

ASSIGNMENT/ASSESSMENT ITEM COVER SHEET

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Student Number: **Email:**

Course Code

Course Title

(Example)

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Campus of Study: (eg Callaghan, Ourimbah, Port Macquarie)

Assessment Item Title: **Due Date/Time:**

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$$a_0 = 2, a_n = a_{n-1} + 3, b_n = 3n + 2$$

Base Case: $n=1$, $a_1 = a_{1-1} + 3 = 5 \therefore a_1 = b_1$
 $b_1 = 3(1) + 2 = 5$

$$\frac{n+1}{(n+1)-1}$$

~~Assume $n=k$, $a_k = a_{k-1} + 3$, $b_k = 3(k) + 2$~~

~~Inductive $n=k+1$, $a_{k+1} = a_k + 3$, $b_{k+1} = 3(k+1) + 2$~~

~~Step~~

~~$a_{k+1} = a_k + 3$, $b_{k+1} = 3k + 3 + 2$~~

Assume, assume that for some arbitrary value of n that $b_n = a_{n-1} + 3$, $a_n = 3n + 2$

Induction, we want to show that $b_{n+1} = a_{n+1-1} + 3$
 Step $a_{k+1} = 3(n+1) + 2$

$$\begin{aligned} b_{n+1} &= 3(n+1) + 2 \\ &= 3n + 3 + 2 \\ &= (a_n) + 3 \\ &= a_n + 3 \end{aligned}$$

$$\begin{aligned} a_{n+1} &= a_{k+1-1} + 3 \\ &= a_n + 3 \end{aligned}$$

Conclusion: therefore $\{a_n = b_n\}$ iff $a_n = b_n, n \in \mathbb{N}_0$

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2A) No, because if 2 ^{odd} ~~even~~ length strings are concatenated, they will be even in length. The closure is all non-empty strings drawn from the alphabet given

B) Yes, it is closed. It doesn't matter if ϵ doesn't end in β because it is no string. Every string ending in β will be closed by Kleene Star

4A) False. If you have $L_1 = \{a\}$, $L_2 = \{a^*\}$, then $L_1 L_2 = L_2 L_1$ and $L_1 \neq L_2$.

Eg: $L_1 L_2 = \{\epsilon, a, aa, \dots\}$, $L_2 L_1 = \{\epsilon, a, aa, \dots\}$

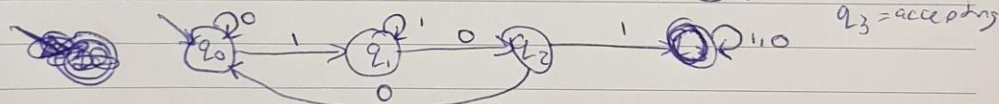
B) False. If you have $L_1 = \{a^*\}$, $L_2 = \{a^+\}$, then $L_1^+ = L_2^+ = \{a^+\} \neq \{a^*\}$. Therefore, false.

C) False.

$L_1 = \{a^*\}$, $L_2 = \{b^*\}$, $L_1 - L_2 = \{\epsilon\}$

Epsilon in $L_1 - L_2$ is finite but the languages L_1 and L_2 are not. Therefore, false.

5A) $\{w \in \{0,1\}^* : w \text{ has } 101 \text{ as substring}\}$



$K = \{q_0, q_1, q_2, q_3\}$

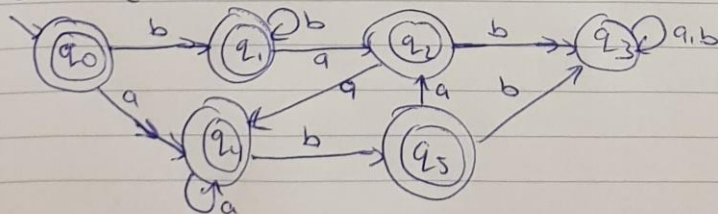
$S = q_0$

$\Sigma = \{0, 1\}$

$A = q_3$

$\delta = \{(q_0, 0) = q_0, (q_0, 1) = q_1, (q_1, 0) = q_2, (q_1, 1) = q_1, (q_2, 0) = q_2, (q_2, 1) = q_3, (q_3, 0) = q_3, (q_3, 1) = q_3\}$

B) $\{w \in \{a,b\}^* : w \text{ has neither } bab \text{ or } abb \text{ as substring}\}$



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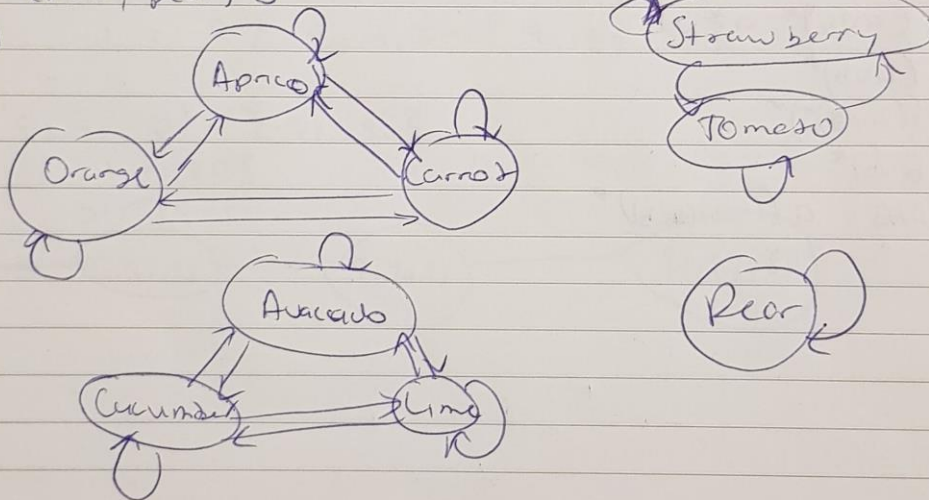
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3) Same colour = $\{ (\text{apricot}, \text{orange}), (\text{avocado}, \text{lime}), (\text{lime}, \text{cucumber}), (\text{strawberry}, \text{tomato}), (\text{carrot}, \text{orange}) \}$

a) Lacks all properties, e.g. it isn't reflexive, symmetric or transitive

b) $\{ (\text{apricot}, \text{orange}), (\text{avocado}, \text{lime}), (\text{lime}, \text{cucumber}), (\text{strawberry}, \text{tomato}), (\text{carrot}, \text{orange}), (\text{apricot}, \text{apricot}), (\text{orange}, \text{orange}), (\text{lime}, \text{lime}), (\text{cucumber}, \text{cucumber}), (\text{strawberry}, \text{strawberry}), (\text{tomato}, \text{tomato}), (\text{carrot}, \text{carrot}), (\text{orange}, \text{orange}), (\text{orange}, \text{apricot}), (\text{lime}, \text{avocado}), (\text{cucumber}, \text{lime}), (\text{tomato}, \text{strawberry}), (\text{orange}, \text{carrot}), (\text{avocado}, \text{cucumber}), (\text{cucumber}, \text{avocado}), (\text{apricot}, \text{carrot}), (\text{carrot}, \text{apricot}), (\text{pear}, \text{pear}) \}$

c)



d) 4

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5B(ONT)

$$K = \{q_0, q_1, q_2, q_3, q_4, q_5\}$$

$$\Sigma = \{a, b\}$$

$$\delta = \{ (q_0, a) = q_4, (q_0, b) = q_1, (q_1, a) = q_2, (q_1, b) = q_1, (q_2, a) = q_4, (q_2, b) = q_3, (q_3, a) = q_3, (q_3, b) = q_3, (q_4, a) = q_4, (q_4, b) = q_5, (q_5, a) = q_2, (q_5, b) = q_3 \}$$

$$S = q_0$$

$$A = \{q_0, q_1, q_2, \text{~~q_3~~, } q_4, q_5\}$$

6) Initial classes = $\{ [6,7], [1,2,3,4,5] \}$

Step 1:

$((6,a), [1,2,3,4,5])$ $((7,a), [1,2,3,4,5])$

$((6,b), [6,7])$ $((7,b), [6,7])$

~~SPLIT 6 and 7~~

DON'T SPLIT

Classes = $\{ [6,7], [1,2,3,4,5] \}$

$((1,a), [1,2,3,4,5])$ $((2,a), [1,2,3,4,5])$

$((1,b), [6,7])$ $((2,b), [6,7])$

$((3,a), [1,2,3,4,5])$ $((4,a), [6,7])$

$((3,b), [1,2,3,4,5])$ $((4,b), [1,2,3,4,5])$

$((5,a), [6,7])$

SPLIT INTO NEW

$((5,b), [1,2,3,4,5])$

CLASSES

Classes: $\{ [6,7], [3], [1,2], [4,5] \}$

$((6,a), [3])$ $((7,a), [3])$

DON'T SPLIT

$((6,b), [6,7])$ $((7,b), [6,7])$

$((1,a), [4,5])$ $((2,a), [1,2])$

SPLIT

$((1,b), [6,7])$ $((2,b), [6,7])$

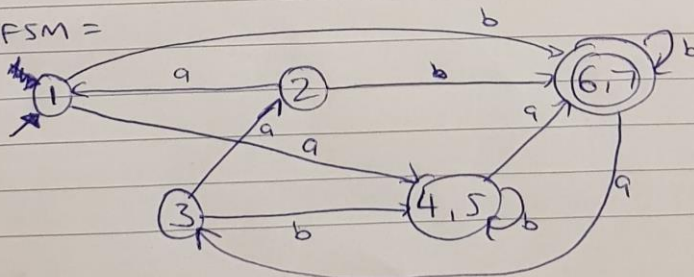
$((4,a), [6,7])$ $((5,a), [6,7])$

DON'T SPLIT

$((4,b), [4,5])$ $((5,b), [4,5])$

Classes: $\{ [6,7], [3], [4,5], [1], [2] \}$

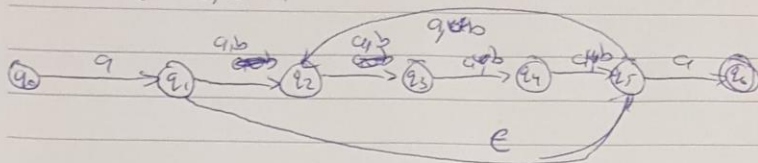
MINDFSM =



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7a) $(a^i b)^*$: two a's separated by whole string is
length is $4i$, for some $i \geq 0$
 $a(a^i b)^* a$



$$K = \{q_0, q_1, q_2, q_3, q_4, q_5, q_6\}$$

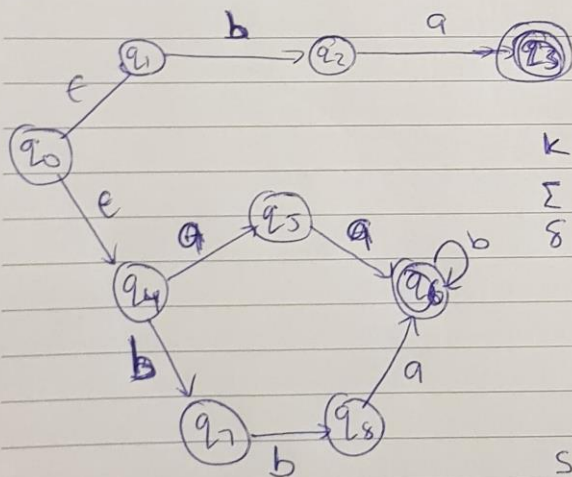
$$\Sigma = \{a, b\}$$

$$\delta = \{(q_0, a) = q_1, (q_1, \epsilon) = q_5, (q_1, a, b) = q_2, (q_2, a, b) = q_3, (q_3, a, b) = q_4, (q_4, a, b) = q_5, (q_5, a) = q_6\}$$

$$S = q_0$$

$$A = q_6$$

7b) $L = \{w \mid \{a, b\}^* ; w = (ba \cup ((a^i b)^* a b^i))\}$



$$K = \{q_0, q_1, q_2, q_3, q_4, q_5, q_6, q_7, q_8\}$$

$$\Sigma = \{a, b\}$$

$$\delta = \{(q_0, \epsilon) = q_1, q_4, (q_1, b) = q_2, (q_2, a) = q_3, (q_4, a) = q_5, (q_5, a) = q_6, (q_6, b) = q_7, (q_7, b) = q_8, (q_8, a) = q_6, (q_6, b) = q_6\}$$

$$S = q_0$$

$$A = \{q_3, q_6\}$$

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Q8)

Step 1: $\text{eps}(1) = \{1, 2\}$ $\text{eps}(4) = \{4\}$
 $\text{eps}(2) = \{2\}$ $\text{eps}(5) = \{4, 5\}$
 $\text{eps}(3) = \{3\}$ $\text{eps}(6) = \{5, 6\}$

Step 2: ~~Start state = $\text{eps}(1) = \{1, 2\}$~~

Step 3: ~~δ' Active states: $\{\{1, 2\}\}$. Consider $\{1, 2\}$
 $\text{eps}(\{1, 2\}, a) = \{3\}$
 $\text{eps}(\{1, 2\}, b) = \{5, 6\}$
Active states: $\{\{1, 2\}, \{3\}, \{5, 6\}\}$. Consider $\{3\}$.
 $\text{eps}(\{3\}, a) =$
 $\text{eps}(\{3\}, b) =$~~

Step 2: Start state = $\text{eps}(5) = \text{eps}(1) = \{1, 2\}$

Step 3: δ' Active = $\{\{1, 2\}\}$. Consider $\{1, 2\}$

$\text{eps}(\{1, 2\}, a) = \{3, 4\}$

$\text{eps}(\{1, 2\}, b) = \{4, 5, 6\}$

Active states: $\{\{1, 2\}, \{3, 4\}, \{4, 5, 6\}\}$ Consider $\{3, 4\}$

$\text{eps}(\{3, 4\}, a) = \{2\}$

$\text{eps}(\{3, 4\}, b) = \{4, 5\}$

Active states: $\{\{1, 2\}, \{3, 4\}, \{4, 5, 6\}, \{2\}, \{4, 5\}\}$

$\text{eps}(\{4, 5, 6\}, a) = \{2, 4, 5\}$

$\text{eps}(\{4, 5, 6\}, b) = \{4, 5, 6\}$

Active states = $\{\{1, 2\}, \{3, 4\}, \{4, 5, 6\}, \{2\}, \{4, 5\}, \{2, 4, 5\}\}$

$\text{eps}(\{2\}, a) = \{4\}$

$\text{eps}(\{2\}, b) = \{4\}$

Active states = $\{\{1, 2\}, \{3, 4\}, \{4, 5, 6\}, \{2\}, \{4, 5\}, \{2, 4, 5\}, \{4\}\}$

$\text{eps}(\{4\}, a) = \{2, 4, 5\}$

$\text{eps}(\{4\}, b) = \{4, 5, 6\}$

Active = $\{\{1, 2\}, \{3, 4\}, \{4, 5, 6\}, \{2\}, \{4, 5\}, \{2, 4, 5\}, \{4\}\}$

$\text{eps}(\{2, 4, 5\}, a) = \{2, 4, 5\}$

$\text{eps}(\{2, 4, 5\}, b) = \{4, 5, 6\}$

Active = Consider $\{4\}$

$\text{eps}(\{4\}, a) = \{2\}$

$\text{eps}(\{4\}, b) = \{4, 5\}$ DONE

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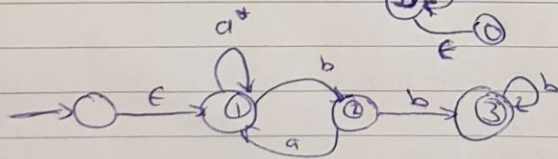
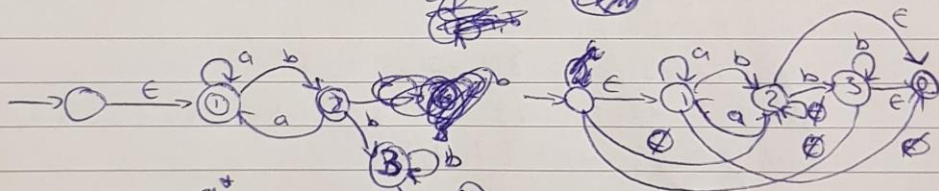
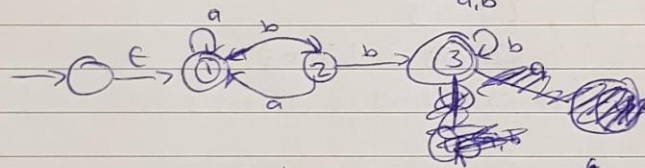
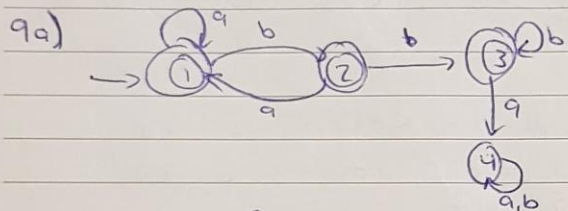
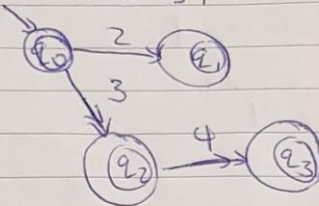
Q8 (cont)

$K' = \{1, 2, 3, \{3, 4\}, \{4, 5, 6\}, \{2\}, \{4, 5\}, \{2, 4, 5\}, \{4\}, \emptyset\}$

$A' = \{1, 2, 3, \{4, 5, 6\}, \{2\}, \{4, 5\}, \{2, 4, 5\}\}$

$\neg L(M) = \{\{3, 4\}, \{2\}, \{4, 5\}\}$

$\neg L(M)$:



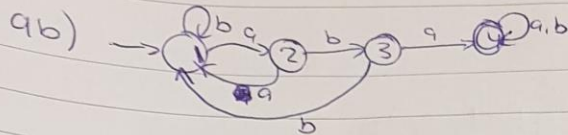
$a^* \quad ba \cup bb \quad b^*$

$a^*(a^*ba \cup bb)b^*$

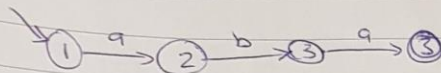
$a^* \cup a^*(ba^* \cup bb)^*b^*$

ANS = $a^* \cup a^*(ba^* \cup bb)b^* \cup a^*(ba^* \cup bb)b^*$

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Get rid of majority of 1, 2, 3 transitions
you get:



regex = aba

~~10a) $\Sigma = \{ 16, 26, 36, 46, 56, 66, 15, 25, 35, 45, 55, 65, 14, 24, 34, 44, 54, 64, 13, 23, 33, 43, 53, 63, 12, 22, 23, 24, 25, 26, 11, 12, 13, 14, 15, 16 \}$~~

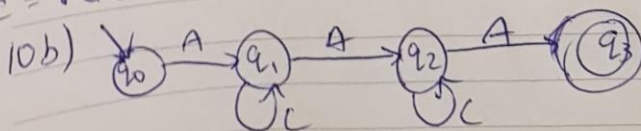
10a) $\Sigma = \{ 11, 21, 31, 41, 51, 61, 12, 22, 32, 42, 52, 62, 13, 23, 33, 43, 53, 63, 14, 24, 34, 44, 54, 64, 15, 25, 35, 45, 55, 65, 16, 26, 36, 46, 56, 66 \}$

Encoding technique would be substituting winning P, winning Q and drawing PQ for letters, ES.

$A = P(w) = \{ 21, 31, 32, 41, 42, 43, 51, 52, 53, 54, 61, 62, 63, 64, 65 \}$

$B = Q(w) = \{ 12, 13, 23, 14, 24, 34, 15, 25, 35, 45, 16, 26, 36, 46, 56 \}$

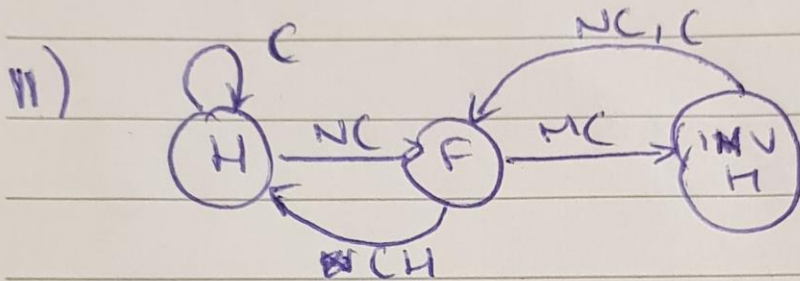
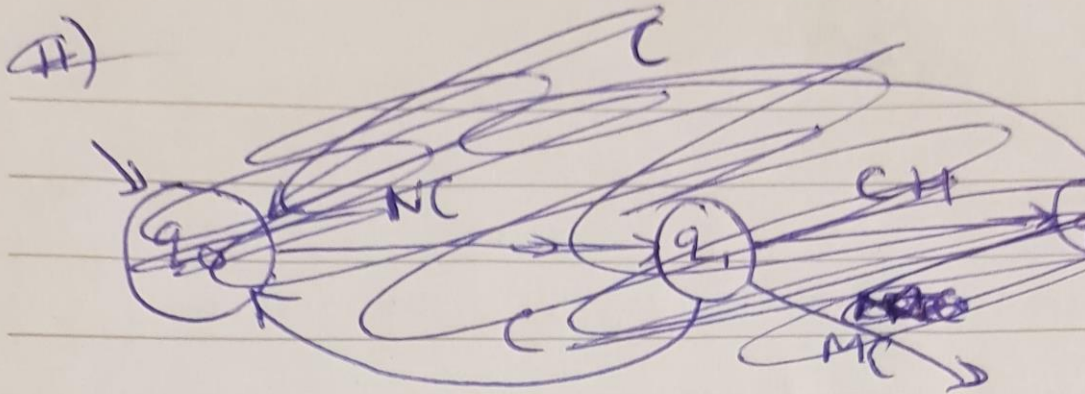
$C = PQ(D) = \{ 11, 22, 33, 44, 55, 66 \}$



$q_1 = (1, 0)$

$q_2 = (2, 0)$

$q_3 = (3, 0)$



H = in hole

F = finding cheese

INV H = invisible in hole