

Analysis

Given:

$$(a_1, password * g^{a_1}) = (324, 11226815350263531814963336315)$$

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we want to find password given elements are multiplicative group Z_p^* where

$p=19807040628566084398385987581$ is prime

$$password * g^{a_1} \equiv A_1(mod \ p) \text{ --- (i)}$$

$$password * g^{a_2} \equiv A_2(mod \ p) \text{ --- (ii)}$$

$$password * g^{a_3} \equiv A_3(mod \ p) \text{ --- (iii)}$$

Multiplying and dividing left hand side of equation 2 and 3 from g^{a_1}

$$(password * g^{a_1}) * g^{(a_2-a_1)} \equiv A_2(mod \ p)$$

$$(password * g^{a_1}) * g^{(a_3-a_1)} \equiv A_3(mod \ p)$$

Using Property: $(a * b)(mod \ m) \equiv a(mod \ m) * b(mod \ m)$

$$((password * g^{a_1}))(mod \ p) * g^{(a_2-a_1)}(mod \ p) \equiv A_2(mod \ p)$$

$$(password * g^{a_1})(mod \ p) * g^{(a_3-a_1)}(mod \ p) \equiv A_3(mod \ p)$$

Using equation (i) we can substitute $((password * g^{a_1})(mod \ p))$

get equation (iv) (v)

$$A_1 g^{(a_3-a_1)} \equiv A_3(mod \ p) \text{ --- (iv)}$$

$$A_1 g^{(a_2-a_1)} \equiv A_2(mod \ p) \text{ --- (v)}$$

$$a_3 - a_1 = 9513 - 324 = 9189$$

$$a_2 - a_1 = 2345 - 324 = 2021$$

we found X and Y. And now apply Extended Euclidean Algorithm

X	Y	q	r	l_1	l_2	l_3	m_1	m_2	m_3
9189	2021	4	1105	1	0	1	0	1	-4
2021	1105	1	916	0	1	-1	1	-4	5
1105	916	1	189	1	-1	2	-4	5	-9
916	189	4	160	-1	2	-9	5	-9	41
189	160	1	29	2	-9	11	-64	41	-50
160	29	5	15	-9	11	-64	41	-50	291
29	15	1	14	11	-64	75	-139	-291	-341
15	14	1	1	-64	75	-139	-291	-341	632
14	1	14	0	75	-139	2021	-341	632	-9189

So, we found the following:

$$gcd(9189, 2021) = 1 \text{ --- (vi)}$$

$$l = -139 \text{ and } m = 632$$

Using Extended Euclid Algorithm

$$X_l + Y_m = gcd(X, Y)$$

From Equation (vi) we Know $gcd(X, Y) = 1$

$$X_l + Y_m = 1 \text{ --- (vii)}$$

$$(A_1^l g^{X_l}) \equiv A_3^l \pmod{p} \text{ --- (viii)}$$

$$A_1^m g^{Y_m} \equiv A_3^m \pmod{p} \text{ --- (xi)}$$

multiply equation (viii) and (xi)

$$A_1^{(l+m)} * g^{(X_l+Y_m)} \equiv (A_3^l * A_3^m) \pmod{p}$$

Using Equation (vii)

$$A_1^{(l+m)} * g \equiv (A_3^l * A_3^m) \pmod{p}$$

$$(A_3^l * A_3^m) \pmod{p} \equiv ((A_3^l \pmod{p}) * (A_3^m \pmod{p})) \pmod{p}$$

$l=139$ so find $(A_3^l \pmod{p})$ for $l=139$ then find the multiplicative inverse of that number. i.e $A_3^l * A_3^{-l} \pmod{p} = 1$

Now if we take RHS as $num = (A_3^l * A_3^m) \pmod{p}$

$$A_3^{139} \pmod{p} = 1438737264732336067040734445$$

$$A_2^{(632)} \pmod{p} = 9086425608952457377582771788$$

Using above values we get value of num:

$$num = 11099199913639351335199364706$$

now equation become

$$A_1^{(l+m)} * g \equiv num \pmod{p} \quad (x)$$

we know that

$$(a * b) \pmod{p} \equiv num$$

$$b \equiv (a^{-1} * num) \pmod{p}$$

now equation (x) becomes

$$g \equiv (A_1^{-(l+m)} * num) \pmod{p}$$

again we have to find modular inverse of $A_1^{(l+m)}$

value of $A_1^{(l+m)} = 2609020618887623880099546994$ lets take

$$a = A_1^{-(l+m)}$$

Value of a = 11763215952453956375186720348

now our equation became

$$g \equiv (a * num) \pmod{p} \quad (xi)$$

now find g from equation (xi) $g = 192847283928500239481729$ then for finding password lets take equation (ii) we can take any equation (i),(ii),(iii)

$$password * g^{a_2} \equiv A_2 \pmod{p} \quad (ii)$$

We know g, a_2 , A_2 and p

$$password \equiv (A_2 * g^{-a_2}) \pmod{p} \quad (xii)$$

find multiplicative inverse of g^{a_2} and put equation (xii)

$$g^{-a_2} = 6507145214719002336566928446$$

we get password from equation (xii)

password = 3608528850368400786036725