

Q203

As we know that it is undirected graph which means we can go from one to second and can go back to city one from second.

So here if we use DFS then we can use backtracking. It involves exhaustive searches of all the nodes by going ahead, if possible, else by backtracking.

It can be implemented by using stacks.

preprocess(g, v)

visited = true

visited[v].distance = 0

stacks

s.push(v)

~~while (s.push)~~

while (s.size() > 0)

$v' = s.top()$

flag = 0

for each edge (v', u) in E

if not visited(u):

visited(u) = true

$u.distance = v'.distance + edge(v', u)$

s.push(u)

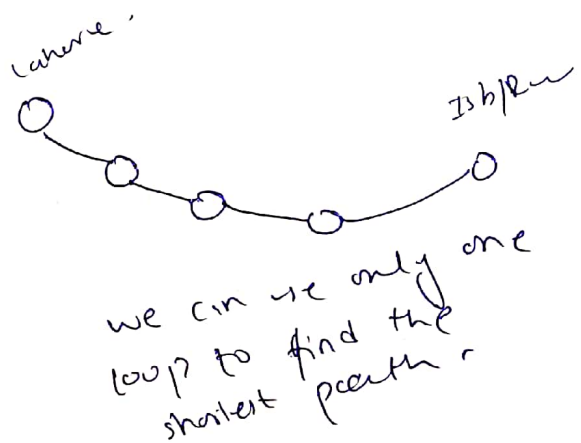
flag = 1

break

if flag == 0

stack.pop()

return $|u.distance - v.distance|$

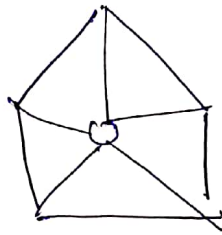


Here the preprocess has run-time of DFS, $O(V+E)$.
thus the best run time must be linear As Distance
have runtime of $O(1)$ since preprocess will
have distance from source.

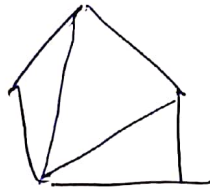
QNO4

if we have 5 cities

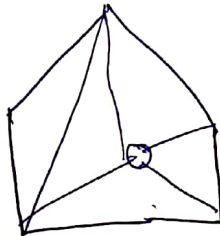
if we introduce bridge b/w them



→ but this not the
requirement



This also cannot
fulfil our requirement



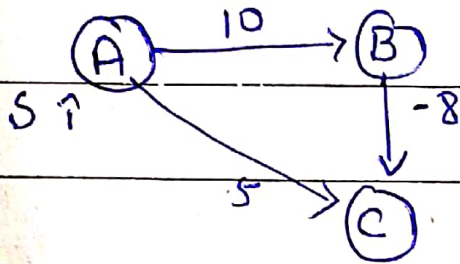
This also cannot
fulfil our requirement

So hence we cannot satisfy NHA requirement
without intersection or building bridges we
cannot connect each city with each other
city.



QNO2

Lets explain with an example suppose we have one graph with negative weight



we know the dijkstra relax the edge only once

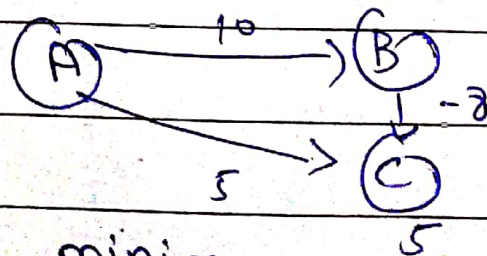
A	B	C
0	∞	∞

initially it will be like this.

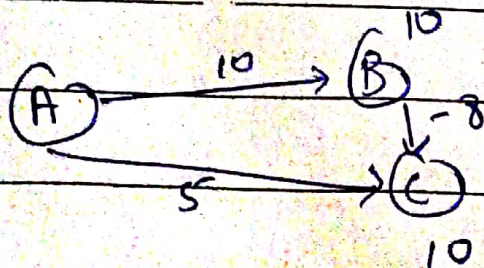
Now when we relax the edges

A	B	C
0	∞	∞
A	10	5
AC	10	5

so here the minimum is 5

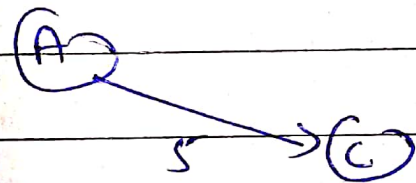


so here the next minimum is 5 10



so at this stage dijkstra will stop because it relax the edge only once. so if we look at this we can clearly see that it is giving wrong shortest path so it fails here!

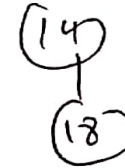
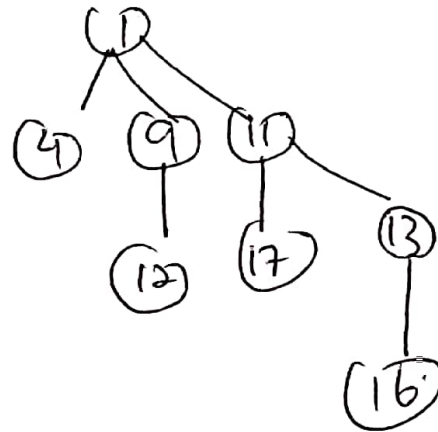
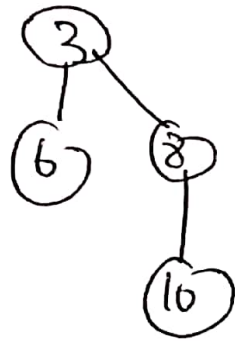
ex. the shortestest here was 2 but it will give 5



For this we use bellmanford to check negative weight edges. because it relax the edge $(v-1)$ times.

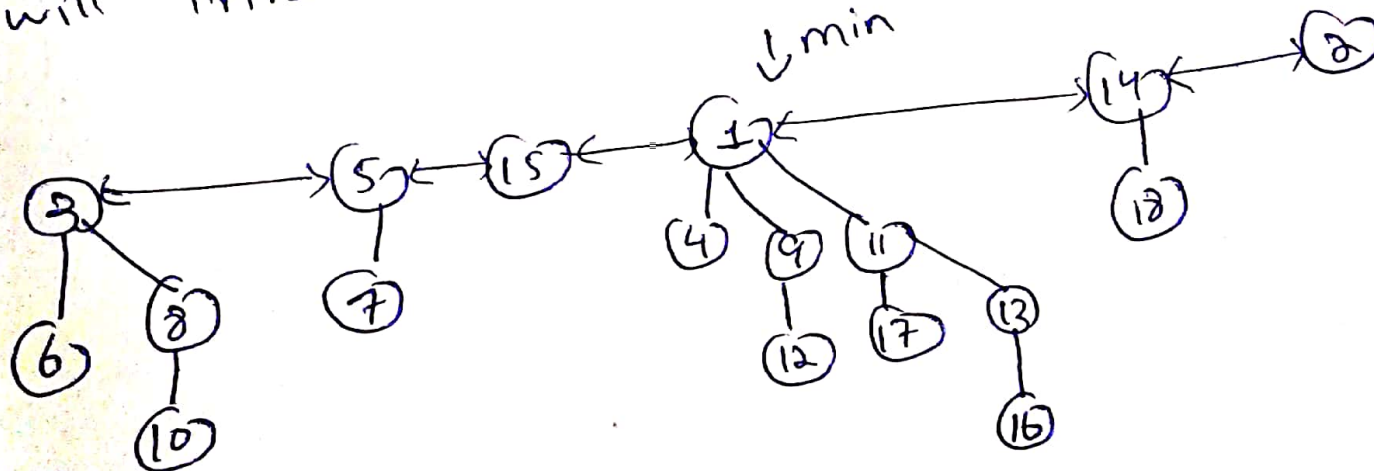
(—————)

Q No 3 (a)



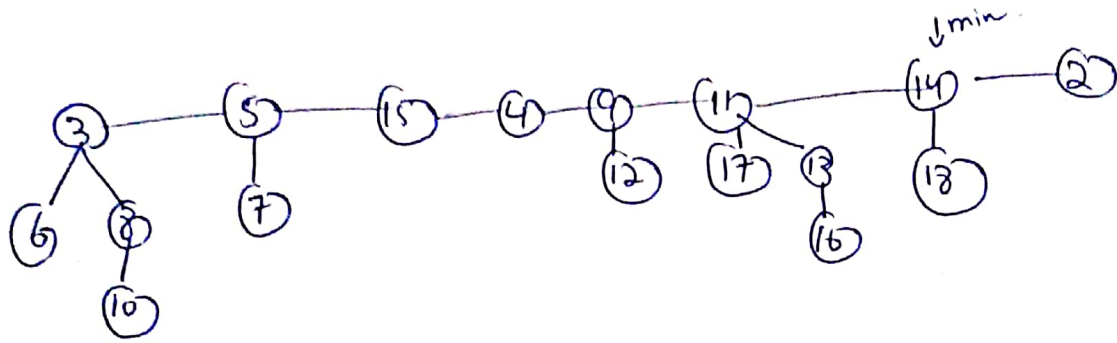
(2)

As we know that these are multiple binomial trees that satisfies the Min-heap property so we will link them to make Fibonacci heap

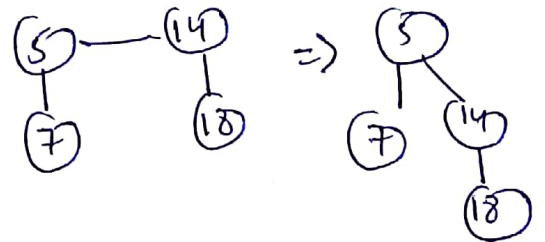


(b) Delete Min

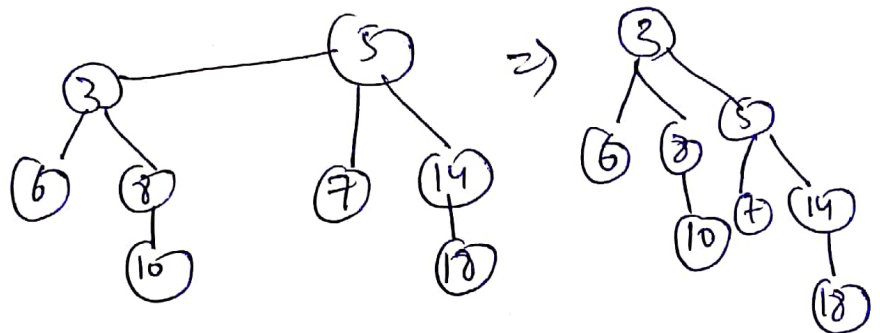
so here we will remove ① because it is minimum here.



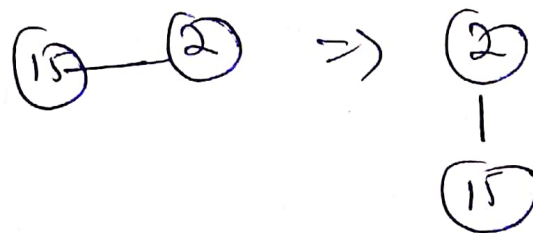
Step 1



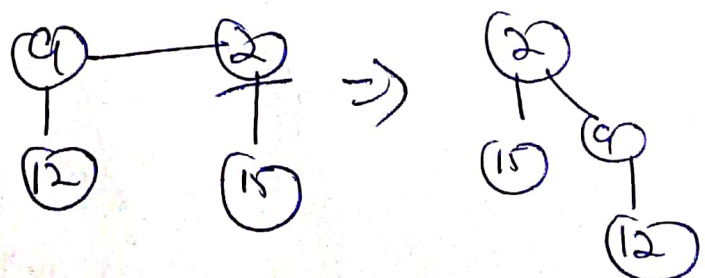
Step 2



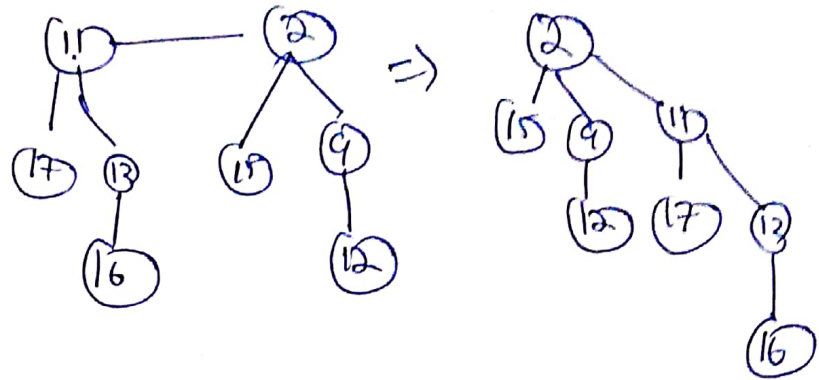
Step 3



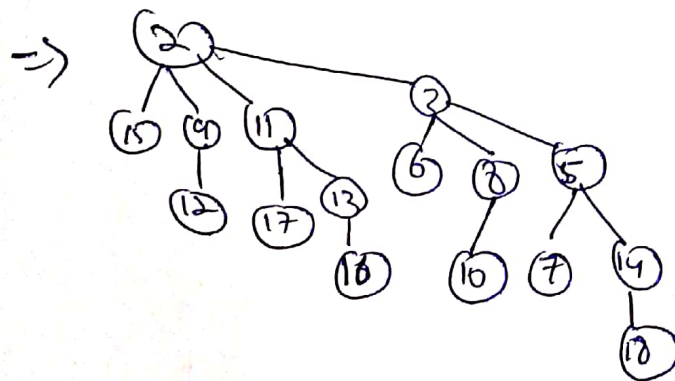
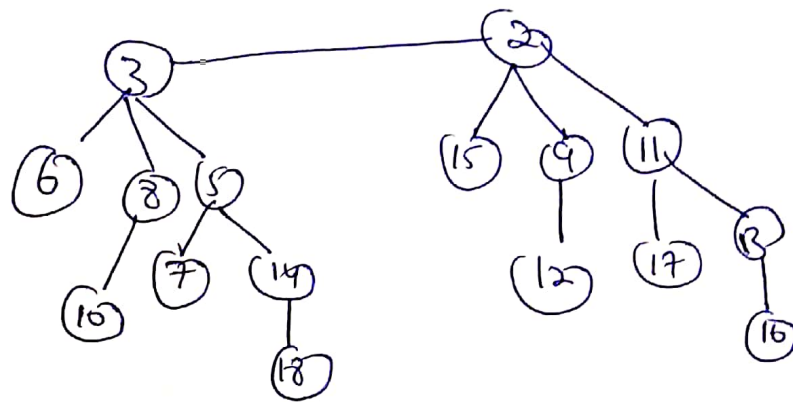
Step 4



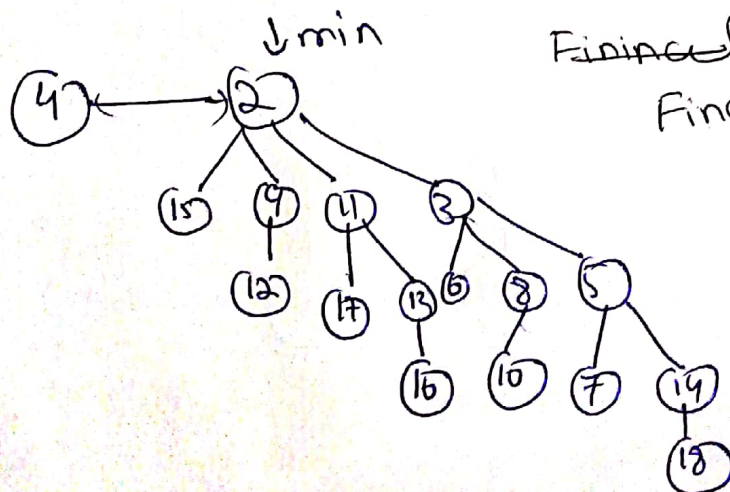
step 5



step 6



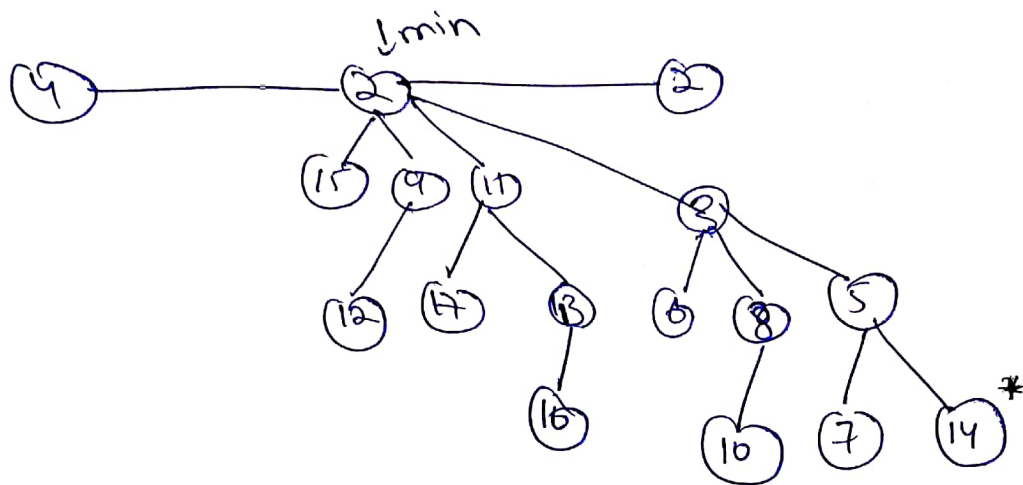
step 7



Final

Final answer

(c) Decrease-key 18 to 2



Here the pointer will be at 2 in the same position of previous. It doesn't matter in case of 2 min (having same values). The pointer will be at previous min.

