

# PS. #2

## Solutions

6.3 Sol1. Write  $A = U \Sigma V^*$  where  $\Sigma = \begin{pmatrix} \sigma_1 & & 0 \\ & \ddots & \\ 0 & & \sigma_n \end{pmatrix}$   $\sigma_1 \geq \sigma_2 \geq \dots \geq \sigma_k > 0$  and  $\sigma_{k+1} = \dots = \sigma_n = 0$ .

Then  $\text{rank}(A) = k$ . Computing  $A^*A = V \Sigma^2 V^*$  we see that  $\text{rank}(A^*A) = \text{rank}(\Sigma^2) = k$ .  
So  $A^*A$  is nonsingular  $\Leftrightarrow \text{rank}(A^*A) = n$   
 $\Leftrightarrow \text{rank}(\Sigma^2) = n \Leftrightarrow \text{rank}(A) = n$ .

Sol2. If  $A$  has full rank then  $A^*A$  is nonsingular since  $A^*A x = 0 \Rightarrow x^* A^* A x = 0 \Rightarrow \|A x\|^2 = 0 \Rightarrow A x = 0 \Rightarrow x = 0$ .  
If  $A^*A$  is nonsingular then  $A$  has full rank since  $A x = 0 \Rightarrow A^* A x = 0$  and then  $x = 0$ .

7.3 Write  $A = QR$ ,  $Q$  unitary ( $Q^*Q = I$ ).  
So  $\det(Q) = \pm 1$  and then  $|\det A| = |\det R|$ .  
But  $R = \begin{pmatrix} r_{11} & & 0 \\ & \ddots & \\ 0 & & r_{nn} \end{pmatrix}$  and  $|\det R| = \prod_j |r_{jj}|$

Now  $r_{jj} = \|P_{\perp} a_j\|_2 \leq \|a_j\|_2$  where  $P_{\perp}$  is the projection on the orthogonal space to  $\langle a_1, \dots, a_{j-1} \rangle$ .

The geometric interpretation comes from the fact that  $|\det A|$  is the "volume" of the parallelepiped determined by the column vectors of  $A$ . The volume is maximized (if we fix the lengths of the vectors) when the vectors are orthogonal to each other, in which case the volume is just the product of the lengths.

7.5 a)  $A = \hat{Q} \hat{R}$ . Then  $A$  has rank  $n$  if and only if  $\hat{R}$  has diagonal entries nonzero. Indeed if  $Ax = 0$  then  $\hat{Q} \hat{R} x = 0$  and then  $\hat{R} x = 0$ . Now let  $\hat{R} = \begin{pmatrix} r_{11} & & 0 \\ & \ddots & \\ 0 & & r_{nn} \end{pmatrix}$  and  $x = \begin{pmatrix} x_1 \\ \vdots \\ x_n \end{pmatrix}$ . If  $r_{nn} \neq 0, \dots, r_{k+1, k+1} \neq 0$  but  $r_{kk} = 0$  we can show that  $\hat{R} x = 0$  implies  $x_n = x_{n-1} = \dots = x_{k+1} = 0$  and then we can choose  $x_k = 1$  and determine the rest of  $x_{k-1}, \dots, x_1$ . So we get a nonzero vector  $x$  such that  $\hat{R} x = 0 \Rightarrow Ax = 0$ , so a contradiction.

If all  $r_{ii}$ 's are nonzero then  $A$  has full rank and if  $A$  has full rank then  $r_{ii}$ 's are nonzero, these follow by the argument given above or the reverse of it (just take and prove  $\text{rank}(\hat{R}) = n$  iff all  $r_{ii}$ 's are nonzero).

b) Using the argument above it can be shown that if  $\hat{R} = \begin{pmatrix} r_{11} & & 0 \\ & \ddots & \\ 0 & & r_{kk} & 0 \end{pmatrix}$  then  $\text{rank}(A) \geq k$ . So  $\text{rank}(A) \in \{k, \dots, n\}$  because  $A$  can not have full rank. There is an example when  $k=1$  and  $\text{rank}(A)$  can be any value in  $\{k, \dots, n-1\}$ .

Take  $\hat{R} = \begin{pmatrix} 1 & 0 & \dots & 0 \\ 0 & 0 & \dots & 0 \\ & & \ddots & \\ 0 & & & 1 & 0 & \dots & 0 \\ & & & 0 & \ddots & & 0 \\ & & & & & \ddots & \\ & & & & & & 1 \end{pmatrix}_{k \times k}$

and  $\hat{Q} = [q_1 | \dots | q_n]$

we get  $A = \hat{Q} \hat{R} = [q_1 | 0 | \dots | q_k | 0 | \dots | q_n]$   
so  $A$  has rank  $k$ .

Problem 8.2

function modif GS = mgs(A)

% compute the reduced QR decomposition of A using the modified GS algorithm.

% compute m and n, the number of rows and columns of A

m = size(A,1); n = size(A,2);

for i = 1:n

    v = A(:,i);

    R(i,i) = norm(v);

    Q(:,i) = v / R(i,i); % set the i<sup>th</sup> orthogonal vector of Q.

% inner loop of the modified GS's

% every time a new orthogonal vector is computed

% project all remaining columns vectors of A onto the space of orthogonal to that vector.

for j = (i+1):n

    R(i,j) = conj(Q(:,i)' \* A(:,j))

    A(:,j) = A(:,j) - R(i,j) \* Q(:,i)

end

end.

R = [R zeros(m-n, n)];

modif GS = [QR]

% Recover Q and R

mQR = clog(A);

m = size(mQR,1); n = size(mQR,2)/2;

QmGS = mQR(:,1:n); RmGS = mQR(1:n, n+1:2\*n).