

M.Tech Computer Science Information security.

Cryptography Basic.(CS6530)

Assignment-4. (C language Program for implementation of ElGamal algorithm for encryption and decryption.)

Roll Number - CS21M515

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Batch Year-2021

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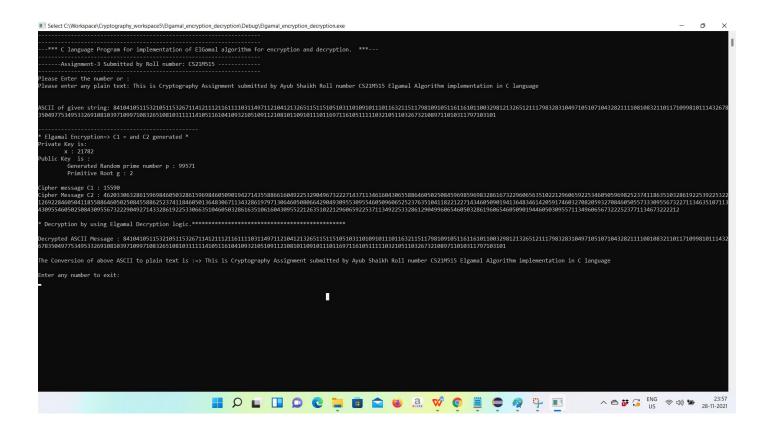
TA: Honorable Venkatakrishnan Sutharsan Sir

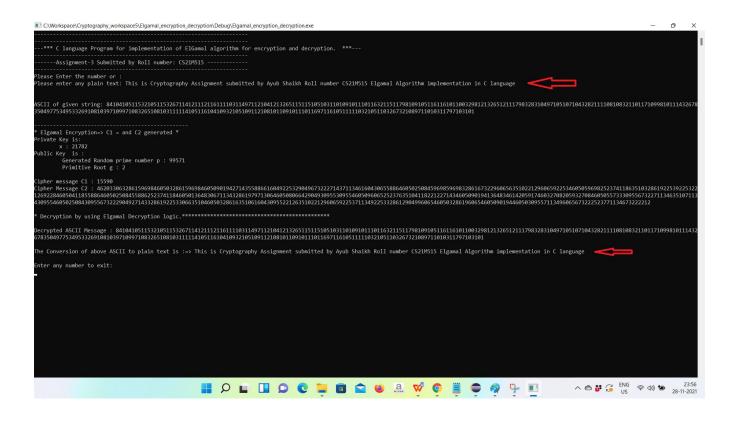
Program Solution and Explanations:-

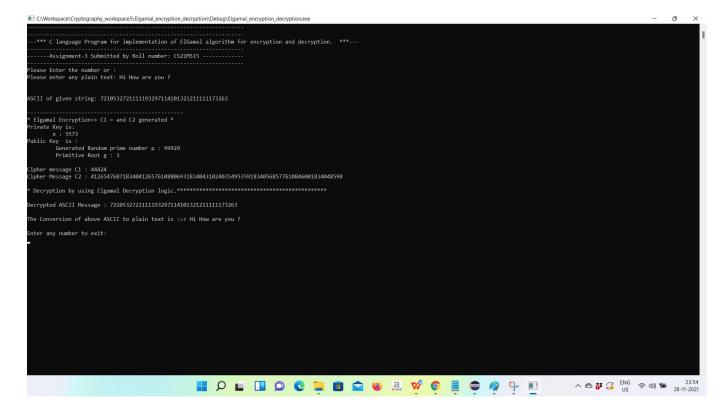
I Have did the above example by using the C program.

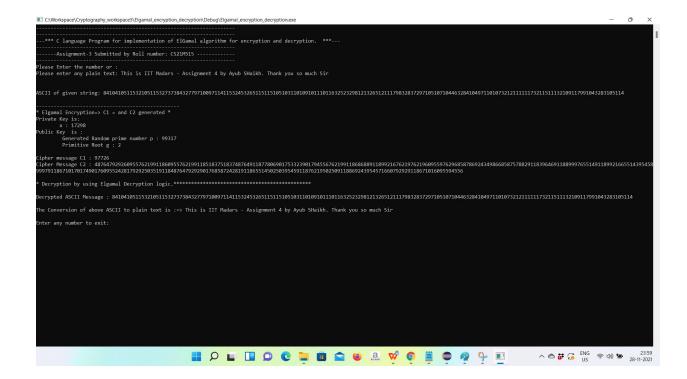
And .exe is shared in the Github and also the program is shared.

Below are the screen shots of the C programm executed .









ElGamal algorithm for encryption and decryption Explanations and Program Logical Steps Explanations:-

Key-Generation Algorithm:

let A be a person creates key!

1. select large number p and large number 9.

p is generator.

2. Seelects random x in Z*p

3. y = 9x mod p.

(P. g. y) -> (publice key)

z -> (secret /private key)

finding or for given 7.9.4 is computationally infemible as of now.

(where x can not be 1 or p-1)

Encryption under Kommande patternation

1) choose a random r 1<x<p-1)

2) C1 = g mod P

(see it is (9 mod p) mod p

that (927) mod p)

m is 1023 bits

3) G = m. y mod p. (mis meisage m<p)

C, C2 are Cipher text.

for RSA, the relation between

c and m are fixed but

for Elgamal, c and m are not

fixed) as all regime ellewated as me

Decryption:

(c, c2 (P, 9, y), x) 1 private Key

seven feet de se de se de se de combination de pla combination de pla combination de se de combination de se de combination de se de combination de combinat

using this decrypting c, = g* mod P, ca = my mod p is impossible for Decryptor) owner of the key too.

So other way is (all are med P) y'= gx = (g') = c,

So Decrypting Ci

- 1 compute Cix mod p
- (2) compute $m = \frac{C_2}{C_1^{\pm}}$

m = (((, x)) mod p

By Euclid Extended Algorithm. we can find m

Program Logical Steps:-

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* Program Logic
 * 1) Take large prime numbers p and g (g <p, g is preferably the prime root of p)
Note: If g is a prime root of the prime number p, then g mod p, g^2 mod p, ..., g^p-1
mod p are permutations from 1 to p-1
2) Randomly select an integer x (2 <= x <= (p-2), (p, g, x) is the private key)
3) Calculate y = g^x (\underline{mod} p) ((p, g, y) is the public key)
2. Encryption process
1) Randomly select an integer k (2 <= k <= (p-2) and k and (p-1) are relatively prime)
2) Calculate a = g^k \mod p, b = m^*y^k \mod p (m is the plaintext to be encrypted)
3) Ciphertext C = (a, b)
3. <u>Decryption</u> process
1. m = b * a^{-}(-x) \mod p
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Note: $b * a^{(-x)} \mod p = m * y^k * g^{(-xk)} = m * g^{(xk)} * g^{(-xk)} = m$

THANK You Sir 🙂 😅 😅