

## Discrete Mathematics

## Mathematical Logic

DPP: 2

**Q1** Let  $p(x)$  and  $q(x)$  denote the following open statements.

$$p(x): x^2 > 0$$

$$q(x): x \text{ is odd}$$

for the universe of all integers, determine the truth or falsity of each of the statement.

$$S_1: \forall x [p(x) \rightarrow q(x)]$$

$$S_2: \exists x [p(x) \rightarrow q(x)]$$

Which of the following is true?

- (A)  $S_1$  only
- (B)  $S_2$  only
- (C) Both  $S_1$  and  $S_2$
- (D) Neither  $S_1$  nor  $S_2$

**Q2** Consider following two First Order Logic Statements:

$$S_1: [\forall x (\sim P(x) \vee Q(x))] \rightarrow [\forall x P(x)] \rightarrow [\forall x Q(x)]$$

$$S_2: [\exists x P(x)] \rightarrow [\exists x Q(x)] \rightarrow [\exists x (P(x) \rightarrow Q(x))]$$

Which of the following is valid?

- (A)  $S_1$  only
- (B)  $S_2$  only
- (C) Both  $S_1$  and  $S_2$
- (D) None of these

**Q3**  $P(y) = \sqrt{y}$  is real in the domain of  $Z^+$ , then which of the following is / are correct?

- (A)  $\forall P(y)$
- (B)  $\exists y P(y)$
- (C)  $\forall y \sim P(y)$
- (D)  $\exists y \sim P(y)$

**Q4** Which of the following is not valid logical expression?

- (A)  $\forall x [P(x) \rightarrow Q(x)] \rightarrow [\forall x P(x)] \rightarrow [\forall x Q(x)]$
- (B)  $\forall x [P(x) \vee Q(x)] \rightarrow [\forall x P(x)] \vee [\forall x Q(x)]$
- (C)  $\exists x [P(x) \wedge Q(x)] \rightarrow [\exists x P(x)] \wedge [\exists x Q(x)]$
- (D)  $\forall x [P(x) \leftrightarrow Q(x)] \rightarrow [\forall x P(x)] \leftrightarrow [\forall x Q(x)]$

**Q5** Consider following logical expressions:

$$\text{I: } \forall y [P(y) \rightarrow Q] \leftrightarrow [\forall y P(y)] \rightarrow Q$$

$$\text{II: } \exists y [P(y) \rightarrow Q] \rightarrow [\exists y P(y)] \rightarrow Q$$

which of the following logical expression is valid?

- (A) I only
- (B) II only
- (C) Both I and II
- (D) None of these

**Q6** Consider

Actor ( $x$ ) =  $x$  is an actor

Smart ( $x$ ) =  $x$  is smart

and the well-formed formula:

$$\exists x (\text{Actor}(x) \wedge \text{Smart}(x))$$

Choose the correct representation of above in English sentence.

- (A) Some Actor is smart
- (B) Some Actor is not smart
- (C) All actors are smart
- (D) All smart are actors

**Q7** Consider the following statement

"There is exactly one apple".

Let  $G(x)$ :  $x$  is an apple.

Now consider the predicate logic statements:

$$\text{I. } \exists x \text{ apple}(x) \wedge \forall y (\text{apple}(y)) \Rightarrow x = y$$

$$\text{II. } \exists x \text{ apple}(x)$$

The correct representation in predicate logic is ?

- (A) Only I
- (B) Only II
- (C) Both I and II
- (D) Neither I and II

**Q8** Consider the following statements:

$$\text{I. } \exists x \{p(x) \rightarrow \{\forall x P(x) \rightarrow \forall x Q(x)\}\}$$

$$\text{II. } \exists x \forall y P(x, y) \rightarrow \forall y \exists x P(x, y)$$

The number of valid statements is/are \_\_\_\_\_.


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**Q9** Choose among the following that are not equivalent to the given first order logic statement:

$$(\exists x) (\forall y) [p(x, y) \wedge q(x, y) \wedge \neg r(x, y)]$$

$$(A) (\forall x) (\exists y) [p(x, y) \wedge q(x, y) \rightarrow r(x, y)]$$

$$(B) (\exists x) (\forall y) [p(x, y) \vee q(x, y) \wedge \neg r(x, y)]$$

$$(C) \neg (\forall x) (\exists y) [p(x, y) \vee q(x, y) \rightarrow r(x, y)]$$

$$(D) \neg (\forall x) (\exists y) [p(x, y) \wedge q(x, y) \rightarrow r(x, y)]$$

**Q10** Choose the correct representation for the below statement:

“Every player is liked by some coach”

$$(A) \forall (x) [\text{player}(x) \rightarrow \exists y [\text{coach}(y) \wedge \text{likes}(y, x)]]$$

$$(B) \forall (x) [\text{player}(x) \rightarrow \exists y [\text{coach}(y) \rightarrow \text{likes}(y, x)]]$$

$$(C) \exists (x) [\text{player}(x) \rightarrow \forall y [\text{coach}(y) \rightarrow \text{likes}(y, x)]]$$

$$(D) \exists (x) [\text{player}(x) \rightarrow \forall y [\text{coach}(y) \wedge \text{likes}(y, x)]]$$



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## Answer Key

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Q1 (B)  
Q2 (C)  
Q3 (A, B)  
Q4 (B)  
Q5 (D)

Q6 (A)  
Q7 (A)  
Q8 2  
Q9 (A, B, C)  
Q10 (A)



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## Hints & Solutions

### Q1 Text Solution:

**Statement  $S_1$ :**  $\forall x[p(x) \rightarrow q(x)]$

As we know the  $\forall x$  connected through ' $\wedge$ ' operator.

So, check the statement for  $x = 3$

$$\therefore [p(3) \rightarrow q(3)] = [(3^2 > 0) \rightarrow (3 \text{ is odd})]$$

$$\equiv [\text{True} \rightarrow \text{True}] \equiv \text{True}$$

Now, check the statement for  $x = 2$

$$\therefore [p(2) \rightarrow q(2)] = [(2^2 > 0) \rightarrow (2 \text{ is odd})]$$

$$\equiv [\text{True} \rightarrow \text{False}] \equiv \text{False}$$

Here  $S_1$  is false.

**Statement  $S_2$ : True**

If  $\exists x$  is true for one value then the overall the truth value of the statement will be true.

So, Check the statement for  $x = 3$

$$\therefore [p(3) \rightarrow q(3)] = [(3^2 > 0) \rightarrow (3 \text{ is odd})]$$

$$\equiv [\text{True} \rightarrow \text{True}] \equiv \text{True}$$

Hence,  $S_2$  is True.

### Q2 Text Solution:

$$S_1: [\forall x (P(x) \rightarrow Q(x)) \rightarrow [\forall x P(x)] \rightarrow [\forall x Q(x)]$$

(Property of Predicate Logic)

$$[\forall x (\sim P(x) \vee Q(x)) \rightarrow [\forall x P(x)] \rightarrow [\forall x Q(x)]$$

$$S_2: [\exists x P(x)] \rightarrow [\exists x Q(x)]$$

$$\rightarrow [\exists x (P(x) \rightarrow Q(x))]$$

**Proof:**

$$(P_1 \vee P_2) \rightarrow (Q_1 \vee Q_2)$$

$$\rightarrow [(P_1 \rightarrow Q_1) \vee (P_2 \rightarrow Q_2)]$$

$$(P_1' P_2' + Q_1 \vee Q_2)$$

$$\rightarrow [(P_1' + Q_1) + (P_2' + Q_2)]$$

$$(P_1 + P_2) \cdot (Q_1' Q_2') + P_1' + Q_1$$

$$+ P_2' + Q_2$$

$$P_1 Q_1' Q_2' + P_2 Q_1' Q_2' + P_1' + Q_1$$

$$+ P_2' + Q_2$$

$$\boxed{A'B + A = A + B}$$

$$P_1 + P_2 + P_1' + Q_1 + P_2' + Q_2$$

$$P_1 + P_1' = 1 \text{ and } 1 + \text{anything} = 1$$

$$1 + P_1 + P_2 + Q_1 + Q_2 + P_2'$$

$$1 \text{ True}$$

Hence both are valid.

### Q3 Text Solution:

$$P(y) = \sqrt{y} \text{ is real}$$

domain = positive integers ( $z^+$ )

$$(a) \forall y P(y) \text{ True}$$

For every values of  $y$ ,  $\sqrt{y}$  is real because domain is positive integer

$$(b) \exists y P(y) \text{ True}$$

For some values of  $y$ ,  $\sqrt{y}$  is real

$$(c) \forall y \sim P(y) \text{ False}$$

$$(d) \exists y \sim P(y) \text{ False}$$

### Q4 Text Solution:

$$(P_1 \vee Q_1) \wedge (P_2 \vee Q_2) \rightarrow (P_1 \wedge P_2) \vee (Q_1 \wedge Q_2)$$

$$(P_1 + Q_1) \wedge (P_2 + Q_2) \rightarrow P_1 P_2 + Q_1 Q_2$$

$$P_1' + Q_1' + P_2' + Q_2' + P_1 P_2 + Q_1 Q_2$$

$$P_1 P_2 + Q_1 Q_2 + P_1' + Q_1' + P_2' + Q_2' \text{ (Invalid)}$$

Remaining all are valid.

### Q5 Text Solution:



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- I:**  $\forall y [P(y) \rightarrow Q] \leftrightarrow [\forall y P(y)] \rightarrow Q$   
 $(P_1 \rightarrow Q) \wedge (P_2 \rightarrow Q) \equiv (P_1 \wedge P_2) \rightarrow Q$   
 $(P_1' + Q) \wedge (P_2' + Q) \equiv P_1' + P_2' + Q$   
 $P_1' P_2' + P_1' Q + P_2' Q + Q \equiv P_1' + P_2' + Q$   
 $P_1' P_2' + Q \neq P_1' + P_2' + Q$  (invalid)
- II:**  $\exists y [P(y) \rightarrow Q] \rightarrow [\exists y P(y)] \rightarrow Q$   
 $(P_1 \rightarrow Q) \vee (P_2 \rightarrow Q) \rightarrow (P_1 \vee P_2) \rightarrow Q$   
 $P_1' + P_2' + Q \rightarrow P_1' P_2' + Q$   
 $P_1 + P_2 Q' + P_1' P_2' + Q$   
 $P_1 P_2 + P_1' P_2' + Q$  (Invalid)  
Hence, option (d) is correct

**Q6 Text Solution:**

- $\exists x$  represents some/any/at least one:
- Actor (x)  $\wedge$  Smart (x) means x is an actor and smart.
- Therefore  $\exists x$  (Actor (x)  $\wedge$  smart (x)) represents some actor is smart.

**Q7 Text Solution:**

I is the correct representation as it reads “there exist an apple x and there exist an apple y and if apple y exists then it is equal to x” that means there is only one apple (exactly one).  
II is absolutely incorrect as it says “some apple or at least one apple” instead of exactly one apple.

**Q8 Text Solution:**

- I:**  $\exists x \{P(x) \rightarrow Q(x)\} = \exists x \{\neg (P(x)) \vee Q(x)\}$   
 $= \{\exists x \neg (P(x)) \vee \exists x Q(x)\}$   
 $= \forall x P(x) \rightarrow \exists x Q(x) \therefore \text{True.}$
- II.**  $\exists x \forall y P(n, y)$   
 $\forall_y P(a, y)$  for same a  
 $P(a, b)$  is true for  $\forall_b = \exists x P(x, b)$   
 $\forall_y \exists x P(x, y)$  is true.

**Q9 Text Solution:**

Two points/rules to solve the question:

- I.**  $\exists x f(x) \equiv \neg \forall (x) \neg f(x)$   
**II.**  $\forall x f(x) \equiv \neg \exists (x) \neg f(x)$

The given statement:

$$(\exists x) (\forall y) [p(x, y) \wedge q(x, y) \wedge \neg r(x, y)]$$

**Can be written as: -**

$$\neg (\forall x) \exists (y) [p(x, y) \wedge q(x, y) \rightarrow r(x, y)]$$

**NOTE:**  $[\neg (x \wedge y) \vee 3 \equiv x \wedge y \rightarrow 3]$

**Q10 Text Solution:**

We write “Every player” as  $\forall x [\text{player}(x) \rightarrow]$   
“there is some coach who likes x” as  $\exists y [\text{coach}(y) p(x)]$   
Where P is the property.  
Therefore, we can write the first order logic for the given statement as  
 $\forall (x) [\text{player}(x) \rightarrow \exists y [\text{coach}(y) \wedge \text{likes}(y, x)]]$



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