OPERATING SYSTEM – Lab 5

Submitted by: Ayesha Javaid

Roll #: BITF22M523

Deadlock Detection and Avoidance

1. Is Deadlock Possible?

Yes, deadlock is possible in this program.

- Each thread (thread1, thread2, thread3) locks resources (m1, m2, m3) in different orders.
- If one thread acquires a lock and another thread acquires a different lock, they can **block each other** indefinitely while waiting for the other lock to be released.

2. Resource Allocation Graph (RAG)

- Threads (T1, T2, T3) represent processes.
- Resources (m1, m2, m3) represent mutex locks.

Graph Representation:

- $\bullet \quad T1 \to m1 \to T1 \to m2 \to T1 \to m3$
- $\bullet \quad T2 \to m2 \to T2 \to m3 \to T2 \to m1$
- $\bullet \quad T3 \to m3 \to T3 \to m1 \to T3 \to m2$

Cycle Detected:

• There is a **circular wait condition** because each thread holds one mutex while waiting for the next mutex held by another thread.

Conclusion: Deadlock exists.

3. Deadlock Conditions (Coffman Conditions)

- 1. **Mutual Exclusion:** Each resource (m1, m2, m3) can only be held by one thread at a time. ✓ Satisfied
- 2. Hold and Wait: Threads hold some resources while waiting for others.✓ Satisfied
- 3. **No Preemption:** Resources cannot be forcibly taken from threads. **Satisfied**
- 4. **Circular Wait:** Threads form a circular chain, where each thread waits for a resource held by the next thread. ✓ Satisfied

Code Reference Example:

- thread1 locks m1 → waits for m2 → waits for m3
- thread2 locks m2 → waits for m3 → waits for m1
- thread3 locks m3 → waits for m1 → waits for m2

4. Deadlock Avoidance Using Banker's Algorithm

Banker's Algorithm can be applied to ensure **safe resource allocation sequences** and prevent deadlock.



