

# CS & IT ENGINEERING

Compiler Design



## Lexical Analysis

One Shot

One Shot



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# Topics to be Covered



GATE



Topic

Lexical Analysis

Topic

Parsing

(Syntax Analysis)

Topic

Syntax Directed Translation

(Semantic Analysis)

Topic

Intermediate Code Generation

Topic

Code Optimisation(Dataflow Analysis)

Topic

Runtime Environment



# Topics to be Covered



Topic

??????

Lexical Analysis.





## Topic : Compiler Design

- Lexical Analysis
- Parsing
- Syntax Directed Translation
- Intermediate Code Generation
- Runtime Environment



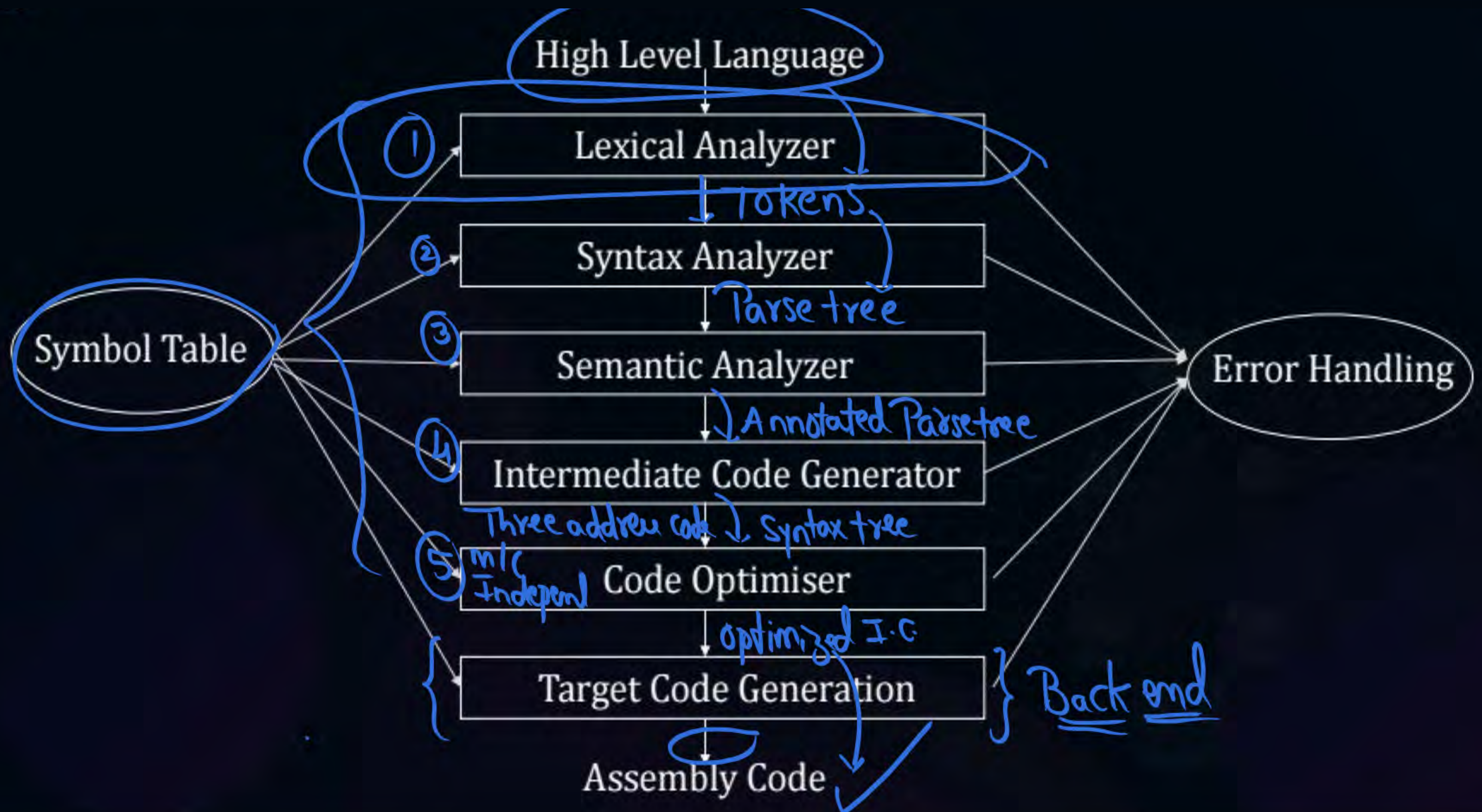


## Topic : Introduction of Compiler Design

- Compiler is a translator which converts a program written in one language (Source Language) to other language (Object/Target Language).











## Topic : Semantic Analysis Phase

- It is a program that takes parse tree as input and produces annotated parse tree as output.
- It also detects semantic errors present in the program.(type checking)



## Topic : Intermediate Code Generation Phase

- It is a program that translates high level code into intermediate code.
- Advantage of generating intermediate code is to perform optimization in simple way.





## Topic : Code Optimization Phase

- It is a program that reduces time and space required by the target machine by removing some unnecessary code .
- There are two types of optimizations performed by compilers known as machine independent optimization and machine dependent optimization.

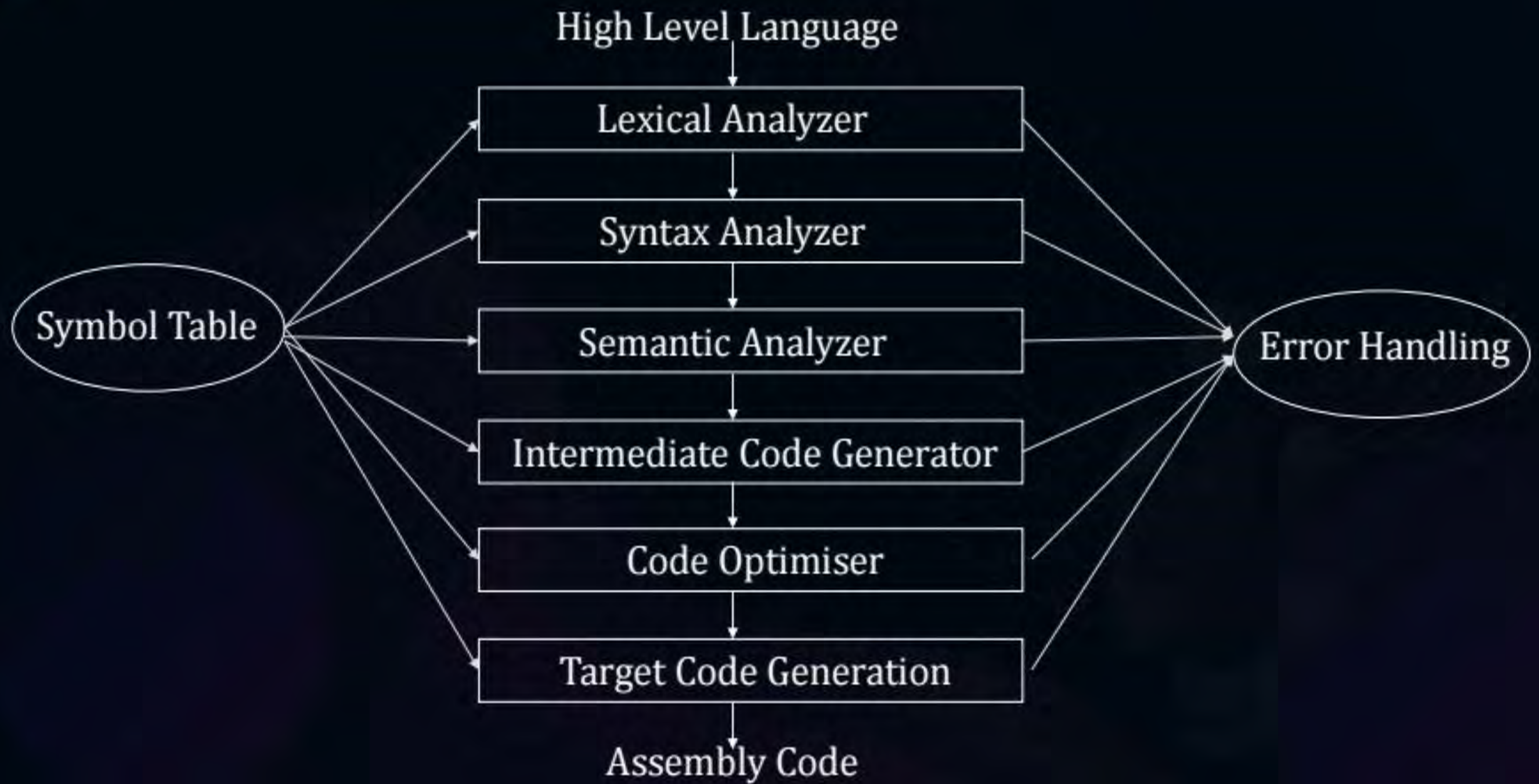




## Topic : Code Generation Phase

- It is a program that translates optimized intermediate code into assembly language or target code.







## Topic : Functionality of Lexical Analysis

➤ ① It reads source program as stream of characters (Scanner) (abc)  
↑ ↑

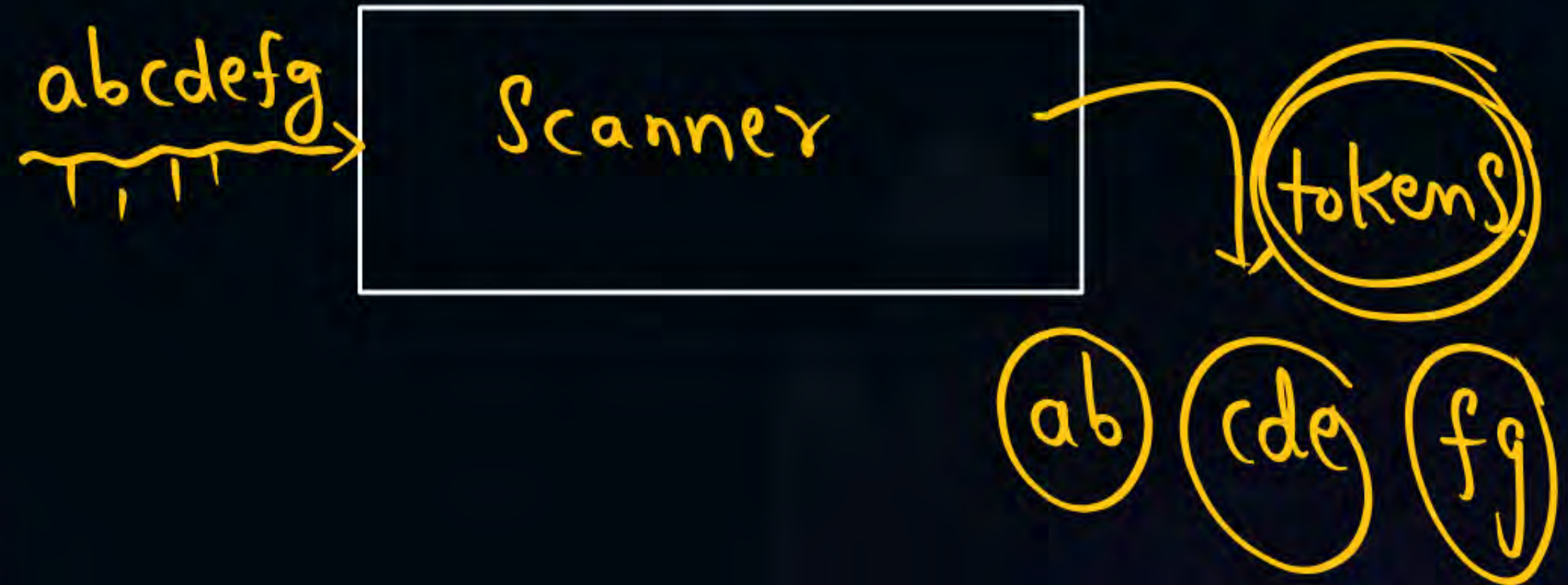
➤ ② It produces tokens as output

➤ ③ It detects lexical errors

➤ ④ It eliminates comment lines

➤ ⑤ It eliminates white space characters

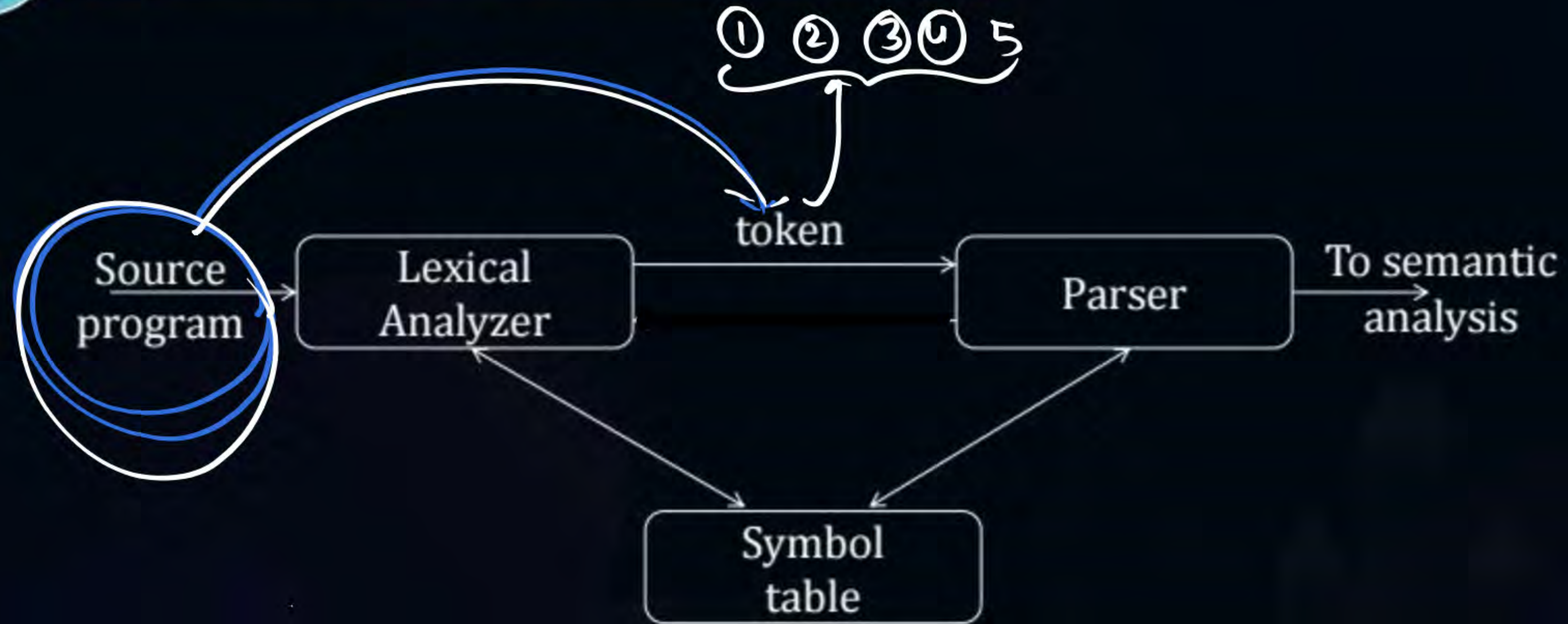
➤ It constructs symbol table







## Topic : The Role of Lexical Analyzer





① Regular Expressions

② Finite Automata



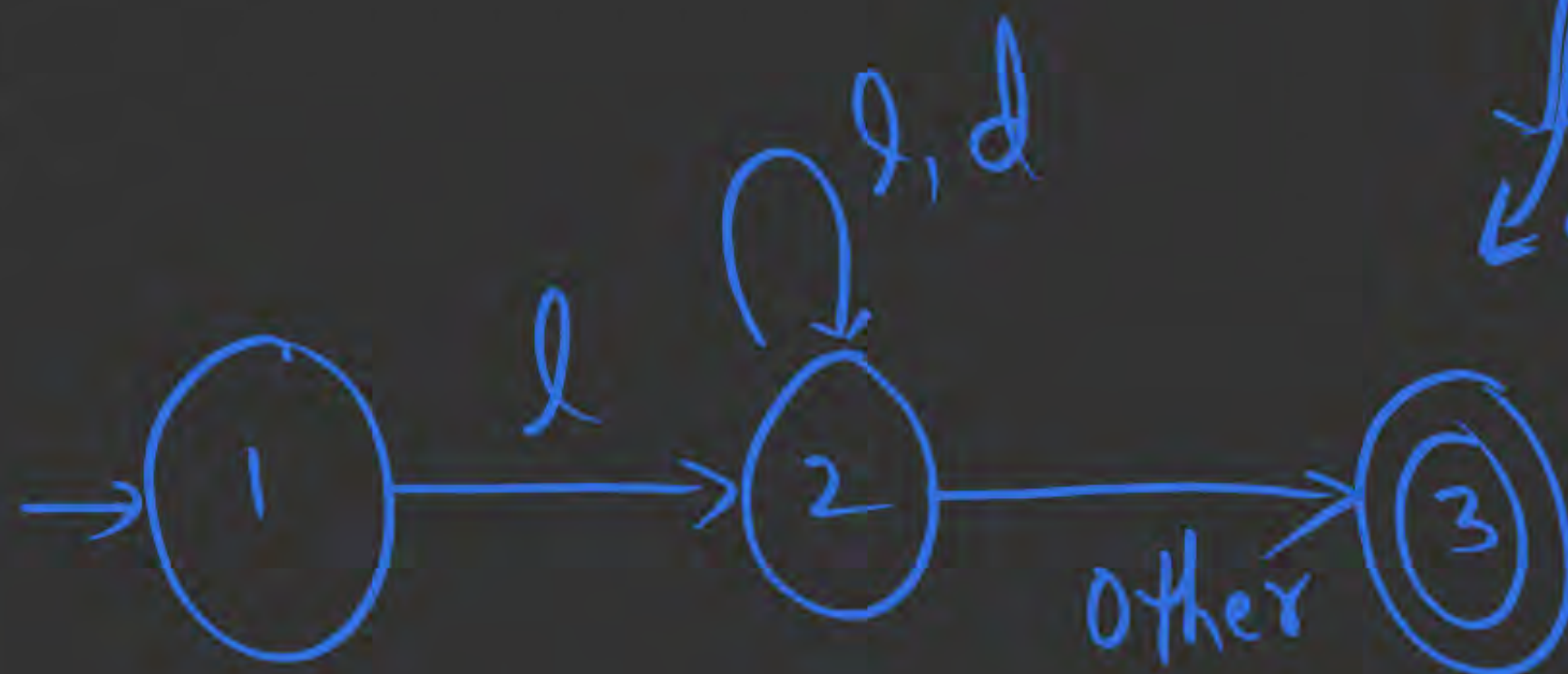
# Identifier

ID

else
if
float
int

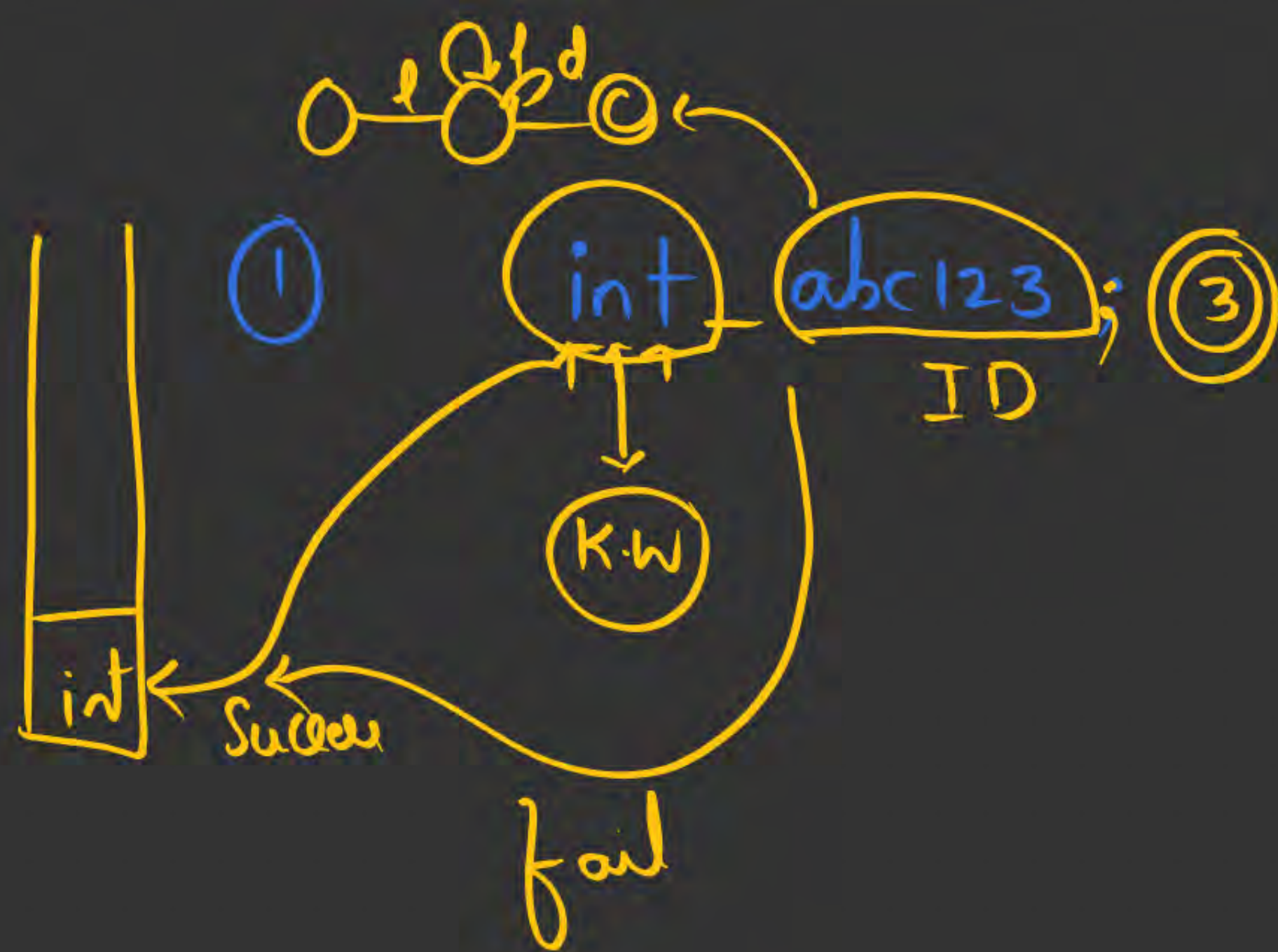
Success K.W

fail ID.



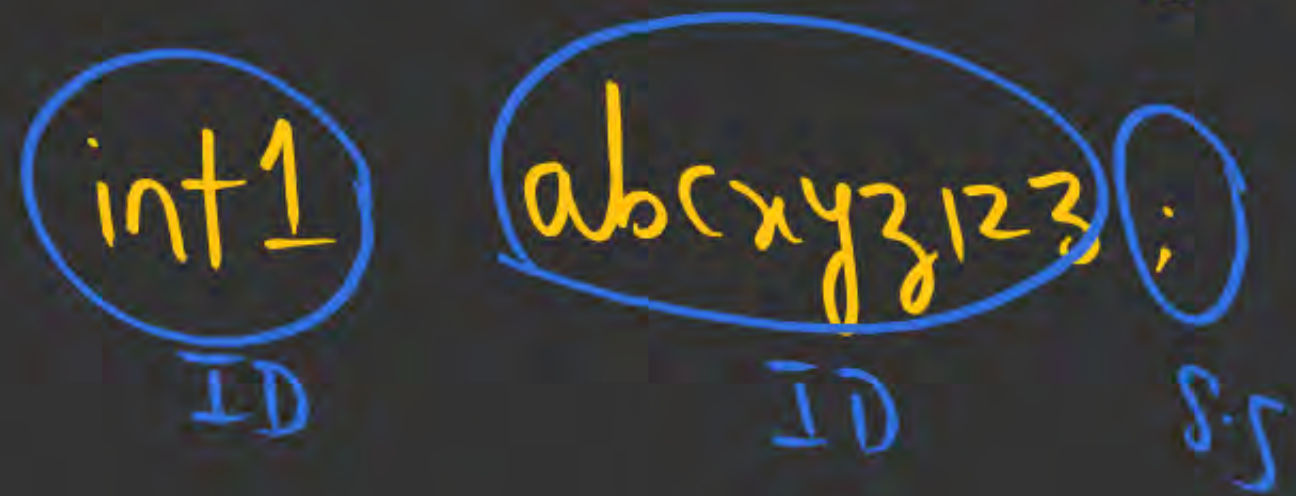
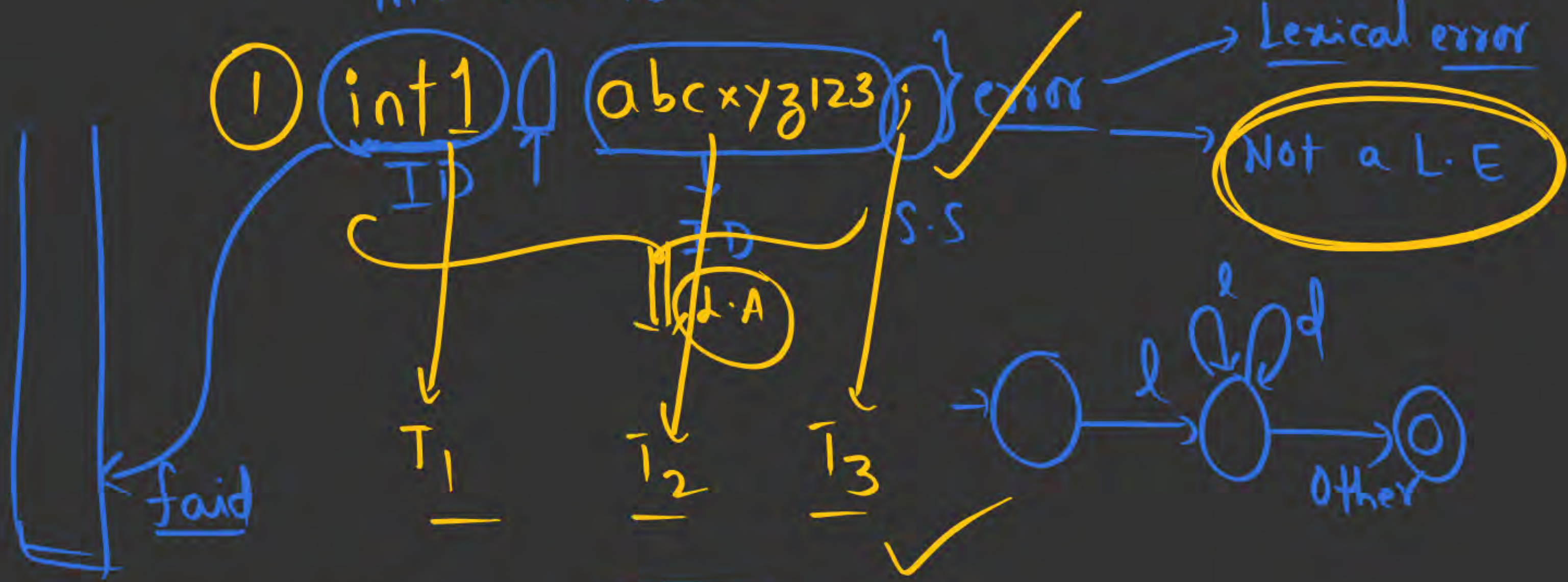
int  
float



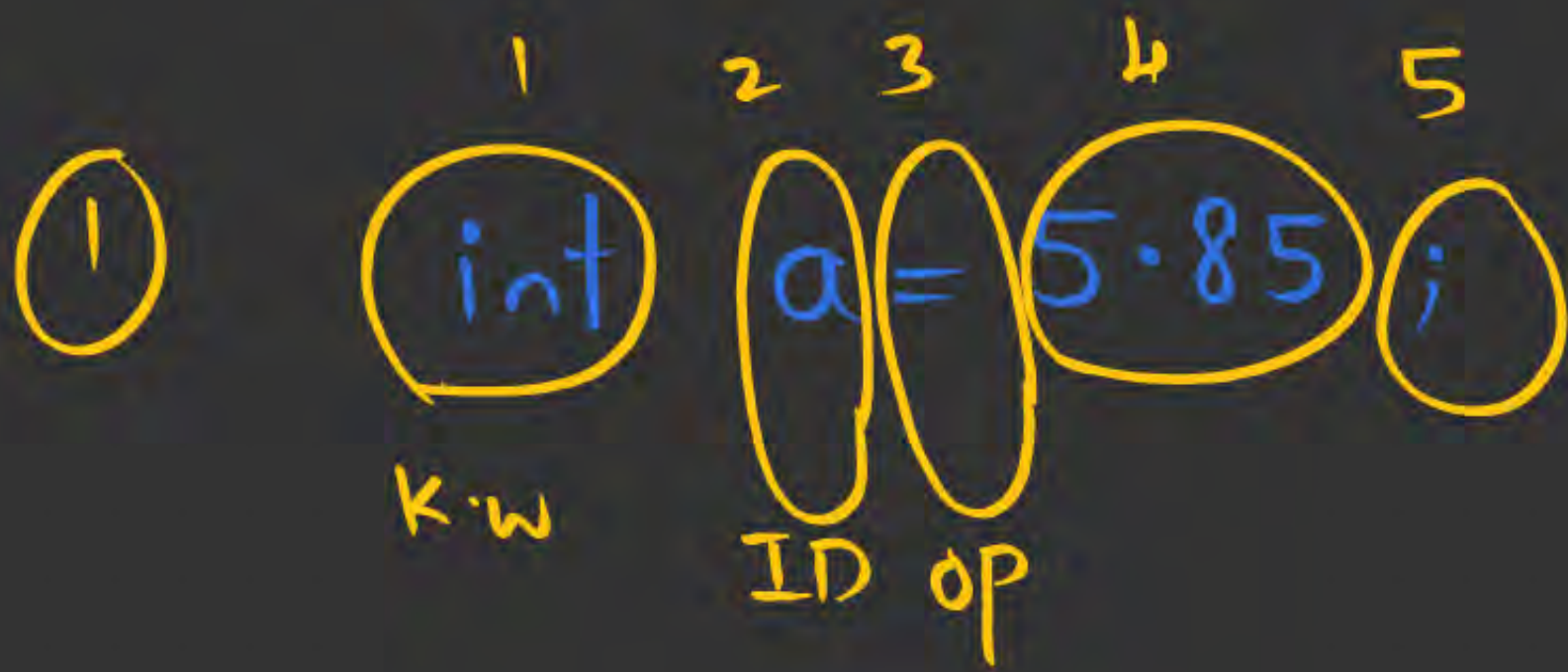




int abcxyz123;











## Topic : The Role of Lexical Analyzer

### **Why to separate Lexical analysis and parsing**

- Simplicity of design
- Improving compiler efficiency
- Enhancing compiler portability



## Topic : The Role of Lexical Analyzer

- **Token:** Token is a group of characters from source code that can be treated as a single logical entity.
- Typical tokens are,
  - 1. Identifiers
  - 2. keywords
  - 3. operators
  - 4. special symbols
  - 5. constants



## Tokens

①

keywords

(int, float, for)...

②

Identifiers

(xyz, abc)

③

Constants

(10, 20, 30)

④

Operators

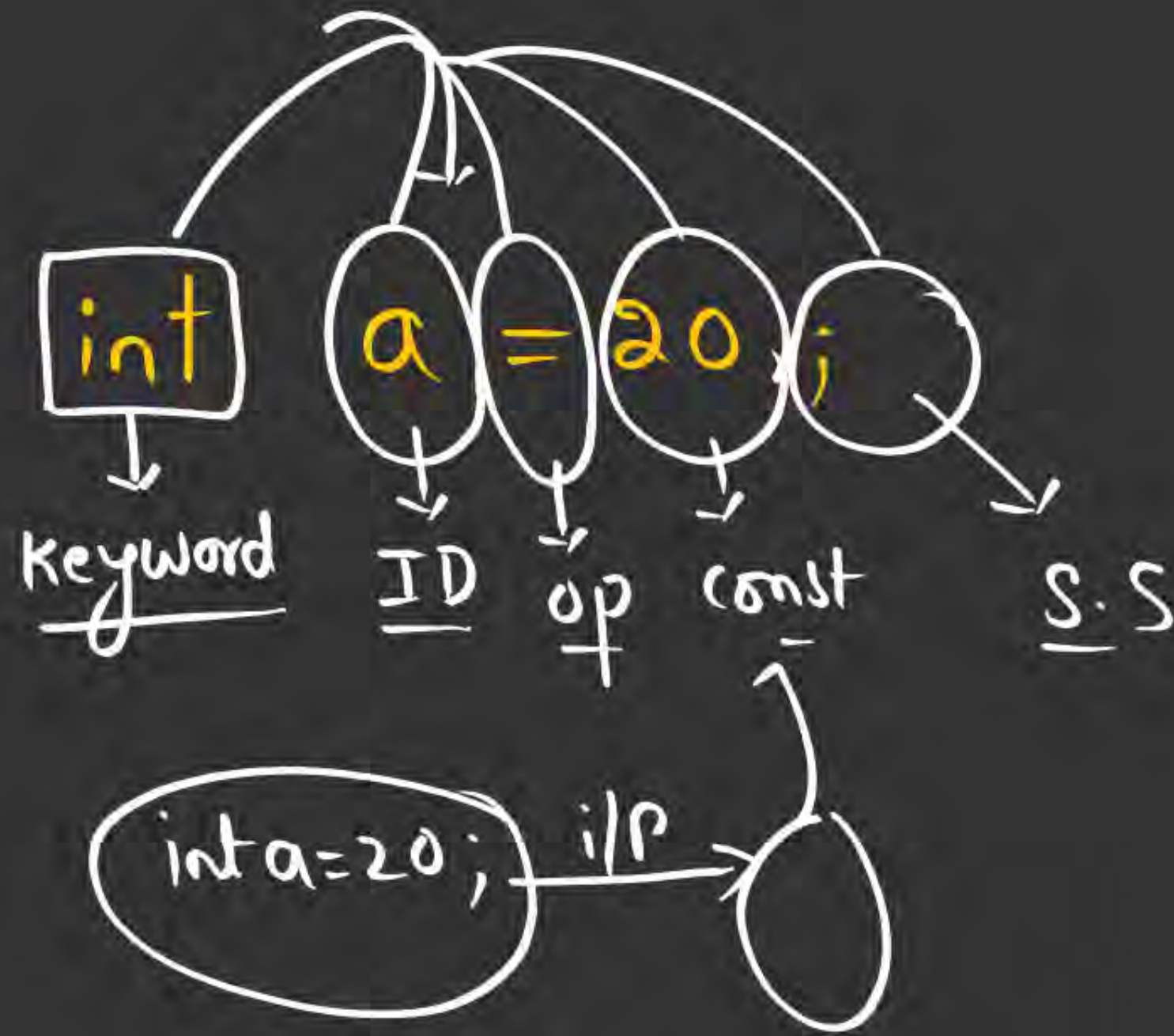
(+, \*)

⑤

Special symbols

{(, ), \*, }, ; }

①



5 tokens



# keywords (ReserveWords)

int

float

char

for

if

goto

do

else

goto

while

short

void.

32



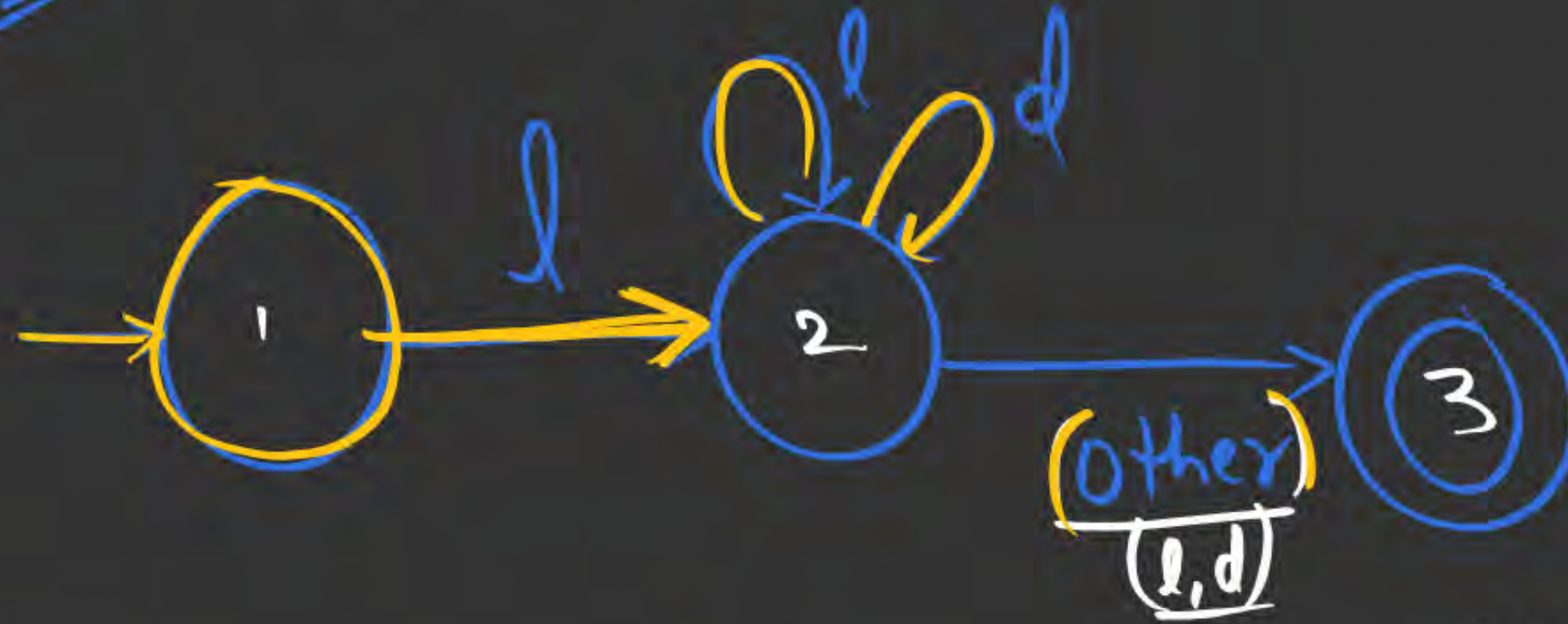
# Identifier

- ① Identifier can be formed using letter and digit
- ② Identifier can start with letter only (can't start with digit)
- ③ keyword name is not allowed as Identifier
- ④ Identifier can't include any special characters (-)

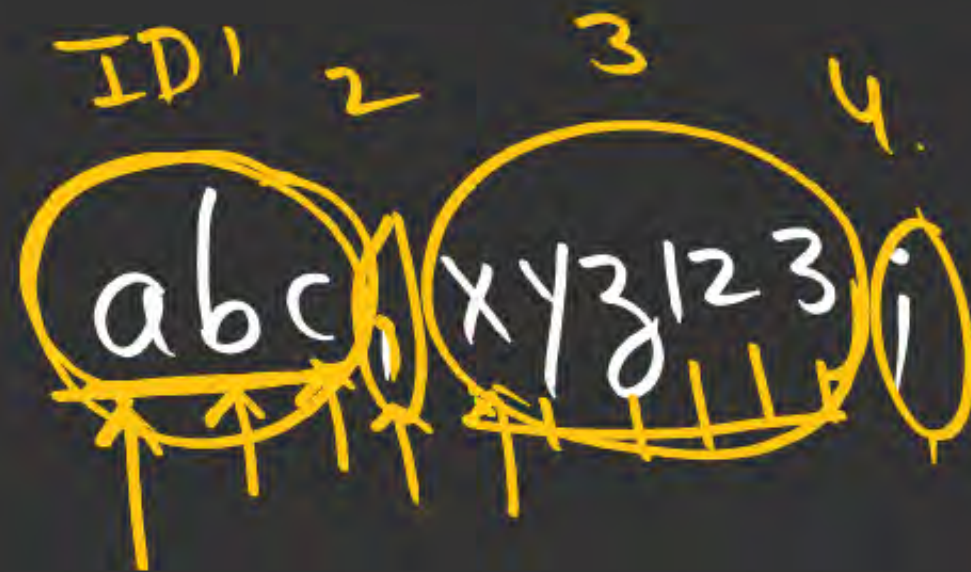


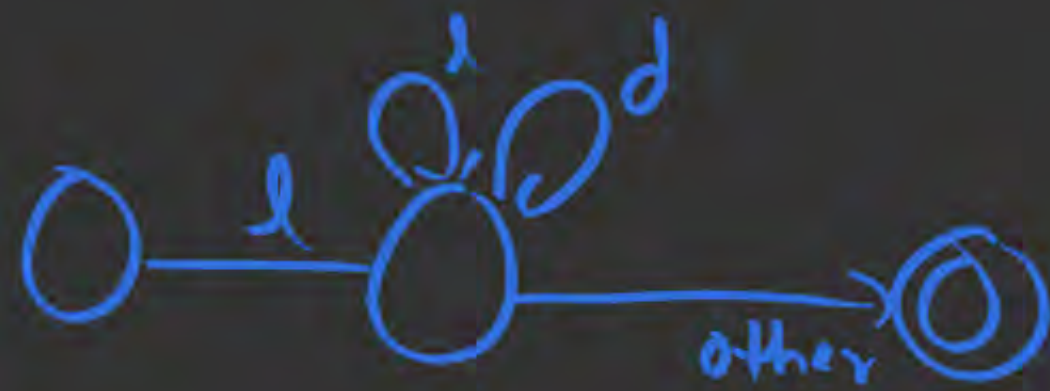
ID F.A

$l(l+d)^*$  other



return ID





intfloatchar; → 2

↑      ↑      2



# constants

d

10000

→ Integer Constant (100, 200)

→ Real Constant (10.5, 20.6)

→ Character Constant 'a' 'b'

→ String Constant "Compiler"

① <sup>1</sup>int<sup>2</sup> <sup>3</sup>a<sup>4</sup>=<sup>5</sup>"Compiler"; } 5  
K.W ID op

② int a="Compiler"; } NO Lexical error

③ <sup>1</sup>int<sup>2</sup> <sup>3</sup>a<sup>4</sup>=<sup>5</sup>20<sup>6</sup>+<sup>7</sup>"gate"; }



①

```
printf("%d %d", x, y);
```

Diagram illustrating the components of the `printf` function call:

- 1: `printf` (function name)
- 2: `(` (opening parenthesis)
- 3: `"%d %d"` (format string)
- 4: `,` (comma separator)
- 5: `x` (first argument)
- 6: `,` (comma separator)
- 7: `y` (second argument)
- 8: `)` (closing parenthesis)
- 9: `;` (semicolon)

main()

{

① int a=10; ✓

② int b; ✓

③ char c="abc"; ✓

④ 

b	=	a	+	c	;
1	2	3	4	5	6

 ✓

Ⓐ Lexical error

Ⓑ No Lexical error



# Operator

Increment/ Decrement	++ , --
Arithmetic Operator	+ - * / %
Relational Operator	< <= >= == !=
Assignment Operator	= , += , -= , *= , /= , %=
Bitwise operator	<< , >>



# Operator (token)

Increment/ Decrement	$++$ , $--$
Arithmetic Operator	$+$ $-$ $*$ $/$ $\%$
Relational Operator	$<$ $<=$ $>=$ $==$ $!=$
Assignment Operator	$=$ , $+=$ , $-=$ , $*=$ , $/=$ , $\%=$
Bitwise operator	$<<$ , $>>$



✓ ①  $\overset{1}{a} \overset{2}{=} \overset{3}{b} \overset{4}{;} \rightarrow 4$

---

②  $\overset{1}{a} \overset{1}{=} \overset{1}{=} \overset{1}{b} \overset{1}{,} \overset{1}{\}$  } 4

---

③  $\overset{1}{a} \overset{2}{+} \overset{3}{b}$

---

④  $\overset{1}{a} \overset{1}{+} \overset{1}{+} \}$  2

---

⑤  $\overset{1}{a} \overset{2}{+} \overset{2}{+} \overset{3}{+} \overset{3}{+} \overset{4}{+} \overset{4}{+} \overset{5}{+} \}$  5

---

⑥  $a \mid < \mid b \mid \}$  3 tokens.

---

⑦  $\overset{1}{a} \overset{2}{<} \overset{3}{=} \overset{3}{b}$

---

⑧  $\overset{1}{a} \overset{2}{+} \overset{3}{=} \overset{3}{b}$

---

⑨  $\overset{1}{a} \overset{2}{-} \overset{3}{=} \overset{3}{b}$

---

⑩  $\overset{1}{a} \overset{2}{>} \overset{3}{=} \overset{3}{b} \}$

1	2	3	4	5	6	7	8	9	10	11	12	13
for	(	i	=	0	;	i	<=	100	;	i	++	)

for      s.s

How many tokens?

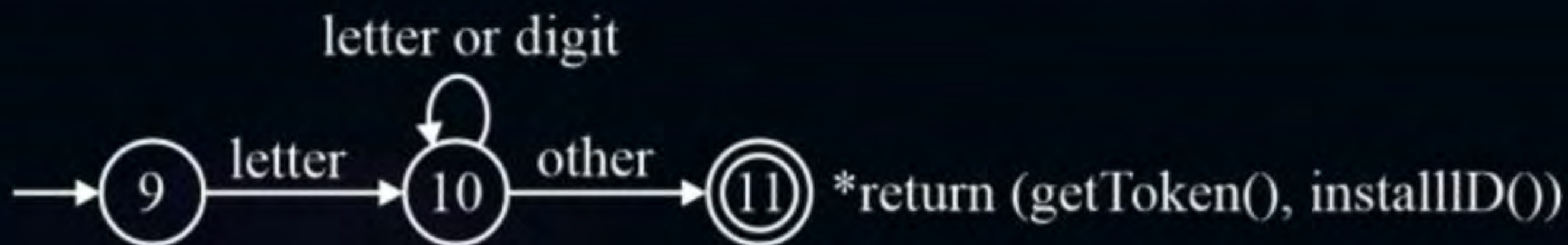
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## Topic : Transition diagrams

- Transition diagram for reserved words and identifiers





## Topic : Lexical Analyzer

① <sup>1</sup>int <sup>2</sup>abc <sup>3.5</sup>; → 3  
Keyword ID

② <sup>1</sup>int <sup>2</sup>a <sup>3</sup>, <sup>4</sup>b <sup>5.</sup>; → 5

③ <sup>1</sup>abc <sup>2</sup>= <sup>3</sup>b <sup>4</sup>+ <sup>5</sup>c <sup>6</sup>; → 6  
ID





## Topic : Lexical Analyzer

④ `if (a < b)`

⑤ `int a=100, b=200;`

⑥ `int abc1234float;`



## Topic : Lexical Analyzer

⑦ `char b = 'a';`

⑧ `char a = "abc";`

⑨ `printf("Compiler Design");`



$$\textcircled{10} \quad x = y / ++ ,$$

$$\textcircled{11} \quad x = y + + + + + z ;$$

$$\textcircled{12} \quad x = * y * z ;$$

$$\textcircled{13} \quad x = y + 125 ;$$

# Comment

→ ①

1	2	3	4	5	6
x	=	y	+	z	;

/\* comment \*/ } 6.

---

②

1	2	3	4	5	6
x	=	/* comment */	y	+	z;

↓ d/p

③

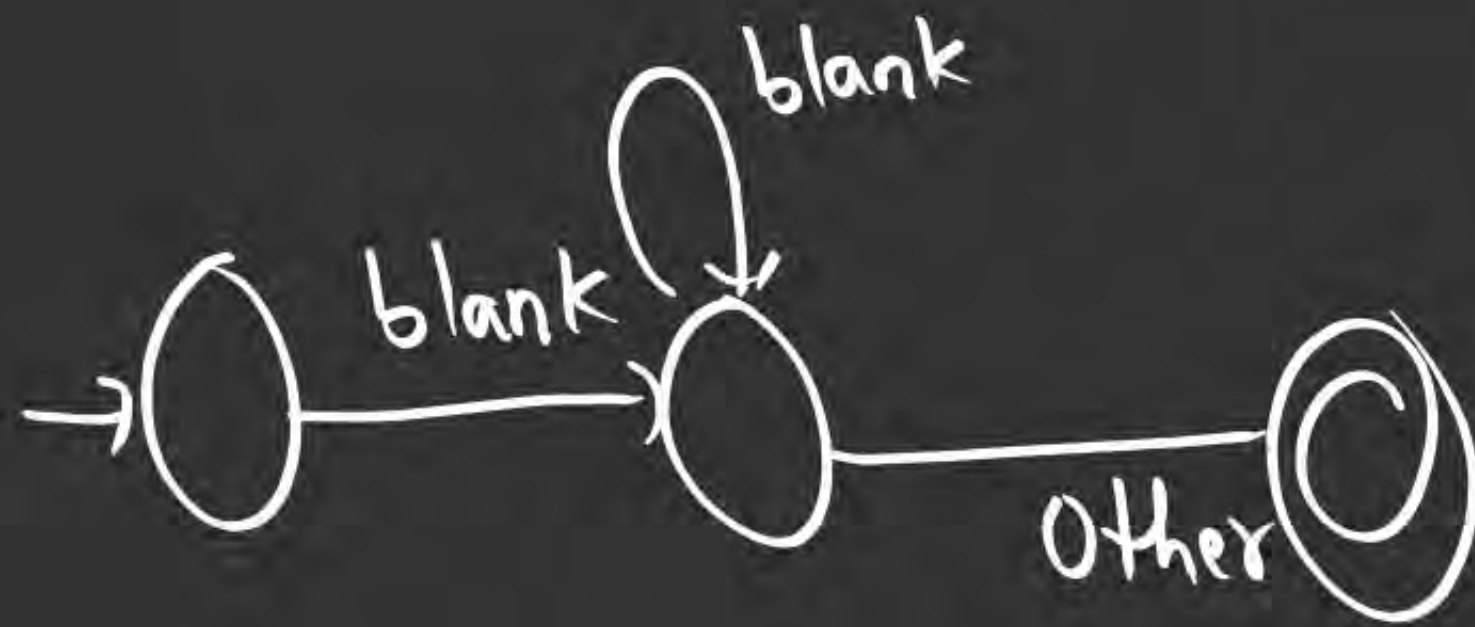
x = y + z;
------------

d/p





return (no token)



return (no token)



## Topic : The Role of Lexical Analyzer

- **Lexeme:** A lexeme is a sequence of characters in the source program that is matched by the pattern for a token.





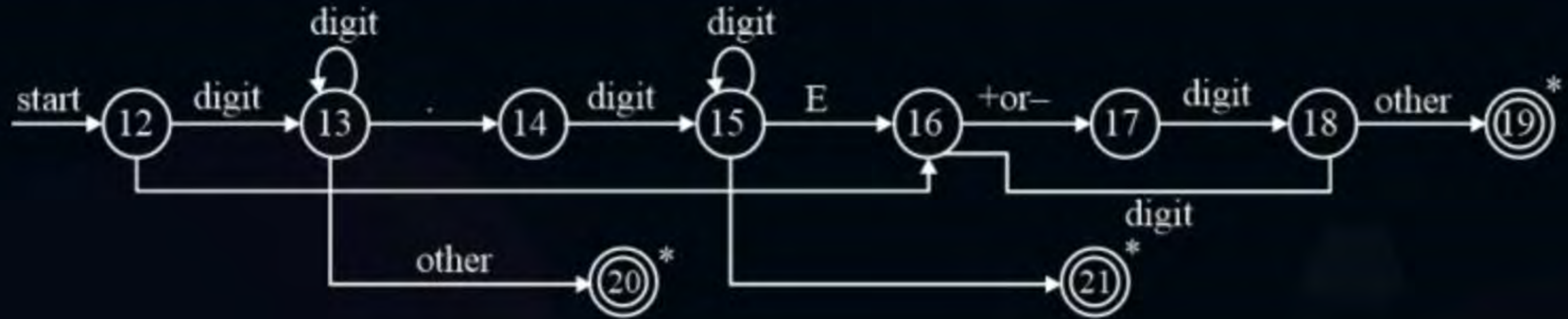
## Topic : Tokens, Patterns and Lexemes

- A token is a pair a token name and an optional token value
- A pattern is a description of the form that the lexemes of a token may take
- A lexeme is a sequence of characters in the source program that matches the pattern for a token



## Topic : Transition diagrams (cont.)

- Transition diagram for unsigned numbers







## Topic : Transition diagrams (cont.)

- Transition diagram for whitespace



(ab c defgg)



#Q. How many tokens are generated by the lexical analyzer, if the following program has no lexical error?  
what type of error

```
main()  
{  
    int x, y;  
    fl/*gate*/oat z;  
    x = /*exam*/10;  
    y = 20;  
}
```

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```
main()  
{  
    int x, y;  
    fl oat z;  
    x = 10;  
    y = 20;  
}
```

(a) Lexical error

(b) No lexical error



# Lexical error

① int 1 a, b; → ✓

② int 1abc; → Lexical error  
k.w

③ char a = "Compiler"

④ char b = 'a' → Lexical error

x ① k.w

x ② I.D

x ③ Const

x ④ Op

x ⑤ S.S

a=100;

#Q. Consider the C program

```

main()
{
    int x = 10;
    x = x + y + z; /* compiler design */
}

```

Handwritten annotations: The number 18 is circled in green above the closing brace. The number 5 is written next to the opening brace. The number 9 is written next to the semicolon in the first line. The number 13 is written next to the semicolon in the second line. The entire second line is circled in green.

How many tokens are identified by lexical analyzer?



#Q. Consider the following program segment:

```
main ()
{
    int a, b;
        a= 5 + 8 +;
        printf("%d", a);
        /* b = 5 ; */
}
```

The number of token present in the above program segment

#Q. Consider the following expression of C program

`abcd + (2 - 5 + × 6/2 - +;`

How many tokens are generated by the above expression during lexical analysis?



#Q. Which of the following is not a token in C language?

- A** Semicolon
- B** Identifier
- C** Keyword
- D** White space

#Q. Consider the following program segment:

```
main ()  
{  
    int a, b, c;  
        a = 50;  
        b = &a;  
    printf("%d", b);  
}
```

The number of tokens in the above C code are \_\_\_\_\_.



#Q. Find the type of error produced by the following C code.

```
main()
{
    in/*comment t x*/;
    floa/*comment*/t gate
}
```

**A**

Lexical error

**B**

Syntax error

**C**

Both (a) and (b)

**D**

None of these

#Q. If (z == a) function1 (10);  
The number of tokens in the above statement are

- A** 9
- B** 11
- C** 10
- D** 12



#Q. Which of the following C program statement having Lexical error.

1. iit a, b;
2. int 1 a,b;
3. int x1,b;
4. int a="hello";
5. int b = 5.89;
6. char a=a;
7. char b="hello";
8. int d="hello"
9. a=b;+c;



**THANK - YOU**