

INDUCTION IS OFTEN EMPLOYED TO PROVE PROPERTIES ABOUT RECURSIVE ALGORITHMS.

→ RECURSION SOLVES A PROBLEM BY SPECIFYING A SOL. OF 1 OR MORE BASE CASES AND DEMONSTRATING HOW TO DERIVE THE SOL TO A PROBLEM OF AN ARBITRARY SIZE FROM SOLS. TO SMALLER PROBLEMS OF THE SAME TYPE.

→ MATHEMATICAL INDUCTION PROVES A PROPERTY ABOUT THE NATURAL NUMBERS BY PROVING THE PROP. ABOUT THE BASE CASE - USUALLY 1 OR 0 - AND THEN PROVING THAT THE PROPERTY MUST BE TRUE FOR ANY ARBITRARY NUM "n". IF IT IS TRUE FOR THE NATURAL #S SMALLER THAN "n".

PROVE FUNCTION FACT RETURNS THE VALUES

$$\text{factorial}(0) = 0! = 1$$

$$\text{factorial}(n) = n! = n \cdot (n-1) \cdot (n-2) \cdot \dots \cdot 1 \text{ if } n > 0$$

PROOF BY INDUCTION ON "n"

BASIS : SHOW THAT PROP IS TRUE FOR $n=0$. SO, SHOW $\text{fact}(0)$ RETURNS 1.

BUT THIS IS SHOWN IN ITS DEF (BASE CASE). MUST ALSO ESTABLISH THAT

PROPERTY IS TRUE FOR AN ARBITRARY $k \Rightarrow$ PROP TRUE FOR $k+1$

INDUCTIVE HYPOTHESIS: ASSUME PROP IS TRUE FOR $n=k$.

$$\text{factorial}(k) = k \cdot (k-1) \cdot (k-2) \cdot \dots \cdot 1 \text{ if } k > 0$$

INDUCTIVE CONCLUSION: SHOW TRUE FOR $n=k+1$

$$\text{factorial}(k+1) = (k+1) \cdot k \cdot (k-1) \cdot \dots \cdot 2 \cdot 1$$

BY DEF $\text{factorial}(k+1)$ RETURNS

$$(k+1) \cdot \text{factorial}(k)$$

\therefore BY INDUCTION $\text{factorial}(k)$ RETURNS

$$k \cdot (k-1) \cdot (k-2) \cdot \dots \cdot 1 \text{ THUS } \text{factorial}(k+1) \text{ RETURNS } (k+1) \cdot k \cdot (k-1) \cdot (k-2) \cdot \dots \cdot 2 \cdot 1 \therefore \text{INDUCTIVE PROOF IS COMPLETE}$$

NOTE: SUBTLETIES OF SOME ALGORITHMS ENCOUNTERED INDICATE THE NEED FOR MATHEMATICAL TECHNIQUES TO PROVE THEIR CORRECTNESS. HELPS TO ELIMINATE ERRORS IN LOGIC & DESIGN OF VARIOUS COMPONENTS OF THE SOL. (E.G. MATHEMATICAL INDUCTION & LOOP INVARIANTS).