Learn RISC-V CPU Implementation and BSV

(BSV: a High-Level Hardware Design Language)

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L16: RISC-V: The Fife pipelined CPU



Reminders

Please git clone: https://github.com/rsnikhil/Learn_Bluespec_and_RISCV_Design (git pull for latest version). Repsitory structure:

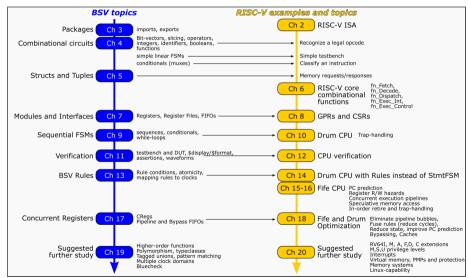
```
./Book_BLang_RISCV.pdf
 Slides/
     Slides 01 Intro.pdf
     Slides_02_ISA.pdf
 Exercises/
     Ex-03-A-Hello-World/
     Ex-03-B-Top-and-DUT/
      . . .
 Code/
     src Top/
     src_Drum/
     src_Fife/
      src Common/
 Doc/Installing_bsc_Verilator_etc.{adoc.html}
```

- Slides and Exercise are numbered in sync with book Chapter numbers.
- For Exercises, please see Appendix E of the book.
 Some (not all) exercises have associated code in the Exercises/ directory.

To compile and run the code for exercises, Drum and Fife, please make sure you have installed:

- bsc compiler (see https://github.com/B-Lang-org/bsc)
- Verilator compiler (see https://www.verilator.org/)

Chapter Roadmap



Flow of information between stages in Drum and Fife

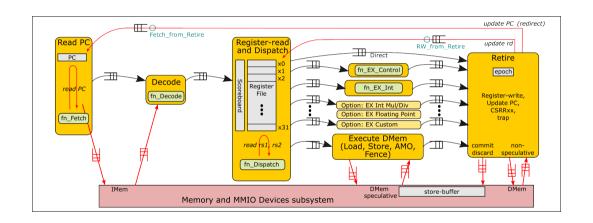


Table of Contents

- Reminders
- 2 Module Interface (same for Drum and Fife)
- 3 Top-level mkCPU module
- 4 Fetch stage
- Decode stage
- 6 Register-Read and Dispatch stage (and Register-Write)
- Execute Control stage
- 8 Execute Int stage
- Retire stage
- Final Comments

Fife CPU: Module Interface (same for Drum and Fife)

6 / 43

Module Interface

```
src_Common/CPU_IFC.bsv: line 27
interface CPU IFC:
   method Action init (Initial Params initial params):
   // IMem
   interface FIFOF O #(Mem Reg) fo IMem reg:
   interface FIFOF_I #(Mem_Rsp) fi_IMem_rsp;
   // DMem, speculative
   interface FIF0F_0 #(Mem_Req) fo_DMem_S_req;
   interface FIFOF I #(Mem Rsp) fi DMem S rsp:
   interface FIFOF O #(Retire to DMem Commit) fo DMem S commit:
   // DMem. non-speculative
   interface FIFOF_O #(Mem_Reg) fo_DMem_reg;
   interface FIFOF_I #(Mem_Rsp) fi_DMem_rsp;
   // Set TIME
   (* always ready, always enabled *)
  method Action set_TIME (Bit #(64) t):
endinterface
```

The speculative DMem sub-interfaces:

 $\label{local_sum} \mbox{fo_DMem_S_req} \qquad \mbox{fi_DMem_S_rsp} \qquad \mbox{fo_DMem_S_commit} \\ \mbox{were unused in Drum, and are used in Fife.}$

Fife CPU: Top-level mkCPU module

Please simultaneously view file: Code/src_Fife/CPU.bsv

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Overall Fife CPU module structure

```
From src_Drum/CPU.bsv _____
(* synthesize *)
module mkCPU (CPU IFC):
   // STATE (sub-modules for pipeline stages)
   . . .
                                                                                  Details in slides that follow
   // STATE (sub-modules for inter-stage connections)
   . . .
   // INTERFACE
   . . .
endmodule
```

Fife's mkCPU module is actually simpler than Drum's mkCPU because the major functionality is now in the separate stage sub-modules (Drum did not have stage sub-modules).

Fife module state: sub-modules for pipeline stages

```
From src_Fife/CPU.bsv .
// Pipeline stages
Fetch_IFC
                stage_F
                                  <- mkFetch:
Decode_IFC
                stage_D
                                  <- mkDecode:
RR RW IFC
                stage_RR_RW
                                  <- mkRR RW:
                                                    // Branch, JAL, JALR
EX_Control_IFC
                stage_EX_Control <- mkEX_Control;</pre>
EX_Int_IFC
                stage_EX_Int
                                  <- mkEX_Int;
                                                    // Integer ops
Retire IFC
                stage_Retire
                                  <- mkRetire:
```

Fife instantiates a submodule for each of the pipeline stages

Fife module state: inter-stage forward-flow connections (sub-modules)

```
From src_Fife/CPU.bsv __
// Forward flow connections
// Fetch->Decode->RR-Dispatch, and direct path RR-Dispatch->Retire
mkConnection (stage_F.fo_Fetch_to_Decode, stage_D.fi_Fetch_to_Decode);
mkConnection (stage_D.fo_Decode_to_RR,
                                           stage_RR_RW.fi_Decode_to_RR);
mkConnection (stage_RR_RW.fo_RR_to_Retire, stage_Retire.fi_RR_to_Retire);
// RR-Dispatch->various EX
mkConnection (stage_RR_RW.fo_RR_to_EX_Control.
              stage EX Control.fi RR to EX Control):
mkConnection (stage_RR_RW.fo_RR_to_EX_Int.
              stage_EX_Int.fi_RR_to_EX_Int);
// Various EX->Retire
mkConnection (stage_EX_Control.fo_EX_Control_to_Retire.
             stage Retire.fi EX Control to Retire):
mkConnection (stage EX Int.fo EX Int to Retire.
              stage Retire.fi EX Int to Retire):
```

Each connection is a FIFO

Fife module state: inter-stage backward-flow connections (sub-modules)

```
From src_Fife/CPU.bsv

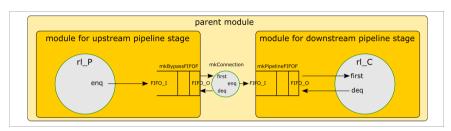
// Backward flow connections

// Fetch<-Retire (redirection)
mkConnection (stage_Retire.fo_Fetch_from_Retire, stage_F.fi_Fetch_from_Retire);
// RR-Dispatch<-Retire (register writeback)
mkConnection (stage_Retire.fo_RW_from_Retire, stage_RR_RW.fi_RW_from_Retire);
```

Each connection is a FIFO

FIFO connections between separately compiled Fife sub-modules

Each inter-stage FIFO mkConnection has this form:



- Data can traverse from producer to consumer in 1 tick, as desired, despite there being two FIFOs, because of the semantics of mkBypassFIFOF and mkPipelineFIFOF (Chapter 17).
- The structure allows the producer and consumer to be compiled independently by bsc, with no "rule-scheduling" constraints leaking across stage boundaries.
- There are no combinational paths crossing the stage boundary (through the two FIFOs).
- The structure allows us to reason about correctness of each stage completely independently of other stages.

Fife module interface definitions

```
From src Fife/CPU bsv
method Action init (Initial Params initial params):
   rg_flog <= initial_params.flog;</pre>
   stage F.init (initial params):
   stage_D.init (initial_params);
   stage_RR_RW.init (initial_params);
   stage_EX_Control.init (initial_params);
   stage_EX_Int.init (initial_params);
   stage Retire.init (initial params):
endmethod
// IMem
interface fo_IMem_req = stage_F.fo_Fetch_to_IMem;
interface fi IMem rsp = stage D.fi IMem to Decode:
// DMem, speculative
interface fo_DMem_S_req
                           = stage_RR_RW.fo_DMem_S_req;
interface fi DMem S rsp
                           = stage Retire.fi DMem S rsp:
interface fo DMem S commit = stage Retire.fo DMem S commit:
// DMem. non-speculative
interface fo DMem reg = stage Retire.fo DMem reg:
interface fi DMem rsp = stage Retire.fi DMem rsp:
// Set TIME
method Action set TIME (Bit #(64) t) = stage Retire.set TIME (t):
```

The init method initializes this module and also each stage sub-module.

The set_TIME method is for updating the time CSR.

Fife CPU Fetch stage

Please simultaneously view file: Code/src_Fife/S1_Fetch.bsv

Fetch stage: interface

```
src_Fife/S1_Fetch.bsv: line 33 ...
interface Fetch_IFC;
method Action init (Initial_Params initial_params);

// Forward out
interface FIFOF_0 #(Fetch_to_Decode) fo_Fetch_to_Decode;
interface FIFOF_0 #(Mem_Req) fo_Fetch_to_IMem;

// Backward in
interface FIFOF_I #(Fetch_from_Retire) fi_Fetch_from_Retire;
endinterface
```

Interface

The FIFO_I from Retire carries redirection information (update PC on misprediction/exception/interrupt).

Fetch stage: implementation module

For module mkFetch let us examine file: Code/src_Fife/S1_Fetch.bsv

- Rule rl_Fetch_req is the forward-flow "fetching rule", repeatedly invoking fn_Fetch to produce IMem requests and information for Decode.
 - The expression "rg_pc+4" is, more generally, "predict(rg_pc)". In future we can substitute new, improved, predictor functions (see book "Section 18.3.6.4 Better next-PC prediction").
- Rule rl_Fetch_from_Retire is the backward-flow. It receives Fetch_from_Retire messages from Retire whenever Retire needs to change the control flow to something other than the predicted PC (due to BRANCH, JAL, JALR, or traps). This rule updates rg_pc to the newly specified PC, and also updates rg_epoch.
- The expression (! f_Fetch_from_Retire.notEmpty) in rl_Fetch_req's condition gives it lower priority compared to rl_Fetch_from_Retire when both are enabled. The sooner we perform redirection, the fewer wrong-path instructions will be fetched.

rg_oiaat is used for "one instruction at a time" mode (for debugging), and can be ignored for now.

Fife CPU Decode stage

Please simultaneously view file: Code/src_Fife/S2_Decode.bsv

Decode stage: interface

```
src_Fife/S2_Decode.bsv: line 33 ...

interface Decode_IFC;
method Action init (Initial_Params initial_params);

// Forward in
interface FIF0F_I #(Fetch_to_Decode) fi_Fetch_to_Decode;
interface FIF0F_I #(Mem_Rsp) fi_IMem_to_Decode;

// Forward out
interface FIF0F_O #(Decode_to_RR) fo_Decode_to_RR;
endinterface
```

Interface

Decode stage: implementation module

For module mkDecode code, let us examine file: Code/src_Fife/S2_Decode.bsv

- FIFO f_Fetch_to_Decode is a mkSizedFIF0(4) instance, aiming to balance the Fetch-to-Decode direct path length with the Fetch-to-IMem-to-Decode path length (equalize the number of items that can be "in flight" on each path).
- Rule rl_Decode is the forward-flow; it repeatedly invokes fn_Decode to produce information for the next stage.

20 / 43

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Fife CPU Register-Read and Dispatch stage (and Register-Write)

Please simultaneously view file: Code/src_Fife/S3_RR_RW.bsv

Register-Read and Dispatch stage (and Register-Write): Interface

```
src Fife/S3_RR_RW.bsv: line 40 ...
interface RR RW IFC:
  method Action init (Initial_Params initial_params);
  // Forward in
  interface FIFOF I #(Decode to RR) fi Decode to RR:
  // Forward out
  interface FIFOF_0 #(RR_to_Retire) fo_RR_to_Retire;
  interface FIFOF_O #(RR_to_EX_Control) fo_RR_to_EX_Control;
  interface FIFOF_O #(RR_to_EX)
                                     fo_RR_to_EX_Int;
  interface FIF0F_0 #(Mem_Req)
                                        fo DMem S reg:
  // Backward in
  interface FIFOF I #(RW from Retire) fi RW from Retire:
endinterface
```

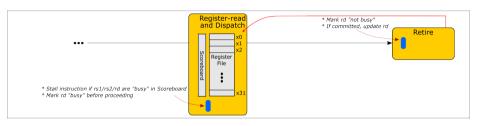
Interface

The four FIFO_O interfaces feed the four execution paths.

The FIF0_I from Retire carries information for registerwrite and release of scoreboard reservations.

Reg-Read and Dispatch (and Reg-Write): GPRs

For module mkRR_RW code, let us examine file: Code/src_Fife/S3_RR_RW.bsv



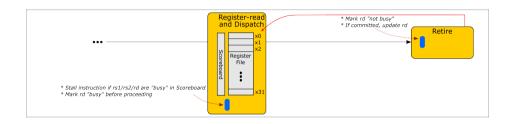
The GPRs (general-purpose registers) are instantiated in this module because we will be reading Rs1 and Rs2 here for each instruction:

```
GPRs_IFC #(XLEN) gprs <- mkGPRs_synth;
```

mkGPRS_synth is just a version of mkGPRs instantiated with a specific XLEN width, and with the (* synthesize *) attribute so that it becomes a Verilog module.

23 / 43

Reg-Read and Dispatch (and Reg-Write): Scoreboard



The Scoreboard accompanies the GPRs. It contains one bit per GPR, indicating whether it is busy (1) or not-busy (0).

```
From src_Fife/CPU.bsv ______

typedef Vector #(32, Bit #(1)) Scoreboard;

Reg #(Scoreboard) rg_scoreboard <- mkReg (replicate (0));
```

(Type Vector#(n,t) represents a vector of n elements, each of type t.)

Reg-Read and Dispatch (forward flow)

Let us examine rl_RR_Dispatch in file: Code/src_Fife/S3_RR_RW.bsv

- It stalls (waits) if the instruction has Rs1, Rs2 or Rd, and these are busy according to the scoreboard.
- Othersize
 - It reads Rs1 and Rs2 registers for the current instruction.
 - It sets the scoreboard for the current instruction's Rd to 1, marking it "busy" (if the instruction has an Rd);
 - It uses information from Decode to dispatch to the four Execute pipes. We always (for every instruction) send an execution tag and additional information on the direct pipe. Depending on the instruction, it may also send information into one of the other Execute pipes: (Execute Control, Execute Integer, and DMem)

Reg-Write (backward flow)

Let us examine rl_RW_from_Retire in file: Code/src_Fife/S3_RR_RW.bsv

- We pop the message x from the f_RW_from_Retire FIFO.
- We perform its specified scoreboard-release for register Rd. If the Rd value is to be committed, we write it into GPR [Rd].

The two rules rl_RR_Dispatch and rl_RW_from_Retire run concurrently.

If rl_RR_Dispatch was stalled on an instruction whose Rs1 or Rs2 was busy, executing rl_RW_from_Retire may "un-stall" it, if its Rd-write is to the pending register.

The explicit condition on rl_RR_Dispatch:

```
(! f_RW_from_Retire.notEmpty)
```

prioritizes the forward rule lower than the backward rule rl_RW_from_Retire, so that this "un-stalling" (if it happens) can happen as soon as possible.

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Fife CPU Execute Control stage

Please simultaneously view file: Code/src_Fife/S4_EX_Control.bsv

This is a very simple, forward-flow-only, one-in, one-out module; it just applies fn_EX_Control() (which we have already seen in Drum) to each item as it transits.

27 / 43

Fife CPU Execute Int stage

Please simultaneously view file: Code/src_Fife/S4_EX_Int.bsv

This is a very simple, forward-flow-only, one-in, one-out module; it just applies fn_EX_Int() (which we have already seen in Drum) to each item as it transits.

28 / 43

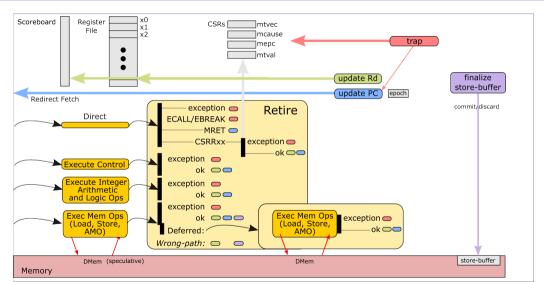
Fife CPU Retire stage

Please simultaneously view file: Code/src_Fife/S5_Retire.bsv

Retire stage module interface

```
src Fife/S5 Retire bsv: line 33 ... -
interface Retire IFC:
  method Action init (Initial_Params initial_params);
   // Forward in
   interface FIFOF I #(RR to Retire)
                                      fi RR to Retire:
   interface FIFOF I #(EX Control to Retire) fi EX Control to Retire:
   interface FIFOF I #(EX to Retire)
                                            fi EX Int to Retire:
   // DMem, speculative
   interface FIFOF_I #(Mem_Rsp)
                                            fi_DMem_S_rsp;
   interface FIFOF O #(Retire to DMem Commit) fo DMem S commit:
   // DMem, non-speculative
   interface FIFOF O #(Mem Reg) fo DMem reg:
   interface FIFOF I #(Mem Rsp) fi DMem rsp:
   // Backward out
   interface FIFOF 0 #(Fetch from Retire) fo Fetch from Retire:
   interface FIFOF_O #(RW_from_Retire)
                                         fo_RW_from_Retire:
   // Set TIME
   (* always_ready, always_enabled *)
   method Action set_TIME (Bit #(64) t):
endinterface
```

Retire stage flows



L16: RISC-V: The Fife pipelined CPU

Retire stage module "mode"

Normally, the Retire stage retires an instruction on every clock (MODE_PIPE). Except

- For DMem "Deferred" instructions (for MMIO). we send the DMem request, and transition to MODE_DMEM_RSP. In MODE DMEM RSP we collect the DMem response. retire the instruction. and return to MODE PIPE.
- For exceptions, we transition to MODE_EXCEPTION. In MODE_EXCEPTION we perform the trap actions. and return to MODE_PIPE.

```
src Fife/S5 Retire.bsv: line 60 ... _____
typedef enum {MODE PIPE.
                         // Normal pipeline operation
             MODE DMEM RSP. // Handle Non-speculative DMem response
             MODE EXCEPTION
} Module Mode
deriving (Bits, Eq. FShow):
```

I 16: RISC-V: The Fife pipelined CPU

Retire stage module state: highlights

```
From Code/src_Fife/S5_Retire.bsv _
module mkRetire (Retire_IFC);
   . . .
   // Control-and-Status Registers (CSRs)
   CSRs_IFC csrs <- mkCSRs;
   // For managing speculation, redirection, traps, etc.
   Reg #(Epoch) rg_epoch <- mkReg (0);</pre>
   . . .
   ... FIFOs for incoming and outgoing flows ...
   . . .
   Reg #(Module_Mode) rg_mode <- mkReg (MODE_PIPE);</pre>
   . . .
endmodule
```

Retire stage module: help-functions and useful predicates

```
From Code/src_Fife/S5_Retire.bsv _
module mkRetire (Retire IFC):
   function Action fa_redirect_Fetch (...);
   function Action fa_update_rd (...);
   . . .
   function Action fa_retire_store_buf (...);
   . . .
   RR_to_Retire x_rr_to_retire = f_RR_to_Retire.first;
   Bool wrong_path = (x_rr_to_retire.epoch != rg_epoch);
   Bool is_Direct = (x_rr_to_retire.exec_tag == EXEC_TAG_DIRECT);
   Bool is_Control = (x_rr_to_retire.exec_tag == EXEC_TAG_CONTROL);
   Bool is Int = (x rr to retire.exec tag == EXEC TAG INT):
   Bool is DMem = (x rr to retire.exec tag == EXEC TAG DMEM):
   . . .
endmodule
```

Retire stage module: rule for wrong-path instructions (mispredicted)

```
src Fife/S5 Retire.bsv: line 190 ... _
rule rl_Retire_wrong_path ((rg_mode == MODE_PIPE)
                           && wrong_path);
  f RR to Retire.deg:
  // Unreserve/commit rd if needed
  fa update rd (x rr to retire, False, ?):
  // Discard related pipe
  if (is_Control) f_EX_Control_to_Retire.deg;
   if (is_Int) f_EX_Int_to_Retire.deg;
  if (is_DMem) begin
      let mem_rsp <- pop_o (to_FIFOF_O (f_DMem_S_rsp));</pre>
     // Send 'discard' (False) to store-buf, if needed
     fa_retire_store_buf (x_rr_to_retire, mem_rsp, False);
   end
   . . .
endrule
```

Retire stage module: rules for Direct path instructions (1/2)

```
From Code/src Fife/S5 Retire.bsv _
rule rl_Retire_CSRRxx ((rg_mode == MODE_PIPE)
                       && (! wrong_path)
                       && is Direct
                       && (! x_rr_to_retire.exception)
                       && is_legal_CSRRxx (x_rr_to_retire.instr));
  match { .exc. .rd val } <- csrs.mav csrrxx (...):
   . . .
  if (! exc)
      . . .
  else begin
     rg_mode <= MODE_EXCEPTION:
   end
endrule
```

```
From Code/src.Fife/S5.Retire.bsv

rule rl_Retire_MRET ((rg_mode == MODE_PIPE)
&& (! wrong_path)
&& is_Direct
&& (! x_rr_to_retire.exception)
&& is_legal_MRET (x_rr_to_retire.instr));
...
endrule
```

I 16: RISC-V: The Fife pipelined CPU

Retire stage module: rules for Direct path instructions (2/2)

```
From Code/src Fife/S5 Retire bsv
rule rl_Retire_ECALL_EBREAK ((rg_mode == MODE_PIPE)
                              && (! wrong_path)
                              && is Direct
                              && (! x rr to retire.exception)
                              && (is_legal_ECALL (x_rr_to_retire.instr)
                                  || is_legal_EBREAK (x_rr_to_retire.instr)));
   . . .
  rg_mode <= MODE_EXCEPTION:
   . . .
endrule
rule rl_Retire_Direct_exception ((rg_mode == MODE_PIPE)
                                  && (! wrong path)
                                  && is Direct
                                  && x_rr_to_retire.exception);
   . . .
  rg mode <= MODE_EXCEPTION:
   . . .
endrule
```

Retire stage module: rules for Execute Control and Int paths instructions

```
From Code/src_Fife/S5_Retire.bsv

rule rl_Retire_EX_Control ((rg_mode == MODE_PIPE)
&& (! wrong_path)
&& is_Control);
...
if (! x2.exception)
...
else begin
...
rg_mode <= MODE_EXCEPTION;
end
endrule
```

Retire stage module: rules for Execute DMem path instructions (1/2)

```
From Code/src_Fife/S5_Retire.bsv

rule rl_Retire_EX_DMem ((rg_mode == MODE_PIPE)
&& (! wrong_path)
&& is_DMem
&& is_DMem
&& (f_DMem_S_rsp.first.rsp_type != MEM_REQ_DEFERRED));
...
if (! exception)
...
else begin
...
rg_mode <= MODE_EXCEPTION;
end
endrule
```

```
From Code/src_Fife/S5_Retire.bsv

rule rl_Retire_DMem_deferred ((rg_mode == MODE_PIPE)
&& (! wrong_path)
&& is_DMem
&& (f_DMem_S_rsp.first.rsp_type
== MEM_REQ_DEFERRED));
...
f_DMem_req.enq (mem_req);
rg_mode <= MODE_DMEM_RSP; // go to await response
endrule
```

Retire stage module: rules for Execute DMem path instructions (2/2)

```
From Code/src_Fife/S5_Retire.bsv

rule rl_Retire_DMem_rsp (rg_mode == MODE_DMEM_RSP);
...
if (exception) begin
...
rg_mode <= MODE_EXCEPTION;
end
else begin
...
rg_mode <= MODE_PIPE;
end
enddendrule
```

Retire stage module: rule for exception-handling

Fife CPU: final comments

- There is very little in this chapter that is RISC-V specific; these pipeline structures are needed for a pipelined CPU for any ISA. All RISC-V-specific issues are the same as in Drum.
- Much scope for performance optimization. E.g.,
 - Faster delivery of redirection to Fetch (less wrong-path wastage)
 - Faster delivery of Rd value to Register-Read (less stalling)
 - In Retire, handle exceptions earlier except perhaps in CSRRxx case.

These are discussed further in Chapter 17.

End

