

Student Information

Full Name : Furkan KARACA

Id Number : 2521698

Answer 1

First, let's draw an adjacency and incidence matrices of the graph.

$$\begin{bmatrix} 0 & 1 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 & 1 \\ 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 1 \\ 1 & 1 & 0 & 1 & 0 \end{bmatrix}$$

Where indexes go a to e from left to right and from up to down. This is the adjacency matrix.

$$\begin{bmatrix} 1 & 1 & 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 1 & 0 & 0 & 1 \end{bmatrix}$$

Is the incidence matrix representation.

a) Sum of all degrees = Sum of all vertex adjacents. For vertex a, count of adjacents = 3

$$b \implies 3$$

$$c \implies 3$$

$$d \implies 2$$

$$e \implies 3$$

$3 + 3 + 3 + 2 + 3 = 14$ is the sum of all degrees of all nodes.

b) We can observe that the sum of non-zero entries is 14 from the adjacency matrix.

c) We can observe that the sum of non-zero entries is 21 from the incidence matrix.

- d) No. There are no such 4 vertices that form a complete graph.
- e) G is not bipartite. From the theorem, it is impossible to paint vertices in two different colors and no two adjacents aren't the same color.
- f) There are 7 undirected edges in the graph. Either of them could have two different ways. So, 2^7 total different directional graphs underlying in the graph G.
- g) $abcde$ is the longest path.
- h) Since the graph is connected. There is only one connected component.
- i) There is not an Euler circuit in G because G includes vertices with odd-numbered adjacents.
- j) Since G does not have an Euler circuit, and G does not have exactly two vertices with odd-numbered adjacents, there is not an Euler path in G.
- k) Yes. $abcdea$ is a Hamilton circuit.
- l) Yes. $abcde$ is a Hamilton path.

Answer 2

Both graphs G and H have 5 vertices and 5 edges. They also have 5 vertices with 2 edges. If we take the symmetric of vertex e from the a-d axis, vertex c from the d-b axis and so on we can obtain the same shape with graph H. So H and G are isomorphic. Also, their adjacency matrix representations are same.

Answer 3

First, we will keep a list that represents the distances with respect to s and another for a previous list to reach that vertex.

All the distances before visiting the vertex are infinity $D=Distance$ and $P=Prev$. So:

$D=(s, 0), (u, \infty), (w, \infty), (v, \infty), (y, \infty), (x, \infty), (z, \infty), (t, \infty)$

$P=(s, -), (u, -), (w, -), (v, -), (y, -), (x, -), (z, -), (t, -)$

We will traverse the graph and not visit the visited vertices. If we find a shorter path we will change the distance. We will traverse starting from vertex s.

The adjacents of vertex s are w, u and v, and their distance in the list (∞) is bigger than 3, 4, and 5, so we change the distance between them as follows:

$D=(s, 0), (u, 4), (w, 3), (v, 5), (y, \infty), (x, \infty), (z, \infty), (t, \infty)$

$P=(s, -), (u, s), (w, s), (v, s), (y, -), (x, -), (z, -), (t, -)$

We repeat the operation for the closest vertex adjacent to s. We add the distance value to the new vertex by adding it the distance of its previous. Closest is the w in this case.

$D=(s, 0), (u, 4), (w, 3), (v, 5), (y, \infty), (x, 11), (z, 15), (t, \infty)$

$P=(s, -), (u, s), (w, s), (v, s), (y, -), (x, w), (z, w), (t, -)$

Note that we didn't change the bigger or equal total distance (v, u and s).

Repeat for u.

$D=(s, 0), (u, 4), (w, 3), (v, 5), (y, 11), (x, 11), (z, 15), (t, \infty)$

$P=(s, -), (u, s), (w, s), (v, s), (y, u), (x, w), (z, w), (t, -)$

Repeat for v.

$D=(s, 0), (u, 4), (w, 3), (v, 5), (y, 11), (x, 7), (z, 15), (t, \infty)$

$P=(s, -), (u, s), (w, s), (v, s), (y, u), (x, v), (z, w), (t, -)$

We changed the distance and prev of x because there is a smaller distance if we go from v instead of w.

Repeat for x

$D=(s, 0), (u, 4), (w, 3), (v, 5), (y, 8), (x, 7), (z, 13), (t, \infty)$

$P=(s, -), (u, s), (w, s), (v, s), (y, x), (x, v), (z, x), (t, -)$ (y,z changed)

Repeat for y

$D=(s, 0), (u, 4), (w, 3), (v, 5), (y, 8), (x, 7), (z, 12), (t, 17)$

$P=(s, -), (u, s), (w, s), (v, s), (y, x), (x, v), (z, y), (t, y)$ (t added, z changed)

Repeat for z

$D=(s, 0), (u, 4), (w, 3), (v, 5), (y, 8), (x, 7), (z, 12), (t, 15)$

$P=(s, -), (u, s), (w, s), (v, s), (y, x), (x, v), (z, y), (t, z)$ (t changed)

Repeat for t

$D=(s, 0), (u, 4), (w, 3), (v, 5), (y, 8), (x, 7), (z, 12), (t, 15)$

$P=(s, -), (u, s), (w, s), (v, s), (y, x), (x, v), (z, y), (t, z)$ (nothing changed)

We traversed across every vertex and tried every possible way. In this graph, using the previous list, we can see the shortest path is s-v-x-y-z-t. And the distance is 15.

Answer 4

a)

Using Prim's algorithm: I pick a for the base point.

Choice 1 is {a,b} We keep choosing the smallest edges among what we have reached.

Choice 2 is $\{c,b\}$
 Choice 3 is $\{f,c\}$
 Choice 4 is $\{d,c\}$
 Choice 5 is $\{k,d\}$
 Choice 6 is $\{j,f\}$
 Choice 7 is $\{g,f\}$
 Choice 8 is $\{i,f\}$
 Choice 9 is $\{h,i\}$
 Choice 10 is $\{e,f\}$

This is the correct order by using Prim's algorithm.

b)

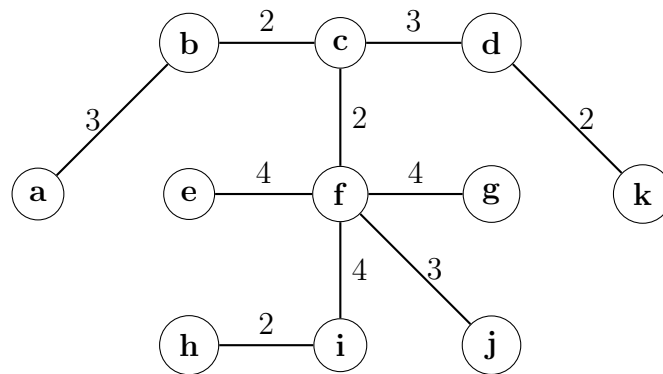


Figure 1: Tree.

c) The minimum spanning tree is not unique because the tree could be formed in different forms such as $\{g,j\}$ edge can be replaced to $\{f,g\}$.