

MATH 680 Computation Intensive Statistics

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Duality

1 Duality

- **Primal problem:**

$$\begin{aligned} \min_x f(x) \\ \text{subject to } h_i(x) \leq 0, i = 1, \dots, m \\ \ell_j(x) = 0, j = 1, \dots, r, \end{aligned}$$

so we have m inequality constraints and r equality constraints. In unconstrained problems, we have $m = r = 0$. The above says that we have a properly defined optimization problem to solve.

- **Primal feasible:** x is primal feasible if it satisfies

$$\begin{aligned} h_i(x) \leq 0, 1 \leq i \leq m \\ \ell_j(x) = 0, 1 \leq j \leq r. \end{aligned}$$

Let C denotes the set of x 's that are primal feasible, then C is called a **primal feasible set**.

- **Primal optimal:** define primal optimal x^* as

$$x^* = \arg \min_{x \in C} f(x)$$

Denote

$$f^* = f(x^*).$$

• **Lagrangian:**

$$L(x, u, v) = f(x) + \sum_{i=1}^m u_i h_i(x) + \sum_{j=1}^r v_j \ell_j(x). \quad u \in \mathbb{R}^m, u \geq 0$$

Important property: for each feasible x

$$f(x) \geq L(x, u, v).$$

• **Lagrangian dual function:**

$$g(u, v) = \min_x L(x, u, v).$$

Proposition 1. $g(u, v)$ is concave.

Proof. For any (u_1, v_1) , (u_2, v_2) and $0 \leq \alpha \leq 1$, let $u^\alpha = \alpha u_1 + (1 - \alpha)u_2$, $v^\alpha = \alpha v_1 + (1 - \alpha)v_2$ and

$$g(u^\alpha, v^\alpha) = \min_x L(x, u^\alpha, v^\alpha).$$

Note that

$$\begin{aligned} L(x, u^\alpha, v^\alpha) &= f(x) + \sum_{i=1}^m u_i^\alpha h_i(x) + \sum_{j=1}^r v_j^\alpha \ell_j(x) \\ &= \alpha L(x, u_1, v_1) + (1 - \alpha)L(x, u_2, v_2), \end{aligned}$$

which implies that

$$\begin{aligned} L(x, u^\alpha, v^\alpha) &\geq \alpha \min_x L(x, u_1, v_1) + (1 - \alpha) \min_x L(x, u_2, v_2) \\ &= \alpha g(u_1, v_1) + (1 - \alpha)g(u_2, v_2). \end{aligned}$$

It follows that

$$g(u^\alpha, v^\alpha) = \min_x L(x, u^\alpha, v^\alpha) \geq \alpha g(u_1, v_1) + (1 - \alpha)g(u_2, v_2).$$

□

Proposition 2. Let C denote primal feasible set. If $u_i \geq 0$, then Lagrange dual function is always a lower bound of f^* . i.e.

$$g(u, v) = \min_x L(x, u, v) \leq \min_{x \in C} L(x, u, v) \leq f^* \leq f(x).$$

- **Dual problem:**

$$\max_{u, v} g(u, v)$$

subject to $u \geq 0$

Dual is a concave maximization problem $\Leftrightarrow \min_{u \geq 0} -g(u, v)$ is a convex minimization problem.

- **Dual feasible:** u is dual feasible if $u \geq 0$.

- **Dual optimal:** define the optimal solution (u^*, v^*) of dual problem as

$$(u^*, v^*) = \arg \max_{u \geq 0, v} g(u, v),$$

Denote

$$g^* = g(u^*, v^*).$$

- **Duality gap:** given primal feasible x and dual feasible u, v , the quantity $f(x) - g(u, v)$ is called the duality gap between x and u, v . Note that since $f^* \geq g(u, v)$

$$f(x) - f^* \leq f(x) - g(u, v)$$

Proposition 3. *If the duality gap $f(x_0) - g(u_0, v_0) = 0$, then x_0 is primal optimal (and similarly, u_0, v_0 are dual optimal).*

Proof. Since

$$f(x_0) - f^* \leq f(x_0) - g(u_0, v_0) = 0$$

So

$$f(x_0) = f^*$$

So x_0 is primal optimal. Similarly by

$$g^* - g(u_0, v_0) \leq f(x_0) - g(u_0, v_0) = 0$$

we know that u_0 and v_0 are dual optimal. □

- **Weak duality:** weak duality $g^* \leq f^*$ is always true by Proposition 2.
- **Slater's condition:** there exists at least one strictly feasible $x_0 \in \mathbb{R}^n$, in other words

$$h_1(x_0) < 0, \dots, h_m(x_0) < 0 \text{ and } \ell_1(x_0) = 0, \dots, \ell_r(x_0) = 0.$$

Slater's condition states that the feasible region must have **an interior point**

- **Strong duality:** $g^* = f^*$ is not always true. It requires two conditions:

Proposition 4. *Strong duality holds if (1) the primal is a convex problem (i.e. f and h_1, \dots, h_m are convex, ℓ_1, \dots, ℓ_r are affine) and (2) Slater's condition is satisfied.*