Learn RISC-V CPU Implementation and BSV

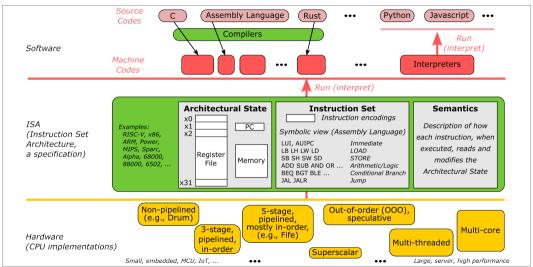
(BSV: a High-Level Hardware Design Language)

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L2: Overview of the RISC-V ISA



What is an ISA?



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Architectural State

The "architectural state" is the state that is visible to instructions. For RV32I, these are:

- The PC (program counter)
- 32 General-Purpose Registers (GPRs, the "register file")
- Memory (byte-addressed)

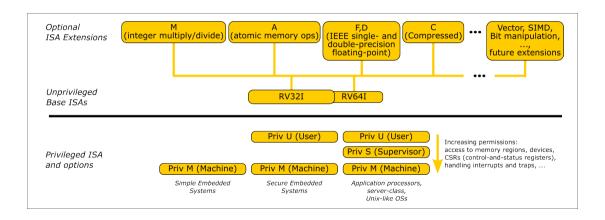
(More architectural state is defined for RV64I, and for most extensions A, F, D, Vector, ...)

The architectural state *does not include* other registers, buffers, FIFOs, memories that may be present in an implementation (they are not visible to instructions).

As such, the architectural state is present in every RISC-V implementation, from tiny CPUs for IoT devices to massive warehouse-scale servers.

Compilers only care about/know about architectural state.

Modularity of the RISC-V ISA



RISC-V Instruction Encodings

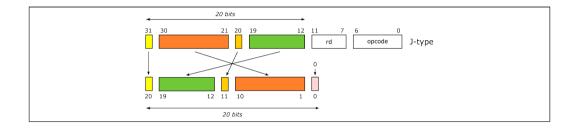
From the RISC-V specification documents:

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31	27	26	25	24		20	19	15	14	12	11	7	6	0	
	funct7				rs2		rs	1	fune	ct3	r	$^{\mathrm{d}}$	opo	code	R-type
	in	nm[11:0)]			rs	1	func	ct3	r	$^{\mathrm{d}}$	opo	code	I-type
	$_{ m imm}[11:5$	5]			rs2		rs	1	fune	ct3	imn	[4:0]	opo	code	S-type
in	nm[12 10	:5]			rs2		rs	1	fune	ct3	imm[4:1 11]	opo	code	B-type
imm[31:12]							r	$^{\mathrm{d}}$	opo	code	U-type				
imm[20 10:1 11 19:12]								r	$^{\mathrm{d}}$	opo	code	J-type			

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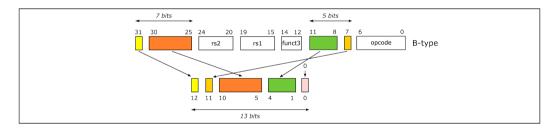
RISC-V Instruction Encodings; J-type immediates



For JAL instruction

RISC-V Instruction Encodings; B-type immediates

For BRANCH set of instructions (BEQ, BNE, BLT, BGE, BLTU, BGEU)



RV32I Instructions

	RV32I	Base Instr	uction S	et			
imm[31:12]				rd	0110111	LUI	↑ Wand (management trade to land annual trade t
		rd	0010111	AUIPC	lack immediate"-kind: to load constant values into a register		
imi	n[20 10:1 11 19		rd	1101111	JAL	f "jump-and-link"-kind: subroutine calls and returns; distant jumps	
imm[11:	imm[11:0] rs1 000			rd	1100111	JALR	Jump-and-link -kind. Subrodeine cans and recurrs, distant jumps
imm[12 10:5]	rs2	rs1	000	imm[4:1 11]	1100011	BEQ	.
imm[12 10:5]	rs2	rs1	001	imm[4:1 11]	1100011	BNE	
imm[12 10:5]	rs2	rs1	100	imm[4:1 11]	1100011	BLT	"conditional branch"-kind: test and possibly jump up to ~0x1000 distance
imm[12 10:5]	rs2	rs1	101	imm[4:1 11]	1100011	BGE	The state of the s
imm[12 10:5]	rs2	rs1	110	imm[4:1 11]	1100011	BLTU	
imm[12]10:5]	rs2	rs1	111	imm[4:1 11]	1100011	BGEU	+
imm[11:	D]	rs1	000	rd	0000011	LB	♦
imm 11:	D)	rs1	001	rd	0000011	LH	
imm 11:	D	rs1	010	rd	0000011	LW	"load data from memory into register" (rs1 and imm specify address)
imm 11:	DÍ	rs1	100	rd	0000011	LBU	
imm[11:	D)	rs1	101	rd	0000011	LHU	↓
imm[11:5]	rs2	rs1	000	imm[4:0]	0100011	SB	*
imm[11:5]	rs2	rs1	001	imm[4:0]	0100011	SH	"store data from register rs2 to memory" (rs1 and imm specify address)
imm[11:5]	rs2	rs1	010	imm[4:0]	0100011	SW	+
imm[11:	0]	rs1	000	rd	0010011	ADDI	A
imm[11:	D)	rs1	010	rd	0010011	SLTI	
imm[11:	D)	rs1	011	rd	0010011	SLTIU	
imm[11:	DÎ	rs1	100	rd	0010011	XORI	
imm 11:	of	rs1	110	rd	0010011	ORI	"integer arithmetic operations (register-immediate)"
imm 11:	DÍ	rs1	111	rd	0010011	ANDI	
0000000	shamt	rs1	001	rd	0010011	SLLI	
0000000	shamt	rs1	101	rd	0010011	SRLI	
0100000	shamt	rs1	101	rd	0010011	SRAI	+
0000000	rs2	rs1	000	rd	0110011	ADD	A
0100000	rs2	rs1	000	rd	0110011	SUB	
0000000	rs2	rs1	001	rd	0110011	SLL	
0000000	rs2	rs1	010	rd	0110011	SLT	
0000000	rs2	rsl	011	rd	0110011	SLTU	
0000000	rs2	rs1	100	rd	0110011	XOR	"integer arithmetic operations (register-register)"
0000000	rs2	rs1	101	rd	0110011	SRL	
0100000	rs2	rs1	101	rd	0110011	SRA	
0000000	rs2	rs1	110	rd	0110011	OR	
0000000	rs2	rs1	111	rd	0110011	AND	↓
fm pre	d succ	rs1	000	rd	0001111	FENCE	<u>.</u>
00000000000		00000 000		00000	1110011	ECALL	"system" operations (ignore FENCE for now)
000000000	001	00000	000	00000	1110011	EBREAK	
							ション・カー・ファイン ション・コー・ファイン

Example specifications

Excerpt from text of ISA specification document for LUI and AUIPC instructions

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SRLI is a logical right shift (zeros are shifted into the upper bits); and SRAI is an arithmetic right shift (the original sign bit is copied into the vacated upper bits).

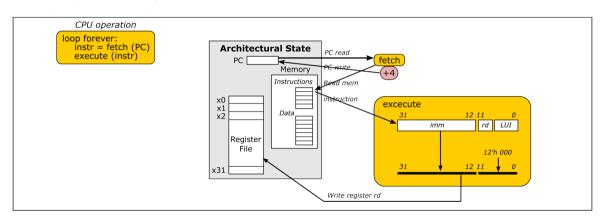
31 12 1	11 7	6	0
imm[31:12]	$^{\mathrm{rd}}$	opcode	
20	5	7	
U-immediate[31:12]	dest	LUI	
U-immediate [31:12]	dest	AUIPC	

LUI (load upper immediate) is used to build 32-bit constants and uses the U-type format. LUI places the U-immediate value in the top 20 bits of the destination register rd, filling in the lowest 12 bits with zeros.

AUIPC (add upper immediate to pc) is used to build pc-relative addresses and uses the U-type format. AUIPC forms a 32-bit offset from the 20-bit U-immediate, filling in the lowest 12 bits with zeros, adds this offset to the address of the AUIPC instruction, then places the result in register rd.

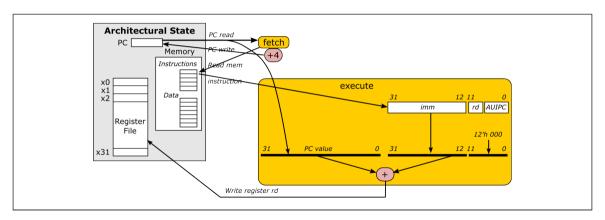
Execution semantics: example

Execution semantics of LUI instruction



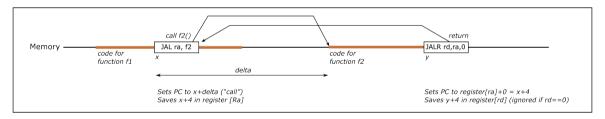
Execution semantics: example

Execution semantics of AUIPC instruction

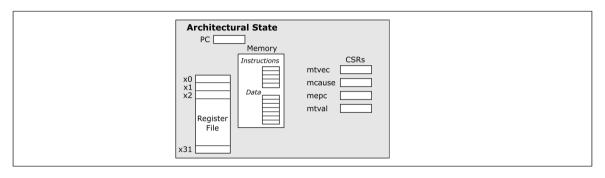


Execution semantics: example

Execution semantics of JAL and JALR instructions (unconditional jumps)



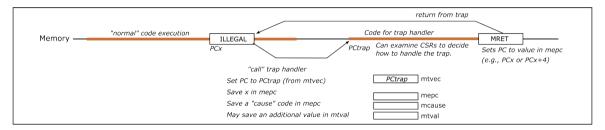
Control and Status Registers (CSRs) for Trap-Handling



These CSRs are implicitly read and written when taking and returning from a trap.

These CSRs can be explicitly read and written by CSRRxx instructions.

Trap and Trap-return flow



There are many possible causes for exceptions and traps; the illustration is for an illegal instruction.

Exception causes

On a trap, a cause-code is written into the mcause CSR:

Exception-Cause code	Description
0	Instruction address misaligned
1	Instruction access fault
2	Illegal instruction
3	Breakpoint
4	Load address misaligned
5	Load access fault
6	Store/AMO address misaligned
7	Store/AMO access fault
11	Environment call M-mode

CSRRxx Instructions

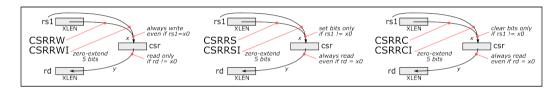
For reading and writing CSRs from RISC-V code (from the RISC-V ISA specifications document)

9.1 CSR Instructions

All CSR instructions atomically read-modify-write a single CSR, whose CSR specifier is encoded in the 12-bit csr field of the instruction held in bits 31–20. The immediate forms use a 5-bit zero-extended immediate encoded in the rs1 field.

31	20 19	15 14 12	11	7 6	0
csr	rs1	funct3	$^{\mathrm{rd}}$	opcode	
12	5	3	5	7	
$\mathrm{source/dest}$	source	CSRRW	dest	SYSTEM	
$\mathrm{source/dest}$	source	CSRRS	dest	SYSTEM	
source/dest	source	CSRRC	dest	SYSTEM	
source/dest	$\operatorname{uimm}[4:0]$	CSRRWI	dest	SYSTEM	
$\mathrm{source/dest}$	$\operatorname{uimm}[4:0]$	CSRRSI	dest	SYSTEM	
source/dest	$\operatorname{uimm}[4:0]$	CSRRCI	dest	SYSTEM	

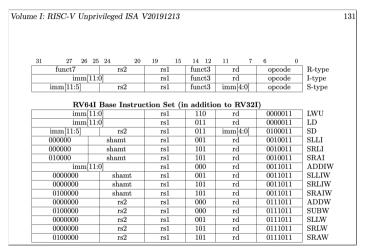
CSRRxx Instruction Semantics



- They all move potentially data into a CSR, and from a CSR to a GPR
- The non "I" variants take input from GPR[rs1] (unless rs1 is zero)
- The "I" variants use rs1 itself as input (unless rs1 is zero)

RV64I instructions

From the RISC-V ISA specifications document



RV64I instructions

In RV64, the 32 GPRs (General Purpose Registers) are each 64-bits wide.

- Most of the RV32I instruction have identical RV64I counterparts (here they operate on 64-bit values).
- A few instructions (SLLI, SRLI, SRAI) are slightly different (allowing 6 bits instead of 5 for the shift amount).
- A few instructions are new, to operate on 32-bit values in the 64-bit registers.
 - LWU to move a 32-bit value from memory into a 64-bit register
 - LD to load a 64-bit value from memory to a 64-bit register
 - SD to store a 64-bit value to memory from 64-bit register
 - ADDIW, SLLIW, ... SRAW to operate on 32-bits of 64-bit registers

End

