MATHEMATICAL CIPHERS Delving into the Cryptographic Complexities of Number Theory and Abstract Algebra

Prove that any prime number $p \equiv 1 \mod 4$ can be expressed as the sum of two squares.

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Key points to note -

- Every square is either O(mod4) or 1(mod4)
 - \circ (2n)^2 = 4n^2 = 0(mod4)
 - \circ (2n+1)^2 = 4n^2 + 4n + 1 = 1(mod4)
- Then what can be the sum of 2 possible squares
 - \circ O + O = O(mod4) Not a prime
 - \circ 1 + 1 = 2(mod4) = 0(mod2) Not a prime
 - \circ O + 1 = 1(mod4) Can be a prime

Prove that any prime number $p \equiv 1 \mod 4$ can be expressed as the sum of two squares.

The exact proof is complicated but if you wanna know how its done exactly -

See this link

Solve the system of congruences:

$$x \equiv 2 \pmod{3}$$

$$x \equiv 3 \pmod{5}$$

$$x \equiv 2 \pmod{7}$$

$$x \equiv 2 \pmod{3}$$

 $x \equiv 3 \pmod{5}$
 $x \equiv 2 \pmod{7}$

This is a typical **Chinese Remainder Theorem** question

- So let, $x = 2 \mod 3$, ie. x = 3k+2
- Now, $3k+2 = 3 \mod 5 => So, 3k = 1 \mod 5$
- So, $k = 2 \mod 5 \Rightarrow k = 5m+2$
- Then, x = 3k+2 = 15m+8
- So, $15m+8 = 2 \mod 7 \Rightarrow 15m = 1 \mod 7 \Rightarrow m = 1 \mod 7 \Rightarrow m = 7n+1$
- So, x = 15(7n+1)+8 = 105n+23
- Hence, $x = 23 \mod 105$

Number Theory

Why its needed?

- **RSA cryptosystem -** relies on the difficulty of factoring large composite numbers into their prime factors.
- Diffie-Hellman Key Exchange: Uses modular exponentiation.
- **Digital Signature Algorithm (DSA)** depends on the difficulty of solving the discrete logarithm problem.
- Error detection Many cryptographic hash functions and error detection codes use number theoretic properties to ensure data integrity.

Cyclic Group

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Let G be a finite group of order m (written multiplicatively).
Let g be in G
Consider the set \{g^0, g^1, g^2, ....\}
Also, g^m = g^0 = 1, so that the set has at max m elements.
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If in this way, we get a set of m elements, we say 'g' is a generator for the set G.

Discrete Logarithm Problem

- Fix cyclic group G of order m, and generator g
- We know that $\{g^0, g^1, ..., g^m-1)\} = G$
- For every $h \in G$, there is a unique $x \in Zm^*$ s.t. $g^x = h$
- Define (base g)logh to be this x the discrete logarithm of h with respect to g (in the group G)

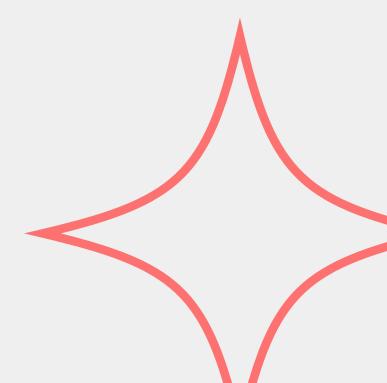
Discrete Logarithm Problem

Dlog problem in G

• given g, h compute log(g)h

And why do we use this?

• Because solving this is hard!!



Example 1

What is log(2)10 in the group Z(11)*?

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What do we need to find?

• 'x' such that $2^x = 10 \mod 11$

What can 'x' be?

- $2^5 = 32$
- and $32 = 10 \mod 11$

So x = 5;

That means -> log(2)10 = 5 in Z(11)* group.

More examples

Suppose $G = \mathbb{F}_{101}^{\times}$. Then $\log_3 37 = 24$, since $3^{24} \equiv 37 \mod 101$.

Things to notice -

- This is a group under multiplication
- So, we need to find 'x' such that $3^x = 37 \mod{101}$
- x = 24
- Hence, log(3)37 = 24

More examples

Suppose $G = \mathbb{F}_{101}^+$. Then $\log_3 37 = 46$, since $46 \cdot 3 \equiv 37 \mod 101$.

Things to notice -

- This is a group under addition
- So, we need to find 'x' such that **3*x = 37mod101**
- And **NOT** 'x' such that 3^x=37mod101
- x = 24
- Hence, log(3)37 = 24

Deffie-Hellman Problem

- Fix group G with generator g
- Define

$$DH_g(h_1, h_2) = DH(g^x, g^y) = g^{xy}$$

So, in essence, given g, h1, h2 we need to find x, y such that x = log(g)h1 and y = log(g)h2

And finally, $DH(g)(h1, h2) = g^(xy)$

Example

In \mathbb{Z}_{11}^* , what is $\mathsf{DH}_2(5,8)$?

Given g, h1, h2 lets first find x and y

So,
$$x = log(2)5 \Rightarrow x = 4$$

and
$$y = log(2)8 \Rightarrow y = 3$$

Thus, $g^{(xy)} \mod 11 = 2^{(12)} \mod 11 = 4 \mod 11$

So,
$$DH(2)(5, 8) = 4$$

Computational Diffie-Hellman (CDH) Problem

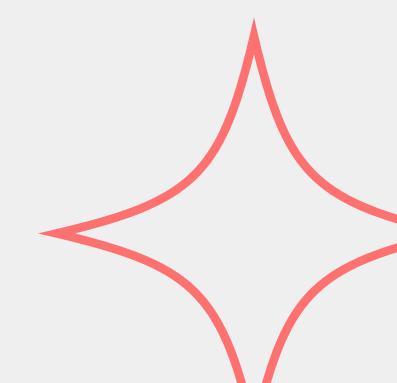
Given g, h1, h2, compute DH(g)(h1, h2)

Decisional Diffie-Hellman (CDH) Problem

Given g, h1, h2, distinguish the correct DH(g)(h1, h2) from a uniform element of G

What does 'uniform' mean?

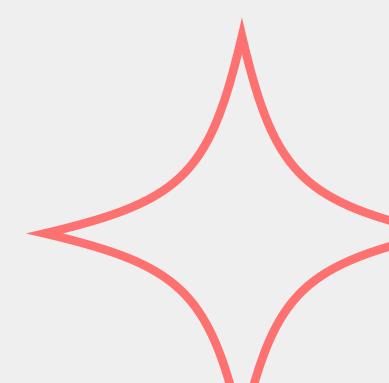
Any random pick from the set G, or more concretely, picking an element in such a way that every element is equally likely to be picked



Group selection

In cryptography, we basically choose **large prime-order groups** as it makes the computation of CDH and DDH problems touher.

Also, it might be worth recalling that in the prev meet, I discussed how prime order groups are always cyclic.



What's ahead?

If you are willing to go deeper in Discrete Logarithm Problem, I've linked notes from MIT that I found interesting

You can take a look at it -

https://math.mit.edu/classes/18.783/2022/LectureNotes9.pdf

Let's review what all we saw in Number Theory

- 1) Shift Cipher
- 2) Some Probability
- 3) Concept of Perfect Secrecy
- 4) Moving to Computational Secrecy
- 5) Modular Arithmetics
- 6) Factoring and **Groups**
- 7) Cyclic groups
- 8) Discrete Log Problems, CDH, DDH

