# CS 173 Study Guide

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## 1 Math Review

#### 1.1 Sets

#### 1.1.1 Important Sets

The following sets are very commonly used when it comes to numerical analysis - make sure to memorize the sets and understand how they work conceptually.

- Integers:  $\mathbb{Z} = \{..., -3, -2, -1, 0, 1, 2, 3, ...\}$
- Natural Numbers (non-negative integers):  $\mathbb{N} = \{0, 1, 2, 3...\}$
- Positive Integers:  $\mathbb{Z}^+ = \{1, 2, 3, ...\}$
- Real Numbers (all rationals and irrationals):  $\mathbb{R}$
- Rational Numbers (all numbers of the form  $\frac{p}{q}$ , where p and q are integers):  $\mathbb{Q}$
- Complex Numbers (of the form a + bi, where a and b are reals and  $i = \sqrt{-1}$ ):  $\mathbb{C}$

#### 1.1.2 Things to Know

- Zero is not included in non-positive or non-negative.
- Natural numbers include zero.
- Real numbers include integers, so not necessary for a real to have a decimal.
- We denote that x is an element of the set A with:  $x \in A$
- MAKE SURE TO LOOK AT THE TYPES OF YOUR VARIABLES.

#### 1.1.3 Intervals

When working with a lot of numbers in a range, intervals help select a large amount of consecutive elements. They're denoted as such:

- Closed Interval (include both endpoints): [a, b]
- Open Interval (include no endpoints): (a, b)
- Half-Open Intervals (include one endpoint): (a, b] or [a, b)

#### 1.2 Pairs of Reals

- A pair of reals is a special way of denoting two real numbers that go together.
- A pair of reals is denoted by (x,y).
- The SET of all pairs of reals is denoted by  $\mathbb{R}^2$ .
- Before working with pairs, you need to understand what each point in the pair represents.

This concept also applies for containers of n dimensions, and the set of all such containers is represented as  $\mathbb{R}^n$ .

## 1.3 Exponentials and Logs

#### 1.3.1 Exponents

- An exponent is an abbreviated notation for repeated multiplication. Eg.  $a \times a \times a = a^3$ .
- Hard to define exponents for fractional or decimal powers, but we assume that it holds.

#### 1.3.2 Special Exponent Cases

The following illustrate some special exponent cases, make sure to memorize these:

- $x^0 = 1$
- $x^0.5 = \sqrt{x}$
- $\bullet \ x^{-n} = \frac{1}{x^n}$

### 1.3.3 Exponent Manipulation

The following illustrate important exponent rules, make sure to memorize these:

- $\bullet \ x^a x^b = x^{a+b}$
- $x^a y^a = (xy)^a$
- $(x^a)^b = x^a b$
- $x^{(a^b)} \neq (x^a)^b$

#### 1.3.4 Logarithms

If we have an exponential equation  $y = x^a$ , then we can invert this to  $y = \log_a x$ . In most cases, a lack of a base indicates that the default base is 2.

#### 1.3.5 Logarithm Manipulation

The following illustrate some important logarithm rules, make sure to memorize these:

- $b^{\log_b x} = x$
- $\log_b xy = \log_b x + \log_b y$
- $\log_b x^y = y \log_b x$

## 1.3.6 Change of Base Formula

- To apply it correctly, choose a and b such that  $\log_b a$  is greater or less than 1.
- Note that both logarithms are merely different by a constant, which gets dropped later on.

## 1.4 Some Handy Functions

#### 1.4.1 Factorials

- Defined as the product of the first n numbers.
- $n! = 1 \times 2 \times 3 \times ... \times (n-1) \times n$
- Note that 0! is predefined as 1.

#### 1.4.2 Permutations

- Defined as the amount of ways to select n objects from a set in any order.
- Permutations ("n choose k"):  $\frac{n!}{k!(n-k)!}$
- Note that this formula requires the set to contain only UNIQUE objects.

## $1.4.3 \quad Max/Min$

- Returns the max and min of their inputs.
- max(2,7) = 7.

## 1.4.4 Floor/Ceiling

- Floor  $(\lfloor x \rfloor)$  rounds a real downwards.
- Ceiling ([x]) rounds upwards.
- Also applies for negative numbers.

## 2 Logic

## 2.1 A Bit About Style

Writing math has two requirements:

- Logical Flow of Ideas
- Express Yourself Fluently

### 2.2 Propositions

- Proposition: A statement that can be true or false, but not both.
- Doesn't deal with variables or complexity, MUST be predefined.
- Eg. "1 < 2" (never changes from true to false)

## 2.3 Complex Propositions

#### 2.3.1 Chaining Propositions

• Can join propositions to get more complex statements that evaluate to true or false.

• Eg. "1 < 2 and Chicago is in Illinois".

#### 2.3.2 Mathematical Notation

We can manipulate propositions using operators like "not", and we can join propositions using operators like "and"/"or". We also have the following mathematical abbreviations for the following operators:

• and:  $\wedge$ 

• or: V

• not: ¬

#### 2.3.3 Truth Tables

We can use truth tables to show the outcome of manipulating propositions. Here's a truth table for the "not" operator:

$$\begin{array}{c|c} A & \neg A \\ \hline T & F \\ F & F \end{array}$$

Here's another truth table for the results of using and/or on two propositions.

Α	В	$A \wedge B$	$A \lor B$
Т	Т	Т	Т
${ m T}$	F	F	$\Gamma$
$\mathbf{F}$	$\Gamma$	F	$\Gamma$
$\mathbf{F}$	F	F	F

Note that the "or" statement is not exclusive - if both of its inputs are true, then the output is also true. In the context of a single exclusive or, there exists an operator called "exclusive or" (XOR).

## 2.4 Implication

- Can join propositions into an "if A, then B" statement.
- Also verbalized as "A implies B".
- Mathematical notation for this:  $A \to B$ .
- Can also be represented as:  $\neg A \lor B$  (not A or B).

Here's the truth table for implication:

A	В	$A \rightarrow B$
Т	Т	Т
${\rm T}$	F	F
$\mathbf{F}$	Т	Т
$\mathbf{F}$	F	F

Note that B's value is only called in if A is true - if A is false, then  $A \to B$  automatically is TRUE.

## 2.5 Converse, Contrapositive, Biconditional

All three of these operate based on the implies statement.

#### 2.5.1 Converse

The converse of  $A \to B$  is  $B \to A$ . To find the converse:

- 1. Find the proposition chains A and B within the original statement.
- 2. Flip B and A, and place the "implies" sign in between.

NOTE THAT THE CONVERSE OF A STATEMENT IS NOT EQUAL TO THE ORIGINAL STATEMENT.

#### 2.5.2 Biconditional

If a problem says "A implies B, and conversely", then this means:  $A \to B \land B \to A$ , which is equivalent to  $A \leftrightarrow B$ . Here's a truth table for the bidirectional operator:

Α	1	В	$A \rightarrow B$	$B \to A$	$A \leftrightarrow B$
П	1	Т	Т	Т	Т
Γ	٦.	F	F	Τ	F
F	7	Т	Т	$\mathbf{F}$	F
F	7	F	Т	Τ	Т

Note that in the case of a biconditional relationship, the converse of a statement is equivalent to the original statement.

## 2.5.3 Contrapositive

The contrapositive of  $A \to B$  is  $\neg B \to \neg A$ . To find the contrapositive:

- 1. Find the proposition chains A and B within the original statement.
- 2. Negate both predicates (A and B).
- 3. Flip the negations of B and A, and place the "implies" sign in between.

Note that the contrapositive of a statement is equal to the original statement.

## 2.6 Complex Statements

- Can combine longer propositions to achieve chains, works the same way.
- Order of operations: parentheses, not, and/or, implications.
- Can also build truth tables, but lot of work as we need more variables.

## 2.7 Logical Equivalence

- Logically Equivalent: values of two propositions A and B are equal for all possible input values.
- Denoted mathematically with  $\equiv$
- Can get to it via truth tables or simplification.

#### 2.7.1 DeMorgan's Laws

Set of laws regarding boolean algebra:

- $\neg (A \lor B) \equiv \neg A \land \neg B$
- $\neg (A \land B) \equiv \neg A \lor \neg B$
- $A \land \neg A \equiv \text{False}$

## 2.8 Some Useful Logical Equivalences

#### 2.8.1 Commutative Rules

- $A \wedge B \equiv B \wedge A$
- $A \lor B \equiv B \lor A$

#### 2.8.2 Distributive Rules

- $A \wedge (B \wedge C) \equiv A \wedge B \wedge C$
- $A \lor (B \lor C) \equiv A \lor B \lor C$
- $A \wedge (B \vee C) \equiv (A \wedge B) \vee (A \wedge C)$
- $A \lor (B \land C) \equiv (A \lor B) \land (A \lor C)$

## 2.9 Negating Propositions

To negate propositons:

- 1. Convert English text into mathematical notation.
- 2. Perform the negation on the implication.
- 3. Convert the mathematical negation back into plain English.

Note that this is dependent on the negation of the implies:  $\neg(A \to B) \equiv A \land \neg B$ 

#### 2.10 Predicates and Variables

- Predicate: A statement that takes on a T/F value based on the variables passed into it.
- Eg.  $x^2 > 10$  is true if x=4, but false if x=3.
- When creating a predicate, need to be explicit about types of variables and assertions.

## 2.11 Other Quantifiers

## 2.11.1 There Exists

- Any single element in the set that fulfills the requirements.
- Denoted with  $\exists$ .

#### 2.11.2 For All

- Every element within the given set.
- Denoted with  $\forall$ .

## 2.11.3 Unique Exists

- A single element within the given set, which is the ONLY one that fulfills the requirements.
- Highly unlikely to show up on an examlet.

#### 2.12 Notation

- Mathematical notation has a strict template: quantifier, variable + domain, predicate.
- Can also be expressed in simple English terms.
- Eg.  $\forall x \in \mathbb{R}, x^2 + 3 \ge 0$  or "for all values of x in the real numbers,  $x^2 + 3 \ge 0$ ".

## 2.13 Useful Notation

- Can also tie together multiple claims if two variables have the same type.
- Eg.  $\forall x, y \in \mathbb{Z}, x + y \ge x$ .
- In this case, x and y don't have to be different both are independent arbitrary values.
- Can also utilize contrapositive in the case of the predicate.

#### 2.14 Notation for 2D Points

Can do any of the two:

- Create a new 2D pair and refer to it later.
- Create two elements.

## 2.15 Negating Statements with Quantifiers

To negate a statement:

- 1. Invert the quantifier (exists becomes for all, and vice versa).
- 2. Negate the predicate/implies statement.

## 2.16 Binding and Scope

- A quantifier binds the variable it defines.
- After a variable is not bound, it's considered free.
- If a variable is free, it should be considered invalid.

## 3 Proofs

NOTE: The approaches defined here change based on universal and existential statements, and based on whether or not the proof is positive or negative. Be sure to take this into account!

### 3.1 Proving a Universal Statement

- 1. Define all vocabulary (eg. rationals, integers, ...).
- 2. Pick a representative value from the set (VARIABLE).
- 3. Go from the value to the claim (hardest part to prove).

#### 3.3 Direct Proof Outline

- Start with known information, move towards final statement.
- Sometimes need to reason backwards, but ALWAYS write forwards.

## 3.4 Proving Existential Statements

- Find a value that matches the claim, done.
- Can choose any value, because the claim is existential.

## 3.5 Disproving a Universal Statement

- Similar process to proving an existential statement.
- Find a value that proves it wrong.

#### 3.6 Disproving an Existential Statement

- Similar process to proving a universal statement.
- Find a representative element, then work from there.

## 3.7 Recap of Proof Methods

$\operatorname{claim}$	prove	disprove
universal	representative element	counterexample
existential	example	representative element

## 3.10 Proof by Cases

- If claim is in the form p or q, then break it up, and prove for p and q individually.
- Combine both statements together afterwards.

## 3.11 Rephrasing Claims

• Depending on the claim, might need to apply negations and DeMorgan's Laws to pull it all together.

## 3.12 Proof by Contrapositive

- Fairly straightforward if you can prove the contrapositive, then you have proven the claim.
- No strict rule as for when to use it.

## 4 Number Theory

## 4.1 Factors and Multiples

- a divides b if a = bn, for any integer n.
- In this case:
  - a is a factor of b
  - b is a multiple of a
- Can express "a divides b" as a|b.
- NOTE: the "smaller" number goes on the left, not the right.

NOTE: An integer p is even iff 2|p.

#### 4.3 Stay in the Set

• Don't introduce rationals if working with ints!

#### 4.4 Prime Numbers

- p  $(p \ge 2)$  is prime iff the only positive factors of p are p and 1, else it's composite.
- Prime Factorization: Expressing any integer p as the product of only prime numbers.

#### 4.5 GCD and LCM

#### 4.5.1 GCD

- Common Divisor: any value that divides 2 integers
- Greatest Common Divisor: largest common divisor of two integers.
- Can express a GCD as gcd(a,b).
- Can calculate the GCD by extracting common factors from the prime factorization.

#### 4.5.2 LCM

- Smallest value possible such that a|c and b|c.
- Can be found with this formula:  $lcm(a,b) = \frac{ab}{acd(a,b)}$ .
- Relatively Prime: When two integers have no shared common divisors.

## 4.6 The Division Algorithm

For any integers a and b, there are unique integers q and r such that a = bq + r.

## 4.7 Euclidean Algorithm

```
Keep doing this algorithm: gcd(a,b):

x = a

y = b

while y > 0:

r = remainder(x, y)

x = y
```

y = r return x

## 4.10 Congruence mod K

- Two integers are congruent mod k if they differ by a value of k.
- If k is any positive integer,  $a \equiv b(modk)$  iff k|(a-b).

## 4.12 Equivalence Classes

- Congruence Class/Equivalence Class: Set of all integers mod k.
- Eg. in congruence mod 7:  $[3] = \{..., -11, -4, 3, 10, 17, ...\}$
- Generally only use integers from 0 to k-1 to name these classes.

## 5 Sets

#### 5.1 Sets

- Sets are unordered collections of objects.
- Items in a set are called elements or members.
- Three ways to define a set:
  - Mathematical English (all integers between 3 and 7, inclusive).
  - List  $({3,4,5,6,7})$
  - Set Builder Notation ( $\{x \in \mathbb{Z}, 3 \le x \le 7\}$ )
- Note that the comma (may be replaced with |) in set builder notation indicates a constant.

## 5.2 Things to be Careful About

- Set is unordered and unique, so  $\{3, 2, 1\} = \{1, 2, 3\}.$
- No duplicates in sets, so  $\{3, 2, 1, 2\} = \{1, 2, 3\}.$
- Note that sets aren't tuples (ordered non-unique collections), so use {} and not ().
- Sets can be empty  $(\emptyset)$  or have a single value.
- Sets can contain objects of multiple types.

## 5.3 Cardinality, Inclusion

- Cardinality: Amount of (unique) objects in a set.
- Subset: All elements in A are also in B, denoted by  $A \subseteq B$ .
- Proper Subset: Subset but both sets must be different, denoted by  $A \subset B$ .

#### 5.4 Vacuous Truth

- Happens when a statement is technically true by definition, but not actually true.
- Eg. is an empty set a subset of A? Yes (all elements in empty set are in A) but also no.

### 5.5 Set Operations

- Intersection(∩): All elements that exist in BOTH sets. If intersection yields empty set, then both sets are disjoint.
- Union( $\cup$ ): Set of all unique elements in both sets combined.
- Set Difference(-): Set of all elements in the first set but not in the second.
- Cartesian Product (×): Set of all 2D-tuples, containing combinations of elements from both sets.

### 5.7 Size of Set Union

- Inclusion-Exclusion Principle:  $|A \cup B| = |A| + |B| |A \cap B|$
- Can extend this idea to multiple sets (2+).

#### 5.8 Product Rule

- Very intuitive method to determine cardinality of Cartesian product.
- $\bullet$   $|A \times B| = |A| \times |B|$

## 5.10 Proving Facts About Set Inclusion

- Start with a representative element from A.
- Perform algebra to show that A exists in B.
- Conclude by saying that the claim has been proven.

## 6 Relations

#### 6.1 Relations

- Set of ordered pairs of elements from A to elements in A.
- If (x, y) exists in the set, we denote it as xRy.
- Note that elements can relate to themselves.
- Can draw directed graphs to represent relations.

#### 6.2 Properties of Relations: Reflexive

- Three main cases:
  - Reflexive (every element relates to itself)
  - Irreflexive (no element relates to itself)
  - Neither (some elements relate to themselves, but not all)

## 6.3 Symmetric and Antisymmetric

- If xRy and yRx, then relation is symmetric.
- Represented in a graph by bidirectional arrows.
- Antisymmetric if no two elements relate to each other in both directions.

#### 6.4 Transitive

- If aRb and bRc, then aRc.
- Whenever theres a n-length path in a graph from a to c, there must also be a length 1 path.
- NOTE: if aRb and bRa, then a must relate to itself in order for graph to be transitive.

## 6.5 Types of Relations

- Partial Order: Reflexive, antisymmetric, transitive.
- Linear Order (total order): Partial order where all elements relate to each other.
- Strict Partial Order: Irreflexive, antisymmetric, transitive.
- Equivalence Relation: Reflexive, symmetric, transitive.

## 7 Functions and Onto

#### 7.1 Functions

- Function F maps items from set A to set B.
- A is the domain, and B is the co-domain.
- This is denoted in the type signature:  $f: A \to B$
- If x is an element of A, then f(x) is the image of x.

## 7.2 When Are Functions Equal

- Must assign the same inputs to the same outputs.
- Must also have the same type signatures.

## 7.3 What Isn't a Function?

- An input has no output.
- An input has multiple outputs.

#### 7.4 Images and Onto

- Image:  $\{f(x), x \in A\}$  (all f(x) for x in A).
- A function is onto if its co-domain IS its image.
- To prove: generate a representative output and find its pre-image.
- Mathematical Def of Onto :  $\forall y \in B, \exists x \in A, f(x) = y$ .

## 7.6 Negating Onto

- Start out with negation:  $\neg(\forall y \in B, \exists x \in A, f(x) = y)$ .
- End up with this statement:  $\exists y \in B, \forall x \in A, f(x) \neq y$ .
- In other words, there's an output that doesn't have an input.

## 7.7 Nested Qualifiers

- Need to be careful with nested qualifiers.
- Easy to mess up order, with existential and universal qualifiers.

## 8 Functions and One-To-One

#### 8.1 One-To-One

- A function is one-to-one if it maps each input to only one output.
- Also depends on the type signature (eg. absolute value of x on integers vs. naturals).
- Mathematical Def:  $\forall x,y \in A, x \neq y \rightarrow f(x) \neq f(y)$  OR  $\forall x,y \in A, f(x) = f(y) \rightarrow x = y$

## 8.2 Bijections

- A function is a bijection if it's one-to-one and onto.
- In this case, both the domain and the co-domain are the same size.

## 8.3 Pidgeonhole Principle

- If we have n objects and k labels, then multiple objects must have the same label if n > k.
- Often tricky to implement, but useful when it comes to "packaging".

#### 8.4 Permutations

- To make an ordered choice of k objects from a set of n:  $P(n,k) = \frac{n!}{(n-k)!}$
- Represents the number of one-to-one functions from a set of k objects to a set of n objects.

## 9 Graphs

#### 9.1 Graphs

- Consists of a set of nodes V and a set of edges E.
- Two nodes connected by an edge are neighbors or adjacent.
- Edges can be directed or undirected (we'll assume undirected).
- Some graphs might have multiple edges or loops.
- Simple Graph: No multiple edges, no loop (will be norm for this class).

### 9.2 Degrees

- deg(v) is the total number of edges with v as an endpoint (self-loops count twice).
- Sum of all degrees  $= 2 \times \text{number of edges}$ .

## 9.3 Complete Graphs

- Graphs in which every node is connected to every other node.
- Denoted by  $K_n$ .
- Total Edges:  $\frac{n(n-1)}{2}$

## 9.4 Cycle Graphs and Wheels

#### 9.4.1 Cycle Graphs

- Circular graph in which each node has only 2 neighbors.
- Denoted by  $C_n$ .
- $\bullet$  Total edges: n

#### **9.4.2** Wheels

- Like a cycle graph, but there's a "hub" in the center.
- Denoted by  $W_n$ .
- Total edges: 2n.

## 9.5 Isomorphism

- Isomorphism: Function from nodes in Graph A to nodes in Graph B.
- Graphs are isomorphic if there's an isomorphism from A to B.
- To prove that graphs aren't isomorphic, find a feature that's different.

## 9.6 Subgraphs

- Subgraph: Nodes are a subset of the main graph, and so are edges.
- If two graphs are isomorphic, then every subgraph must have a counter.

## 9.7 Paths, Walks, and Cycles

- Walk is a sequence of nodes and edges (generally just one) from node a to node b.
  - Closed: Starting node and ending node is the same, otherwise open.
- Path: A walk in which no nodes are re-used.
- Cycle: Closed walk with 3+ nodes, in which no node is used more than once.
- Can create a path from a walk by pruning nodes.
- Acyclic: Graph has no cycles in it.

### 9.8 Connectivity

- Connected: Walk exists between every pair of nodes in a graph.
- Connected Components: Amount of largest possible sets of connected nodes.
- Cut Edge: Edge that would break a connected component into 2 if removed.

#### 9.9 Distances

- d(a,b) is the length of the shortest path from a to b.
- Diameter: maximum distance between any pair of nodes in the graph.

#### 9.10 Euler Circuit

- Closed walk that uses each edge of the graph once.
- Requires a connected graph and each node to have even degree.

#### 9.11 Bipartite Graph

- Bipartite if we can split V into two subsets such that every node in  $V_1$  connects only to nodes in  $V_2$ .
- No edges connecting nodes of 1 set to nodes in the same set.

## 10 Two-Way Bounding

Use two subproofs, one for lower bound and one for upper.

## 10.1 Marker Making

- Being able to cut out as many shapes as you can from fabric.
- If lower and upper bounds meet, then that's the answer.

## 10.2 Pigeonhole Point Placement

- Break up a problem into smaller ones, then use one of those as a case.
- Identify why it wouldn't work in this case.

### 10.3 Graph Coloring

- Chromatic Color: Minimum colors needed to color a graph without neigbors being the same color.
- Need to show an upper bound (can be colored with n colors) and a lower bound (cannot be colored with less than n colors).

## 11 Induction

Induction

## 12 Recursive Definition

## 12.1 Recursive Definition

- Two parts:
  - Base case(s)
  - Recursive formula
- Need to have both, but can't label them differently.
- Inputs are assumed to be integers.

## 12.2 Finding Closed Forms

- Can sometimes guess, but often need to use unrolling.
- Keep substituting a definition into itself until a pattern's clear.
- Can also apply summation formulas to make it cleaner.
- Find value of k at base case, then plug and chug.

## 13 Trees

Trees

## 14 Big-O

Big-O

# 15 Algorithms

Algorithms

## 16 NP

NP

## 17 Proof by Contradiction

Proof by Contradiction

## 18 Collections of Sets

- Most sets previously contained atomic elements (numbers, strings, tuples).
- Sets can also contain other sets.
- Collection: A set that contains other sets.

## 18.1 Sets Containing Sets

- Happens when we need to get subsets of another set.
- Can divide a set into an amount of (non)overlapping subsets.
- If a non-overlapping set of subsets is the domain of a function, each subset results in an an output.

#### 18.1.1 Cardinality

- The cardinality of a set is the amount of elements in the set itself, not in its subsets.
- The empty set can be put in another set, and it counts as an element.

#### 18.2 Powersets and Set-Valued Functions

- Powerset of A contains ALL subsets of A (including empty set).
- Denote a powerset with  $\mathbb{P}(A)$ .
- Powerset of A contains  $2^n$  elements, if A has n elements.
- Use powersets when a function returns MULTIPLE values, and hence consistency is important.
- Are useful for defining the domains of above-mentioned functions.

#### 18.3 Partitions

- Division of a base set A into non-overlapping subsets.
- Corresponds to equivalence reactions (and vice-versa).
- Three requirements for a partition of A:
  - 1. Elements put together cover all of A.
  - 2. No empty set in partition.
  - 3. No overlap within elements of partition sets.

### 18.4 Combinations

#### 18.4.1 Combinations

- Combinations happen when we have an n-element set, from which we need all subsets of size k.
- K-Combination: Subset of size k.
- Care about order in a permutation, but not in a combination.

#### 18.4.2 Equations

- Choose k elements without an order:  $\frac{n!}{(n-k)!}$
- Choose k elements in order (AKA binomial coefficient):  $C(n,k) = \binom{n}{k} = \frac{n!}{k!(n-k)!}$
- Note that this is defined if  $n \ge k \ge 0$

## 18.5 Applying the Combinations Formula

- Used to select a set of locations and assign values.
- Might need to apply it multiple times to generate the correct answer.

### 18.6 Combinations with Repetition

- Need a clever way to count the amount of possibilities with multiple groups.
- Replace each item with a star, and find amount of places you can put the separator.
- To choose k objects from a list of size n:  $\binom{k+n-1}{n-1} = \binom{k+n-1}{k}$

## 18.7 Identities for Binomial Coefficients

- $\bullet \ \binom{n}{k} = \binom{n}{n-k}$
- $\bullet \ \binom{n+1}{k} = \binom{n}{k} + \binom{n}{k-1}$

## 19 State Diagrams

## 19.1 Introduction

- Directed graph, nodes represent states and edges represent actions.
- Label on an edge indicates what happens as a system moves across states.
- Walks must follow arrows.
- NOTE: INCLUDE STATES AND ACTIONS WHEN WRITING IT OUT.

## 20 Countability

## 20.1 The Rationals and the Reals

- Three sets of numbers: integers, rationals, reals.
- Integers are discrete whole numbers
- Majority of real numbers are irrational, and only a few are rational.

## 20.2 Completeness

- Reals have completeness, rationals don't.
- Completeness: Any subset of reals with upper bound has a SMALLEST upper bound.

#### 20.3 Cardinality

- Two sets have the same cardinality iff there's a bijection from A to B.
- Eg.  $f: \mathbb{R} \to \mathbb{Z}, f(n)$   $\begin{cases} \frac{n}{2}, n \equiv 0 \pmod{2} \\ \frac{-(n+1)}{2}, n \equiv 1 \pmod{2} \end{cases}$
- Countably Infinite: Bijection exists from  $\mathbb{N}$  or  $\mathbb{Z}$  onto an infinite set A.
- Countable: A set is countably infinite or finite. (Includes all subsets of integers).

#### 20.4 Cantor Schroeder Bernstein Theorem

- $|A| \leq |B|$  iff there's a one-to-one function from A to B.
- Can reverse this and do it twice (A to B and then B to A), and show that a bijection exists.