#### Exercises: Artificial Intelligence

Automated Reasoning: Movable Objects

#### Solution: Movable Objects

- English to logic
- Logic to implicative normal form
  - Model
  - Assumption to prove
- Apply resolution
  - Derive inconsistency:
  - Model + negated assumption

#### Solution: Model to logic

- If all movable objects are blue, then all non-movable objects are green.
  - $-(∀x mov(x) \rightarrow blue(x)) \rightarrow (∀y \neg mov(y) \rightarrow green(y))$
- If there exists a non-movable object, then all movable objects are blue.
  - $-(\exists x \neg mov(x)) \rightarrow (\forall y mov(y) \rightarrow blue(y))$
- D is a non-movable object.
  - $-\neg mov(D)$

- $(\forall x \text{ mov}(x) \rightarrow \text{blue}(x)) \rightarrow (\forall y \neg \text{mov}(y) \rightarrow \text{green}(y))$ 
  - mov(A) ∨ mov(y) ∨ green(y) (← true) mov(y) ∨ green(y) ← blue(A)
- $(\exists x \neg mov(x)) \rightarrow (\forall y mov(y) \rightarrow blue(y))$ 
  - mov (x)  $\vee$  blue(y)  $\leftarrow$  mov(y)
- ¬mov(D)
  - false  $\leftarrow$  mov(D)
- Negated assumption:  $\neg \exists x \ green(x) \leftrightarrow \forall x \ \neg green(x)$ 
  - false  $\leftarrow$  green(x)

- Prove using resolution:
  - Assumption: false  $\leftarrow$  green(x)
- Model:
  - $mov(A) \lor mov(y) \lor green(y) (\leftarrow true)$
  - $mov(y) \lor green(y) \leftarrow blue(A)$
  - $-\operatorname{mov}(x) \vee \operatorname{blue}(y) \leftarrow \operatorname{mov}(y)$
  - false  $\leftarrow$  mov(D)

#### Solution: Resolution

```
mov(x) \lor blue(y) \leftarrow mov(y)
                                                                           false \leftarrow mov(D)
                        blue(y) \leftarrow mov(y)
                                                                           mov(y1) \vee green(y1) \leftarrow blue(A)
               mov(y1) \vee green(y1) \leftarrow mov(A)
                                                                          mov(A) \vee mov(y2) \vee green(y2)
mov(y1) \vee green(y1) \vee mov(y2) \vee green(y2) \leftarrow true
                mov(y1) \vee green(y1) \leftarrow true
                                                                           false \leftarrow mov(D)
                                                                           false \leftarrow green(x)
                        green(D) \leftarrow true
                           false ← true
```

#### Exercises: Artificial Intelligence

Automated Reasoning: Politicians

#### **Problem: Politicians**

#### Given:

- If a poor politician exists, then all politicians are male.
- If people are friends with a politician, then this politician is poor and female.
- Lazy people have no friends.
- People are either male or female, but not both.
- If Joel is not lazy, then he is a politician.

#### Proof by resolution:

There exists no person who is a friend of Joel.

#### Solution: English to logic

- $(\exists x \text{ pol}(x) \land \text{poor}(x)) \rightarrow (\forall y \text{ pol}(y) \rightarrow \text{male}(y)).$
- $\forall x (pol(x) \land (\exists y fr(y,x))) \rightarrow poor(x) \land fem(x).$
- $\forall x | \text{lazy}(x) \rightarrow (\neg(\exists y | \text{fr}(y,x))).$
- $\forall x (male(x) \lor fem(x)) \land (\neg(male(x) \land fem(x))).$
- $\neg$ lazy(Joel)  $\rightarrow$  pol(Joel).

- male(y) ← pol(x) ∧ poor(x) ∧ pol(y)
- $poor(x) \leftarrow pol(x) \land fr(y,x)$
- $fem(x) \leftarrow pol(x) \land fr(y,x)$
- false  $\leftarrow$  lazy(x)  $\wedge$  fr(y,x)
- male(x)  $\vee$  fem(x)
- false  $\leftarrow$  male(x) $\land$ fem(x)
- lazy(Joel) v pol(Joel)

#### • Prove:

- There exists no person who is a friend of Joel
  - $\neg \exists x \ fr(x,Joel) \leftrightarrow \forall x \ \neg fr(x,Joel)$
- Negate assumption:
  - There exists a person who is a friend of Joel
    - ∃x fr(x,Joel)
  - Call the friend S
    - fr(S,Joel)

- male(y)  $\leftarrow$  pol(x)  $\land$  poor(x)  $\land$  pol(y)
- $poor(x) \leftarrow pol(x) \land fr(y,x)$
- $fem(x) \leftarrow pol(x) \land fr(y,x)$
- false  $\leftarrow$  lazy(x)  $\land$  fr(y,x)
- male(x) \( \times \) fem(x)
- false ← male(x) ∧ fem(x)
- lazy(Joel) v pol(Joel)
- fr(S,Joel)

- male(y1) ← pol(x1) ∧ poor(x1) ∧ pol(y1)
  - poor(x2)  $\leftarrow$  pol(x2)  $\land$  fr(y2,x2)
  - RESOLUTION: {x2/x1}
- male(y1) ← pol(x1) ∧ pol(y1) ∧ fr(y2,x1)
  - FACTORING: {y1/x1}
- male(x1) ← pol(x1) ∧ fr(y2,x1)
  - 'Politicians who have friends must be male'

- male(x1) ← pol(x1) ∧ fr(y2,x1)
  - false  $\leftarrow$  male(x3)  $\land$  fem(x3)
  - RESOLUTION: {x3/x1}
- false ← pol(x1) ∧ fr(y2,x1) ∧ fem(x1)
  - Politicians who have friends cannot be female'

- false ← pol(x1) ∧ fr(y2,x1) ∧ fem(x1)
  - fem(x4) ← pol(x4)  $\wedge$  fr(y4,x4)
  - RESOLUTION: {x4/x1}
- false ← pol(x1) ∧ fr(y2,x1) ∧ pol(x1) ∧ fr(y4,x1)
  - FACTORING: {}
- false ← pol(x1) ∧ <u>fr(y2,x1)</u> ∧ <u>fr(y4,x1)</u>
  - FACTORING: {y4/y2}
- false ← pol(x1) ∧ fr(y2,x1)
  - 'Politicians do not have friends'

- false ← pol(x1) ∧ fr(y2,x1)
  - lazy(Joel) \times pol(Joel)
  - RESOLUTION: {x1/Joel}
- lazy(Joel) ← fr(y2,Joel)
  - 'If Joel has friend, then he must be lazy'

- <u>lazy(Joel)</u> ← fr(y2,Joel)
  - false ← lazy(x5) ∧ fr(y5,x5)
  - RESOLUTION: {x5/Joel}
- false ← <u>fr(y2,Joel)</u> ∧ <u>fr(y5,Joel)</u>
  - FACTORING: {y5/y2}
- false ← fr(y2,Joel)
  - 'Joel does not have any friends'

- false ← <u>fr(y2,Joel)</u>
  - fr(S,Joel)
  - RESOLUTION: {y2/\$}
- false ← true