

# Exercise session 8: LTL and CTL

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## 1 Check equivalence

Which of the following pairs of CTL formulas are equivalent? For those which are not, exhibit a model of one of the pair which is not a model of the other:

1.  $EF \phi$  and  $EG \phi$
2.  $EF \phi \vee EF \psi$  and  $EF (\phi \vee \psi)$
3.  $AF \phi \vee AF \psi$  and  $AF (\phi \vee \psi)$
4.  $AF \neg \phi$  and  $\neg EG \phi$
5.  $EF \neg \phi$  and  $\neg AF \phi$
6.  $A (\phi_1 \cup A (\phi_2 \cup \phi_3))$  and  $A (A (\phi_1 \cup \phi_2) \cup \phi_3)$ , hint: it might make it simpler if you think first about models that have just one path
7.  $\top$  and  $AG \phi \rightarrow EG \phi$
8.  $\top$  and  $EG \phi \rightarrow AG \phi$
9.  $A [\phi \cup \psi]$  and  $\phi \wedge AF \psi$
10.  $A [\phi \cup \psi] \vee A [\tau \cup \psi]$  and  $A [(\tau \vee \phi) \cup \psi]$

## 2 Express in CTL and LTL

Express the following properties in CTL and LTL whenever possible. If neither is possible, try to express the property in CTL\*:

1. Whenever  $p$  is followed by  $q$  (after finitely many steps), then the system enters an 'interval' in which no  $r$  occurs until  $t$ .
2. Event  $p$  precedes  $s$  and  $t$  on all computation paths. (You may find it easier to code the negation of that specification first)
3. After  $p$ ,  $q$  is never true. (Where this constraint is meant to apply on all computation paths.)
4. Between the events  $q$  and  $r$ , event  $p$  is never true.
5. Transitions to states satisfying  $p$  occur at most twice.
6. Property  $p$  is true for every second state along a path.

### 3 Expressable in ...

1. Give example of an LTL-formula for which equivalent translation in CTL does not exist.
2. Give example of an CTL-formula for which equivalent translation in LTL does not exist.
3. Give example of an CTL\*-formula for which equivalent translation in LTL either in CTL does not exist.

### 4 Proof the equivalence

Given the definitions:

- $\pi \models \psi \cup \phi$  iff there is some  $i \geq 1$  such that  $\pi^i \models \phi$  and for all  $j = 1, \dots, i - 1$  we have  $\pi^j \models \psi$
- $\pi \models \psi \text{ R } \phi$  iff either there is some  $i \geq 1$  such that  $\pi^i \models \psi$  and for all  $j = 1, \dots, i$  we have  $\pi^j \models \phi$ , or for all  $k \geq 1$  we have  $\pi^k \models \phi$

Proof the following theorem:

$$\neg(\psi \cup \phi) \equiv \neg\psi \text{ R } \neg\phi \quad (1)$$

### 5 Nim game

*If you have time left, and haven't made the Nim game yet last session, complete this exercise. The assignment from last week can still be found on Toledo.*