

# CENG 280

## Formal Languages and Abstract Machines

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### Homework 2

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### Answer for Q1

a.  $(a(b+c)^*a + b + aa)(a+b)^*$

b.

A= 0:

B= 1:

C= 0,1:

D=2:

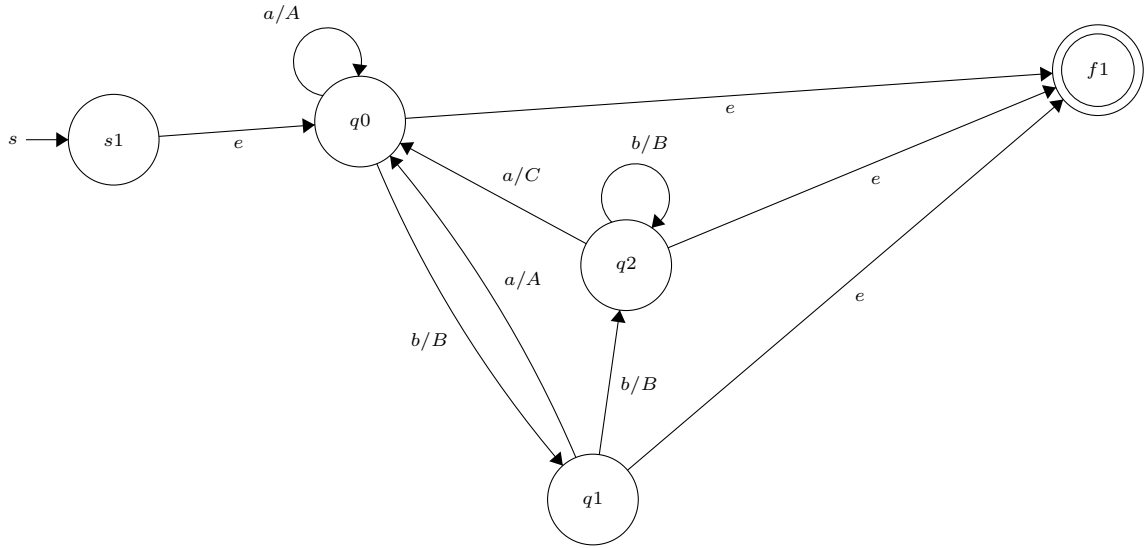
E=1:

F=0,2:

### Answer for Q2

a. We should use state elimination algorithm to find the output language.

b. We pick  $q_0$  as starting state and Final states as  $[q_0, q_1, q_2]$ , then we create a new start state ( $s_1$ ) and connect it to  $q_0$  with an empty transition. After that, we create a final state ( $f_1$ ) and connect all final states to  $f_1$  with empty transitions and make final states non-final. Now, after these updates, we are ready to use our algorithm. Also, we should change the output strings as given input-output pattern. Our new alphabet is  $\Sigma = [A, B, C]$  It will look like this:



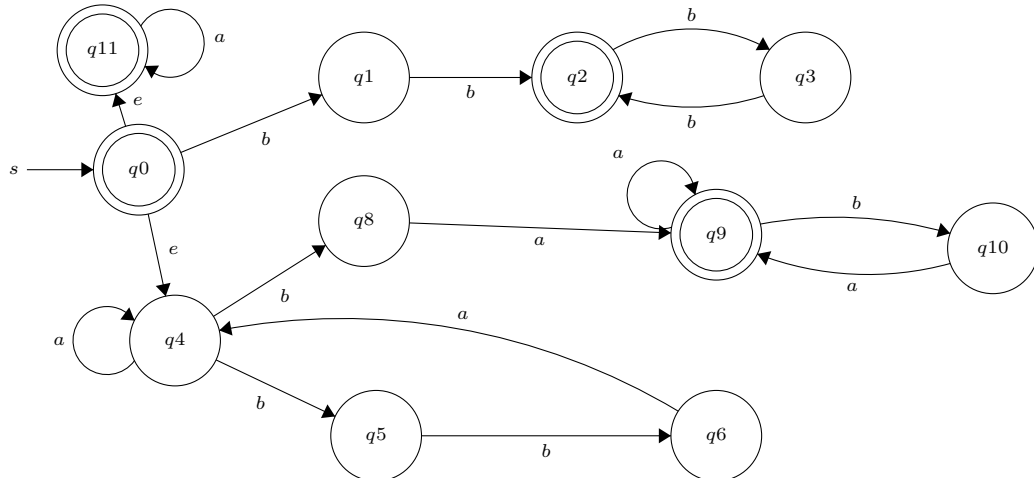
c. After making the changes in part 2, we can construct the desired part of the language which ends with C. We will follow the path  $q_0, q_1, q_2, q_0$  path to reach the strings that end with C. We first eliminate the  $q_0$ .  $A^*(BA)^*$  is the way of reaching  $q_1$  considering the loops. To reach  $q_2$  from  $q_1$  we have only one option  $b/B$ . To reach  $q_0$  from  $q_2$  we can form  $B^*C$  strings and it is the desired final state. Putting it all together, our output language is  $(A^*(BA)^*BBB^*C)^+$

## Answer for Q3

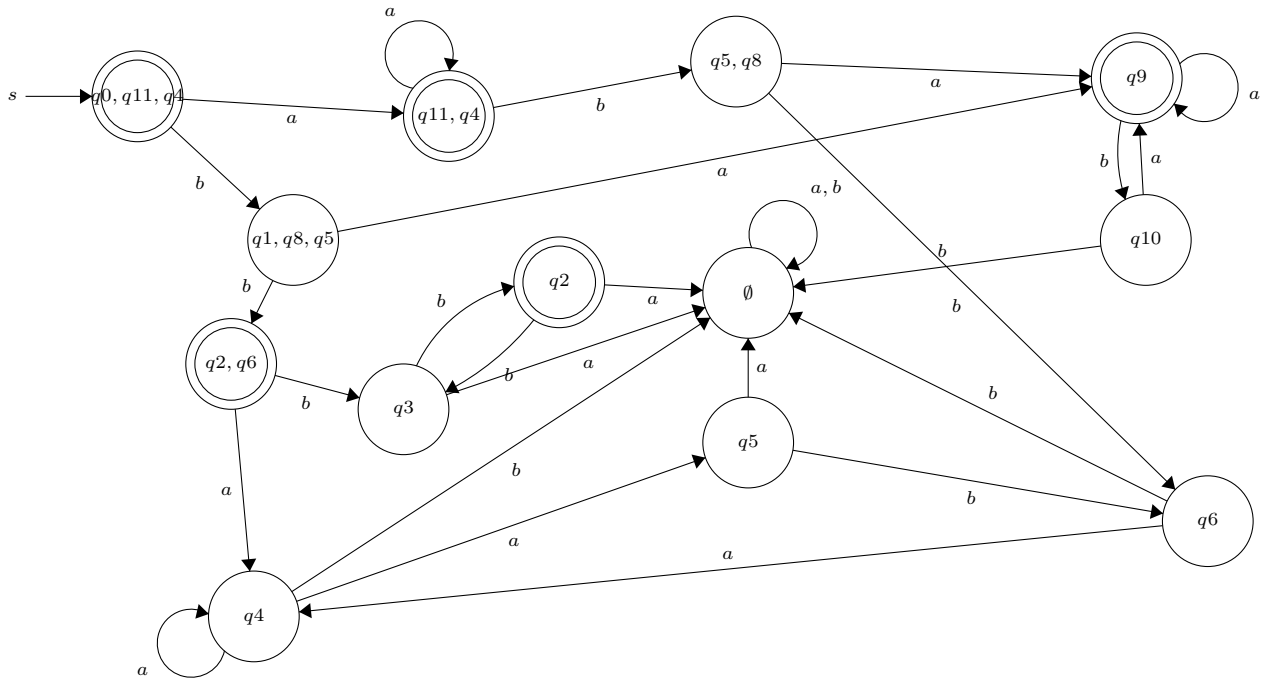
We first find a regular expression both we can get by the given regular expression and can be accepted by one of the NFA's.

Checking  $N_2$ , we can eliminate the upper loop (namely  $q_0, q_4, q_3$  and  $q_0, q_3$ ) since we cannot reach the final state from this loop since it ends with a trap state. Also, we can eliminate the  $q_0, q_1, q_2$  path since our output regular expression does not contain strings with " $BBA^*$ ". Only string sequence can reach the final state of  $N_2$  is  $A^*$  and  $BB$ .

Now let's look at  $N_3$ . In this NFA, we can eliminate the state  $q_4$  since there is no way to go through it, no strings end with " $CC$ ". Also we can eliminate  $q_1$  and  $q_2$  since we cannot reach the final state using this path. And the only way to reach the final state is possible with  $q_0, q_3$  since our RE can generate it. For the  $q_0, q_3$  path, we can pass by using  $(A + BBA)^*BA(A + BA)^* + (BB)^*$  path. Getting all together and in combination with  $N_1$  our final RE is  $(A + BBA)^*BA(A + BA)^* + (BB)^* + A^*$ . NFA for the given system is:



We will create a DFA according to this NFA. Using the algorithm we learned in class;



is the DFA