



- 6.1 匿名的基础知识
- 6.2 如何对比特币去匿名化
- 6.3 混币
- 6.4 分布式混币
- 6.5 零币和零钞



匿名(Anonymity): 无关联性的化名

- 比特币系统中,使用者不需要使用真实的姓名
- 需要使用公钥哈希值作为交易标识
- 一个用户可以随机创建出任意多个比特币地址



■ 比特币具有化名性,但是*不能达到绝对 隐私* 

●使用数字货币(如比特币)支付时,在 真实的物理世界里容易暴露身份,进而 关联到地址,以及其他所有的交易



#### 无关联性

- ■同一个用户的不同地址应该不易关联
- ■同一个用户的不同交易应该不易关联
- 一个交易的交易双方应该不易关联



区块链货币中,所有交易都记录在一个 公开账本上,也就是说相关交易信息可 以*永久追踪* 

- ✓ 希望能够达到传统银行能够达到的隐私保护级别,降低公共区块链带来的信息暴露风险
- ✓ 超越传统银行给我们的隐私保护级别



匿名化和去中心化

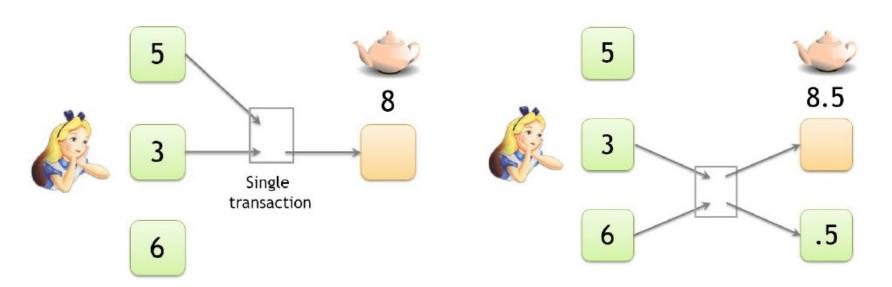
• Chaum的电子现金系统,采用了*盲签名* 技术,但还需一个*中央权威机构* 

Zerocoin, Zerocash: 匿名化&去中心化 的加密数字货币系统



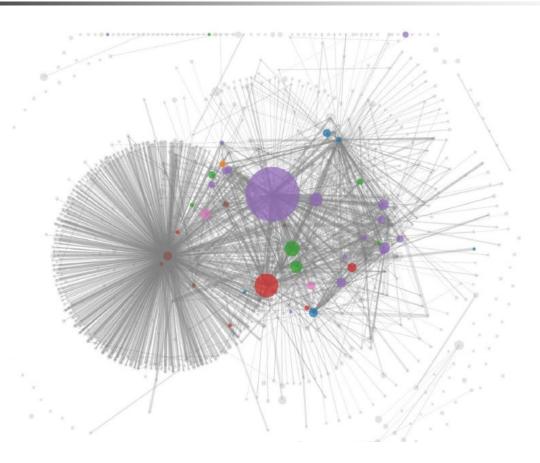
# 6.2 如何对比特币去匿名化

■ 关联性



容易暴露零钱地址

# 6.2 如何对比特币去匿名化

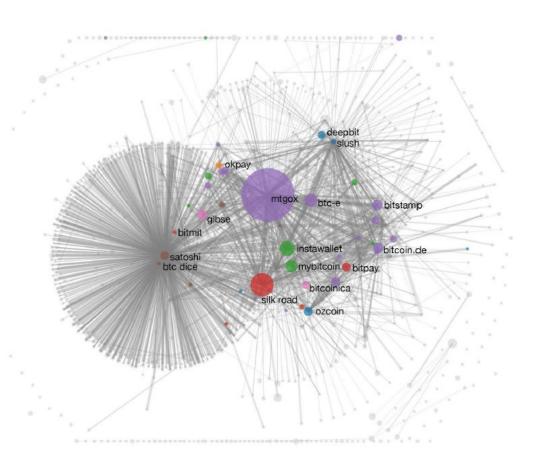


Clustering of addresses: Characterizing Payments Among Men with No Names

# 交易图谱分析

# 6.2 如何对比特币去匿名化

■利用交易进行标记



# 6.2 如何对比特币去匿名化

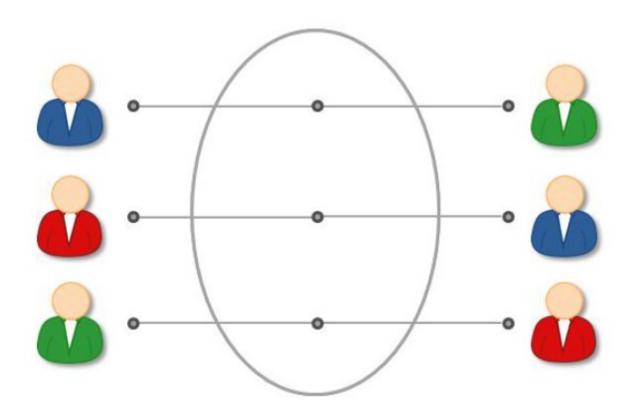
■网络层的去匿名化



"the first node to inform you of a transaction is probably the source of it."

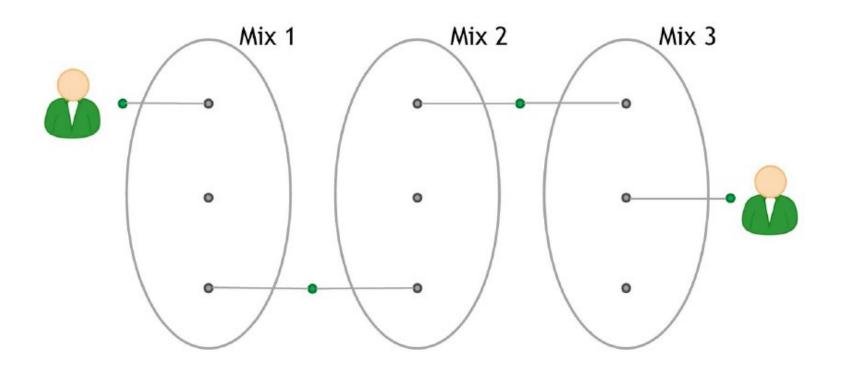


■ 想要匿名化,使用一个中介媒体



# 6.3 混币

多重混币





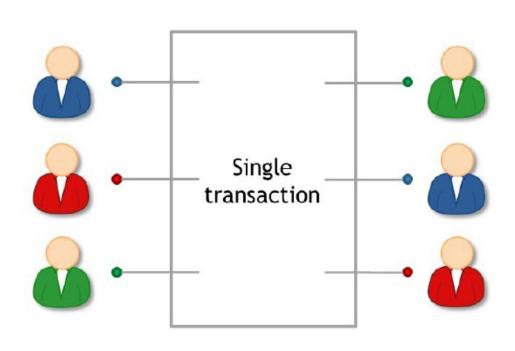
# 6.4 分布式混币

■ 分布式混币(Decentralized Mixing)

■ 采用一种用户之间的**点对点**模式实现混 币交易的协议

# 6.4 分布式混币

攻击者无法建立输入和输出的匹配关系



A Coinjoin transaction



# 6.4 分布式混币

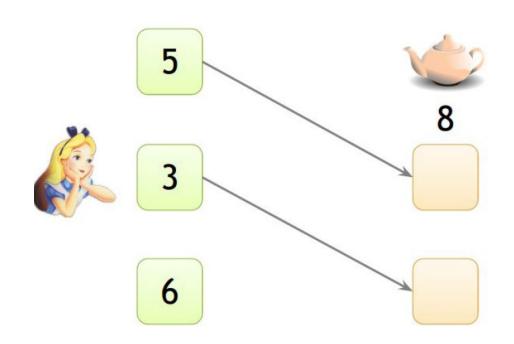
#### 高交易风险流

为了完成一笔支付,用户通常会组合所 拥有的数字货币,这样便有足够数额可 以支付到单一接收地址

■ 规避: 所有输入地址被关联在一起

# 4

# 6.4 分布式混币



Merge avoidance



# 零知识证明

Zero Knowledge Proof: 证明者(Prover)要让验证者(Verifier)相信自己拥有某种知识,但又不泄漏它

■ 很多场景中有着广泛的应用,比如金融 交易中,保护支付方、接收方、交易金 额的**隐私** 

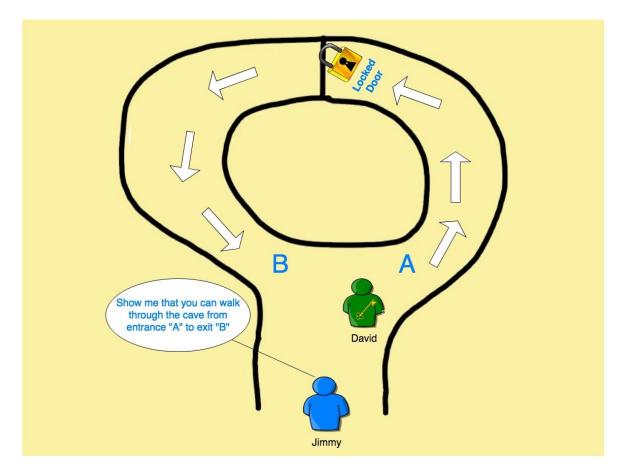


Example 1. A Key to a Door

Example 2. Coloring Problem



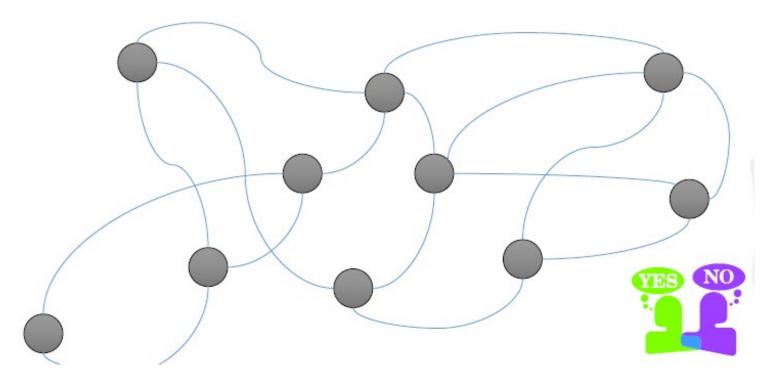
# Example 1





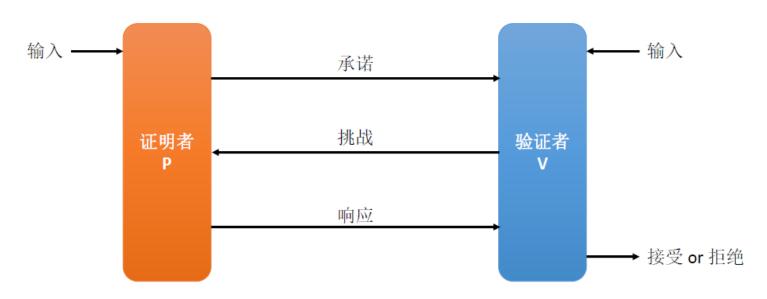
# Example 2

#### Can we label this graph with {r, g, b}?



# 交互式零知识证明的一般模型

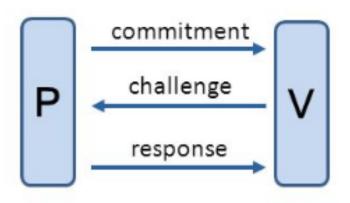
#### 交互式零知识证明的一般模型



- 证明者和验证者共享一个公共输入,证明者可能拥有某个秘密输入。
- 如果验证者认可证明者的响应,则输入接收(Accept); 否则,输出拒绝(Reject)。



#### Sigma-protocols

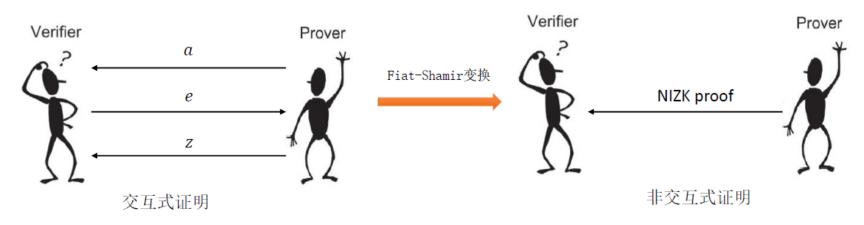


- P sends V a message a
- V sends P a random t-bit string e
- P sends a reply z, and V decides to accept or reject based solely on the data it has seen; i.e., based only on the values (x, a, e, z).

# 非交互式零知识证明

Fiat-Shamir变换是一种可以将Sigma协议变成非交互证明的技术。它能够让证明者Prover可以通过给验证者Verifier发送一个证明信息即可完成证明(无需交互,无需返回挑战)。

而且,它能把任何一个Simga协议变成一个数字签名,签名的含义就是"知道这个Sigma协议的秘密的人已经签署了这个消息"。Prover能够创造一个证明,然后分发给很多个验证者,验证者可以不必联系Prover即可验证证明有效性。同时零知识也变得容易了,因为验证者或者其他敌手不能做任何事情。



个密码学安全的 Hash 函数可以近似地模拟传说中的「随机预言机



#### Zerocoin:

```
I know x such that H(x \mid | \langle other known inputs \rangle) < \langle target \rangle.
```

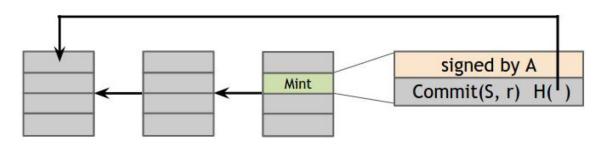
"I know x such that H(x) belongs to the following set:  $\{...\}$ ".

The proof would reveal nothing about x



#### Zerocoin

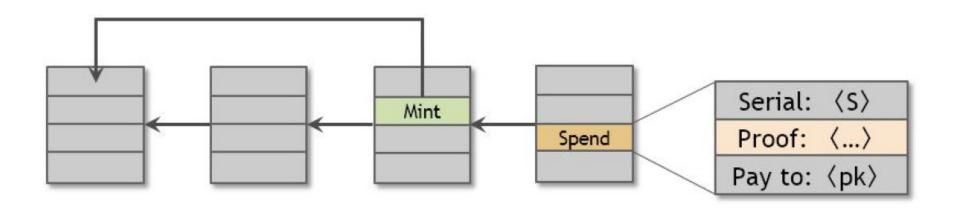




Putting a zerocoin on the block chain

Committing to a serial number





Spending a zerocoin



# Spending a zerocoin with serial number S to redeem a new basecoin

 Create a special "spend" transaction that contains S, along with a zero-knowledge proof of the statement:

"I know r such that Commit(S, r) is in the set  $\{c_1, c_2, ..., c_n\}$ ".

- Miners will verify your zero-knowledge proof which establishes your ability to open one of the zerocoin commitments on the block chain, without actually opening it.
- Miners will also check that the serial number S has never been used in any previous spend transaction (since that would be a double-spend).
- The output of your spend transaction will now act as a new basecoin. For the output address, you should use an address that you own.



Zerocoin: 匿名性

- ■铸币交易或者花费交易中没有展示过r
- 无人知道序列号对应哪一个具体的零币



#### Zerocash

- zk-SNARK (Zeroknowledge Succinct Noninteractive Arguments of Knowledge)
- DAP (Decentralized Anonymous Payment Scheme)
- Hiding user identities, transaction amounts, and account balances from public view

# Summary

System	Туре	Anonymity attacks	Deployability
Bitcoin	pseudonymous	transaction graph analysis	default
Manual mixing	mix	transaction graph analysis, bad mixes/peers	usable today
Chain of mixes or coinjoins	mix	side channels, bad mixes/peers	bitcoin-compatible
Zerocoin	cryptographic mix	side channels (possibly)	altcoin, trusted setup
Zerocash	untraceable	none known	altcoin, trusted setup

A comparison of the anonymity technologies