

# THEORY OF COMPUTATION ASSIGNMENT

Q2.2 Given a PDA, convert it to CFG

## **BASIC STRUCTURE**

1. **PDA\_CFG.h** is the header file that contains all the variables and the function prototypes
2. **PDA\_CFG.cpp** is where the conversion of PDA to CFG takes place, all the permutations are printed
3. **pda\_cfg\_main.cpp** is the driver function where the file input\_pda.txt is read and the corresponding PDA is loaded onto the program which the PDA\_CFG.cpp to convert PDA to CFG

## **VARIABLES**

TYPE	NAME	DESCRIPTION
vector<string>	states	States in the PDA
vector<string>	input_symbols	Input symbol
vector<string>	init_state	Stores the initial state before transition
vector<string>	final_state	Stores the final state after transition
vector<string>	stack_top	Top of the stack
vector<string>	push_stack	Symbol to be pushed on the stack
ifstream	fin	Input Stream object to operate on file

## **FUNCTIONS**

1. **FileRead()** - Reads a file input.txt which contains the grammar
  - Checks error in opening file.
  - Pushes the states in PDA present on Line 1 in states vector.
  - Pushes the initial state, input symbol and top of stack of LHS to corresponding vectors, trimming the first three characters  $\delta$ ( and last )
  - Pushes the final state and symbol to be pushed onto the stack of RHS to corresponding vectors trimming the parenthesis.
2. **DisplayPDA()** - Display the PDA read from the file

3. **Permute()** - The construction of an equivalent grammar uses variables each of which represents an "event" consisting of :
1. The net popping of some symbol X from the stack, and
  2. A change in state from some p at the beginning to q when X has finally been replaced by  $\epsilon$  on the stack

Such a variable is represented by a composite symbol [pXq]

This function permutes all the variables in the Grammar (composite [pXq]) using standard recursive permutation technique.

4. **Print\_Permute()** -Print all the Permutations found in Premute() function

5. **PDA\_CFG()** - The PDA changes state as it pops stack symbols, so we note the state that it enters when it finally pops a level off its stack. The computation proceeds until each of the symbols on the stack is removed. All the variables generated from the states and stack symbols are printed

## INPUT

The screenshot shows a text editor window titled "input\_pda.txt" with the following transition rules for a Pushdown Automaton (PDA):

```
q
δ(q, 0, Z) = (q, AZ)
δ(q, 1, Z) = (q, BZ)
δ(q, 0, A) = (q, AA)
δ(q, 0, B) = (q, ε)
δ(q, 1, A) = (q, ε)
δ(q, 1, B) = (q, BB)
δ(q, ε, Z) = (q, ε)
```

Overlaid on the text is a state transition diagram with two states represented by circles. The bottom circle is the start state, indicated by an incoming arrow from the left. It has two outgoing transitions: one to the top circle labeled "0, B/ε" and "1, A/ε", and another to itself labeled "1, Z/ε". The top circle has two outgoing transitions: one back to the bottom circle labeled "0, Z/AZ" and "1, Z/BZ", and another to itself labeled "0, A/AA" and "1, B/BB".

## OUTPUT

```
MacBook:PDA_TO_CFG ayushsingh$ g++ pda_cfg_main.cpp
MacBook:PDA_TO_CFG ayushsingh$ ./a.out

-----
THE INPUT PDA IS :
STATES = q
 $\delta(q, 0, Z) = (q, AZ)$ 
 $\delta(q, 1, Z) = (q, BZ)$ 
 $\delta(q, 0, A) = (q, AA)$ 
 $\delta(q, 0, B) = (q, \epsilon)$ 
 $\delta(q, 1, A) = (q, \epsilon)$ 
 $\delta(q, 1, B) = (q, BB)$ 
 $\delta(q, \epsilon, Z) = (q, \epsilon)$ 
-----

Corresponding CFG :
START STATE = S -> [qXq]

 $\delta(q, 0, Z) = (q, AZ)$ 
[qZq] -> 0 [qAq] [qZq]

 $\delta(q, 1, Z) = (q, BZ)$ 
[qZq] -> 1 [qBq] [qZq]

 $\delta(q, 0, A) = (q, AA)$ 
[qAq] -> 0 [qAq] [qAq]

 $\delta(q, 0, B) = (q, \epsilon)$ 
[qBq] -> 0

 $\delta(q, 1, A) = (q, \epsilon)$ 
[qAq] -> 1

 $\delta(q, 1, B) = (q, BB)$ 
[qBq] -> 1 [qBq] [qBq]

 $\delta(q, \epsilon, Z) = (q, \epsilon)$ 
[qZq] ->  $\epsilon$ 

[qXq] ->  $\epsilon$ 
MacBook:PDA_TO_CFG ayushsingh$
```