 've\_case\_ix' on <case\_id, ve\_num> on entity 'party\_involve', and 'case\_loc\_ix' on <case\_id, loc\_num> on entity 'case'. In this case, when doing hash join, the system can do a fast full scan on 'party involve' and 'case' instead of a full scan.

The index 've\_case\_ix' is relevant when we want to group the collisions according to the type of vehicles, and the index 'case\_loc\_ix' is used for the grouping of collisions according to the county\_city code. In this case it is indeed useful to index the 'case\_id' according to 've\_num' and 'loc\_num' as a B-tree. It could indeed search the desired result binarily.

We can see that the index 've\_case\_ix' has been used twice for a fast full scan for the matching of vehicle and case, which reduces the loss greatly, and improves the quality of the query.

#### **Distribution of cost:**

Now for the optimized plan, the majority cost is spent on doing group by based on [ve\_type, county\_city] for 'vehicle' entity and the join of case and party\_involve table. After applying the three indexes, we could find that the total cost is reduced from 48986 to 24096. The cost of access of party\_involve is reduced from 19124 to 6855, and cost of accessing the case table is reduced from 16023 to 3406, then the cost of the join of them is also reduced. Due to the case\_loc\_ix index, the join of the location and the case is also reduced. The county\_city\_ix index plays a small role in the optimization.

#### Initial plan

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# Improved plan>

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OPERATION	OBJECT_NAME	OPTIONS	CARDINALITY	COST
⊟ ® SELECT STATEMENT			1	24096
- GORT		AGGREGATE	1	
□ MESTED LOOPS			1	24096
			1	. 6
⊟ • • SORT		AGGREGATE	1	
Ė∎ VIEW	SYS. VM NWVW 1		540	6
⊟ hash		GROUP BY	540	6
O€ INDEX	COUNTY CITY IX	FAST FULL SCAN	4362	5
□ VIEW			1	24090
□ OV Filter Predicates				
CITY_NUM>=TOTAL_CITY/2				
⊟ • • SORT		GROUP BY	15	24090
□···■ VIEW			287	24090
⊟ ● FILTER				
□ <b>○ ♥</b> Filter Predicates				
COUNT (*)>=10				
⊟ SORT		GROUP BY	287	24090
😑 🔀 hash join			7251269	23898
- On Access Pred	licates			
LOCATION.	LOC_NUM=CASE, LOC_NUM			
·····□를 INDEX	COUNTY CITY IX	FAST FULL SCAN	436	2 5
⊨ Mash join			7251269	23875
Access	Predicates			
	LE. VE_NUM=PARTY_INVOLVE. VI	E_NUM		
TABLE ACC	CEVEHICLE	FULL	26260	68
□ MASH JOI			7286600	23788
₽ On Acce	ess Predicates			
Т РА	ARTY_INVOLVE.CASE_ID=CASE.	CASE_ID		
₽- <b>M</b> neste	D		7286600	23788
	'A'			
	CASE LOC IX	FAST FULL SCAN	3678063	3406
<b></b>	DIVE CASE IX	RANGE SCAN		6855
- □	Access Predicates			
	PARTY_INVOLVE. CASE_II	D=CASE.CASE_ID		
o∉ INDEX	VE CASE IX	FAST FULL SCAN	7286600	6855

# Query 6:

< Initial Running time/IO: 5 measurements shows initial average run time is 3641.6 ms / 47735

Optimized Running time/IO: 5 measurements shows optimized running time is 2980 ms / 37102

# **Explain the improvement:**

## Added the following index:

# create index loc\_in\_case on case(loc\_num)

Necessities: From the initial plan, we can find there are five tables joining: COL\_CITY LOCATION, CASE, PARTY\_INVOLVE, ASSOCIATE\_VICTIM and each of them involves a full file scan. We only need [case\_id, loc\_num] of the case table, [[party\_id, case\_id] of the party\_involve table, [party\_id, vic\_age] of the associate\_victim table, [loc\_num, county\_city] of the location table, and we have selection predicates on these fields. So we try to build indexes on these columns to reduce full scan and optimize.

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We tried several combinations of the indexes of these columns, but finally we found this one is useful in the timing, though other keys like location(county\_city) (for selection during the join), party\_involve(case\_id, loc\_num) (for utilizing the clustered index) will also reduce the cost and change the plan by using index.

With the loc\_in\_case of the case table, the hash join of table case and location is completed by index, and it also utilizes the clustered case\_id row number as index, which is very nice. While joining and accessing the needed columns, they can just use the index range scan to check if the predicates are satisfied, thus the IO is reduced.

Influence on the cost: Cost is related to the CPU and IO. From the optimized plan, we can find that the total cost is reduced from 47735 to 37104, and the cost of case table access is reduced from 16023 to 225, and the cost of location access is reduced from 7 to 2, and the total cost of their join is reduced from 16048 to 5417 (this is why the total cost is reduced). There are also two hash joins becoming nested join, both because of the index scan. Maybe the optimizer chooses finds the indexed nested loop more efficient.

However, the cost of party\_involve full scan is still large and plays an important role in the total cost. We can reduce its cost by creating an index on party\_involve(party\_id, case\_id) and it indeed helps to reduce the cost, but because the running time suffers, we did not use it. From the plan we can also find that the group cost and join cost if dominated by the file access cost.

Initial plan

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ERATION	OBJECT_NAME	OPTIONS	CARDINALITY COST	
SELECT STATEMENT			21893	4773
⇒ ◆ SORT		ORDER BY	21893	4773
UIEW VIEW			21893	4737
□ <b>58</b> Filter Predicates				
ROW_NUMBER<=10				
⊕ WINDOW		SORT PUSHED RANK	21893	4737
⊕ <b>σ∀</b> Filter Predicates				
ROW_NUMBER() OVER ( PARTITION BY from\$_subquery\$_006.	.COUNTY_CITY ORDER BY SUM(	VIC_AGE)/COUNT(VIC_AGE))		
⊟- ● HASH		GROUP BY	21893	4737
⊨ MHASH JOIN			21893	4667
□ On Access Predicates				
ASSOCIATE_VICTIM. PARTY_ID=PARTY_INVOLVE. PARTY_I	D			
⇒ MHASH JOIN			40481	3514
□ Om Access Predicates				
PARTY_INVOLVE. CASE_ID=CASE. CASE_ID				
□ MHASH JOIN			20434	1604
□ On Access Predicates				
CASE. LOC_NUM=LOCATION. LOC_NUM				
⊟-MHASH JOIN			24	1
□ On Access Predicates				
from\$_subquery\$_006.COUNTY_CITY=LOCATION	_			
⇒ ■ VIEW	SYS. null		3	
□ OF Filter Predicates				
from\$_subquery\$_006.rowlimit_\$\$_rownu	mber<=3			
□ ■ WINDOW		SORT PUSHED RANK	3055	
⇒ ov Filter Predicates				
ROW_NUMBER() OVER ( ORDER BY from \$_		0) <=3		
⊟ VIEW	SYS. null		3055	
⊟ HASH		UNIQUE	3055	
TABLE ACCESS	LOCATION	FULL	4362	
TABLE ACCESS	LOCATION	FULL	4362	
TABLE ACCESS	CASE	FULL	3678063	1602
TABLE ACCESS	CASE	FULL	3678063	1602
TABLE ACCESS	PARTY INVOLVE	FULL	7286606	1908
□ ■ TABLE ACCESS	ASSOCIATE VICTIM	FULL	3940663	746

# Improved plan>

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OPERATION	OBJECT_NAME	OPTIONS	CARDINALITY	COST
⊞ SELECT STATEMENT			2189	93 37104
SORT		ORDER BY	2189	93 37104
□ ■ VIEW			2189	93 36745
⇔ σ♥ Filter Predicates				
ROW_NUMBER<=10				
= WINDOW		SORT PUSHED RANK	2189	36745
GV Filter Predicates				
ROW_NUMBER() OVER ( PARTITION BY from\$_subquery\$_006.	COUNTY_CITY ORDER BY SUM	(VIC_AGE) /COUNT (VIC_AGE)	<=10	
⊟- ● HASH		GROUP BY	2189	36745
⊟-MHASH JOIN			2189	36042
- <b>O</b> Access Predicates				
ASSOCIATE_VICTIM. PARTY_ID=PARTY_INVOLVE. PARTY_II	)			
HASH JOIN			4048	24516
PARTY_INVOLVE. CASE_ID=CASE. CASE_ID				
□ NESTED LOOPS			2043	5417
□ M NESTED LOOPS			2043	5417
⇒ MHASH JOIN			2	24 16
GN Access Predicates				
from \$_subquery \$_006. COUNTY_CITY=LOCATI	ON. COUNTY_CITY			
. □ VIEW	SYS. null			3
Ģ <b>σ∀</b> Filter Predicates				
from\$_subquery\$_006.rowlimit_\$\$_row	number<=3			
⊟-● WINDOW		SORT PUSHED RANK	305	55
ভ তেই Filter Predicates				
ROW_NUMBER() OVER ( ORDER BY from	\$_subquery\$_005.rowlimit_	.\$_0) <=3		
□ ■ VIEW	SYS. null		305	55
⊟- ● HASH		UNIQUE	305	
TABLE ACCESS	LOCATION	FULL	436	-
TABLE ACCESS	LOCATION	FULL	436	32
⊟-o∉ INDEX	LOC IN CASE	RANGE SCAN	84	43 2
□ O <sub>N</sub> Access Predicates				
➡ <b>O</b> ™ Access Predicates				
CASE, LOC_NUM=LOCATION, LOC_NUM				
TABLE ACCESS	CASE	BY INDEX ROWID	84	3 225
TABLE ACCESS	PARTY INVOLVE	FULL	728660	6 19080
□-⊞ TABLE ACCESS	ASSOCIATE VICTIM	FULL	394066	3 7460
□ <b>♂♥</b> Filter Predicates				
VIC AGE IS NOT NULL				

#### Query 8:

<Initial Running time/IO: 2.478s / 19386</p>

Optimized Running time/IO: 1.492s / 5998

# **Necessity of indexing:**

This query requires us to a group by operation on the 've\_num' attribute on 'party\_involve' entity. However, there are no indexes for the 've\_num' attribute for 'party\_involve', and group by in this case can consume much cost. Therefore, it is necessary to add an index to 've\_num' for 'party\_involve'.

# **Explain the improvement:**

# Added the following index:

create index ve\_party\_ix on party\_involve(party\_id, ve\_num)

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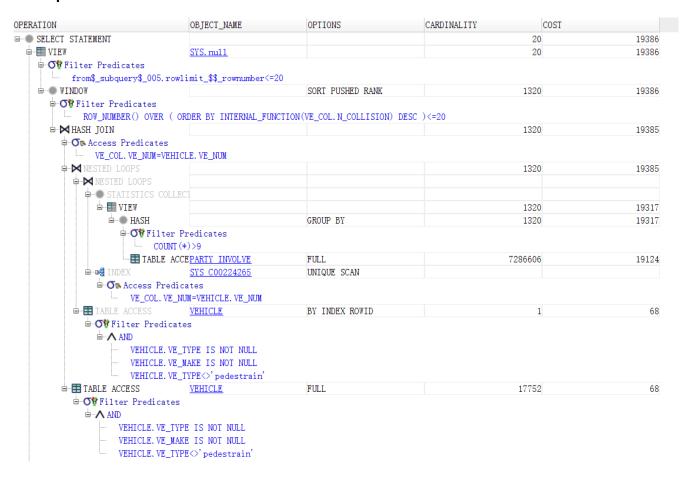


For this query, we added a joint index 've\_party\_ix' on <party\_id, ve\_num> on entity 'party\_involve'. In this case, when doing group by, the system can use the 've\_party\_ix', which reduces the cost for group by from 19317 to 5929.

#### **Distribution of cost:**

Now for the optimized plan, the majority cost is spent on doing group by based on 've\_num' for 'party\_involve' entity. Little cost is spent on accessing the 'vehicle' entity (because it can access by its primary key, the index).

# Initial plan



#### Improved plan>

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PERATION	OBJECT_NAME	OPTIONS	CARDINALITY COST	
∍ ● SELECT STATEMENT			20	5998
Ģ…≣ VIEW	SYS. nu11		20	5998
⊕ <b>σ</b> ♥ Filter Predicates				
from\$_subquery\$_005.r	owlimit_\$\$_rownumber<=	<del>:</del> 20		
- WINDOW		SORT PUSHED RANK	1320	5998
⊖ <b>O∀</b> Filter Predicates				
ROW_NUMBER() OVER	ORDER BY INTERNAL_FU	NCTION(VE_COL.N_COLLISION)	DESC )<=20	
⊨ MHASH JOIN			1320	5997
₽ <b>७</b> ⋒ Access Predicates				
VE_COL. VE_NUM=VE	HICLE. VE_NUM			
₽ MESTED LOOPS			1320	5997
₽ MESTED LOOPS				
	LECI			
⇒ ■ VIEW			1320	5929
⊟ ● HASH		GROUP BY	1320	5929
	r Predicates			
COUN				
		FAST FULL SCAN	7286606	5736
⊟o INDEX		UNIQUE SCAN		
□ <b>O</b> ™ Access Pre				
	E_NUM=VEHICLE.VE_NUM			
□ TABLE ACCESS		BY INDEX ROWID	1	68
□ <b>O</b> ♥ Filter Predic	cates			
□ ∧ AND	TE TUDE TO NOT MILL			
	/E_TYPE IS NOT NULL			
	/E_MAKE IS NOT NULL			
TABLE ACCESS	/E_TYPE<>'pedestrain' VEHICLE	FULL	17752	68
□ OV Filter Predicat		FULL	17792	08
□ AND	es			
* * * * * * * * * * * * * * * * * * * *	TYPE IS NOT NULL			
_	MAKE IS NOT NULL			
	MARE IS NOT NOLL  TYPE<>'pedestrain'			
VEHICLE. VE_	TIPEN/ pedestrain			

# Query 9:

<Initial Running time/IO: 4.343s / 16227</p>

Optimized Running time/IO: 1.457s / 2260

# **Necessity of indexes:**

Here in the query, we require the system to select from two different entities, given an equality condition 'loc\_num'. Hence, the system has to do a join on the two entities. However, it only has a primary key 'loc\_num' for 'location' entity, and has no indexes for the 'case' entity. In this case, adding an index to 'loc num' for 'case' entity is necessary to reduce the costs for the join.

# **Explain the improvement:**

# Added the following indexes:

create index loc\_in\_case on case(loc\_num)

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For this query, we added an index 'loc\_in\_case' on 'loc\_num' for 'case' entity.

We can see from the query plan that the index 'loc\_in\_case' has been used twice.

Based on these indexes, when the system is doing hash join for 'location' entity and 'case' entity on 'location.loc\_num = case.loc\_num' (to find the name of the cities respective to the location numbers), it can do a fast full scan on 'case' entity, instead of a full scan, which greatly reduces the cost from 16023 to 2056.

#### Distribution of cost:

From the optimized query plan, we can see very little cost is spent on scanning the table 'location', the majority cost is spent on doing range scan on 'case' tabel, fulfilling the equality constraint. Despite this, it still saves much cost compared to the original plan.

#### Initial plan

OPERATION	OBJECT_NAME	OPTIONS	CARDINALITY	COST
□ SELECT STATEMENT			10	16227
SORT   SORT		ORDER BY	10	16227
Ė≣ VIEW	SYS. null		10	16226
from\$_subquery\$_003.rd	owlimit_\$\$_rownumber	<=10		
WINDOW		SORT PUSHED RANK	540	16226
🗦 🔿 Filter Predicates				
ROW_NUMBER() OVER (	ORDER BY COUNT (*)	ESC )<=10		
⊟ ● HASH		GROUP BY	540	16226
⊨ ⋈ HASH JOIN			3678063	16039
□ <b>O</b> ® Access Predicate	es			
LOCATION. LOC_N	UM=CASE.LOC_NUM			
□ ■ TABLE ACCESS	<u>LOCATION</u>	FULL	4362	7
Ġ <b>σ∀</b> Filter Predic	ates			
LOCATION. COU	JNTY_CITY IS NOT NUL	L		
TABLE ACCESS	<u>CASE</u>	FULL	3678063	16023

Improved plan>

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OPERATION	OBJECT_NAME	OPTIONS	CARDINALITY	COST
■ SELECT STATEMENT			10	2258
		ORDER BY	10	2258
Ė…∰ VIEW	SYS. nu11		10	2257
from\$_subquery\$_	003.rowlimit_\$\$_rownumber	<=10		
□···· ● WINDOW		SORT PUSHED RANK	540	2257
⇔ <b>σÿ</b> Filter Predicat	es			
ROW_NUMBER() C	OVER ( ORDER BY COUNT(*) D	DESC )<=10		
⊟ HASH		GROUP BY	540	2257
⊨ MHASH JOIN			3678063	2070
₽ <b>O</b> ™ Access Pre	dicates			
LOCATION.	LOC_NUM=CASE.LOC_NUM			
₽ MESTED LOOP	S		3678063	2070
- STATISTIC	CS COLL			
i od INDEX	COUNTY CITY IX	FAST FULL SCAN	4362	5
⊟ <b>o</b> †Fi1	lter Predicates			
L	OCATION. COUNTY_CITY IS NO	T NULL		
≐• <b>₫</b> INDEX	LOC IN CASE	RANGE SCAN	843	2056
⊕ <b>o</b> Acces	s Predicates			
LOCA	ATION.LOC_NUM=CASE.LOC_NUM	I		
	LOC IN CASE	FAST FULL SCAN	3678063	2056

# **General Comments**

In this project, we finished the design of our database system, from designing the ER diagram, preparing data to implement the queries, and optimizing the execution. We successfully compelted the tasks requried in the project, a very comprehensive and practical project.

Many thanks for the help from the professors and TAs!